



**HACETTEPE UNIVERSITY**

**Department of Computer Engineering**

**BBM204**

**Programming Lab.**

**Programming Assignment 1**

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# Analysis of Algorithms

## 1.Introduction

In computer science, a sorting algorithm is an algorithm that puts elements of a list in a certain order. In computer science, the computational complexity or simply complexity of an algorithm is the amount of resources required to run it. <sup>1</sup>

The number of (machine) instructions which a program executes during its running time is called its time complexity in computer science. This number depends primarily on the size of the program's input, that is approximately on the number of the strings to be sorted (and their length) and the algorithm used. So approximately, the time complexity of the program “sort an array of  $n$  strings by minimum search” is described by the expression  $c \cdot n^2$ .

“ $c$ ” is a constant which depends on the programming language used, on the quality of the compiler or interpreter, on the CPU, on the size of the main memory and the access time to it. <sup>2</sup>

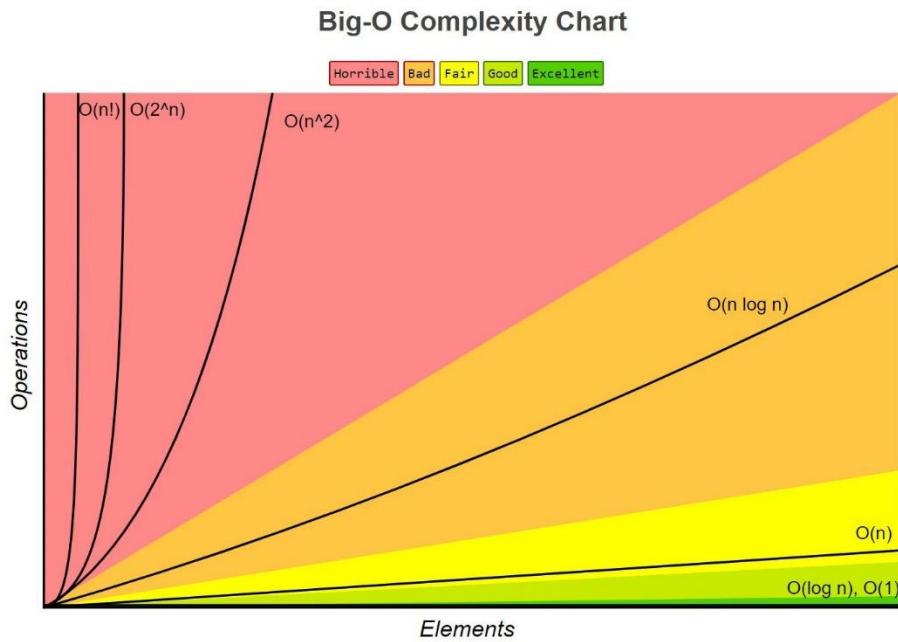
Generally, we perform the following types of analysis;

- Worst-case– The maximum number of steps taken on any instance of size  $a$ .
- Best-case– The minimum number of steps taken on any instance of size  $a$ .
- Average case– An average number of steps taken on any instance of size  $a$ .

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<sup>1</sup> [https://en.wikipedia.org/wiki/Sorting\\_algorithm](https://en.wikipedia.org/wiki/Sorting_algorithm)

<sup>2</sup> <https://www.leda-tutorial.org/en/official/ch02s02s03.html>



*Figure 1 Time Complexity Graph*

## 2.Problem

In this experiment, 5 sequencing algorithms will be analyzed:

- Comb Sort Algorithm
- Gnome Sort Algorithm
- Shaker Sort Algorithm
- Stooge Sort Algorithm
- Bitonic Sort Algorithm

- 1) A brief information will be given about algorithms.
- 2) The time we obtained will be given according to the number of inputs.
- 3) It will be proved whether it is stable or not.
- 4) The data will be visualized with graphics.
- 5) The final step will be the results and benefits from this experiment.

## 3. Sorting Algorithm

### 3.1 Comb Sort Algorithm

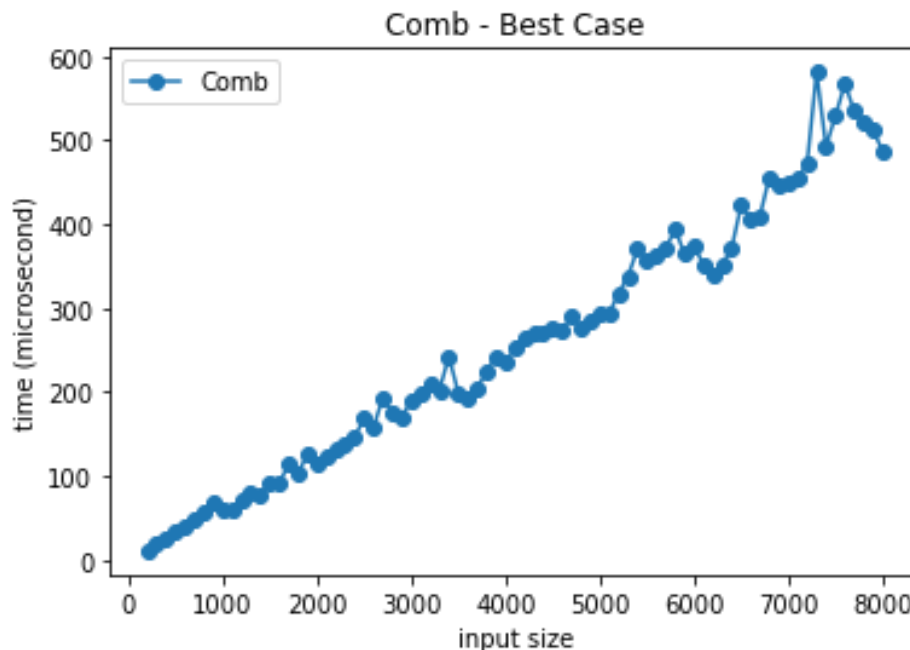
Comb sort is a relatively simple sorting algorithm originally designed by Włodzimierz Dobosiewicz and Artur Borowy in 1980, later rediscovered by Stephen Lacey and Richard Box in 1991. Comb Sort improves on Bubble Sort by using gap of size more than 1. The gap starts with a large value and shrinks by a factor of 1.3 in every iteration until it reaches the value 1.<sup>3</sup>

- Worst-case performance  $O(n^2)$
- Best-case performance  $\Theta(n \cdot \log(n))$
- Average performance  $\Omega(n^2 / 2^p)$ , where  $p$  is the number of increments
  - Space Complexity  $O(1)$

Data: (Unit microsecond)

(Sorted Array) => Best Case

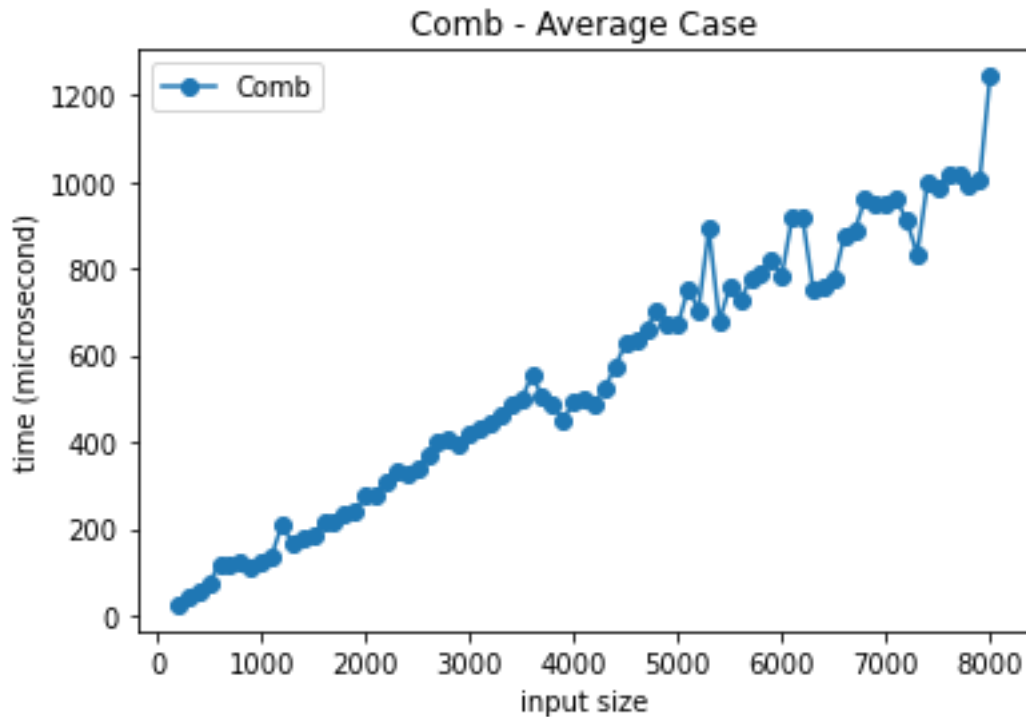
Input	200	900	1600	2300	3000	3700	4400	5100	5800	6500	7200	7900
Comb	10.971429	69.757143	90.814286	138.157143	189.507143	204.914286	271.071429	293.835714	394.485714	422.535714	472.85	512.5



<sup>3</sup> [https://en.wikipedia.org/wiki/Comb\\_sort](https://en.wikipedia.org/wiki/Comb_sort)

(Unsorted Array) => Average Case

Input	200	900	1600	2300	3000	3700	4400	5100	5800	6500	7200	7900
Comb	25.185714	111.971429	214.521429	332.807143	416.25	504.657143	573.464286	752.107143	790.742857	773.985714	910.121429	1001.564286



Stable: No

```
Unsorted Array
[(9,a), (9,b), (8,a), (8,b), (12,a), (6,a), (9,c), (20,a), (18,a), (9,d), (8,c), (18,b), (1,a), (1,b), (18,c), (3,b)]
Comb Sort
[(1,a), (1,b), (3,b), (6,a), (8,c), (8,a), (8,b), (9,b), (9,a), (9,d), (9,c), (12,a), (18,a), (18,b), (18,c), (20,a)]
```

Note:

When we look at Figure 1, it can be said that my average graph looks like the " $n * \log(n)$ " graph. Again, this can be said for our best graph. It surprised me that the Bubble sort algorithm was used for this algorithm and the gap value was determined to be 1.3. I observed that the experiment result was not stable. I can say that it is efficient compared to others. In addition, space complexity is efficient as memory usage is low.

## 3.2 Gnome Sort Algorithm

Gnome sort (dubbed stupid sort) is a sorting algorithm originally proposed by Iranian computer scientist Hamid Sarbazi-Azad (professor of Computer Science at Sharif University of Technology) in 2000.<sup>4</sup>

The gnome sort is a sorting algorithm which is similar to insertion sort in that it works with one item at a time but gets the item to the proper place by a series of swaps, similar to a bubble sort. It is conceptually simple, requiring no nested loops.

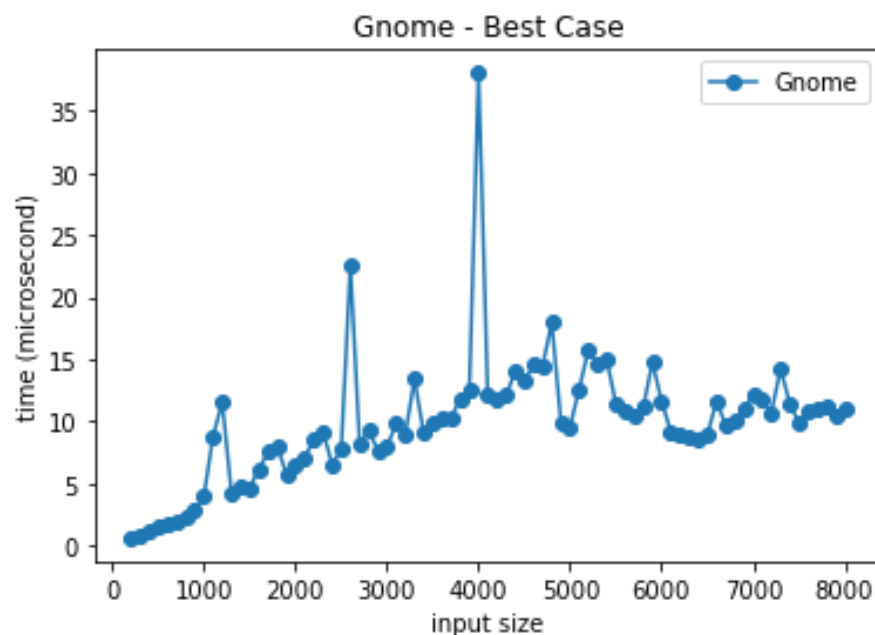
<sup>5</sup>

- Worst-case performance  $O(n^2)$
- Best-case performance  $\Theta(n)$
- Average performance  $\Omega(n^2)$ 
  - Space Complexity  $O(1)$

Data: (Unit microsecond)

(Sorted Array) => Best Case

Input	200	900	1600	2300	3000	3700	4400	5100	5800	6500	7200	7900
Gnome	0.621429	2.9	6.178571	9.171429	7.964286	10.25	14.0	12.614286	11.135714	8.907143	10.728571	10.478571

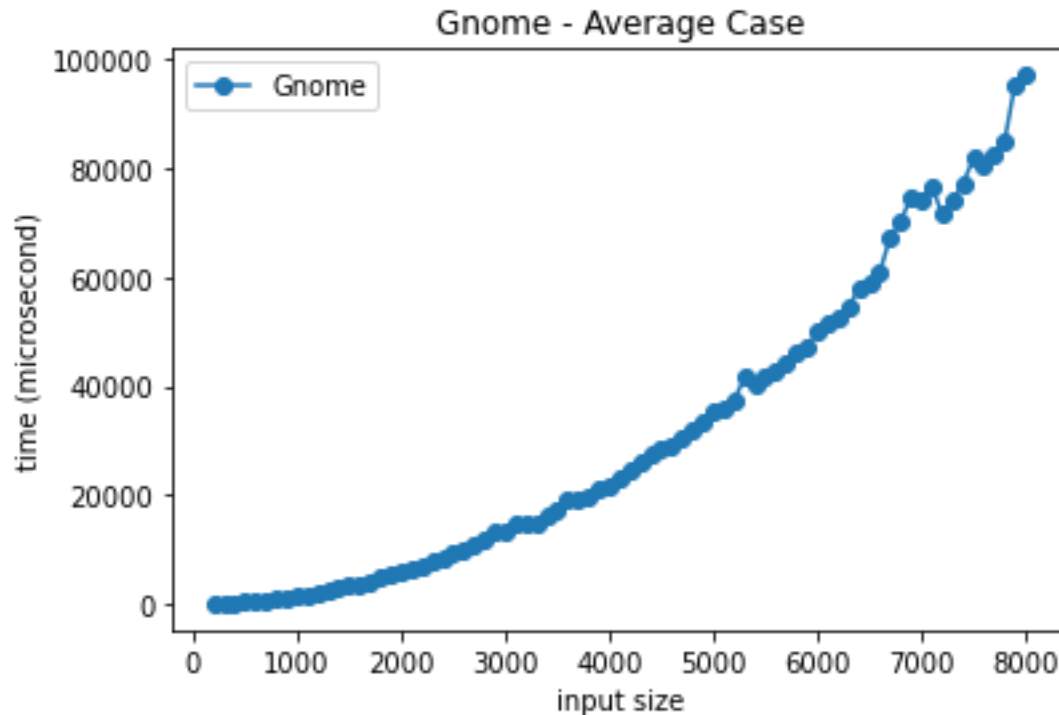


<sup>4</sup> [https://en.wikipedia.org/wiki/Gnome\\_sort](https://en.wikipedia.org/wiki/Gnome_sort)

<sup>5</sup> <https://www.geeksforgeeks.org/gnome-sort-a-stupid-one/>

(Unsorted Array) => Average Case

Input	200	900	1600	2300	3000	3700	4400	5100	5800	6500	7200
Gnome	119.357143	1188.442857	3725.521429	7742.3	13551.071429	19000.564286	27587.992857	35849.421429	45918.564286	59067.45	71414.685714



Stable: Yes

```
Unsorted Array
[(9,a), (9,b), (8,a), (8,b), (12,a), (6,a), (9,c), (20,a), (18,a), (9,d), (8,c), (18,b), (1,a), (1,b), (18,c), (3,b)]
Gnome Sort
[(1,a), (1,b), (3,b), (6,a), (8,a), (8,b), (8,c), (9,a), (9,b), (9,c), (9,d), (12,a), (18,a), (18,b), (18,c), (20,a)]
```

Note:

When we look at Figure 1, it can be said that the average graph looks like " $n * \log(n)$ " or " $n^2$ " graph. It can be said that there is " $n$ " graph for our Best graph. I would say it has a simple algorithm. I observed that the experiment result was stable. I can say that it is moderately efficient compared to others. I use it for a simple series, but not for a mixed series. It can be said that it is efficient with  $O(1)$  in terms of memory.

### 3.3 Shaker Sort Algorithm

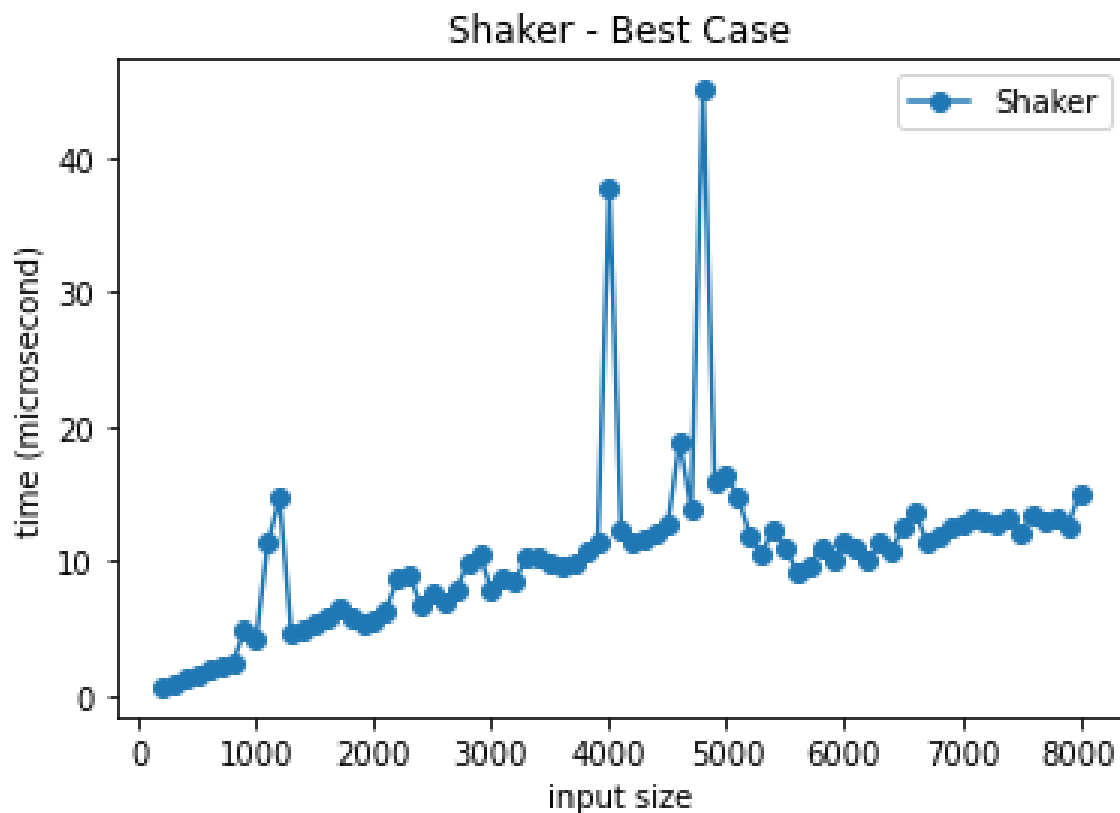
Cocktail shaker sort, also known as bidirectional bubble sort, cocktail sort, shaker sort (which can also refer to a variant of selection sort), ripple sort, shuffle sort, or shuttle sort, is an extension of bubble sort. The algorithm extends bubble sort by operating in two directions.<sup>6</sup>

- Worst-case performance  $O(n^2)$
- Best-case performance  $\Theta(n)$
- Average performance  $\Omega(n^2)$ 
  - Space Complexity  $O(1)$

Data: (Unit microsecond)

(Sorted Array) => Best Case

Input	200	900	1600	2300	3000	3700	4400	5100	5800	6500	7200	7900
Shaker	0.707143	4.9	5.878571	9.014286	7.935714	9.935714	12.092857	14.8	11.1	12.507143	12.978571	12.635714

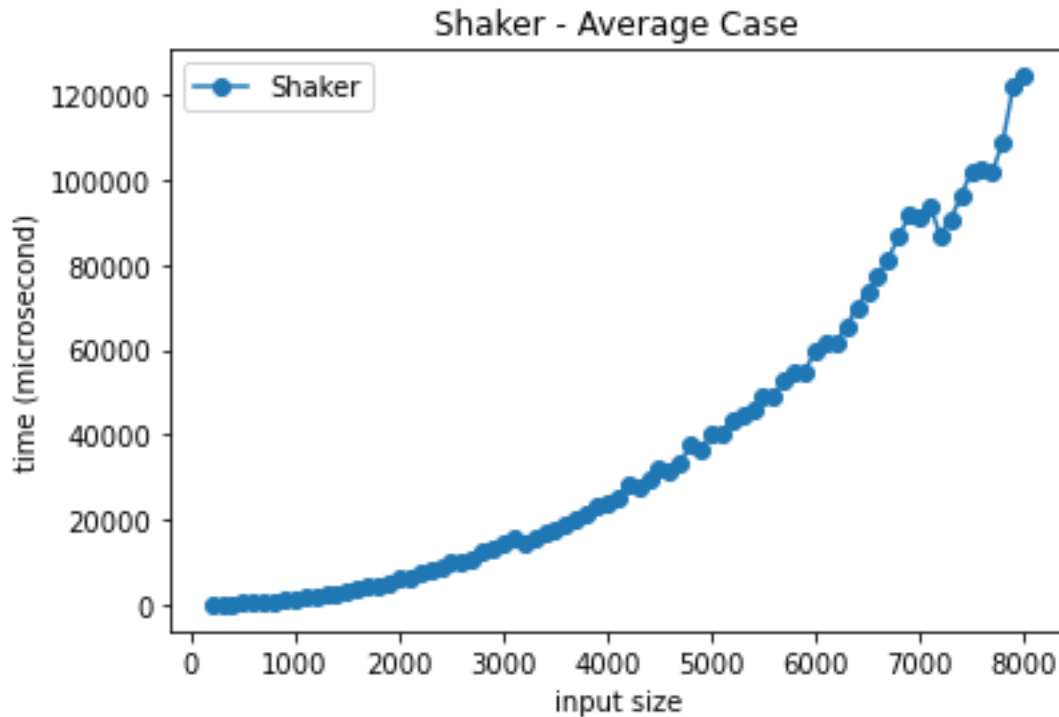


<sup>6</sup> [https://en.wikipedia.org/wiki/Cocktail\\_shaker\\_sort](https://en.wikipedia.org/wiki/Cocktail_shaker_sort)



(Unsorted Array) => Average Case

Input	200	900	1600	2300	3000	3700	4400	5100	5800	6500	7200
Shaker	132.771429	1386.985714	3802.0	7932.742857	14235.671429	20407.721429	29315.285714	40228.45	54835.778571	73661.078571	87028.535714



Stable: Yes

```
Unsorted Array
[(9,a), (9,b), (8,a), (8,b), (12,a), (6,a), (9,c), (20,a), (18,a), (9,d), (8,c), (18,b), (1,a), (1,b), (18,c), (3,b)]
Shaker Sort
[(1,a), (1,b), (3,b), (6,a), (8,a), (8,b), (8,c), (9,a), (9,b), (9,c), (9,d), (12,a), (18,a), (18,b), (18,c), (20,a)]
```

Note:

When we look at Figure 1, it can be said that the average graph looks like " $n * \log(n)$ " or " $n^2$ " graph. It can be said that there is a " $n$ " graph for our Best graph. This algorithm is an extension of bubble sort. Effective for small values at the end of the list. I observed that the experiment result was stable. I can say that it is medium-high efficient compared to others. I might prefer bubble sort instead. I would say that the algorithm is memory-efficient.

### 3.4 Stooge Sort Algorithm

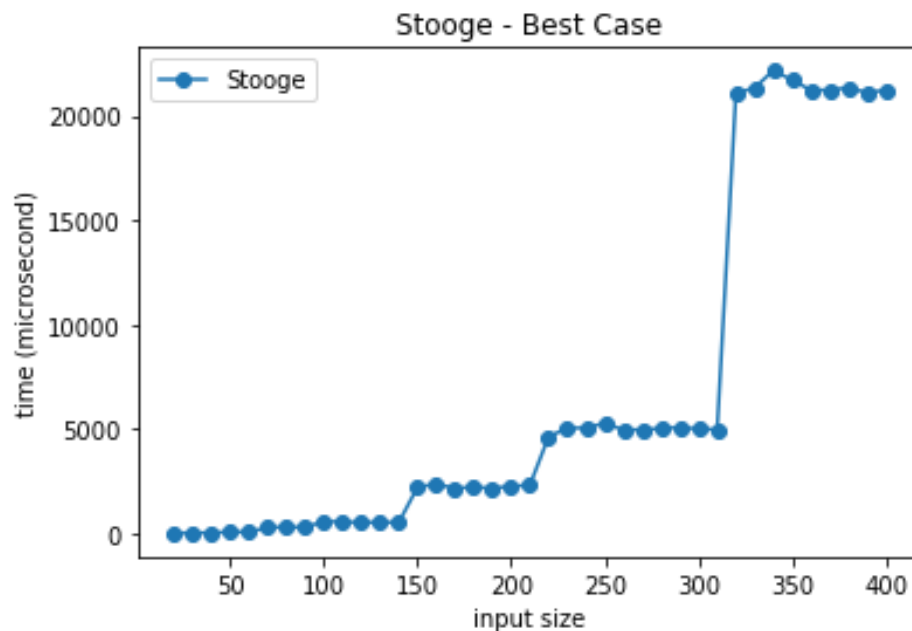
Stooge sort is a recursive sorting algorithm. The algorithm is defined as follows:

- If the value at the start is larger than the value at the end, swap them.
- If there are 3 or more elements in the list, then:
  - Stooge sort the initial 2/3 of the list
  - Stooge sort the final 2/3 of the list
  - Stooge sort the initial 2/3 of the list again<sup>7</sup>
- Worst-case performance  $O(n^{\log_3 \log 1.5})$  /  $O(n^{2.709})$ 
  - Space Complexity  $O(n)$

Data: (Unit microsecond)

(Sorted Array) => Best Case

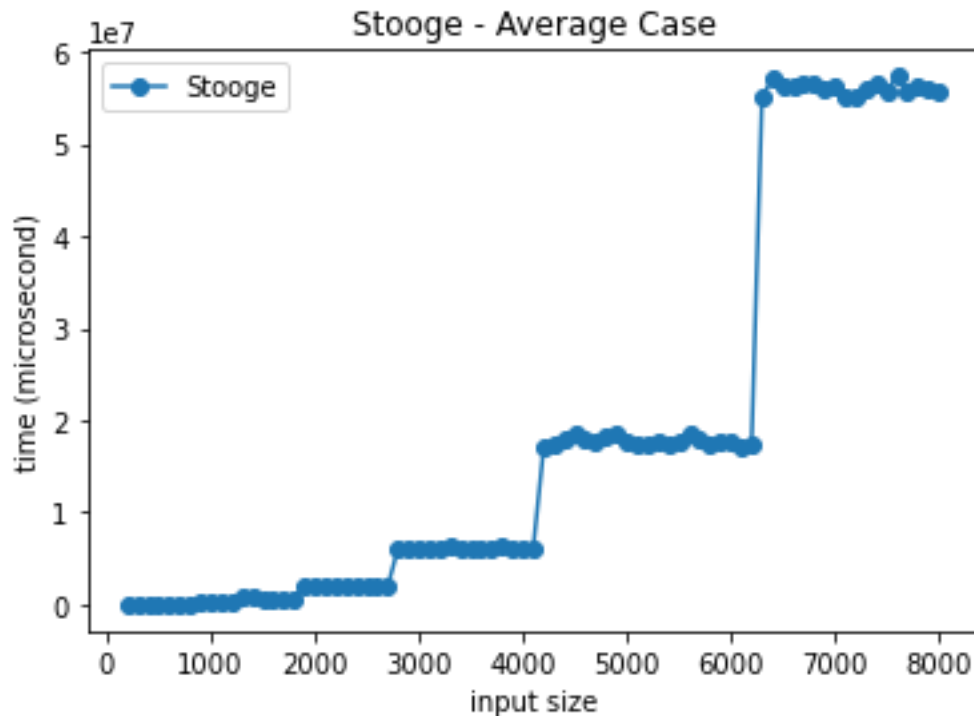
	20	60	100	140	180	220	260	300	340	380
<b>Stooge</b>	15.057143	82.585714	594.478571	538.007143	2226.914286	4648.557143	4964.935714	5016.285714	22172.807143	21352.142857



<sup>7</sup> [https://en.wikipedia.org/wiki/Stooge\\_sort](https://en.wikipedia.org/wiki/Stooge_sort)

(Unsorted Array) => Average Case

	200	900	1600	2300	3000	3700	4400	5100	5800	6500	7200	7900
Stooge	4650.8	221789.4	696608.9	1957236.9	6178074.2	6154738.9	17918993.6	17423569.3	17437037.0	56153075.4	55027805.7	55919418.2



Stable: No

```
Unsorted Array
[(9,a), (9,b), (8,a), (8,b), (12,a), (6,a), (9,c), (20,a), (18,a), (9,d), (8,c), (18,b), (1,a), (1,b), (18,c), (3,b)]
Stooge Sort
[(1,a), (1,b), (3,b), (6,a), (8,b), (8,a), (8,c), (9,d), (9,c), (9,b), (9,a), (12,a), (18,a), (18,c), (18,b), (20,a)]
```

Note:

When we look at Figure 1, it can be said that my average graph looks like the " $n^2$ " graph. It is observed that he makes some jump due to algorithm and computer. It can be said that there is a " $n^2$  or  $n^{2.7}$ " graph for our Best graph. This algorithm is a recursive algorithm. It can be said that it is quite inefficient. I observed that the experiment result was not stable. I can say that it is very inefficient compared to the others. I would not choose. Again, it is one of the reasons why I do not prefer it because of space complexity. It is quite inefficient.

### 3.5 Bitonic Sort Algorithm

Bitonic mergesort is a parallel algorithm for sorting. It is also used as a construction method for building a sorting network. The algorithm was devised by Ken Batcher.<sup>8</sup>

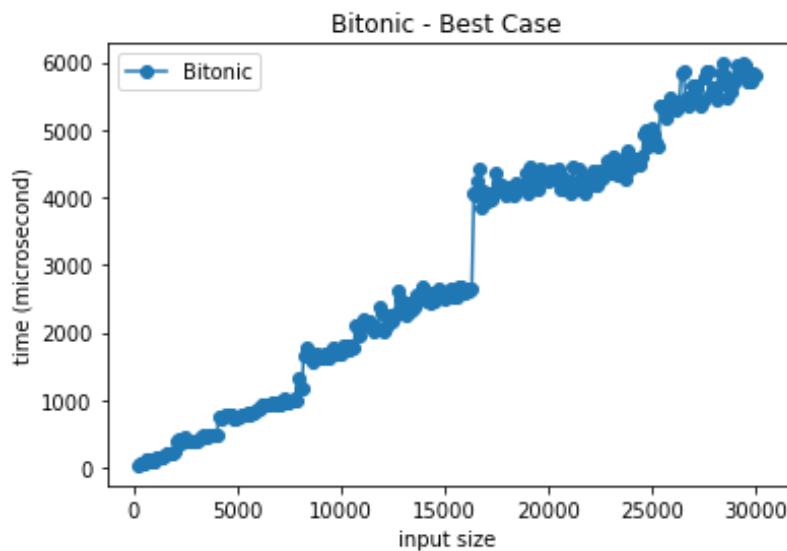
The number of comparisons done by Bitonic sort are more than popular sorting algorithms like Merge Sort [ does  $O(n \log n)$  comparisons], but Bitonic sort is better for parallel implementation because we always compare elements in predefined sequence and the sequence of comparison doesn't depend on data. Therefore it is suitable for implementation in hardware and parallel processor array.<sup>9</sup>

- Worst-case performance  $O(\log^2(n))$
- Best-case performance  $\Theta(\log^2(n))$
- Average performance  $\Omega(\log^2(n))$ 
  - Space Complexity  $O(n \cdot \log(n)^2)$

Data: (Unit microsecond)

Input	200	3200	6200	9200	12200	15200	18200	21200	24200	27200
Bitonic	36.807143	453.328571	938.821429	1683.892857	2068.192857	2601.4	4072.621429	4445.135714	4565.664286	5460.778571

(Sorted Array) => Best Case

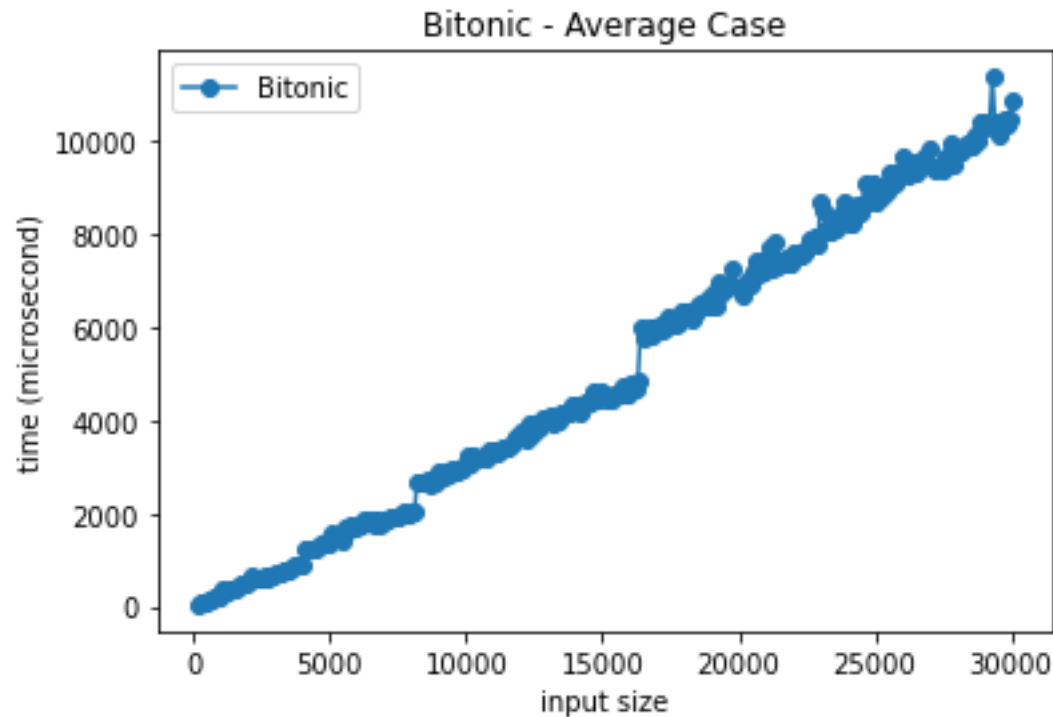


<sup>8</sup> [https://en.wikipedia.org/wiki/Bitonic\\_sorter](https://en.wikipedia.org/wiki/Bitonic_sorter)

<sup>9</sup> <https://www.geeksforgeeks.org/bitonic-sort/>

(Unsorted Array) => Average Case

Input	200	3200	6200	9200	12200	15200	18200	21200	24200	27200
Bitonic	39.25	725.542857	1886.771429	2807.885714	3589.471429	4439.271429	6372.042857	7287.971429	8236.978571	9371.421429



Stable: No

```
Unsorted Array
[(9,a), (9,b), (8,a), (8,b), (12,a), (6,a), (9,c), (20,a), (18,a), (9,d), (8,c), (18,b), (1,a), (1,b), (18,c), (3,b)]
Bitonic Sort
[(1,a), (1,b), (3,b), (6,a), (8,c), (8,a), (8,b), (9,a), (9,c), (9,b), (9,d), (12,a), (18,c), (18,a), (18,b), (20,a)]
```

Note:

When we look at Figure 1, it can be said that my average and best graph is similar to the "n" graph. An algorithm that only works at powers of 2. It can be said that it is very efficient. I observed that the experiment result was not stable. The algorithm needs more memory than others, I see the reason for this at the speed of the algorithm. I expected it to be more efficient in my experiment, but I still prefer it because of its efficiency.

## 4. Discussion and Result

time unit = ns = nanosecond					
Input_size	Comb	Gnome	Shaker	Stooge	Bitonic
2	902700	662000	568700	621900	790800
4	4800	1300	1700	3800	2900
8	5100	5600	5400	22800	5500
16	21700	16300	18100	99600	14000
32	37000	41400	44200	135600	66000
64	50300	111800	194800	924300	31800
128	233000	581700	473900	2111800	68600
256	146900	976700	983000	9832400	133200
512	295900	1772600	2072200	59336500	185700
1024	1266500	10512600	6580500	199616700	246400
2048	1104000	5688300	7278600	1825346300	1098000
4096	1664400	46448400	65661500	18439338400	1038400

Figure 2

*time comparison each of the algorithms with the same input*

In this assignment, I observed that the knowledge of theory does not make sense unless it is tested in experiment and that differences may arise between theory and practice. These are weak but noticeable differences. The reason for this may be the programming language we use, our computer speed, random sequences, numerical errors. In order to minimize these errors, the following can be done; choosing Java as the programming language, stopping every other program while the program is running, taking the average of a few test results.

In this experimental analysis, I saved the outputs to the ".csv" file. I processed these with the help of Python's Numpy, Pandas and Matplotlib libraries. I converted the data from nanoseconds to microseconds for convenience. I used my data by taking the average of 15 different data sets. The most accurate result is observed with too many input values and too many data sets.

As a result of the experiment, I observed that the Stooge algorithm was the most inefficient. Shaker and Gnome were close to each other. Comb and Bitonic algorithms were the most efficient of all. Bitonic may be preferred for powers of 2 and Comb for others. Of course, Shaker or Gnome should be used for stable state.

Stable status is important in ordering our data with more than one variable. For this reason, I tried all of them using the same array. As a result I observed that the Shaker and Gnome algorithms are stable.

As a result of this experiment, I had the opportunity to learn 5 different sorting algorithms with the subject of time complexity. I had the chance to research the properties of these algorithms and test them with each other.

Last Note:

I would say that the worst-case in this experiment is almost similar to the average-case. Although it is difficult to create the worst case scenario specifically for algorithms, we have the opportunity to learn worst-cases from algorithm analysis. As seen in the experiment, there are places where we can think of average-cases as worst-cases.

## *5. References*

- Wikipedia - *[https://en.wikipedia.org/wiki/Main\\_Page](https://en.wikipedia.org/wiki/Main_Page)*
- GeeksforGeeks - *<https://www.geeksforgeeks.org/>*
- Algorithms, 4th Edition by Robert Sedgewick and Kevin Wayne - *<https://algs4.cs.princeton.edu/home/>*
- Tudorcodes - *<https://www.tudorcodes.com/blog/stooge-sort/>*