

1. 课上测试完善:

```
def partition(a, p, r): # 以一个确定的基准元素a[p]对子数组a[p:r]进行划分
    i = p
    j = r+1
    x = a[p]
    while 1:
        while 1:
            i += 1
            if a[i] < x and i < r:
                pass
            else:
                break
        while 1:
            j -= 1
            if a[j] > x:
                pass
            else:
                break
        if i >= j:
            break
        mid = a[i]
        a[i] = a[j]
        a[j] = mid
    a[p] = a[j]
    a[j] = x
    return j
```

```
def randpartition(a, p, r): # 随机化划分
    i = random.randrange(p, r)
    mid = a[i]
    a[i] = a[p]
    a[p] = mid
    return partition(a, p, r)
```

```
def randselect(a, p, r, k): # 选择
    if p == r:
        return a[p]
    i = randpartition(a, p, r)
    j = i - p + 1
    if k <= j:
        return randselect(a, p, i, k)
    else:
        return randselect(a, i + 1, r, k-j)
```

```
a = [4, 5, 1, 6, 2, 7, 3, 8]
for i in range(1, 5):
    print(randselect(a, 0, 7, i))
```

2.

要使得 $P(n_1) = P(n_2) = \dots = P(n_d) = 0$ 且最高次系数为1
 则该多项式为 $(x-n_1)(x-n_2)\dots(x-n_d)$
 要求该多项式:
 ① 使用算法1:
 ~~$f(n[d], d) = \{$~~
~~如果 $d=1$, 返回 $x-n[0]$~~
~~否则 返回算法1 ($n[d-1]$) 和 $f(n, d-1)$ 的结果~~
 ~~$\}$~~
 效率 $1+2+\dots+(d-1) = \frac{(d+1)d}{2} = O(d^2)$
 ② 使用算法2: (假设 d 为2的幂次, 不足则补足)
 $f(n[d], d) \{$
 如果 $d=1$, 返回 $x-n[0]$
 否则 返回算法2计算的 $f(n, \frac{d}{2})$ 和 $f(n+\frac{d}{2}, \frac{d}{2})$ 的乘积)
 $\}$
 效率 $\frac{d}{2} \cdot \log \frac{d}{2} + (\frac{d}{4} \cdot \log \frac{d}{4}) \cdot 2 + (\frac{d}{8} \cdot \log \frac{d}{8}) \cdot 4 + \dots + \frac{d}{2} \cdot \log 1$
 $= \frac{d}{2} (\log \frac{d}{2} + \log \frac{d}{4} + \log \frac{d}{4} + \dots + \log \frac{d}{d})$
 $= \frac{d}{2} \cdot \log_2 d (\log_2 d - \log_2 2 - \log_2 2 - \dots - \log_2 d \log_2 d)$
 $= \frac{d \log_2 d}{4} (\log_2 d - \log_2 2)$
 $= O(d \log_2 d)$
 综合比较: 算法2的时间复杂度更低, 效率更高

3.

代码如下:

```
#include <iostream>
#include <fstream>
using namespace std;
```

```
int factorial(int n)
```

```
{
    int C = 1;

    while (n >= 1)
    {
        C = C * n;
        --n;
    }
}
```

```
    return C;
```

```
}//计算 n 的阶乘
```

```
int dic(int arrange[], int n)
```

```
{
    int head, tail, count = 0;

    if (n == 1)
        return 1;
    for (int i = 0; i < n - 1; ++i)
    {
        if (arrange[n - 1] > arrange[i])
            ++count;
    }
    head = factorial(n - 1) * count;
    tail = dic(arrange, n - 1);
```

```
    return tail + head;
```

```
}//递归计算字典序
```

```
void redic(int arrange[], int n, int order)
```

```
{
    int num, mid, * array, count;

    array = (int*)malloc(n * sizeof(int));
    if (!array)
        exit(-1);
    for (int i = 0; i < n; ++i)
```

```

        *(array + i) = 1;
for (int i = 0; i < n; ++i)
{
    num = factorial(n - i - 1);
    mid = order / num;
    order = order - mid * num;
    ++mid;
    count = 0;
    for (int k = 0; k < n; ++k)
    {
        if (*(array + k) == 1)
            ++count;
        if (count == mid)
        {
            arrange[i] = k + 1;
            array[k] = 0;
            break;
        }
    }
}
} //循环计算下一个排列

int main()
{
    int n, * arrange, * mid, order;
    ifstream in;
    ofstream out;

    in.open("input.txt", ios::in);
    if (!in)
    {
        cout << "文件打开失败! " << endl;
        exit(-1);
    }
    in >> n;
    arrange = (int*)malloc(n * sizeof(int));
    if (!arrange)
        exit(-1);
    mid = (int*)malloc(n * sizeof(int));
    if (!mid)
        exit(-1);
    for (int i = 0; i < n; ++i)
        in >> *(mid + i);
    in.close();

```

```

for (int i = 0; i < n; ++i)
    *(arrange + i) = *(mid + (n - i - 1));
out.open("output.txt", ios::out);
if (!out)
{
    cout << "文件打开失败! " << endl;
    exit(-1);
}
order = dic(arrange, n) - 1;
out << order << endl;
redic(arrange, n, order + 1);
for (int i = 0; i < n; ++i)
    out << *(arrange + i) << " ";

return 0;
}

```

运行结果如下图：

```

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2 6 4 5 8 3 1 7

```