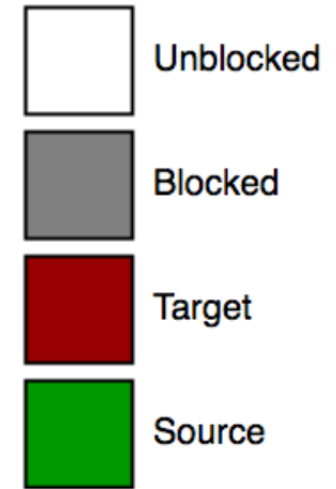
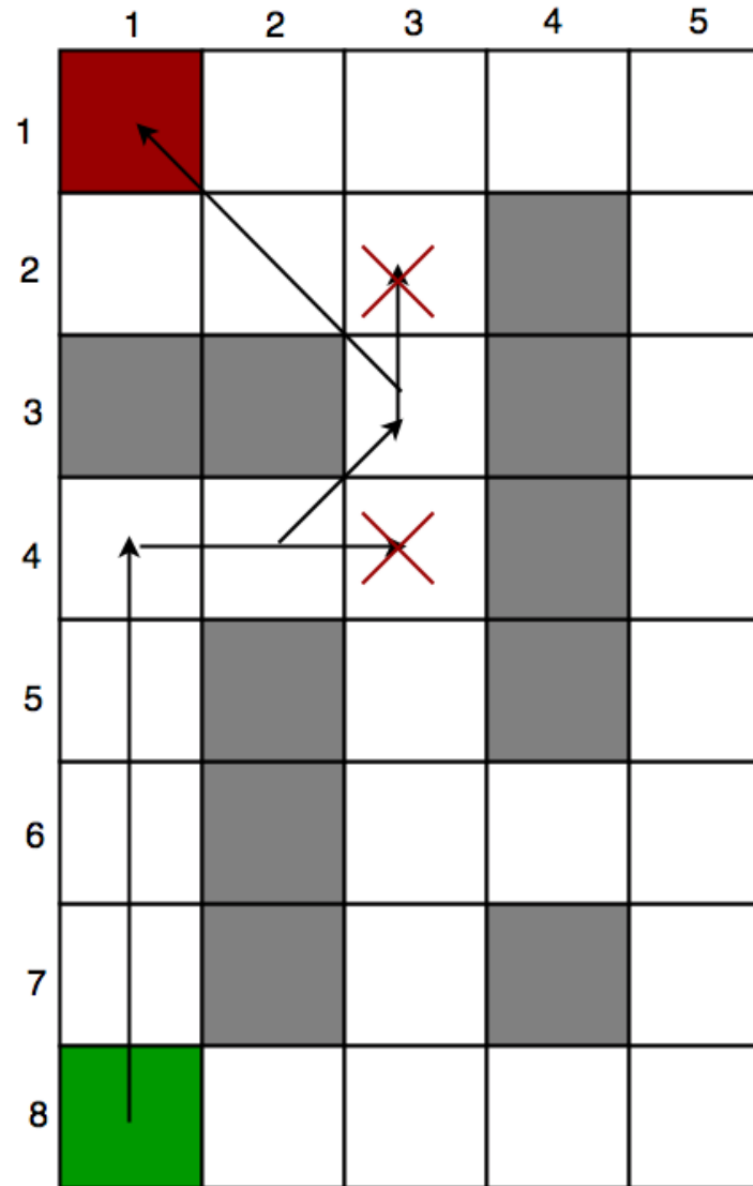


# CSCI 3202: Intro to Artificial Intelligence

## Lecture 7: A\* Search and Heuristics

Rachel Cox  
Department of  
Computer Science



A\* Search Algorithm makes the most intelligent choice at each step. Hence you can see that algorithm goes from (4,2) to (3,3) and not (4,3) (shown by cross).

Similarly the algorithm goes from (3,3) to (2,2) and not (2,3) (shown by cross).

[Source](#)

## **Review: Uniform-cost Search (UCS)**

---

- Expand out in contours, where least cost dictates which nodes we explore.
- Eventually, we will find a path to the goal - but the search is not directed

# Search Algorithms

---

❖ Search algorithms are fundamentally the same except for their frontier strategies.

**Uninformed Search**: e.g. Uniform Cost Search

- the good: UCS is complete and optimal → if a solution exists, it will find it with the least cost path
- the bad: explores in every direction

**Informed Search**: include information about where the goal is

- what do we need to have? A heuristic.

**heuristic**: A function that estimates how close a state is to a goal.

# Greedy best-first search

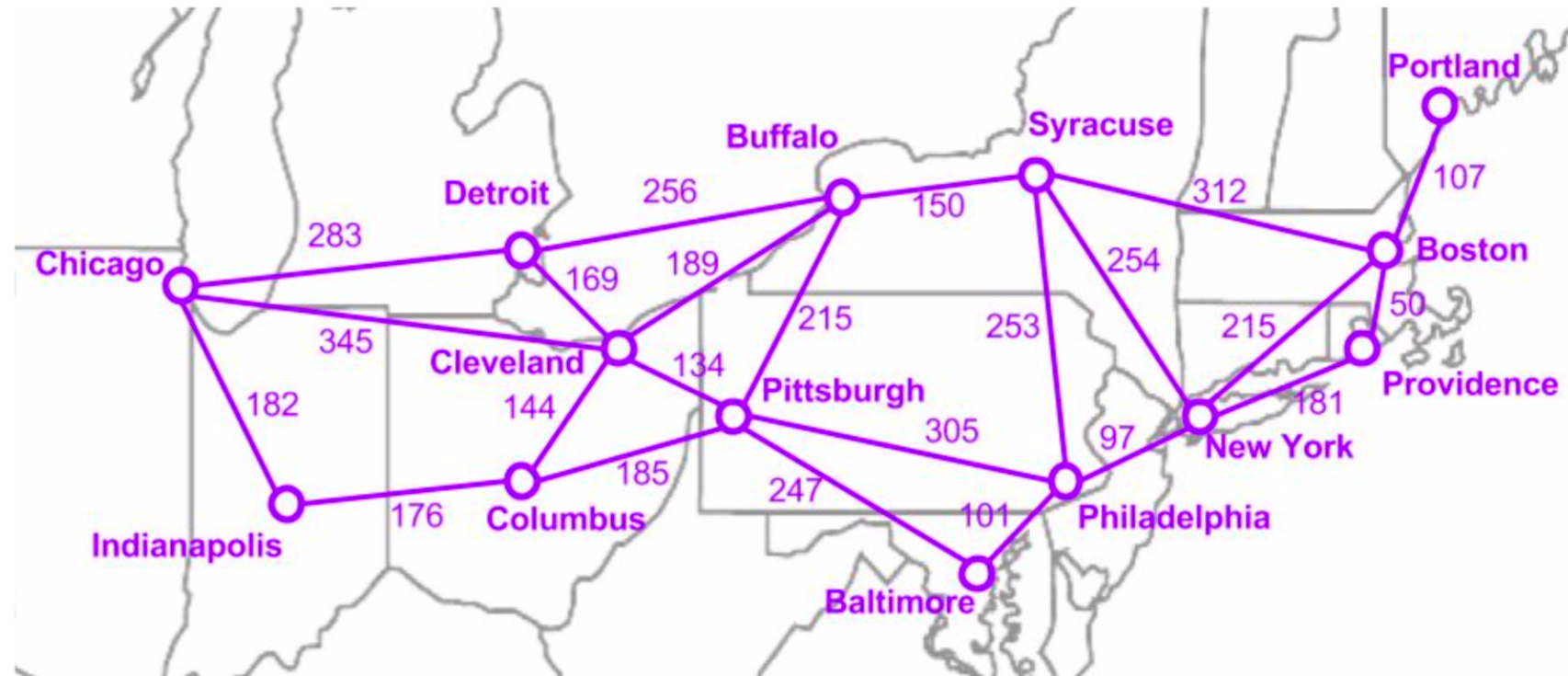
---

❖ First expand the path that's closest to the goal.

To determine what's closest to the goal, we need to define a heuristic function.

**Example:** For the traveling in the northeast problem, let's estimate the distance to the goal as the straight-line distance between city and the goal city.

Step costs: miles between cities along major highways

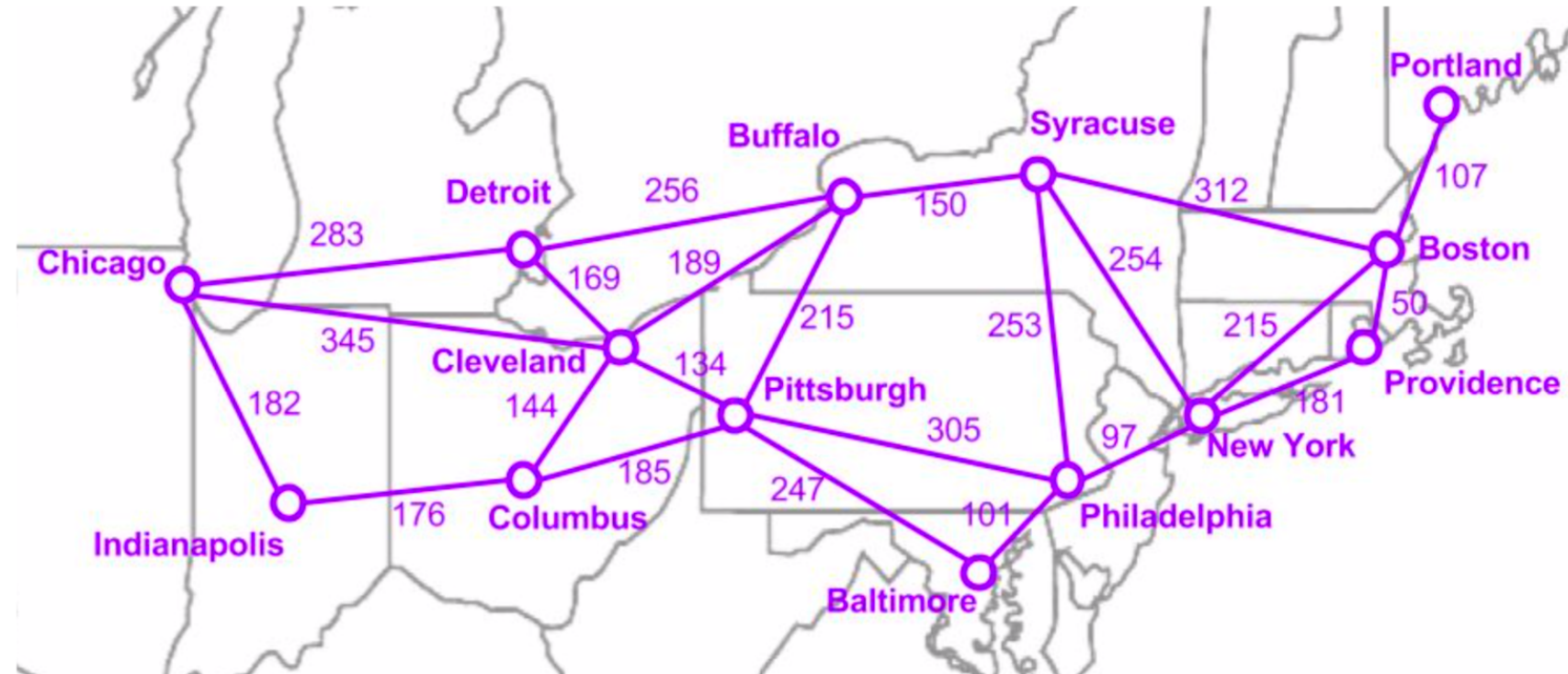


# Greedy best-first search

---

**Example:** Use the greedy best-first search to find a route from Chicago to Providence.

Heuristic:  $h(n)$  = straight-line distance to Providence



# Greedy best-first search

---

Possible Issue: Won't necessarily find the optimal path



# A\* Search

---

**Uniform-cost search:**

$$f(n) = g(n) \quad (\text{cost to get to } n)$$

**Greedy:**

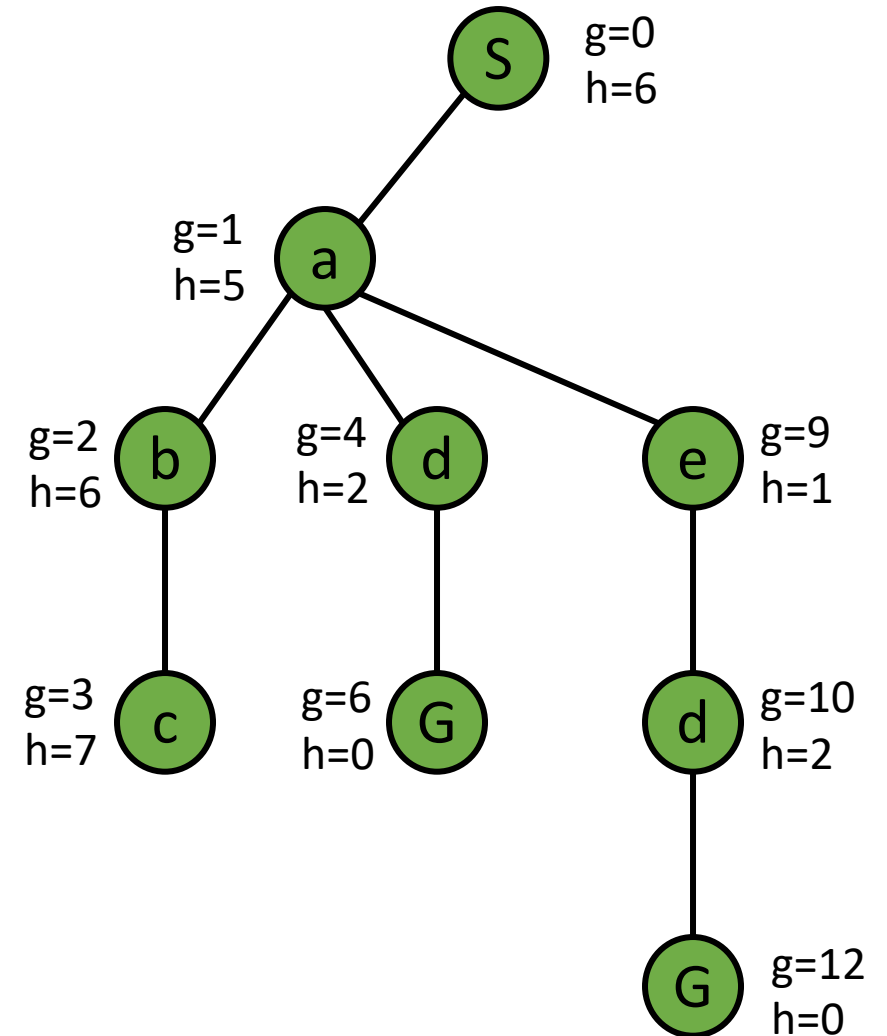
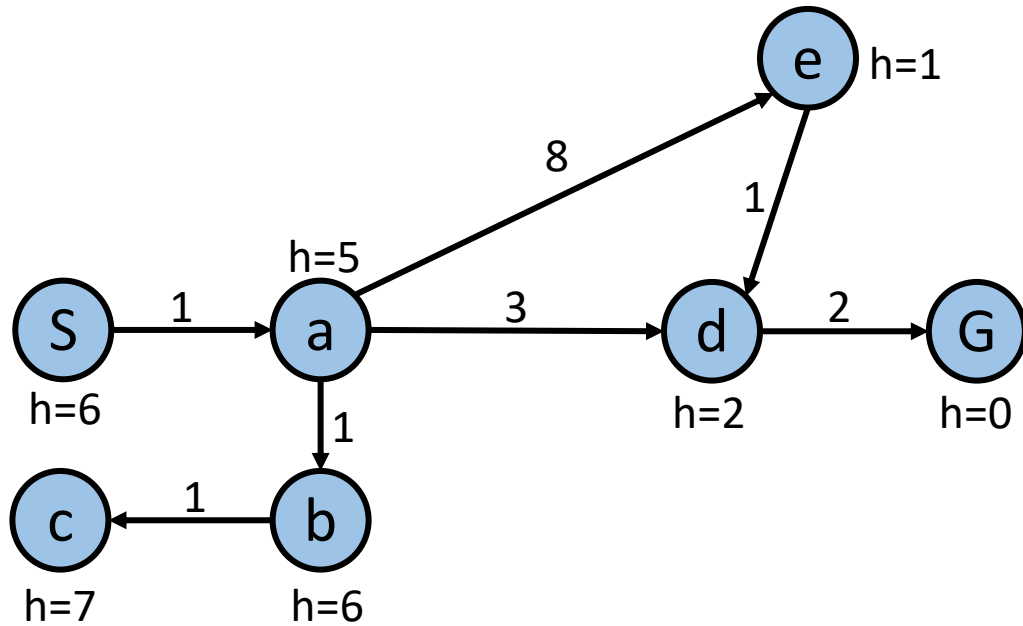
$$f(n) = h(n) \quad (\text{estimated cost to get from } n \text{ to goal})$$

**A\*:**

$$f(n) = g(n) + h(n) \quad (\text{estimated total cost of cheapest solution through } n)$$

# A\* Search

**Example:** Compare Uniform Cost, Greedy Search, and A\* on the graph below.

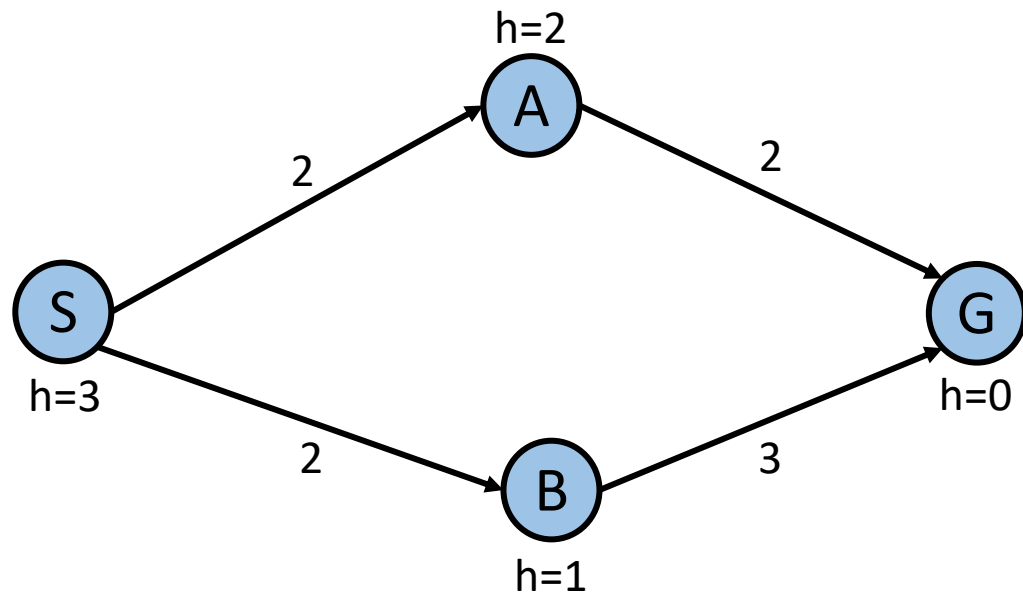




# A\* Search

---

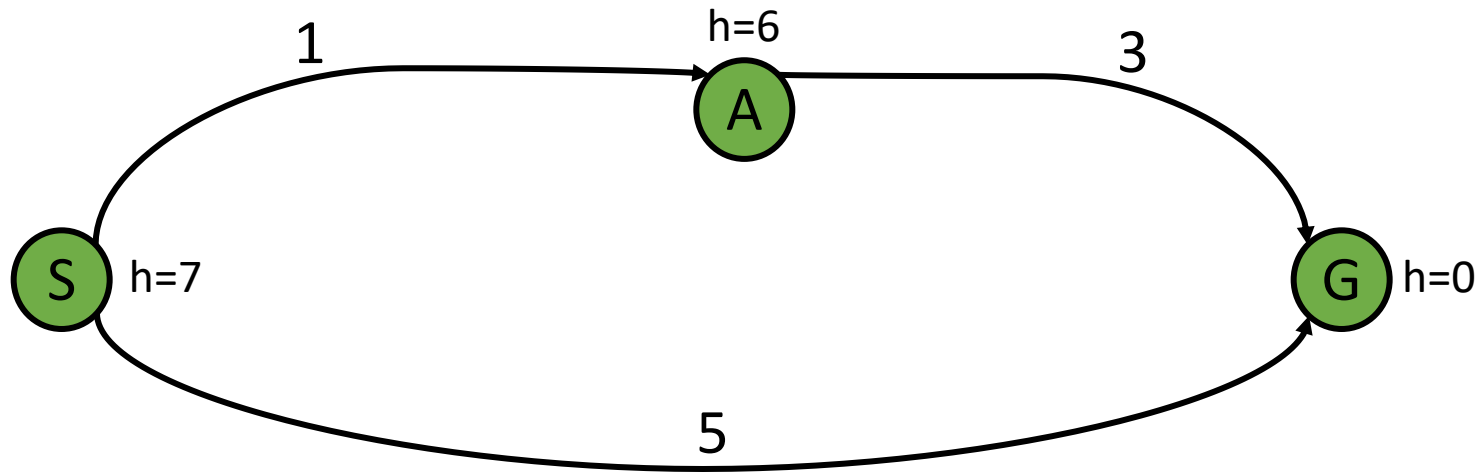
**Example:** When should A\* search terminate?



# A\* Search

---

Is A\* optimal?



# A\* Search

---

**Consistent:** for every node  $n$  and successor  $n'$  of  $n$ , generated by some action  $a$ , the estimated cost of reaching the goal from  $n$  is no greater than the step cost from  $n$  to  $n'$ , plus the estimated cost of reaching the goal from  $n'$

- That is:  $h(n) \leq c(n, a, n') + h(n')$
- General **triangle inequality** between  $n$ ,  $n'$ , and the goal

A heuristic  $h$  is **admissible** (optimistic) if  $0 \leq h(n) \leq h^*(n)$ , where  $h^*(n)$  is the true cost to the nearest goal.

# A\* Search

---

Search only works when:

- domain is fully observable
- domain must be known
- domain must be deterministic
- domain must be static

implementation: use a **node**

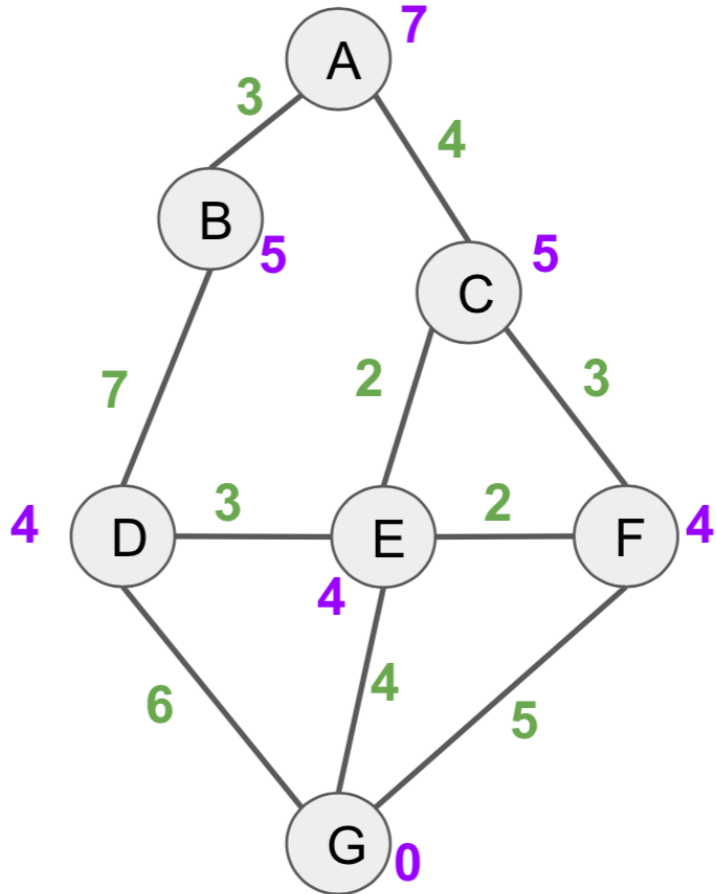
- state - indicates state at end of path
- action - action taken to get here
- cost - total cost
- parent - pointer to another node

# A\* Search

---

## A\* Search:

- Find the cheapest path from A to G
- $h(n)$  values are given in **purple**
- Step costs are given in **green**



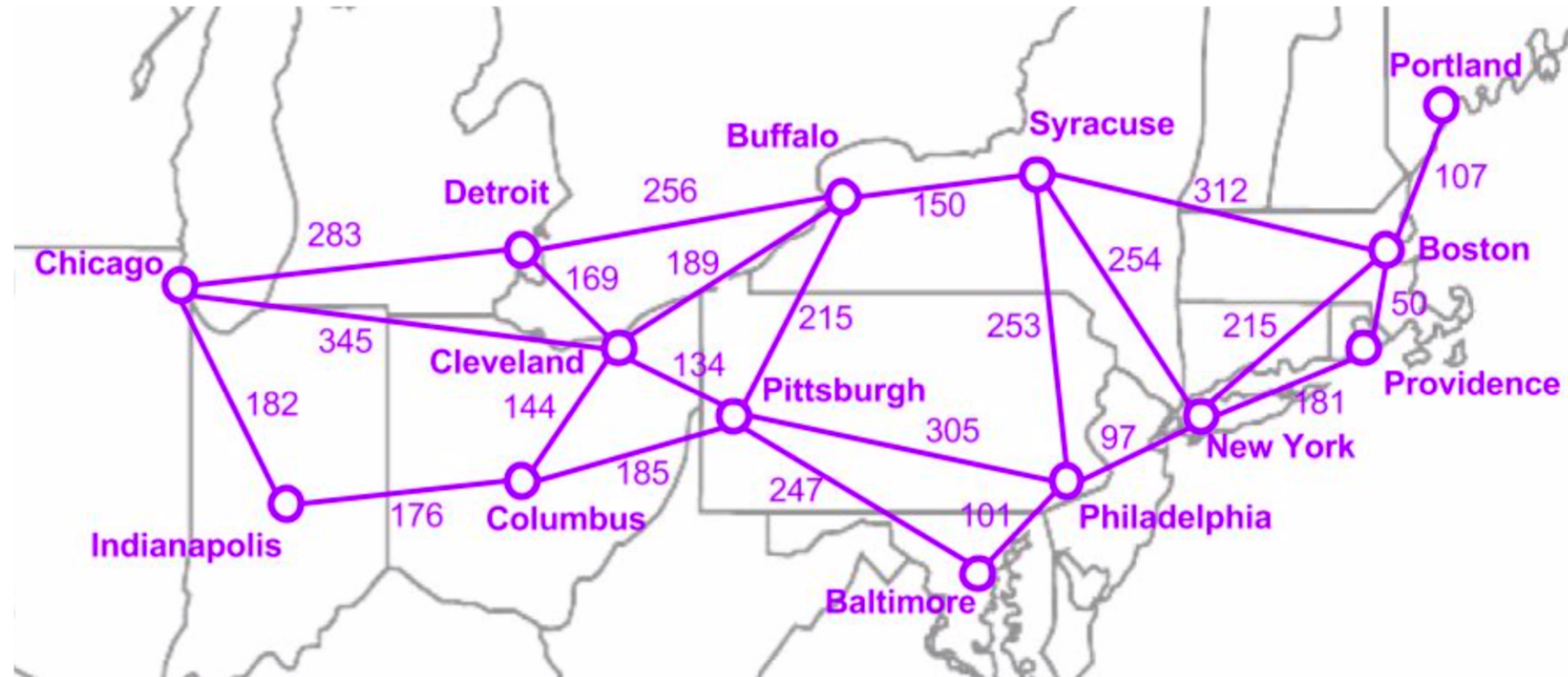
# A\* Search

---

**Example:** Use A\* search to find a route from Chicago to Providence.

$h(n)$  = straight-line distance to Providence

$g(n)$  = Path cost so far



# A\* Search

---

**Any consistent heuristic is also admissible (but not the other way around).**

**Example:** Prove the above statement by induction.

# Next Time

---

Optimality and Variants of A\*