

STACKS, QUEUES, and DEQUEUES

Stacks : A stack is a collection of objects that are inserted and removed according to the LIFO (Last-in-First-out) principle.

Def A stack is a collection of objects that are inserted and removed according to the LIFO (Last-in-First-out) principle.

We insert objects at the top of the stack

To its bottom by pushing them and we can remove from the top by popping the element and that technique called last-in-first-out.

The ADT Stack interface is an interface which is implemented by The Stack Class and you can show the interface as following Syntax :

Codes

```
public interface Stack<E> {  
    int size(); // return stack number of elements  
    boolean isEmpty(); // return true if size = 0  
    void push(E e); // insert object in stack  
    E top(); // return the top element in the stack  
    E pop(); // remove and return the top element in the stack  
}
```

now we will implement the Stack Data Structure based on Array as following Syntax :

Code:

$O(N)$ →

Space-
Complexity

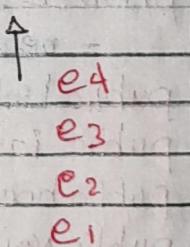
```
public class ArrayStack<E> implements Stack<E> {
    public static final int CAPACITY = 1000; //N
    private E[] data;
    private int t = -1; // index of The Top element
    public ArrayStack() { this(CAPACITY); }
    public ArrayStack(int capacity) {
        data = (E[]) new Object[capacity]; //safe cast
    }
    public int size() { return (t + 1); }
    public boolean isEmpty() { return (t == -1); }
    public void push(E e) throws IllegalStateException {
        if (size() == data.length)
            throw new IllegalStateException("stack is full");
        data[t + 1] = e;
    }
}
```

$O(1)$ →

Time
Complexity

```
public E top() {
    if (isEmpty()) return null;
    return data[t];
}

public E pop() {
    if (isEmpty()) return null;
    E answer = data[t];
    data[t] = null;
    t--;
    return answer;
}
```



Of Course, you notice that there is one drawback for this implementation which is based on fixed capacity so there are limits of the ultimate size.

now we will use **LinkedList** Datastructure To solve
The problem of fixed Capacity in Array using **Adapter**
Design Pattern.

Adapter pattern is to define a new Class in such a way
That it Contains an instance of The **existing** Class as
a hidden Field, and Then to implement each method of
the new Class using methods of This **hidden** instance
Variable.

we **Specify** we will implement the Stack using
SinglyLinkedList Specifically as shown Syntax:

Codes

```
public class LinkedStack <E> implements Stack <E> {  
    private SinglyLinkedList <E> list  
    = new SinglyLinkedList <E>();  
    public LinkedStack() {}  
    public int size() { return list.size(); }  
    public boolean isEmpty() { return list.isEmpty(); }  
    public void push (E element) { list.addFirst(element); }  
    public E top() { return list.getFirst(); }  
    public E pop () { return list.removeFirst(); }  
}
```

now we want to use the stack DataStructure
in ~~the~~ something like Matching Tags in a Markup language
Like **HTML** and **XML**.

we want To make an Algorithm That Take a string Html Syntax and Test that every tag has a closing tag or not and we write it as following Syntax 8

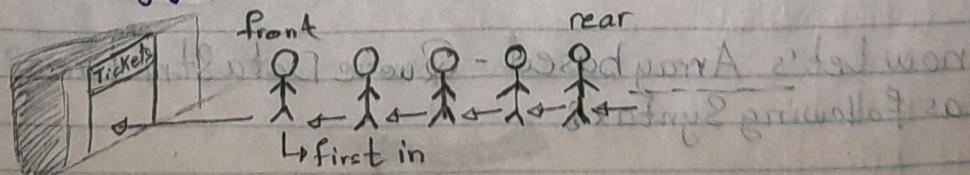
Code:

```
public static boolean isHtmlMatched ( String html ) {  
    Stack < String > buffer = new LinkedStack <> ();  
    int j = html . indexOf ( '<' );  
    while ( j != -1 ) {  
        int k = html . indexOf ( '>', j + 1 );  
        if ( k == -1 ) return false; // invalid tag  
        String tag = html . substring ( j + 1, k );  
        opening & tag if ( ! tag . startsWith ( "/" )) buffer . push ( tag );  
        else { // closing tag  
            if ( buffer . isEmpty ()) return false;  
            to match if ( ! tag . substring ( 1 ). equals ( buffer . pop ()) )  
                return false; // mismatch tag  
        }  
        j = html . indexOf ( '<', k + 1 );  
    }  
    return buffer . isEmpty (); // were all opening tags matched ?  
}
```

Queues 8

Def:

It is a Close "Cousin" of the Stack, but a Queue is a collection of objects that are inserted and removed according to the (first-in-first-out) FIFO principle.



`java.util.Queue;`

The Queue Class implements ADT Queue interface and This Interface is written as following
Syntax :

Code:

```
public interface Queue <E> {  
    int size(); // return number of elements  
    boolean isEmpty(); // return true if size == 0  
    void enqueue(E e); // insert element at rear  
    E first(); // return first element  
    E dequeue(); // return and remove first element  
}
```

now we will implement the Queue Data Structure based on fixed Array as a first development step, but we want to use array Circularly.
we will use The modulo operator % using The following formula :

$$f = (f + 1) \% N$$

↑ new position ↑ front index → length of Array

0	1	2	f	4	5	6	7	8	9
data				F	G	H	I	J	

M	N			F	G	H	I	J	K	L
data										

now Let's Array based - Queue Data Structure as following Syntax :

Code:

$O(N)$

Space complexity

```
public class ArrayQueue<E> implements Queue<E> {
```

```
    private E[] data; public static int final N = 1000;
```

```
    private int f = 0; // front index
```

```
    private int sz = 0; // size
```

```
    public ArrayQueue() { this(N); }
```

```
    public ArrayQueue(int capacity) {
```

```
        data = (E[]) new Object[capacity];
```

```
}
```

```
    public int size() { return sz; }
```

```
    public boolean isEmpty() { return (sz == 0); }
```

```
    public void enqueue(E e) throws IllegalStateException {
```

```
        if (sz == data.length)
```

```
            throw new IllegalStateException("Queue is full");
```

```
        int avail = (f + sz) % data.length; // Available index
```

```
        data[avail] = e;
```

```
        sz++;
```

```
    public E first() {
```

```
        if (isEmpty()) return null;
```

```
        return data[f];
```

```
}
```

```
    public E dequeue() {
```

```
        if (isEmpty()) return null;
```

```
        E answer = data[f];
```

```
        data[f] = null;
```

```
        f = (f + 1) % data.length; // update front index
```

```
        sz--;
```

```
        return answer;
```

```
}
```

```
Time complexity
```

Now let's implement **LinkedQueue** class using **Adapter Design Pattern** by **LinkedList** as following Syntax :

Code:

```
public class LinkedQueue <E> implements Queue <E> {  
    private SinglyLinkedList <E> list  
        = new SinglyLinkedList <E>;  
    public LinkedQueue () {}  
    public int size() { return list.size(); }  
    public boolean isEmpty() { return list.isEmpty(); }  
    public void enqueue(E element) { list.addFirst(element); }  
    public E first() { return list.first(); }  
    public E dequeue() { return list.removeFirst(); }  
}
```

If we want to make a **Circular Queue** interface To rotate The elements in the Queue, we need To make **CircularQueue** extends **Queue** interface as following

Syntax :

Code:

```
public interface CircularQueue <E> extends Queue <E> {  
    void rotate();  
} // Rotates The front element of The Queue to the back of it
```

OK, let's use The **Queue** DataStructure To solve The **josephus Problem** which is a children's game called **Hot Potato**. That we have n children sitting in a Circular way and each of them pass the potato to his next child. When the leader rings the bell, The child who has the potato at that moment will leave the Game and continue the process up to be 1 remaining child. we want to know who is the last survivor when the leader rings bell at kth child.

Let's solve This problem as following Syntax 8

Code:

```
public class Josephus {  
    public static <E> E josephus(CircularQueue<E> queue,  
        int K) {  
        if (queue.isEmpty()) return null;  
        while (queue.size() > 1) {  
            Rotate by K-1 index  
            for (int i = 0; i < K - 1) queue.rotate();  
            E e = queue.dequeue();  
            delete front K-1 element  
            System.out.println("ps." + e + " is out");  
            queue.enqueue(e);  
        }  
        return queue.dequeue(); // winner  
    }  
  
    public static <E> CircularQueue<E> buildQueue  
        (E a[]) {  
        CircularQueue<E> queue  
        = new LinkedCircularQueue<E>();  
        for (int i = 0; i < a.length; i++)  
            queue.enqueue(a[i]);  
        return queue;  
    }  
  
    public static void main(String[] args) {  
        String[] a1 = {"Alice", "Bob", "Cindy", "Doug"};  
        System.out.println("First Winner is " +  
            josephus(buildQueue(a1), 3));  
    }  
}
```

Double Ended - Queues

Defs

it's a datastructure like a Queue That Supports insertion and deletion at both the front and the back of The Queue.

The Deque class implements ADTDeque interface and we can write this interface as following Syntax

Codes

```
public interface Deque <E> {  
    int size();  
    boolean isEmpty();  
    E first();  
    E last();  
    void addFirst(E e);  
    void addLast(E e);  
    E removeFirst();  
    E removeLast();  
}
```

Time Complexity

now we should make it Circular and use the following 2 formulas

$$\text{frontAdd, } f = (f - 1 + N) \% N$$

we can implement The Deque using DoublyLinkedList by Adapter DesignPattern

notes

Any method in deque ADT interface that insert or remove element may throw a NoSuchElementException or accessing an element and return it