

Artificial Intelligence (CS-401)

Chapter 6: Constraint Satisfaction Problem

Objectives

- Backtracking
- MRV and degree heuristic
- Forward checking
- Improvement in forward check (through heuristic combination)
- MAC (maintaining arc consistency)
- Back jumping
- Intelligent backtracking (constraint learning)

Constraint Satisfaction Problem

• So far in Search space, states evaluated by heuristics and goal.

 A constraint satisfaction problem (CSP) is a problem that requires its solution within some limitations/conditions also known as constraints.

For which the typical problems can be converted to CSP and solved.

Converting problems to CSPs

- A problem to be converted to CSP requires the following steps:
- **Step 1:** Create a variable set.(Representing state)
- **Step 2:** Create a domain set.(values)
- **Step 3:** Create a constraint set with variables and domains after considering the constraints.(limitations)

CSP Example1-Sudoku

Rule: Fill all empty Squares so that the numbers 1 to 9 appear once in each row, column and 3x3 box

	1	2	3	4	5	6	7	8	9
A			3		2		6		
В	9			3		5			1
C			1	8		6	4		
D			8	1		2	9		
E	7								8
F			6	7		8	2		
G			2	6		9	5		
Н	8			2		3			9
1			5		1		3		

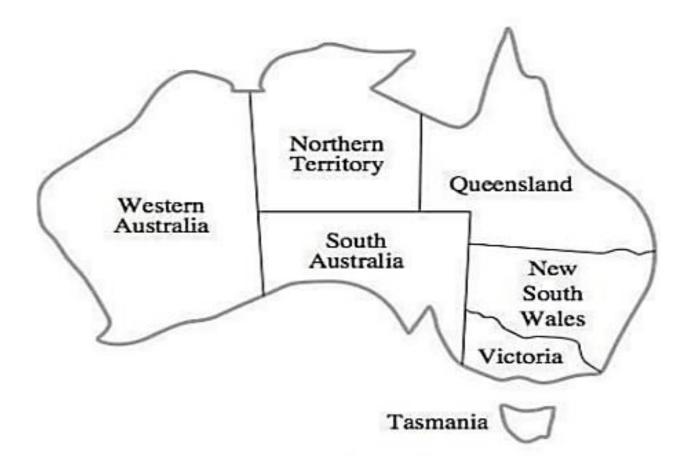
VARIABLES – variables are cells (A1-I9)

DOMAIN – domain of each variable is

{1,2,3,4,5,6,7,8,9}

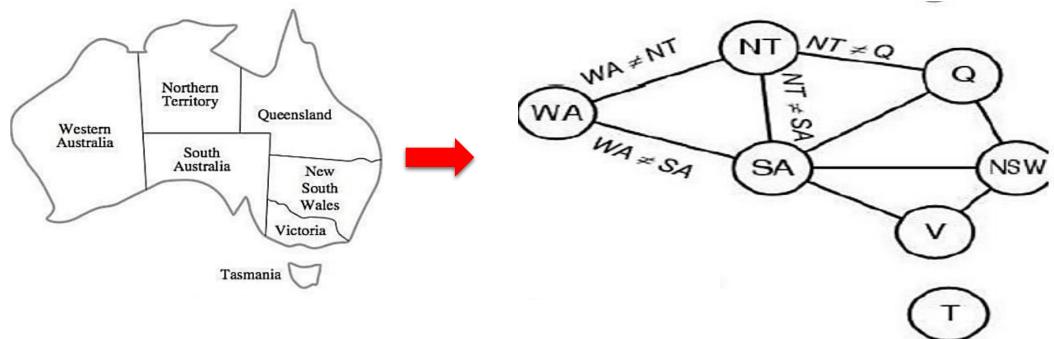
CONSTRAINT – constraints: rows,
columns, boxes contain all different
numbers

Example2-Map Coloring



- Rule Color each region either red, green or blue
 - No adjacent region can have the same color

CSP Representation of Map Coloring



Variables: WA, NT, Q, NSW, V, SA, T

Domains: Di = {red, green, blue}

Constraints: adjacent regions must have

different colors

e.g., WA ≠ NT, or (WA,NT)

Constraint Graph

- Nodes are Variables
- Arcs are Constraint

Varieties of constraints

- Unary constraints involve a single variable, e.g. SA ≠ green
- Binary constraints involve pairs of variables (relates two variables),e.g., SA ≠ WA
- Higher-order constraints involve 3 or more variables, e.g. cryptarithmetic puzzles (column constraints)

Example: Cryptarithmetic

Each letter stands for a distinct digit; the aim is to find a substitution of digits for letters such that the resulting sum is arithmetically correct, with the added restriction that no leading zeroes are allowed



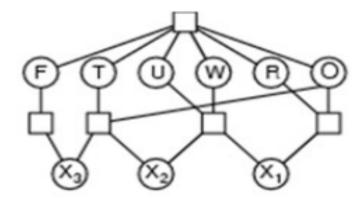
- Variables: FTUWROX₁X₂X₃
- Domains: {0,1,2,3,4,5,6,7,8,9}
- Constraints: Alldiff (F,T,U,W,R,O)
- where XI, X2, and X3 are auxiliary variables representing the digit (0 or 1) carried over into the next column

$$-O+O=R+10\cdot X_1$$

$$- X_1 + W + W = U + 10 \cdot X_2$$

$$- X_2 + T + T = O + 10 \cdot X_3$$

$$-X_3 = F, T \neq 0, F \neq 0$$

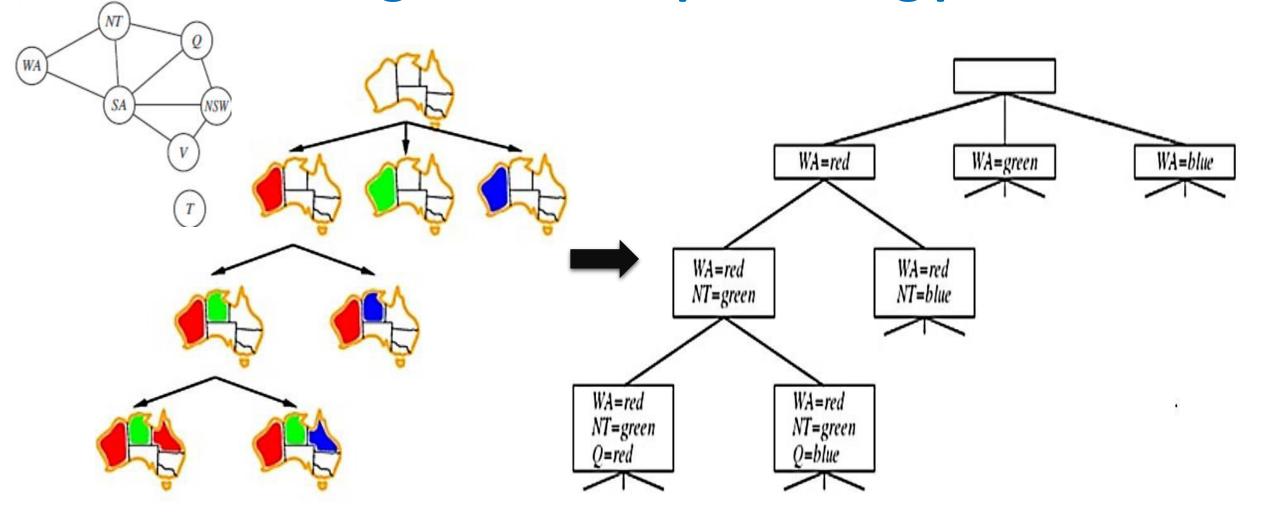


The constraint hyper graph for the cryptarithmetic problem, showing the Alldzff constraint as well as the column addition constraints

Backtracking search

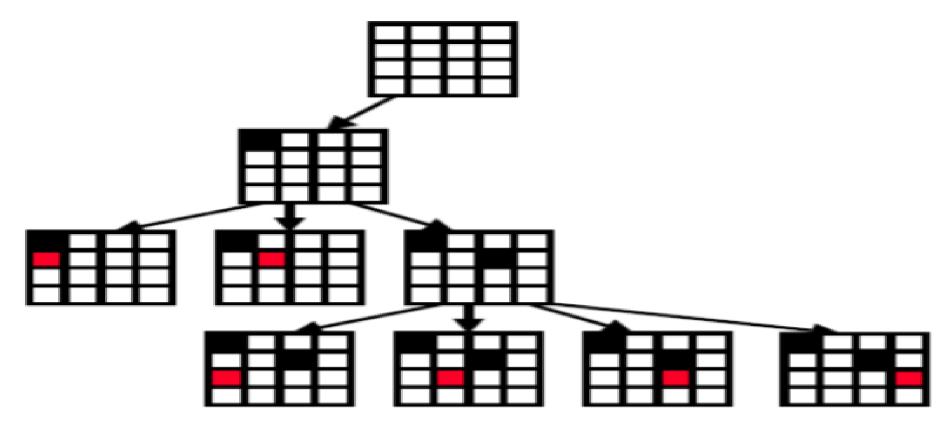
- Depth-first search for CSPs with single-variable assignments is called backtracking search
- Variable assignments are commutative: [WA = red then NT = green]
 same as [NT = green then WA = red]
- Only need to consider assignments to a single variable at each node
- Backtracks when variable has no legal value

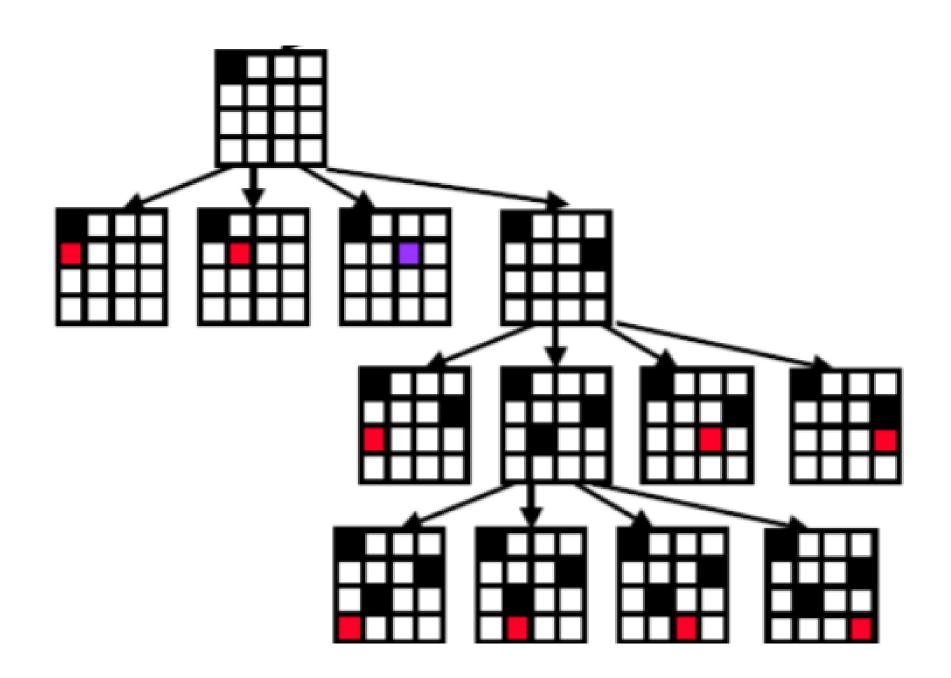
Backtracking for the map-coloring problem



Backtracking Example:2

•4X4 Queens





Solution!

•4X4 Queen

Improving backtracking efficiency (Variable and Value Ordering)

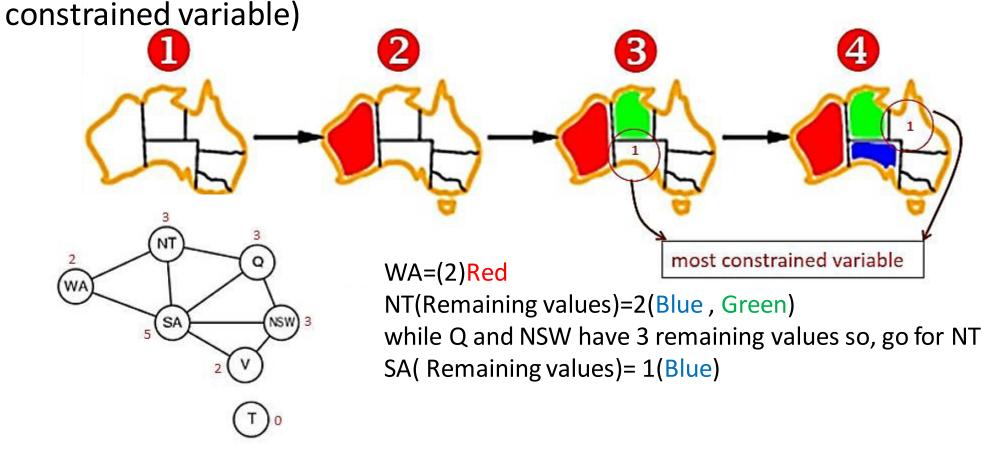
By default, Backtrack algorithm simply selects the next unassigned variable in the order given by the list VARIABLES[csp].

General-purpose methods can give huge gains in speed:

- Which variable should be assigned next?
- In what order should its values be tried?
- Can we detect inevitable failure early?

Minimum remaining values (MRV)

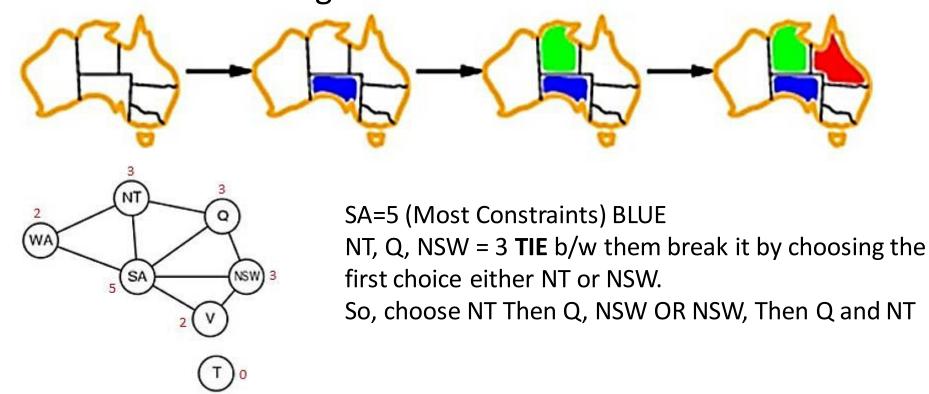
• This intuitive idea-choosing the variable with the fewest "legal" values is called the minimum remaining values (MRV) heuristic (also known most



after the assignments for WA = red and NT = green, there is only one possible value for SA, so it makes sense to assign **SA** = blue next rather than assigning Q, after SA is assigned, the choices for Q, NS W, and V are all forced.

Most-Constraining Variable (Degree Heuristics)

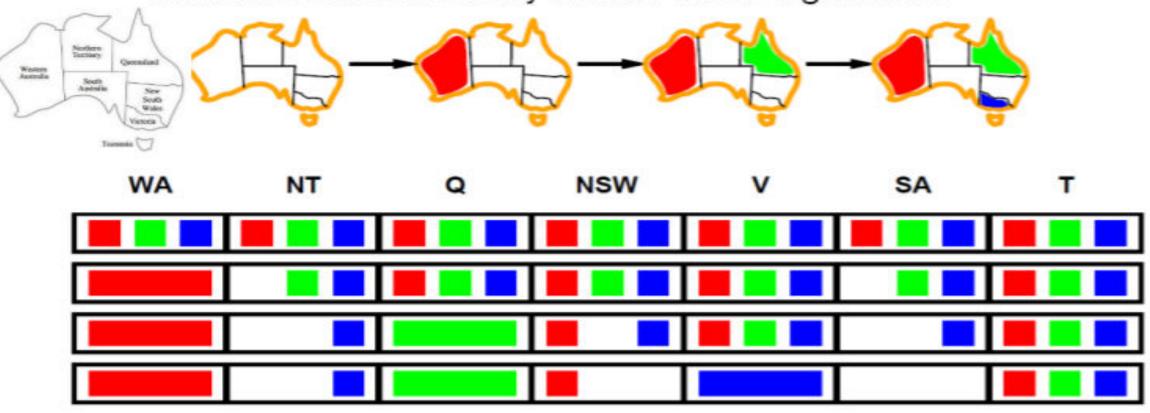
 Most constraining variable: choose the variable with the most constraints on remaining variables



Forward Checking

Idea: Keep track of remaining legal values for unassigned variables

Terminate search when any variable has no legal values



No possible assignments for SA, we try other assignments

Forward Checking Example:2

4X4 Queens

Q1,Q2,Q3,Q4 with domain {1..4}

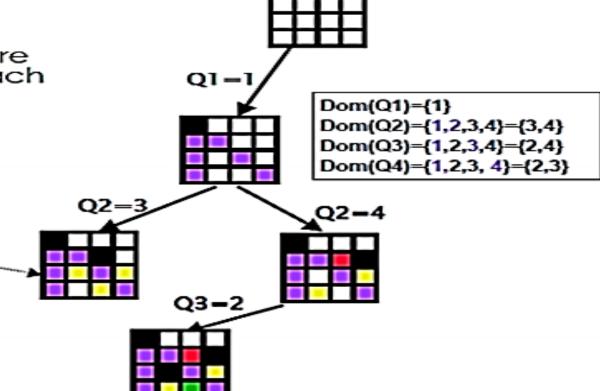
All binary constraints: C(Qi,Qj)

 FC illustration: color values are removed from domain of each row (blue, then yellow, then green)

DWO happens for Q3

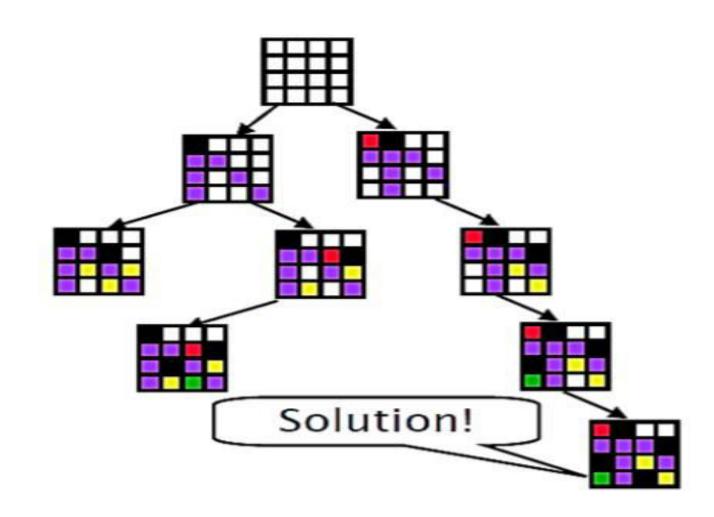
vlaue for Q2

So backtrack, try another



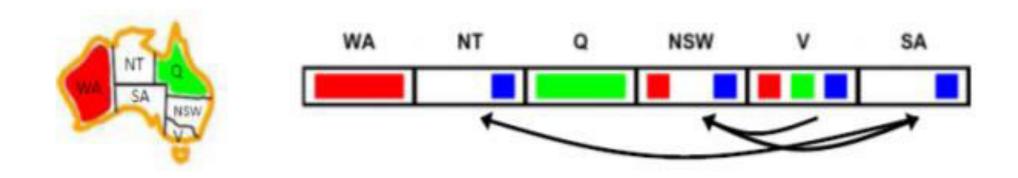
Forward Checking Example:2.....

4X4 Queens
 continue...



ARC Consistency

A simple form of propagation makes sure all arcs are simultaneously consistent:



- Arc consistency detects failure earlier than forward checking
- Important: If X loses a value, neighbors of X need to be rechecked!
- Must rerun after each assignment!

Remember: Delete from the tail!

Intelligent Backtracking: Looking Backward

- Consider a fixed variable ordering Q, NSW, V, T, SA, WA, NT.
 Suppose we have generated the partial assignment
 {Q=red, NSW = green, V = blue, T = red}.
 When we try the next variable,
- SA, we see that every value violates a constraint(chronological backtracking)
- Recoloring Tasmania cannot possibly resolve the problem with South Australia.
- A more intelligent approach to backtracking is to backtrack to a variable that might fix the problem
- The set (in this case {Q=red,NSW =green, V =blue, }), is called the conflict set for SA.
- In this case, back jumping would jump over Tasmania and try a new value for <u>V</u>.