

MSE Database Seminar - Fall 2017

GPU Databases

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Abstract

The content of this paper is splitted into three parts, it begins with an overview about GPU Database Systems. Followed by a port about MapD, that is an example of the available GPU database products. And finally a benchmark comparing the performance of MapD and PostgreSQL. The data of the benchmark is based on a part of the Unified New York City Taxi and Uber data.

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1. Introduction

During the Master of Science in Engineering the students have to participate at two seminars. The goal of these are to elaborate a theme on their own, discuss the result in group and write a paper about the topic.

The Databasesystems Seminar does a focus on GPU Database Systems. The students do have a closer look at a certain GPU Database product and have to do a benchmark.

The benchmark is based on the public NYC Taxi Rides dataset and the queries are predetermined by Prof. Stefan Keller.

2. GPU Databases

Related to the massive amount of data that is collected nowadays, the stagnation of CPU speed and the trend to use the GPU for tasks like machine learning, the database developers have discovered the GPU to improve the performance of their products, too. Hence the main idea of GPU databases is to perform some operations on the GPU for acceleration purposes.

Strengths GPUs do have their strengths on parallelable tasks. This is due to the fact that GPUs can have thousands of cores and high bandwidth memory on each card. Thus most of the GPU databases products focus on analytics.

Weaknesses Beside of these strengths, GPU database host several pitfalls like transfer of the data from the CPU to the GPU, the memory limitations and the massiv costs of GPU servers.

2.1. Components of a GPU database

The current section will give you an overview of the components a GPU database consists of. The paper [Bre+14] of Sebastian Bress et al. does split the components into Functional and Non-Functional properties. Since those categories are reasonable they are used in this paper as well. The interesting Non-Functional properties of GPU database are performance and portability. The next sub section will explain and list the Functional Properties.

2.1.1. Functional properties

Storage system

If we talk about storage systems there are several scenarios conceivable. First we the Video Random Access Memory (VRAM) of the GPUs we have an additional storage medium next to the Random-Access Memory (RAM) and the hard disk. The different mediums provides different advantages and disadvantages. Hard disks are persistent, cheap and have a huge capacity. Unfortunately they are pretty slow. Random-Access Memory is very fast compared to hard disks but the transfer to the GPU is still a bottleneck and it isn't a persistent storage, either. With Video Random Access Memory the bottleneck of the transfer disappears, but most of the time the storage capacity is highly limited.

2. GPU Databases

Storage model

In terms of storage model, there are row stores or column stores [AMH08]. Row stores store table records in a sequence of rows. Column stores store table records in a sequence of columns, the entries of a column is stored in contiguous memory locations [Bre14]. The advantages of column stores are tasks like aggregations, though row stores are more efficient if the result of a query returns multiple rows.

Processing model

There are two processing models in modern databases tuple-at-a-time and operator-at-a-time. The tuple-at-a-time is similar to an iterator, which iterates over the relevant tuples and applies the operations. The operator-at-a-time fits to GPUs, since it applies the same operation on a bulk of data.

Query processing

GPU database systems are able to use the GPU and the CPU, hence it is necessary to choose the right processor for the right task.

Query placement It is a extremely difficult task to decide which processing device is the most accurate for the current query.

Optimization To optimize the performance in therms of query execution time, several factors are include. For example the operations, the data, hardware specifications and even more.

Transaction support

An other problem are transactions and consistency on GPUs, until now there is almost no research done in this area. Thus most of the GPU database don't support transactions.

2.1.2. Proposed architecture

The paper [Bre+14] proposes an architecture with an in-memory storage, using a column store, an operator-at-a-time processing model, cross device processing and no transaction support. With regard to the portability a hardware oblivious architecture is suggested. To my point of view a hardware aware GPU database would make more sense. Since GPU databases are all about speed, use hardware specific tuning would result in more acceleration.

3. MapD

MapD is a GPU database with the goal to speed up queries and analytic tasks with the power of GPU's and their massive parallel architectures consisting of thousands of cores. The first prototype of MapD was develop in 2012 by Todd Mostak. A year leater MapD was incubated at the MIT's Computer Science and Artificial Intelligence Laboratory (CSAIL) database group and in September 2013 Todd Mostak founded MapD Technologies, Inc. They have got two products called MapD Core [Map17c] and MapD Immerse [Map17d]. The MapD Core SQL engine is an open source in-memory, SQL, GPU database and MapD Immerse is a tool for visual analytics on top of MapD Core SQL engine.



Figure 3.1.: *MapD*
Logo

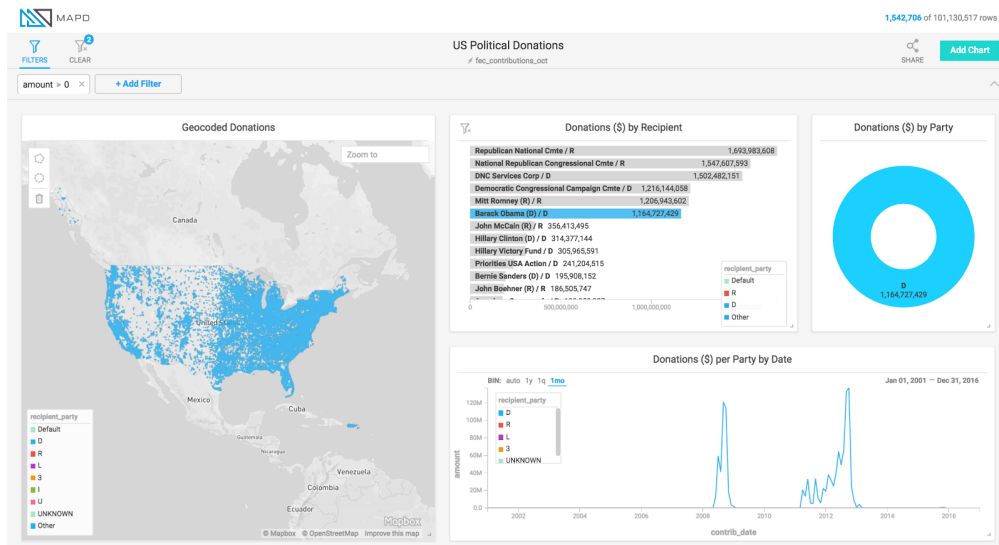


Figure 3.2.: *MapD Immerse* [Dup17]

3. MapD

3.0.1. Functional Properties

Related to the excellent Survey of Sebastian Bress et al. [Bre+14] MapD has the following functional properties.

Storage system

MapD has got a relational DBMS that is able to handle data amounts bigger than the memory space. But tries to hold as much data as possible in-memory to improve the performance.

Storage model

MapD uses a columnar layout to store the data and uses so called chunks, which split the columns in smaller pieces. The chunks are the basic units of the memory manager.

Processing model MapD processes on operator-at-a-time or one chunk per operation. Thus it is a block-oriented processing. The queries are compiled for the CPU and the GPU.

Query processing

Query placement In contrast to [Bre+14] the gained experience with MapD showed that MapD tries to run the queries on the GPU even if there isn't enough space and isn't able to handle such queries on his own. Hence the user has to switch the execution mode from GPU to CPU.

Optimization MapD's optimizer tries to execute the queries on the most suitable device, like text searching using an index on the CPU and table scans on the GPU.

Transactions support

MapD doesn't support transactions and doesn't support UPDATE or DELETE operations on the inserted data, either. That could be related to the difficult task of handling persistence between the CPU and GPU or even distributed systems at all.

3.1. Overview

The following section will give you an overview about the handling of MapD. Often there will be a comparison between MapD and PostgreSQL to point out the differences and similarities of these two databases.

3.1.1. Data Definition Language (DDL)

The DDL seems familiar since it uses SQL syntax [Map17a]. The syntax to handle users, databases, tables, and views is listed below.

User

- CREATE USER
- DROP USER
- ALTER USER

Database

- CREATE DATABASE
- DROP DATABASE

Table

- CREATE TABLE
- CREATE TABLE AS SELECT
- ALTER TABLE
- DROP TABLE
- TRUNCATE TABLE

View

- CREATE VIEW
- DROP VIEW

3. MapD

Datatypes

MapD supports the data types shown in table 3.1. To get a better intuition the corresponding PostgreSQL data types are listed as well.

MapD		PostgreSQL	
Data type	Size [bytes]	Data type	Size [bytes]
TEXT	Variable	text	Variable
TIMESTAMP	8	timestamp	8
TIME	8	time	8
DATE	8	date	4
FLOAT	4	real	4
DOUBLE	8	double precision	8
INTEGER	4	integer	4
SMALLINT	2	smallint	2
BIGINT	8	bigint	8
BOOLEAN	1	boolean	1
DECIMAL	8	numeric	variable

Table 3.1.: *Data types [Map17b] [Gro17a]*

As you can see MapD supports the common data types. And if you compare them to the corresponding PostgreSQL types they have got nearly the same names. In addition PostgreSQL provides further data types like json or box which allow extended possibilities of use.

3.1.2. Data Manipulation Language (DML)

As well as the DDL of MapD the DML uses the SQL syntax too. MapD currently supports the instructions:

- INSERT INTO
- SELECT

Until now there is no support for the operations:

- DELETE
- UPDATE

But they are in development, since MapD don't want to compromising the speed much with these instructions it may take a while.

Furthermore, MapD provides operations like EXPLAIN, LIKELY/UNLIKELY, Aggregate Functions and Conditional Expressions to improve the DML operations and extend the functionality.

3. MapD

3.1.3. Data import

MapD allows to import data from different sources as you can see in the following section.

COPY FROM

The COPY FROM operation is callable from the mapdql terminal and imports data from a local CSV or related format file into the database.

SQL Importer

The SQL Importer is a java tool that allows to run queries on other database via JDBC and stores the results in to MapD.

StreamInsert

The StreamInsert program could be attached at the end of a real-time stream processing engine like Kafka or a similar product to import stream data into MapD for further analytic tasks.

HDFS

The tool sqoop-export offers the possibility to import CSV or Parquet files from a HDFS file system into MapD database.

3.1.4. Data export

To export data from MapD, mapdql provides the command COPY TO that allows to export the result of a SELECT statement to a file. For example like:

- COPY (SELECT * FROM tweets) TO '/tmp/tweets.csv';

3. MapD

3.1.5. Client interfaces

MapD provides a tool called `mapdql` [Map17e] as a client-side SQL console that displays query results you submit to the MapD Core Server. The counterpart of PostgreSQL is `psql` [Gro17b].

Database connection

To following commands compares the connection to MapD respectively PostgreSQL with `mapdql` and `psql`.

`mapdql`: `mapdql <database> -u <user> -p <password> -port <port> -s <host>`

`psql`: `psql -h <host> -p <port> -U <user> -W <password> <database>`

Commands

The table 3.2 lists some basic commands of `mapdql` and `psql`. It is only a slight slice of all possible commands, but another good example to point out how much in common those two products have.

Command	mapdql	psql
List databases	<code>\l</code>	<code>\l</code>
List tables in database	<code>\t</code>	<code>\d</code>
Describe a table	<code>\d <table></code>	<code>\t <table></code>
Connect to a database	<code>\c</code>	<code>\c</code>
Print timing information	<code>\timing</code>	<code>\timing</code>
Switch to GPU mode	<code>\gpu</code>	-
Switch to CPU mode	<code>\cpu</code>	-
Quit	<code>\q</code>	<code>\q</code>

Table 3.2.: *Commands*

Alternative interfaces

Beside of `mapdql` MapD provides the following interfaces:

- JDBC
- ODBC
- `pymapd`
- Python JayDeBeApi
- Squirrel SQL
- RJDBC
- Apache Thrift

4. Benchmark

This section does concerned with a benchmark between the two databases PostgreSQL 10.0 and MapD 3.3.1. The dataset is based on a part of the New York City Taxi and Limousine Commission (TLC) Trip Record Data [New17]. All queries were executed on the same server.

4.1. Dataset



Figure 4.1.: *TLC Trip Record Data*

The New York City Taxi and Limousine Commission (TLC) Trip Record Data [New17] consists of trip records from the yellow and green taxis. Furthermore there are data from the For-Hire Vehicle (FHV). The data includes information about capturing pick-up and drop-off locations, times, trip distances, fares, rate types, and driver-reported passenger counts.

To simplify the handling with the huge amount of data we used the scripts from the Github repository [Sch17b], that is able to download all the data, provides the schema for PostgreSQL and consists of scripts to import the data.

4. Benchmark

4.1.1. Specification

To get a brief overview of the specification have a look at the following list.

Format: CSV
Yellow: January 2009 - June 2017
Green: August 2013 - June 2017
FHV: January 2015 - June 2017
Taxi trips: Over 1.1 billion [Sch17a]
Size: 267 GB [Sch17a]

4.1.2. Data subset

Due to the fact that not all GPU databases are able to handle queries based on data larger than the GPU memory, we had to shrink the dataset fitting into the available memory.

The GPU used on the server for the benchmark is a Nvidia Tesla K40m with 12 GB memory. Hence we used the yellow and green taxi trip data of the year 2015, that has a size of 12 GB, too. Consequently all queries fit into the GPU memory.

4. Benchmark

4.2. Queries

The Queries were predefined by Prof. Stefan Keller and are inspired by the Google BigQuery examples [LLC17]. All the queries are listed below, the used syntax is able to run on MapD.

```
1 SELECT passenger_count , Avg(total_amount)
2 FROM   trips
3 GROUP BY 1;
```

Listing 4.1: Query 1, Counts all the yellow taxi trips

```
1 SELECT cab_type_id , Count(*)
2 FROM   trips
3 GROUP BY 1;
```

Listing 4.2: Query 2, Calculates the average passenger amount per trip

```
1 SELECT passenger_count , Extract(year FROM pickup_datetime) , Count(*)
2 FROM   trips
3 GROUP BY 1, 2;
```

Listing 4.3: Query 3, Sums the yearly amount of passengers

```
1 SELECT passenger_count , Extract(year FROM pickup_datetime) , Cast(
   trip_distance AS INT) , Count(*)
2 FROM   trips
3 GROUP BY 1, 2, 3
4 ORDER BY 2, 4 DESC;
```

Listing 4.4: Query 4, Groups the amount of passenger by year regarding the trip distance

```
1 SELECT *
2 FROM   trips
3 WHERE  ( pickup_longitude BETWEEN -74.007511 AND -73.983479 )
4        AND ( pickup_latitude BETWEEN 40.7105 AND 40.731071 ) LIMIT 10;
```

Listing 4.5: Query 5, Queries all trips in a certain bounding box

```
1 SELECT Extract(HOUR FROM pickup_datetime) AS h, AVG(trip_distance / NULLIF(
   TIMESTAMPDIFF(HOUR,pickup_datetime , dropoff_datetime),0)) AS speed
2 FROM   trips
3 WHERE  ( pickup_longitude BETWEEN -74.007511 AND -73.983479 )
4        AND ( pickup_latitude BETWEEN 40.7105 AND 40.731071 )
5        AND trip_distance > 0
6        AND fare_amount / trip_distance BETWEEN 2 AND 10
7        AND dropoff_datetime > pickup_datetime
8        AND cab_type_id = 1
9 GROUP BY h
10 ORDER BY h;
```

Listing 4.6: Query 6, Determines the average speed of the yellow taxi trips by hour of the day in a bounding box

4. Benchmark

```
1 SELECT Extract(HOUR FROM pickup_datetime) AS h, Avg(trip_distance / NULLIF(
   TIMESTAMPDIFF(HOUR, pickup_datetime, dropoff_datetime), 0)) AS speed
2 FROM   trips
3 WHERE  trip_distance > 0
4        AND fare_amount / trip_distance BETWEEN 2 AND 10
5        AND dropoff_datetime > pickup_datetime
6        AND cab_type_id = 1
7 GROUP BY h
8 ORDER BY h;
```

Listing 4.7: Query 7, Computes the average speed of the yellow taxi trips by hour of the day

```
1 SELECT Extract(DOW FROM pickup_datetime) AS dow, Avg(trip_distance / NULLIF(
   TIMESTAMPDIFF(HOUR, pickup_datetime, dropoff_datetime), 0)) AS speed
2 FROM   trips
3 WHERE  trip_distance > 0
4        AND fare_amount / trip_distance BETWEEN 2 AND 10
5        AND dropoff_datetime > pickup_datetime
6        AND cab_type_id = 1
7 GROUP BY dow
8 ORDER BY dow;
```

Listing 4.8: Query 8, Calculates the average speed of the yellow taxi trips by day of the week

```
1 SELECT Extract(DOW FROM pickup_datetime) AS dow, Avg(trip_distance / NULLIF(
   TIMESTAMPDIFF(HOUR, pickup_datetime, dropoff_datetime), 0)) AS speed
2 FROM   trips
3 WHERE  ( pickup_longitude BETWEEN -74.007511 AND -73.983479 )
4        AND ( pickup_latitude BETWEEN 40.7105 AND 40.731071 )
5        AND trip_distance > 0
6        AND fare_amount / trip_distance BETWEEN 2 AND 10
7        AND dropoff_datetime > pickup_datetime
8        AND cab_type_id = 1
9 GROUP BY dow
10 ORDER BY dow;
```

Listing 4.9: Query 9, Determines the average speed of the yellow taxi trips by day of the week in a bounding box

4. Benchmark

4.3. Results

As already mentioned the queries of the benchmark were applied on a database without GPU acceleration, in our case PostgreSQL 10.0 and a on MapD 3.3.1 a GPU powered database. The dataset was introduced in section 4.1 and the queries are listed in the section 4.2.

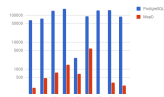


Figure 4.2.: A table beside a figure

Raw acceleration			Processed acceleration		
x	y	z	x	x	z
-6.160	7.185	3.235	-0.021	-0.094	10.001
-6.153	7.204	3.220	-0.042	-0.106	10.006
-6.174	7.170	3.356	-0.023	-0.015	10.041
-6.174	7.173	3.416	-0.033	-0.069	10.062

Table 4.1.: A table beside a figure

5. Conclusion

A. Architecture

Appendices

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