CS317
Information Retrieval
Week 03

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Dictionaries & Tolerant Retrieval

Review Chapter No. 2

- We developed idea of inverted indexes for handling Boolean and proximity queries.
- We discussed positional indexes for supporting general phrase queries.
- We modify intersection of posting list to speedup in generating result-set.
- What about the dictionary? Large Dictionary are still a challenge?

Chapter No. 3

- In this chapter we will develop techniques that are robust to typographical errors in the query, as well as alternative spellings.
- We also develop data structures that help search for terms in the vocabulary in an inverted index.
- We explore the idea of a wildcard query: a query such as *a*e*i*o*u*, which seeks documents containing any term that includes all the five vowels in sequence.

Data Structures for Dictionary

- There are two choices
 - □ Trees
 - Hashtable
- IR systems can use either of the approach.

Sec. 3.1

Hashtables

- Each vocabulary term is hashed to an integer
 - □ (We assume you've seen hashtables before)
- Pros:
 - □ Lookup is faster than for a tree: O(1)
- Cons:
 - No easy way to find minor variants:
 - judgment/judgement
 - □ No prefix search [tolerant retrieval]
 - If vocabulary keeps growing, need to occasionally do the expensive operation of rehashing everything

Sec. 3.1

Tree: B-tree

a-hu
hy-m

□ Definition: Every internal nodel has a number of children in the interval [a,b] where a, b are appropriate natural numbers, e.g., [2,4].

Sec. 3.1

Trees

- Simplest: binary tree
- More usual: B-trees
- Trees require a standard ordering of characters and hence strings ... but we typically have one
- Pros:
 - □ Solves the prefix problem (terms starting with *hyp*)
- Cons:
 - □ Slower: O(log *M*) [and this requires *balanced* tree]
 - □ Rebalancing binary trees is expensive
 - But B-trees mitigate the rebalancing problem

| Wild Card Queries (*)

- Wildcard queries are used in any of the following situations:
 - the user is uncertain of the spelling of a query term (e.g., Sydney vs. Sidney, which leads to the wildcard query S*dney);
 - the user is aware of multiple variants of spelling a term and (consciously) seeks documents containing any of the variants (e.g., color vs. colour);

| Wild Card Queries (*)

- Wildcard queries are used in any of the following situations:
 - the user seeks documents containing variants of a term that would be caught by stemming, but is unsure whether the search engine performs stemming (e.g., judicial vs. judiciary, leading to the wildcard query judicia*);
 - the user is uncertain of the correct rendition of a foreign word or phrase (e.g., the query Universit* Stuttgart).

Sec. 3.2

| Wild-card queries: *

- mon*: find all docs containing any word beginning with "mon".
- Easy with binary tree (or B-tree) lexicon: retrieve all words in range: mon ≤ w < moo</p>
- *mon: find words ending in "mon": harder
 - Maintain an additional B-tree for terms backwards.

Can retrieve all words in range: *nom* ≤ *w* < *non*.

Exercise: from this, how can we enumerate all terms meeting the wild-card query **pro*cent**?

Sec. 3.2

Query processing

- At this point, we have an enumeration of all terms in the dictionary that match the wild-card query.
- We still have to look up the postings for each enumerated term.
- E.g., consider the query:

se*ate AND fil*er

This may result in the execution of many Boolean *AND* queries.

Sec. 3.2

B-trees handle *'s at the end of a query term

- How can we handle *'s in the middle of query term?
 - □ co*tion
- We could look up **co*** AND ***tion** in a B-tree and intersect the two term sets
 - Expensive
- The solution: transform wild-card queries so that the *'s occur at the end
- This gives rise to the **Permuterm** Index.

| General Wild Card Query

- We now study two techniques for handling general wildcard queries.
 - Both techniques share a common strategy: express the given wildcard query q_w as a Boolean query Q on a specially constructed index, such that the answer to Q is a superset of the set of vocabulary terms matching q_w.
 - Then, we check each term in the answer to Q against q_w, discarding those vocabulary terms that do not match q_w. At this point we have the vocabulary terms matching q_w and can resort to the standard inverted index.

Sec. 3.2.1

Permuterm index

- For term *hello*, index under:
 - □ *hello\$, ello\$h, llo\$he, lo\$hel, o\$hell* where \$ is a special symbol.
- Queries:
 - X lookup on X\$ X* lookup on \$X*
 - □ *X lookup on X\$* *X* lookup on X*
 - □ X*Y lookup on Y\$X* X*Y*Z ??? Exercise!

Query = hel*o X=hel, Y=o Lookup o\$hel*

Sec. 3.2.1

| Permuterm query processing

- Rotate query wild-card to the right
- Now use B-tree lookup as before.

Sec. 3.2.2

Sec. 3.2.2

| Bigram (k-gram) indexes

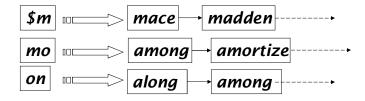
- Enumerate all *k*-grams (sequence of *k* chars) occurring in any term
- e.g., from text "April is the cruelest month" we get the 2-grams (bigrams)

```
$a,ap,pr,ri,il,l$,$i,is,s$,$t,th,he,e$,$c,cr,ru,
ue,el,le,es,st,t$, $m,mo,on,nt,h$
```

- □ \$ is a special word boundary symbol
- Maintain a <u>second</u> inverted index <u>from</u> <u>bigrams to</u> <u>dictionary terms</u> that match each bigram.

Bigram index example

■ The *k*-gram index finds *terms* based on a query consisting of *k*-grams (here *k*=2).



Sec. 3.2.2

Processing wild-cards

- Query *mon** can now be run as
 - □ \$m AND mo AND on
- Gets terms that match AND version of our wildcard query.
- But we'd enumerate *moon*.
- Must post-filter these terms against query.
- Surviving enumerated terms are then looked up in the term-document inverted index.
- Fast, space efficient (compared to permuterm).

Sec. 3.2.2

| Processing wild-card queries

- As before, we must execute a Boolean query for each enumerated, filtered term.
- Wild-cards can result in expensive query execution (very large disjunctions...)
 - pyth* AND prog*
- If you encourage "laziness" people will respond!

Search

Type your search terms, use '*' if you need to. E.g., Alex* will match Alexander.