

Markov chains

We will consider discrete time discrete space M.C.s.

Stochastic process $\{X_n, n \in T\}$ where T is a countable set, usually $\{\dots, -2, -1, 0, 1, 2, \dots\}$ or $\{0, 1, 2, \dots\}$

and $X_n \in \Omega$ where Ω is a countable (~~finite~~) set.

A discrete time discrete space ~~stochastic~~ stochastic process $\{X_n, n \in T\}$ is a Markov chain if $P(X_{n+1}=j | X_n=i, X_{n-1}=i_{n-1}, X_{n-2}=i_{n-2}, \dots) = P_{ij}$ for all $i, i_{n-1}, i_{n-2}, \dots, j \in \Omega$ and all $n \geq 0$, where $\{P_{ij}\}$ is a fixed set of transition probabilities.

Note that since P_{ij} is assumed to not vary (say w/n) this is a homogeneous M.C.

Obviously, $P_{ij} \geq 0 \quad \forall i, j \in \Omega$, $\sum_{j \in \Omega} P_{ij} = 1 \quad \forall i \in \Omega$

Let $P = \{P_{ij}\}$ denote the set of onestep transition probabilities. If $T = \{0, 1, 2, \dots\}$

$$P = \begin{matrix} & \begin{matrix} \text{next state} \\ P_{00} & P_{01} & P_{02} & \dots \end{matrix} \\ \begin{matrix} \text{current state} \\ P_{10} & P_{11} & P_{12} & \dots \\ \vdots & \vdots & \vdots & \vdots \end{matrix} & \end{matrix}$$

possibly infinite dimensional transition probability matrix (t.p.m.)

Markov chain idea: Future (X_{n+1}, \dots)
and past $(X_{n-1}, X_{n-2}, \dots)$ are conditionally
independent given the present (X_n)

Can generate a realization of chain if
we have an initial distr. (Q) on state space.
and t.p.m. P .

$$X_0 \sim Q \quad X_{n+1} | X_n = x \text{ based on t.p.m. } P \text{ for } n=0,1,\dots$$

E.g. 1. Random walk. $\Omega = \{\dots, -2, -1, 0, 1, 2, \dots\} = \mathbb{Z}$

$$P_{i,i+1} = p$$

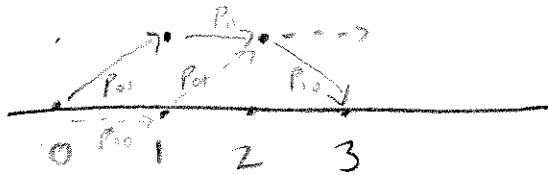
$$P_{i,i-1} = 1-p$$

Walking on straight line 1st step to right w/ prob. p
1 step " left " " $1-p$.

— end 1/25/08
lec #5

E.g. 2. 2 state M.C. w/ states = $\{0, 1\}$

Markovian coin toss: use different coins if you are in state 0 or state 1.



$$P = \begin{bmatrix} P_{00} & P_{01} \\ P_{10} & P_{11} \end{bmatrix}$$

E.g. 3. Let Z_1, Z_2, \dots be iid r.v. on space of non-neg. integers,

so $P(Z_i = j) = p_j$ for $j \in \{0, 1, \dots\}$

w/ $p_j \geq 0$, $\sum_{j=0}^{\infty} p_j = 1$

Let $S_n = \sum_{i=1}^n Z_i$, $n = 1, 2, \dots$

S_n is an M.C. w/ t.p.m.

$$\begin{aligned} P(S_{n+1} = j | S_n = i) &= P\left(\sum_{i=1}^n Z_i + Z_{n+1} = j \mid \sum_{i=1}^n Z_i = i\right) \\ &= P(Z_{n+1} = j - i) = \begin{cases} p_{j-i} & j \geq i \\ 0 & \text{else} \end{cases} \end{aligned}$$

$$P = \begin{bmatrix} p_0 & p_1 & p_2 & p_3 & \dots \\ 0 & p_0 & p_1 & p_2 & \dots \\ 0 & 0 & & & \dots \\ \vdots & \vdots & & & \ddots \end{bmatrix}$$

Gambler's ruin

Suppose gambler plays game where he wins \$1 w/prob p
loses \$1 " " $1-p$
Let X_t = fortune at time t .

Gambler quits when he wins \$ N (maximum) or
when he has \$0 left. (he is 'ruined').

This is a modified random walk.

$$\text{M.C. w/ t.p.m. } \begin{aligned} P_{i,i+1} &= p \\ P_{i,i-1} &= 1-p \end{aligned} \quad \text{for } i=1, \dots, N-1$$

$$\text{and } P_{0,0} = 1 = P_{N,N} \quad 0, N \text{ are 'absorbing states'}$$

Of interest: \Pr gambler gets to goal (\$ N) before
(going broke) 'ruin' (\$0).

Let $f_i = \Pr$ (reach \$ N when already have \$ i) , $i=0, 1, \dots, N$.
 $f_0 = 0$, $f_N = 1$.

Condition on outcome of 1st bet.

$$f_i = \underbrace{f_{i+1}}_{\text{win \$1}} p + \underbrace{f_{i-1}}_{\text{lose \$1}} (1-p) \quad i=1, 2, \dots, N-1$$

Let $q=1-p$ for convenience

Now, since $p+q=1$.

$$p f_i + q f_i = p f_{i+1} + q f_{i-1}$$

$$\Rightarrow q (f_i - f_{i-1}) = p (f_{i+1} - f_i)$$

$$\Rightarrow f_{i+1} - f_i = \frac{q}{p} (f_i - f_{i-1}) \quad i=1, \dots, N-1.$$

$$\text{Hence, } f_2 - f_1 = \frac{q}{p} (f_1 - f_0) = \frac{q}{p} f_1$$

$$f_3 - f_2 = \frac{q}{p} (f_2 - f_1) = \left(\frac{q}{p}\right)^2 f_1$$

$$\vdots$$

$$f_N - f_{N-1} = \left(\frac{q}{p}\right)^{N-1} f_1$$

$$\text{Add first } i-1 \text{ eqns: } f_i - f_1 = \sum_{k=1}^{i-1} \left(\frac{q}{p}\right)^k f_1$$

$$\Rightarrow f_i = \sum_{k=0}^{i-1} \left(\frac{q}{p}\right)^k f_1$$

sum of geometric series. $\frac{1-r^{n+1}}{1-r}$

$$\Rightarrow f_i = \begin{cases} \left[\frac{1 - (q/p)^i}{1 - (q/p)} \right] f_1 & q/p \neq 1 \\ i f_1 & q/p = 1 \end{cases}$$

Need to find f_1 .

Since $f_N = 1$

$$1 = f_N = \begin{cases} \frac{1 - (q/p)^N}{1 - q/p} f_1 & q/p \neq 1 \\ N f_1 & q/p = 1 \end{cases}$$

$$\Rightarrow f_1 = \begin{cases} \frac{1 - q/p}{1 - (q/p)^N} & \text{if } q/p \neq 1 \\ \frac{1}{N} & \text{if } q/p = 1 \end{cases}$$

$$\text{Hence, } f_i = \begin{cases} \frac{(1-(q/p)^i)}{1-q/p} \cdot \frac{1-q/p}{1-(q/p)^N} & \text{if } q/p \neq 1 \\ i/N & \text{if } q/p = 1 \end{cases}$$

$$\text{Note: As } N \rightarrow \infty \quad f_i \rightarrow \begin{cases} 1-(q/p)^i & \text{if } p > 1/2 \\ 0 & \text{if } p = 1/2 \\ 0 & \text{if } p < 1/2 \end{cases}$$

Implication: $p > 1/2$ means true prob. gambler can increase fortune indefinitely

$p \leq 1/2$ means gambler will be ruined w/ prob 1 if playing against infinitely wealthy opponent (casinos!)

— end of lec. 6

Chapman-Kolmogorov Equations

t.p.m. P describes 1 step transitions.

$$P_{ij} = P(X_{t+1}=j | X_t=i) \quad i, j \in \Omega$$

More generally, n -step transitions:

$$P_{ij}^{(n)} = P(X_{t+n}=j | X_t=i) \quad n \in \{1, 2, 3, \dots\}$$

Note: $P^{(n)} = P^n$ so all information is in 1-step t.p.m.

How do we calculate n -step transition probabilities?

Condition on intermediate steps using

C-K eqns.

$$P_{ij}^{n+m} = \sum_{\text{all } k} P_{ik}^n P_{kj}^m, \quad \forall n, m \geq 0, \quad i, j \in \Omega$$

$$\begin{aligned} \text{Pf: } P_{ij}^{n+m} &= P(X_{n+m}=j | X_0=i) \\ &= \sum_k P(X_{n+m}=j, X_n=k | X_0=i) \\ &= \sum_k P(X_{n+m}=j | X_n=k, X_0=i) P(X_n=k | X_0=i) \\ &= \sum_k P(X_{n+m}=j | X_n=k) P(X_n=k | X_0=i) \\ &= \sum_k P_{kj}^m P_{ik}^n \end{aligned}$$

Classification of states of an M.C.

'Accessible': state j is accessible from state i if for some integer $n \geq 0$, $P_{ij}^n > 0$

Can reach j from i in finite # of transitions w/ +ve probability.

Communicate: Two states i and j , accessible to each other, are said to communicate. ($i \leftrightarrow j$)

If i and j do not communicate:

$$\text{either } P_{ij}^n = 0 \quad \forall n \geq 0$$

$$\text{or } P_{ji}^n = 0 \quad \forall n \geq 0$$

or both are true.

Communication is an equivalence relation:

$$(1) \quad i \leftrightarrow i \quad \therefore P_{ij}^0 = \begin{cases} 1 & i=j \\ 0 & i \neq j \end{cases} \quad \text{by defn.}$$

(reflexivity)

$$(2) \quad i \leftrightarrow j \Rightarrow j \leftrightarrow i \quad (\text{symmetry})$$

$$(3) \quad i \leftrightarrow j \text{ and } j \leftrightarrow k \Rightarrow i \leftrightarrow k \quad (\text{transitivity})$$

$$\text{Pf: } i \leftrightarrow j \text{ and } j \leftrightarrow k \Rightarrow \exists n, m \text{ st. } P_{ij}^n > 0 \text{ and } P_{jk}^m > 0$$

$$\Rightarrow P_{ik}^{n+m} = \sum_r P_{ir}^n P_{rk}^m \geq P_{ij}^n P_{jk}^m > 0$$

$$\Rightarrow i \rightarrow k \quad \text{Similarly, } k \rightarrow i.$$

Hence. Can find equivalence classes (of states that communicate w/ each other) that partition the state-space.

Two states that communicate are said to be in the same class.

~~So two classes of states are either identical or disjoint~~

\Rightarrow Communication ~~is~~ partitions state space into a number of ~~different~~ separate classes.

Irreducible: M.C. is irreducible if there is only one class, i.e., all states communicate with each other.

E.g. 1: $P = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{matrix} & \left[\begin{array}{cc|cc} 1/2 & 1/2 & 0 & 0 & 0 \\ 1/4 & 3/4 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 \end{array} \right] \end{matrix} = \begin{pmatrix} P_1 & 0 \\ 0 & P_2 \end{pmatrix}$

Two classes: $\{1, 2\}$ and $\{3, 4, 5\}$ (not irreducible)

E.g. 2: Random walk w/ absorbing barriers (like gambler's ruin problem)

$$\begin{matrix} 0 \\ 1 \\ 2 \\ \vdots \\ (a-1) \\ a \end{matrix} \begin{bmatrix} 1 & 0 & 0 & 0 & \dots & 0 & 0 & 0 \\ (1-p) & 0 & p & 0 & \dots & 0 & 0 & 0 \\ 0 & (1-p) & 0 & p & \dots & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & \dots & (1-p) & 0 & p \\ 0 & 0 & 0 & 0 & \dots & 0 & 0 & 1 \end{bmatrix}$$

— end of lec 7

3 classes: $\{0\}$, $\{1, 2, \dots, a-1\}$, $\{a\}$
Can go from B to A or C but not A to B or C to B (21)

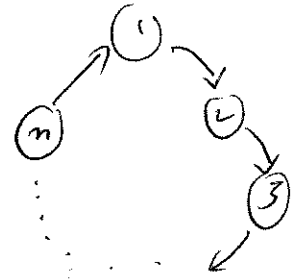
Periodicity

Period of state i ($d(i)$) is the greatest common divisor (g.c.d.) of all integers $n \geq 1$ for which

$$P_{ii}^n > 0 \quad \text{If } P_{ii}^n = 0 \quad \forall n \geq 1, \text{ define } d(i) = 0.$$

E.g.

	1	2	3	4	...	n
1	0	1	0	0	...	0
2	0	0	1	0	...	0
3
4	0	0	1
...
n	1	0	0



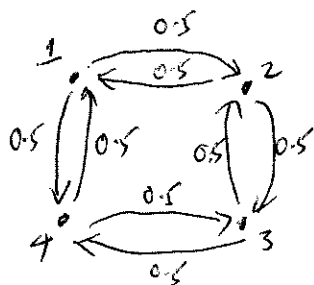
Each state has period n

\therefore starting at i it is possible to only re-enter i at $\{n, 2n, 3n, \dots\}$ g.c.d. is n .

Periodicity is a class property, i.e. if $i \leftrightarrow j$ then $d(i) = d(j)$.

Aperiodic: M.C. in which each state has period 1.
is called aperiodic. For irreducible M. chain
only need to show period = 1 for any one state.

E.g. Random walk on a square



t.p.m. $P = \begin{bmatrix} 0 & 0.5 & 0 & 0.5 \\ 0.5 & 0 & 0.5 & 0 \\ 0 & 0.5 & 0 & 0.5 \\ 0.5 & 0 & 0.5 & 0 \end{bmatrix}$

Start in 1: in 1-step can get to 2) or 4, not return to 1

✓ 2-steps: " " " 3 or return to 1

~~3-steps: " " " 2 or 4~~

from 3, need two steps to get back to 1

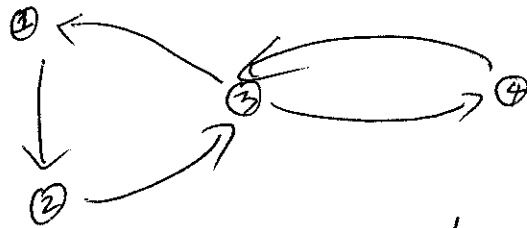
3-steps: cannot get back to 1

✓ 4-steps: can get back to 1

1 has period of 2.

Since 1, 2, 3, 4 are in same class, all have period of 2.

E.g. 2.



Is this M. chain aperiodic?

It is irreducible so only need to consider 1 state
say (3). $d(3) = \gcd\{2, 3, \dots\}$
 $= 1$.

Yes, ~~irre~~ aperiodic. even though at 1st glance, ^{it seems periodic} (1, 3 seems to have periods > 1)

For any states i & j , define
 prob. that, starting in i , the 1st transition
 into j occurs at time n .

Define $f_{ij}^0 = 0$

$$f_{ij}^n = P(X_n = j, X_k \neq j, k = 1, \dots, n-1 | X_0 = i)$$

Now define $f_{ij}^* = \sum_{n=1}^{\infty} f_{ij}^n$

f_{ij}^* = Prob of ever making a transition into j ,
 given that process starts in state i .

$$f_{ij}^* > 0 \Leftrightarrow i \rightarrow j \quad (i \neq j)$$

Defn: State i is recurrent if $f_{ii}^* = 1$
 State i is transient if $f_{ii}^* < 1$

Note the difference between f_{ij}^n and P_{ij}^n

f_{ij}^n = prob. of 1st transition to j at time n .

P_{ij}^n = prob. of a " " " " " " " " " " " "

Thm: All states of a finite, irreducible M.C. are recurrent.

Pf: Since M.C. is irreducible ~~all~~ states are either ^{all} recurrent or all transient.

Suppose all states are transient.

For pt. by contradiction we need the following Lemma.

~~Lemma~~ Define: $Q_{ii} = \text{Prob}(\text{M.C. returns } \infty \text{ often} \mid X_0 = i)$

Lemma: State is recurrent or transient according to whether $Q_{ii} = 1$ or 0 respectively.

Let $Q_{ii}^N = \text{Prob}(\text{M.C. returns to } i \text{ at least } N \text{ times} \mid X_0 = i)$

$$\begin{aligned} Q_{ii}^N &= \sum_{k=1}^{\infty} \text{Prob}(\text{M.C. returns 1st time in } k^{\text{th}} \text{ step}) \\ &\quad \times \text{Prob}(\text{M.C. returns to } i \text{ at least } N-1 \text{ times} \mid X_k = i) \\ &= \sum_{k=1}^{\infty} f_{ii}^k Q_{ii}^{N-1} = Q_{ii}^{N-1} f_{ii}^* \end{aligned}$$

Proceeding recursively, $Q_{ii}^N = (f_{ii}^*)^2 Q_{ii}^{N-2} = \dots = (f_{ii}^*)^{N-1} Q_{ii}^1$

$$= (f_{ii}^*)^N$$

$$Q_{ii} = \lim_{N \rightarrow \infty} Q_{ii}^N = \lim_{N \rightarrow \infty} (f_{ii}^*)^N = \begin{cases} 1 & \text{if recurrent} \\ 0 & \text{if transient} \end{cases}$$

Hence, ^(from Lemma) if state i is transient, M.C. does not return to state i after some finite time T_i . ^(w/ prob 1) True for all states i

~~Simultaneously~~

So, M.C. does not return to any state after time $T = \sup \max \{T_i\}$

w/ prob. 1. Contradiction!

Hence all states must be \S recurrent.

Thm: State i is recurrent iff

$$\sum_{n=1}^{\infty} P_{ii}^n = \infty$$

Sketch of

Pf: If state i is recurrent, prob. (M.C. eventually returns to i) = 1.

(\Rightarrow) Process probabilistically restarts itself when it enters state i so prob(it returns twice to state i) = 1.

So Prob (returns ∞ # of times to state i) = 1

$$E(\# \text{ returns}) = \infty$$

$$\Rightarrow E\left(\sum_{n=1}^{\infty} I(X_n = i | X_0 = i)\right) = \infty$$

(careful about this!) $\Rightarrow \sum_{n=1}^{\infty} E(I(X_n = i | X_0 = i)) = \infty$

$$\Rightarrow \sum_{n=1}^{\infty} P(X_n = i | X_0 = i) = \infty$$

$$\Rightarrow \sum_{n=1}^{\infty} P_{ii}^n = \infty$$

(need Borel-Cantelli etc. for rigorous proof)

(\Leftarrow) If state i is transient, prob. (M.C. never returns to i) = $1 - f_{ii}^* > 0$ ($\because f_{ii}^* < 1$)

$$\text{Prob (MC stays in state } i \text{ for a total } k \text{ times } | X_0 = i) = (f_{ii}^*)^{k-1} (1 - f_{ii}^*) > 0$$

Prob of getting to state i $k-1$ times

Prob of never returning to i

(\Leftarrow) Suppose state i is transient.

Each time process returns to i , there is positive probability, $1 - f_{ii}^* > 0$ that it will never return.

Hence, # visits \sim Geometric ($p = 1 - f_{ii}^*$)

wait until 1st success where "success" = leave and never return

$$E(\# \text{ visits}) = \frac{1}{1 - f_{ii}^*} < \infty$$

$$\sum_{n=1}^{\infty} P_{ii}^n < \infty$$

Hence i is recurrent $\Leftrightarrow \sum_{n=1}^{\infty} P_{ii}^n = \infty$

i is transient $\Leftrightarrow \sum_{n=1}^{\infty} P_{ii}^n < \infty$

Thm: Recurrence, ^{transience} is a class property

Pf: Suppose i is recurrent.

Assume $j \leftrightarrow i$ so j and i are in the same class

Remains to show that j is recurrent

Now, $P_{ij}^n > 0$ and $P_{ji}^m > 0$ for some n, m (by $i \leftrightarrow j$ assumption)

For any $s \geq 0$:

$$P_{jj}^{m+s+n} \geq P_{ji}^m P_{ii}^s P_{ij}^n \quad (\text{from Chapman-Kolmogorov})$$

$$\Rightarrow \sum_{s=0}^{\infty} P_{jj}^{m+s+n} \geq \left(\sum_{s=0}^{\infty} P_{ii}^s \right) P_{ji}^m P_{ij}^n$$

$$\text{But } \sum_{s=0}^{\infty} P_{ii}^s = \infty \quad (\text{by previous thm, since } i \text{ is recurrent})$$

$$\text{Hence } \sum_{s=0}^{\infty} P_{jj}^{m+s+n} = \infty$$

and j is also recurrent.

Follows that transience is also a class property