

Lecture 7 FORMAL POWER SERIES.

Some algebraic systems we have experience with:

Group:	- assoc./unity/inverse	$\left[\begin{array}{l} \text{assoc.} \\ \text{commutative} \end{array} \right] \quad \mathbb{Z}, +$
Abelian group:	- Commutative	$\left[\begin{array}{l} \text{commutative} \\ \text{assoc./unity law} \end{array} \right] \quad \mathbb{Z}_m^n$
Ring: (commutative with unity)	+ : Abelian group	$\left[\begin{array}{l} \text{assoc./unity law} \\ \text{dist. law} \end{array} \right] \quad (\mathbb{Z}, +, \cdot) \quad \mathbb{Z}_m^n$
Field: 2 operations	+ : Abelian group	$\left[\begin{array}{l} \text{assoc./unity law} \\ \text{dist. law} \end{array} \right] \quad \mathbb{R}, \mathbb{C}$
	$\therefore \mathbb{R} \setminus \{0\}$ is abelian group	$\left[\begin{array}{l} \text{assoc./unity law} \\ \text{dist. law} \end{array} \right] \quad \mathbb{Q}$

We learnt one more today. Let R be a ring and q be some symbol

$$R[[q]] := \left\{ \sum_{i=0}^{\infty} a_i q^i = a_0 + a_1 q + a_2 q^2 + a_3 q^3 + \dots \mid a_i \in R \right\}$$

$\sum a_i q^i \hookrightarrow (a_0, a_1, a_2, a_3, \dots) \leftarrow$ infinite string.
Since we calculate only a finite part of this we use the symbol $0(q^m)$ to denote when we have calculated until.
Ex: $1+q+q^2+0(q^3)$
 \uparrow big 0.
 $= a = \sum a_i q^i = \sum b_i q^i \quad \text{if } a_i = b_i \text{ for all } i$
 $+ a+b = \sum_{i=0}^{\infty} (a_i + b_i) q^i \quad (\text{either is defined in ring), add - assoc/comm/0/inverse}$

Prop: Let $a = \sum a_i q^i$. Then a has a multiplicative inverse if and only if a_0 is invertible in R .

Proof: Sp a is invertible. Let $b = \sum b_i q^i$ be the inverse then $ab = 1 \Rightarrow (a_0 + a_1 q + \dots)(b_0 + b_1 q + \dots) = 1 \Rightarrow a_0 b_0 = 1 \Rightarrow a_0$ is invertible in R .

Conversely, let a be given and let $b = \sum_{i=0}^{\infty} b_i q^i$. Consider the equation $ab = 1$

$$\begin{aligned} \Rightarrow a_0 b_0 = 1 &\quad \Rightarrow b_0 = a_0^{-1} \\ a_0 b_1 + a_1 b_0 = 0 &\quad \Rightarrow a_0, a_1, b_0 \text{ known } b_1 \\ a_0 b_2 + a_1 b_1 + a_2 b_0 = 0 &\quad \Rightarrow b_2 \text{ can be calculated} \\ a_0 b_3 + a_1 b_2 + a_2 b_1 + a_3 b_0 = 0 &\quad \Rightarrow b_3 \dots \end{aligned}$$

This is a 'triangle' system and can be solved.
Ex. Write an inductive proof that $b_0, b_1, \dots, b_n, \dots$ can be found from this system.

Remark: If $R = \mathbb{Z}$, a_0 is invertible iff $a_0 = \pm 1$.

(2) We can define $R[[x, y]]$ as $R[x][y]$ etc.
Same for $R[x, y] \leftarrow$ polynomials in 2 variables

(3) Can think of calculating inverses as follows:

$$\frac{1}{1+a_1 q + a_2 q^2 + \dots} = 1 - (a_1 q + a_2 q^2 + \dots)^{-1} + (a_1 q + a_2 q^2 + \dots)^2 + \dots$$

$$x \quad a \cdot b = (\sum a_i q^i)(\sum b_j q^j) = \sum_{i+j=0}^n c_i q^i, \text{ where } c_n = a_0 b_0 + a_1 b_1 + \dots + a_n b_n$$

$$\begin{aligned} \text{Example: } (0, 1, 0, 0, \dots) \times (0, 0, 1, 0, 0, \dots) &= (0, 0, 1, 0, 0, \dots) \\ c_0 = 0 \cdot 0 & \\ c_1 = 1 \cdot 0 + 0 \cdot 1 &= 0 \\ c_2 = 0 \cdot 0 + 1 \cdot 1 + 0 \cdot 0 &= 1 \\ c_3 = 0 \cdot 0 + 1 \cdot 0 + 0 \cdot 1 + 0 \cdot 0 &= 0 \\ q \rightsquigarrow (0, 1, 0, 0, \dots) & \\ q^2 \rightsquigarrow (0, 0, 1, 0, 0, \dots) & \end{aligned}$$

$$\text{Ex: Show } q^k \cdot q^l = q^{k+l} \text{ in this.}$$

Third operation $\cdot : R \times R[[q]] \rightarrow R[[q]]$
Scalar multiplication.

$$\begin{aligned} c a(q) &= \sum_{i=0}^{\infty} c a_i q^i \quad \text{for } c \in R. \\ - (c_1 c_2) a(q) &= c_1 (c_2 a(q)) \\ - \text{Dist laws} & \quad - c (a(q) + b(q)) = (c a(q)) + (c b(q)) \text{ etc.} \\ - 1 \cdot a = a & \quad (c_1 + c_2) a(q) = c_1 a + c_2 a. \\ \text{We say } (R[[q]], +, \text{ scalar mult}) & \text{ is a MODULE} \\ (R[[q]], +, \times, \text{ scalar}) & \text{ is an algebra.} \end{aligned}$$

Can compute coeffs of q^m for $m=0, 1, 2, \dots$ from here.

Def: Let $a(q) = \sum a_i q^i$. The formal derivative is defined as $D(a(q)) = \sum_{k=1}^{\infty} k a_k q^{k-1}$.

Remarks (1) Can prove binet rules for diff from here.

(2) D^k is defined by iteration. $D^k = D(D^{k-1})$ etc.

$$\begin{aligned} (3) \quad \text{if } D a(q) = 0 &\Rightarrow k a_k = 0 \text{ for } k \geq 1 \\ &\Rightarrow a_{k+1} = 0 \text{ for } k \geq 1 \\ \text{so } a(q) &= a_0 + a_1 q + a_2 q^2 + \dots \\ &= a_0. \end{aligned}$$

$$(4) \quad \exp(x) \stackrel{def}{=} \sum_{k=0}^{\infty} \frac{x^k}{k!}$$

$$\begin{aligned} D(\exp(x)) &= \sum_{k=1}^{\infty} \frac{k x^{k-1}}{k!} = \sum_{k=1}^{\infty} \frac{x^{k-1}}{(k-1)!} \\ &= \exp(x). \end{aligned}$$

$e \leftarrow$ meaningless here, but we still use it.

modules are called
the main object of study in Linear Algebra

- $R[[q]] \leftarrow$ Ring structure
- $R[[q]] \leftarrow$ Ring of polynomials $\sum_{i=0}^n a_i q^i$, $a_i \in R$, $n = \text{degree}$
also algebra.
- Usually we use $\mathbb{R}[[q]]$ in combinatorics, on occasion \mathbb{Q} .

Ex: Sage: try $R<q> = \text{PowerSeriesRing}(27)$ in this.
(with power series context)
 $(1+q)^2 = \dots$
 $(1+q)^{100} = \dots$
Default calc.

$$\text{Theorem: (Imp)} \quad (1-q)^{-1} = 1 + q + q^2 + \dots \quad (\text{in FPPS.})$$

Proof: Consider $(1-q)(1+q+q^2+\dots)$
Coeff of q^m : $1 + (-1)(1) + \dots + (-1)^m$
coeff of q^m : $1 \cdot 1 = 1$.
so $(1-q)(1+q+q^2+\dots) = 1$. \square

Prop: (Taylor expansion) $a = \sum a_k q^k$
 $a(q) = a_0 + D(a(0))q + D^2 a(0) \frac{q^2}{2!} + \dots$

Here if $b(q) = \sum b_i q^i$, $b(q) = b_0$.

Proof: Ex (induction).

Remark: Sp $b(q) = 0$
then $a(b(q)) = \sum_{i=0}^{\infty} a_i (b(q))^i$
makes sense.