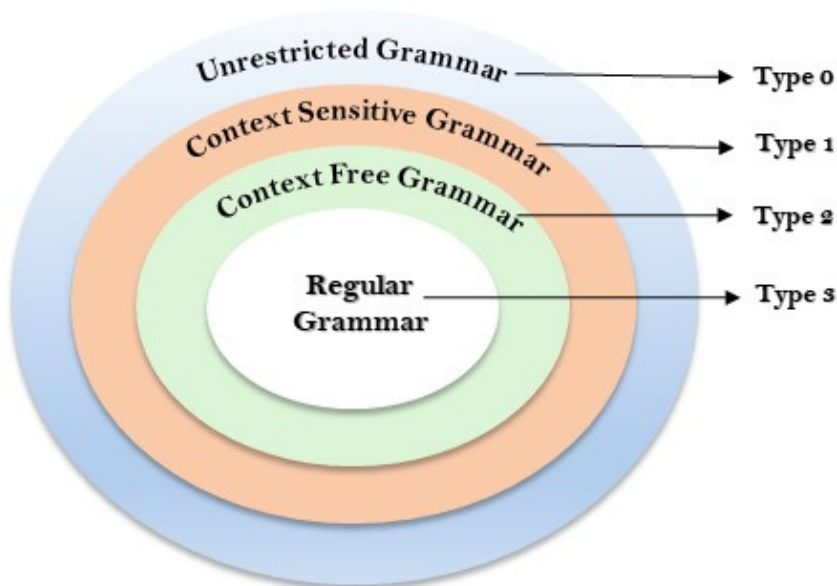


# Chomsky Hierarchy

Chomsky Hierarchy represents the class of languages that are accepted by the different machine. The category of language in Chomsky's Hierarchy is as given below:

1. Type 0 known as Unrestricted Grammar.
2. Type 1 known as Context Sensitive Grammar.
3. Type 2 known as Context Free Grammar.
4. Type 3 Regular Grammar.



**Fig: Chomsky Hierarchy**

This is a hierarchy. Therefore every language of type 3 is also of type 2, 1 and 0. Similarly, every language of type 2 is also of type 1 and type 0, etc.

## Type 0 Grammar:

Type 0 grammar is known as Unrestricted grammar. There is no restriction on the grammar rules of these types of languages. These languages can be efficiently modeled by Turing machines.

**For example:**

$bAa \rightarrow aa$

$S \rightarrow s$

## Type 1 Grammar:

Type 1 grammar is known as Context Sensitive Grammar. The context sensitive grammar is used to represent context sensitive language. The context sensitive grammar follows the following rules:

- The context sensitive grammar may have more than one symbol on the left hand side of their production rules.
- The number of symbols on the left-hand side must not exceed the number of symbols on the right-hand side.

The rule of the form  $A \rightarrow \epsilon$  is not allowed unless  $A$  is a start symbol. It does not occur on the

- right-hand side of any rule.

The Type 1 grammar should be Type 0. In type 1, Production is in the form of  $V \rightarrow T$

◦

Where the count of symbol in  $V$  is less than or equal to  $T$ .

### For example:

$S \rightarrow AT$

$T \rightarrow xy$

$A \rightarrow a$

## Type 2 Grammar:

Type 2 Grammar is known as Context Free Grammar. Context free languages are the languages which can be represented by the context free grammar (CFG). Type 2 should be type 1. The production rule is of the form

$A \rightarrow \alpha$

Where  $A$  is any single non-terminal and  $\alpha$  is any combination of terminals and non-terminals.

**For example:**

$$A \rightarrow aBb$$
$$A \rightarrow b$$
$$B \rightarrow a$$

### Type 3 Grammar:

Type 3 Grammar is known as Regular Grammar. Regular languages are those languages which can be described using regular expressions. These languages can be modeled by NFA or DFA.

Type 3 is most restricted form of grammar. The Type 3 grammar should be Type 2 and Type 1. Type 3 should be in the form of

$$V \rightarrow T^* V / T^*$$

**For example:**

$$A \rightarrow xy$$