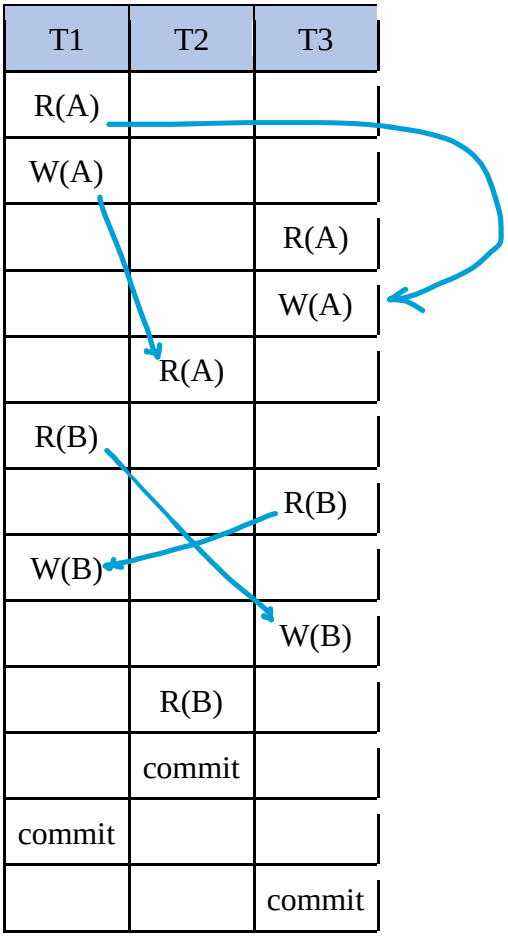
Part 2:



T1	T2	T3
R(A)		
W(A)		
		R(A)
		W(A)
	R(A)	
R(B)		
		R(B)
W(B)		
		W(B)
	R(B)	
	commit	
commit		
		commit

The transaction order here is as follows:

$$T3 \leftarrow T1 \rightarrow T2$$

And since we have a loop between T1 and T3 this means the schedule is NOT conflict-serializable.