

TCP SYN Flood Denial of Service Attack: Implementation and Analysis

Mushfiqur Rahman

Student ID: 2005107

Arnab Dey Kabbya

Student ID: 2005112

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1 Introduction

A TCP SYN flood attack is a type of Denial of Service (DOS) attack that exploits vulnerabilities in the TCP three-way handshake process. By overwhelming a target server with half-open connections, the attack consumes system resources, rendering the service unavailable to legitimate users. This project implements a TCP SYN flood attack in Python, showcasing both non-spoofed and IP-spoofed strategies, and analyzes their effectiveness through empirical results.

2 Overview

The TCP SYN flood attack targets the connection establishment phase of the TCP protocol. By sending a large volume of SYN packets without completing the handshake, the attacker forces the server to allocate resources for connections that never finalize, leading to resource exhaustion.

2.1 TCP Three-Way Handshake

The TCP connection establishment involves three steps:

1. **SYN:** The client sends a SYN packet to the server to initiate a connection.
2. **SYN-ACK:** The server responds with a SYN-ACK packet, acknowledging the request.
3. **ACK:** The client sends an ACK packet to complete the handshake.

2.2 SYN Flood Attack Process

In a SYN flood attack, the attacker:

1. Sends numerous SYN packets to the target server.
2. Fails to send the final ACK packet, leaving connections half-open.
3. Overwhelms the server's connection table, exhausting memory and CPU resources.
4. Prevents legitimate users from establishing connections.

3 Implementation Details

3.1 Core Components

The implementation consists of several key components:

- **Packet Building Utilities:** Functions to compute checksums and construct IP and TCP headers.
- **IP Packet Class:** Configurable IPv4 headers with fields like source and destination IPs.
- **TCP Packet Class:** Constructs TCP headers with SYN flags and random sequence numbers.

3.2 Attack Strategies

The project supports two attack modes:

- **Non-Spoofed Attack:** Uses the attacker's real IP address, making it easier to trace but simpler to implement.
- **IP Spoofed Attack:** Uses randomized source IP addresses, increasing attack effectiveness and bypassing basic defenses.

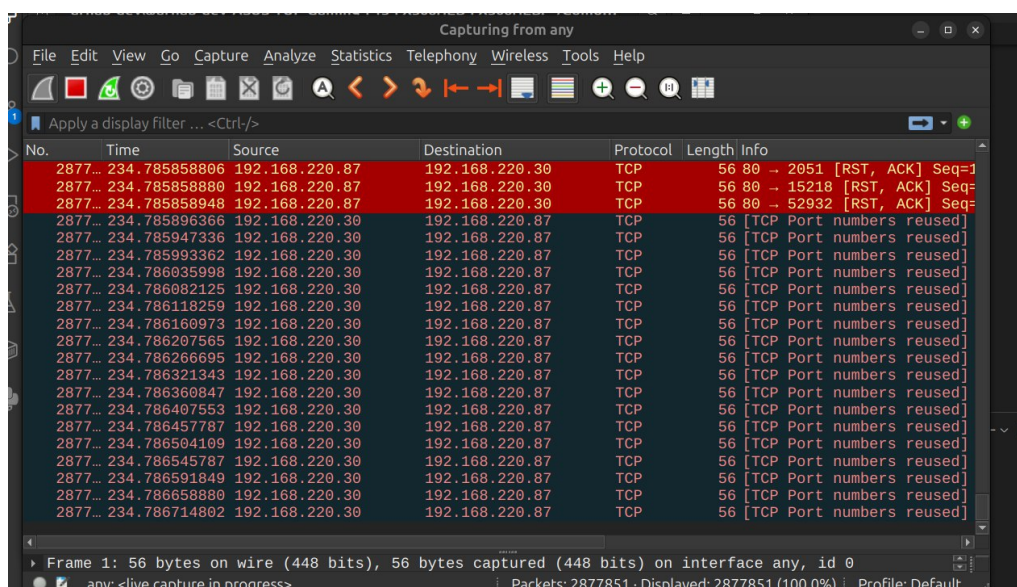
3.3 Key Features

- **Raw Socket Implementation:** Allows direct packet crafting for precise control.
- **Response Monitoring:** Captures SYN-ACK responses to verify target connectivity.
- **Configurable Parameters:** Adjustable target IP, port, attack rate, and thread count.
- **IP Spoofing Option:** Enhances attack effectiveness with randomized source addresses.
- **Rate Control:** Precise control over packets per second (PPS).

4 Attack Results Analysis

The attack's effectiveness was evaluated with and without IP spoofing, as shown in the following figures.

4.1 Without IP Spoofing



No.	Time	Source	Destination	Protocol	Length	Info
2877...	234.785858806	192.168.220.87	192.168.220.30	TCP	56	80 → 2051 [RST, ACK] Seq=1
2877...	234.785858880	192.168.220.87	192.168.220.30	TCP	56	80 → 15218 [RST, ACK] Seq=
2877...	234.785858948	192.168.220.87	192.168.220.30	TCP	56	80 → 52932 [RST, ACK] Seq=
2877...	234.785896366	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.785947336	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.785993362	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786035998	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786082125	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786118259	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786160973	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786207565	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786266695	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786321343	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786360847	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786407553	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786457787	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786504109	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786545787	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786591849	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786658880	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]
2877...	234.786714802	192.168.220.30	192.168.220.87	TCP	56	[TCP Port numbers reused]

Figure 1: Network traffic during a non-spoofed TCP SYN flood attack. The attacker uses their real IP address, resulting in easily traceable traffic.

- **Attack Method:** Uses the attacker's actual IP address without any IP randomization.
- **Detection Risk:** High, since the traffic originates from a single source.

- **Effect on Target:** Causes moderate resource consumption; blocking is easier.

4.2 With IP Spoofing

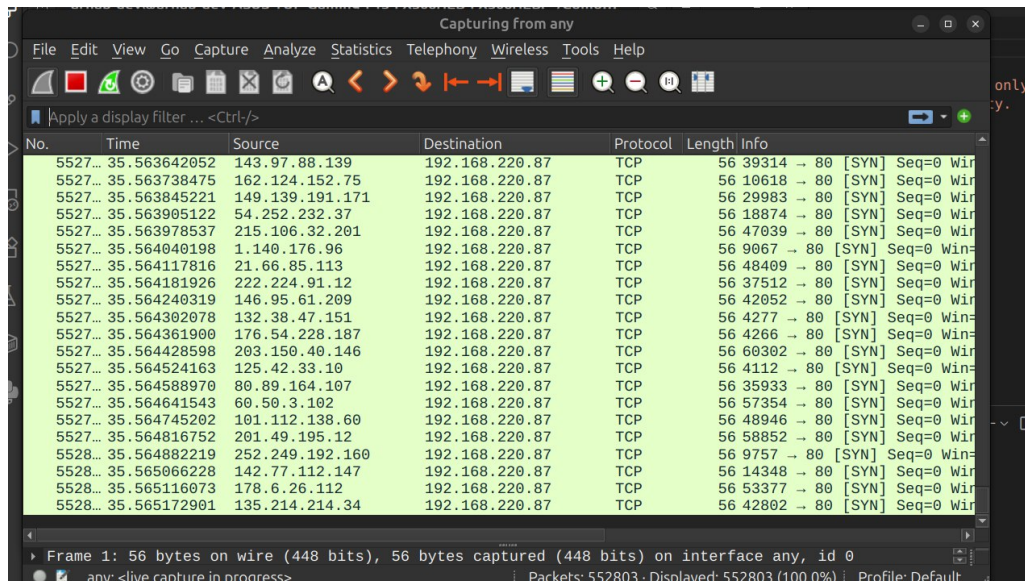


Figure 2: Network traffic during an IP-spoofed TCP SYN flood attack. The attacker sends SYN packets using randomized source IP addresses.

- **Attack Method:** Randomized spoofed IP addresses for each SYN packet.
- **Detection Risk:** Low, because packets appear to come from multiple sources.
- **Effect on Target:** Significantly higher resource exhaustion; filtering becomes complex.

4.3 Victim Impact

The resource utilization on the victim server differs significantly based on whether IP spoofing is used. Each scenario is explained below with its corresponding visualization.

Without IP Spoofing

- The server experiences moderate resource strain due to repeated SYN packets from a single IP.
- Blocking the attacker's IP using firewall rules or rate-limiting is straightforward.
- Network traffic patterns are consistent and easily identifiable.

With IP Spoofing

- The server faces severe resource depletion because each SYN packet appears to originate from a different IP.
- Filtering becomes extremely difficult, making the attack far more effective.
- Connection queues fill rapidly, leading to service denial for legitimate users.

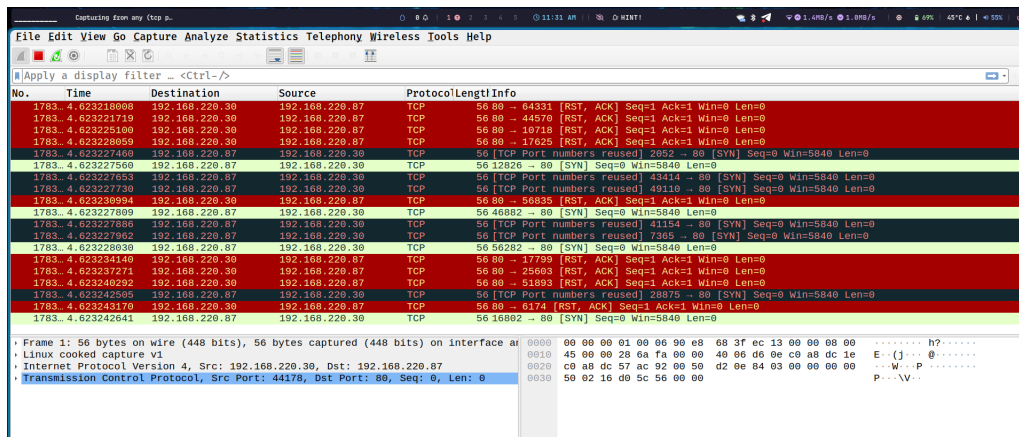


Figure 3: Server resource usage during a non-spoofed TCP SYN flood attack. Traffic originates from a single IP, making detection and blocking easier.

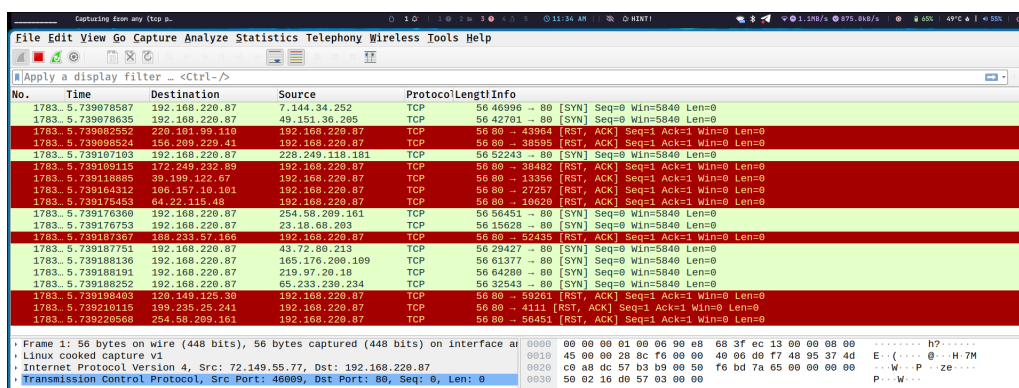


Figure 4: Server resource usage during an IP-spoofed TCP SYN flood attack. The system struggles with diversified spoofed IP traffic, making mitigation challenging.

5 Technical Implementation

5.1 Packet Structure

The TCP SYN flood attack constructs raw packets by manually crafting IP and TCP headers. The key fields used in the attack are summarized below:

IP Header Fields

Field	Description
Version	4 (IPv4)
Header Length	5 (20 bytes)
Type of Service	0
Total Length	40 bytes (IP + TCP)
Identification	Randomly generated
Flags	Don't Fragment
TTL	64
Protocol	TCP (6)
Source IP	Attacker IP (or spoofed)
Destination IP	Victim server IP

Table 1: IP Header Fields used in the SYN flood packet

TCP Header Fields

Field	Description
Source Port	Randomized
Destination Port	Target service port
Sequence Number	Randomized
Acknowledgment Number	0
Header Length	5 (20 bytes)
Flags	SYN (0x02)
Window Size	5840
Checksum	Calculated dynamically

Table 2: TCP Header Fields used in the SYN flood packet

The packet construction ensures compliance with the TCP/IP protocol while allowing flexibility for attack customization. The fields highlighted above are dynamically modified during the flood to avoid detection and improve attack efficiency.

5.2 Attack Flow

The attack follows these steps:

1. **Initialization:** Sets up raw sockets with IP header inclusion.
2. **Target Validation:** Verifies the target IP address format.
3. **Connectivity Test:** Sends an initial SYN packet to confirm reachability.
4. **Thread Deployment:** Launches multiple worker threads for parallel packet transmission.
5. **Flood Execution:** Continuously sends SYN packets at the configured rate.
6. **Monitoring:** Tracks SYN-ACK responses to adjust attack parameters.

6 Defense Mechanisms

To mitigate TCP SYN flood attacks, the following defenses can be implemented:

6.1 Network-Level Defenses

- **SYN Cookies:** Use stateless connection handling to avoid resource allocation for half-open connections.
- **Connection Rate Limiting:** Restrict the number of new connections per IP address.
- **Firewall Rules:** Block traffic exhibiting suspicious patterns.
- **Load Balancing:** Distribute incoming connections across multiple servers.

6.2 System-Level Defenses

- **TCP Backlog Tuning:** Increase the connection queue size to handle more half-open connections.
- **Timeout Reduction:** Accelerate the cleanup of half-open connections.

- **Resource Monitoring:** Continuously track system resource usage.
- **Intrusion Detection:** Implement automated systems to recognize attack patterns.

7 Conclusion

The TCP SYN flood attack remains a potent threat to network availability. This implementation highlights the simplicity of executing such attacks and the enhanced effectiveness of IP spoofing. By understanding these attack vectors, network administrators can develop robust countermeasures to ensure service availability and security.

8 References

- Eddy, W. (2007). TCP SYN Flooding Attacks and Common Mitigations. RFC 4987, IETF.
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- Postel, J. (1981). Transmission Control Protocol. RFC 793, IETF.