

# **NATIONAL UNIVERSITY OF COMPUTER AND EMERGING SCIENCES**

## **CS 201 – DATA STRUCTURES LAB**

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### **LAB SESSION # 06**

#### **Outline**

- Deletion in BST
- Searching in AVL
- Insertion in AVL
- Exercise

# DELETION IN BST

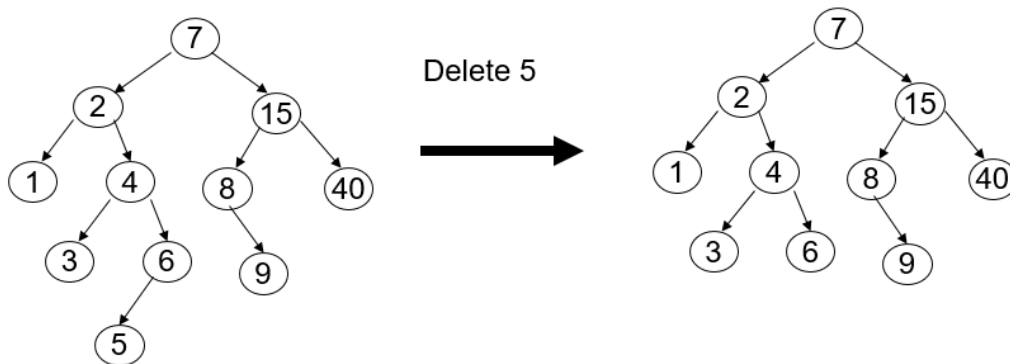
There are three cases:

1. The node to be deleted is a leaf node.
2. The node to be deleted has one non-empty child.
3. The node to be deleted has two non-empty children.

## CASE 1: DELETING A LEAF NODE

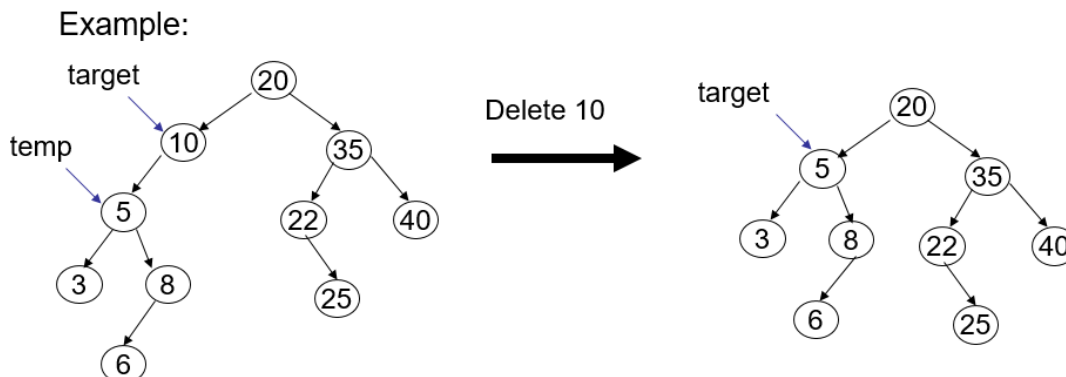
Convert the leaf node into an empty tree by using the detachKey method:

Example: Delete 5 in the tree below:

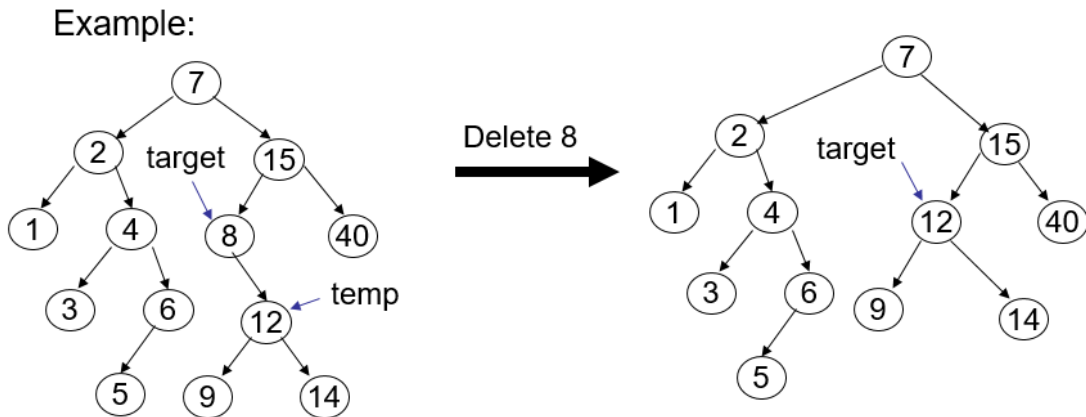


## CASE 2: THE NODE TO BE DELETED HAS ONE NON-EMPTY CHILD

- a. The right subtree of the node x to be deleted is empty.



b. The left subtree of the node x to be deleted is empty.



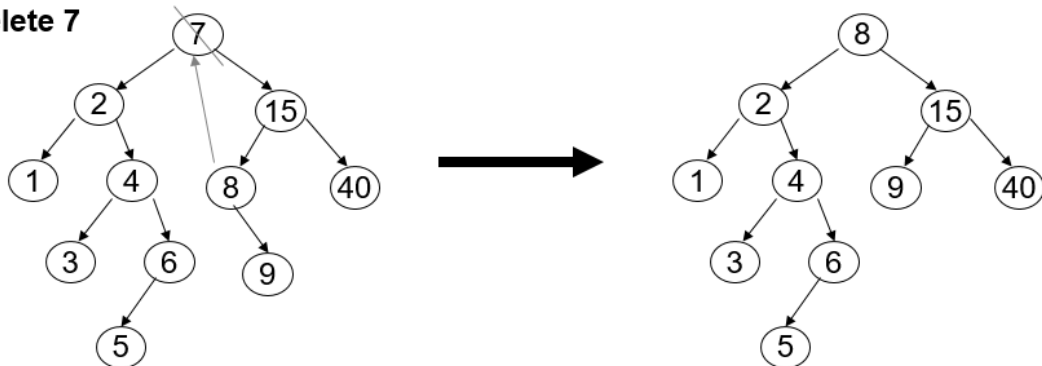
### CASE 3: DELETING A NODE THAT HAS TWO CHILDREN

#### a. Method # 1: DELETION BY COPYING

Copy the minimum key in the right subtree of x to the node x, then delete the one-child or leaf-node with this minimum key.

Example:

Delete 7



#### b. Method # 2: DELETION BY COPYING:

Copy the maximum key in the left subtree of x to the node x, then delete the one-child or leaf-node with this maximum key.

Example:

Delete 7



Deletion by Copying Code

```
// find the minimum key in the right subtree of the target node
Comparable min = target.getRightBST().findMin();

// copy the minimum value to the target
target.key = min;

// delete the one-child or leaf node having the min
target.getRightBST().withdraw(min);
```

All the different cases for deleting a node are handled in the **withdraw (Comparable key)** method of **BinarySearchTree** class

## AVL TREE

An AVL tree is a binary search tree with a height balance property:  
For each node  $v$ , the heights of the subtrees of  $v$  differ by at most 1.

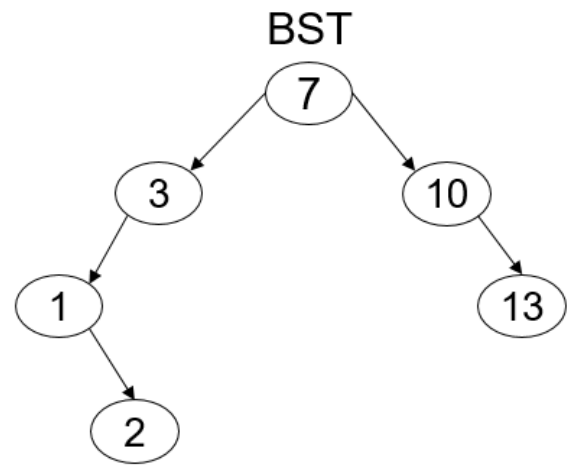
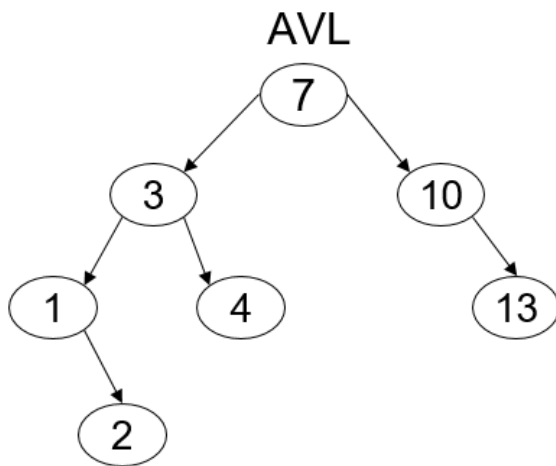
A subtree of an AVL tree is also an AVL tree.

For each node of an AVL tree can have a balance factor: Balance factor = height (right subtree) - height (left subtree)

An AVL tree is rebalanced after each insertion or deletion.

The height-balance property ensures that the height of an AVL tree with  $n$  nodes is  $O(\log n)$ .

Searching, insertion, and deletion are all  $O(\log n)$ .



## AVL TREE IMPLEMENTATION:

```
public class AVLTree extends BinarySearchTree{
    protected int height;
    public AVLTree(){ height = -1;}

    public int getHeight(){ return height } ;

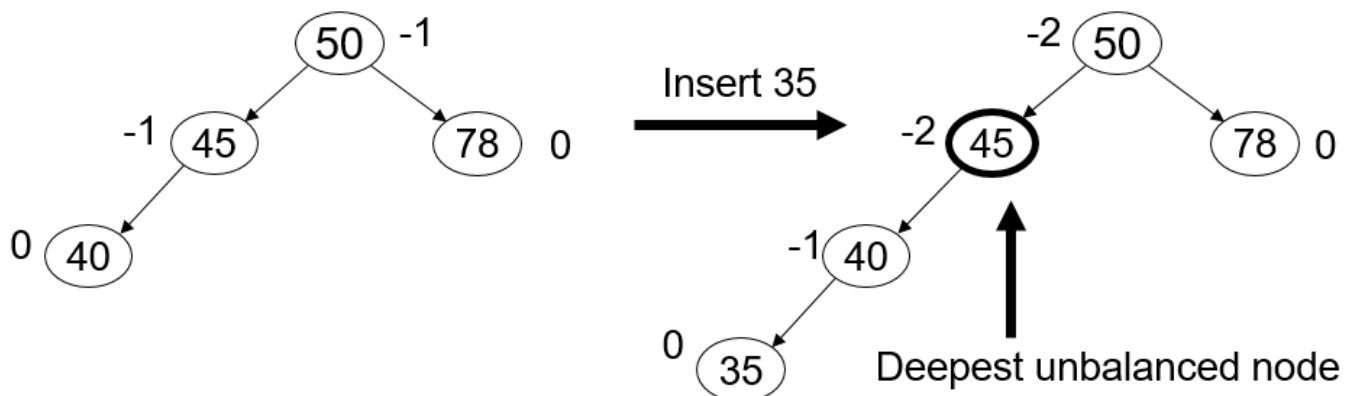
    protected void adjustHeight(){
        if( isEmpty())
            height = -1;
        else
            height = 1 + Math.max( left.getHeight() , right.getHeight());
    }

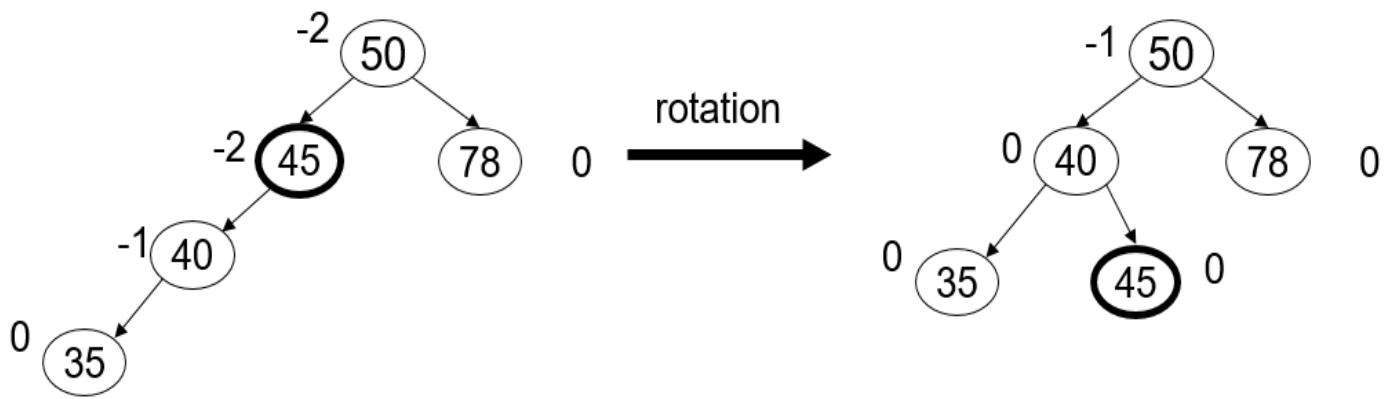
    protected int getBalanceFactor(){
        if( isEmpty())
            return 0;
        else
            return right.getHeight() - left.getHeight();
    }
    // . . .
}
```

## AVL TREE ROTATION:

- A rotation is a process of switching children and parents among two or three adjacent nodes to restore balance to a tree.
- **An insertion or deletion may cause an imbalance in an AVL tree.**

The deepest node, which is an ancestor of a deleted or an inserted node, and whose balance factor has changed to -2 or +2 requires rotation to rebalance the tree.





### Tree Imbalance

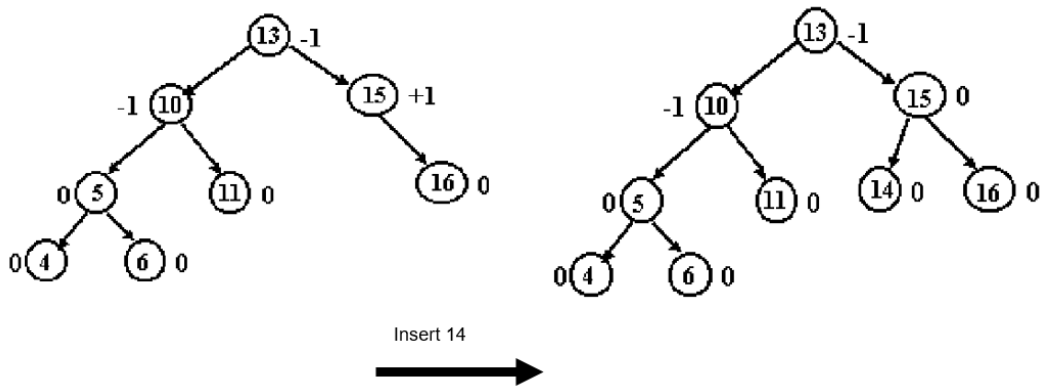
- Left Left Imbalance (Single Rotation)
- Right Right Imbalance (Single Rotation)
- Left-Right Imbalance (Double Rotation)
- Right-Left Imbalance (Double Rotation)

### AVL INSERTION:

- Insert using a BST insertion algorithm.
- Rebalance the tree if an imbalance occurs.
- An imbalance occurs if a node's balance factor changes from -1 to -2 or from +1 to +2.
- Rebalancing is done at the deepest or lowest unbalanced ancestor of the inserted node.
- There are three insertion cases:
  1. Insertion that does not cause an imbalance.
  2. Same side (left-left or right-right) insertion that causes an imbalance.
    - Requires a single rotation to rebalance.
  3. Opposite side (left-right or right-left) insertion that causes an imbalance.
    - Requires a double rotation to rebalance.

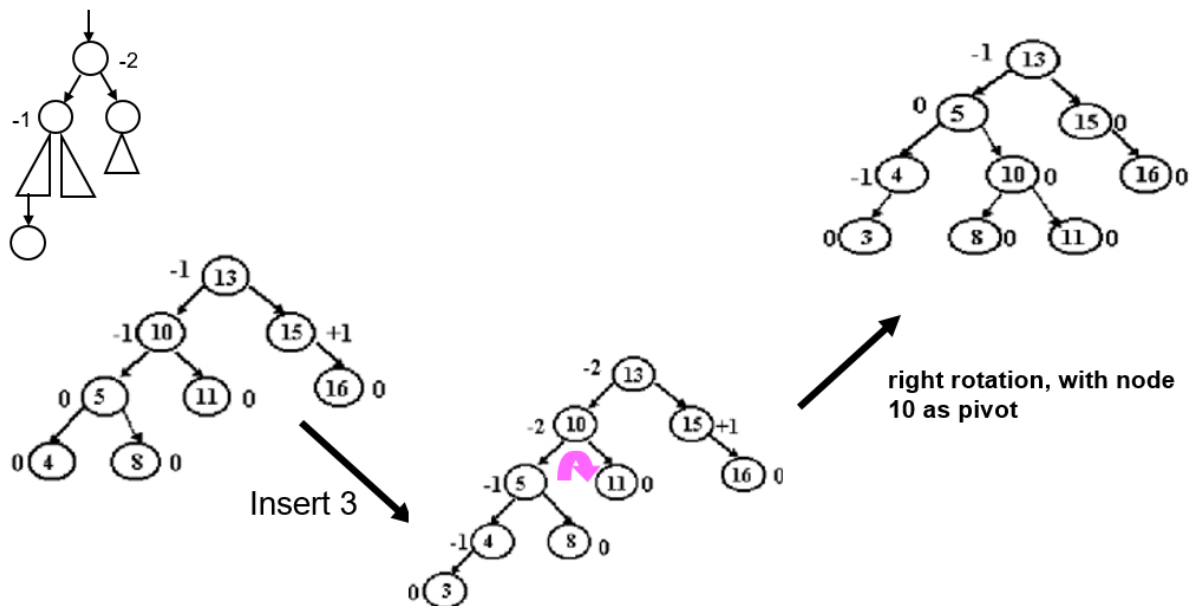
## Insertion Case 1:

- Example: An insertion that does not cause an imbalance.

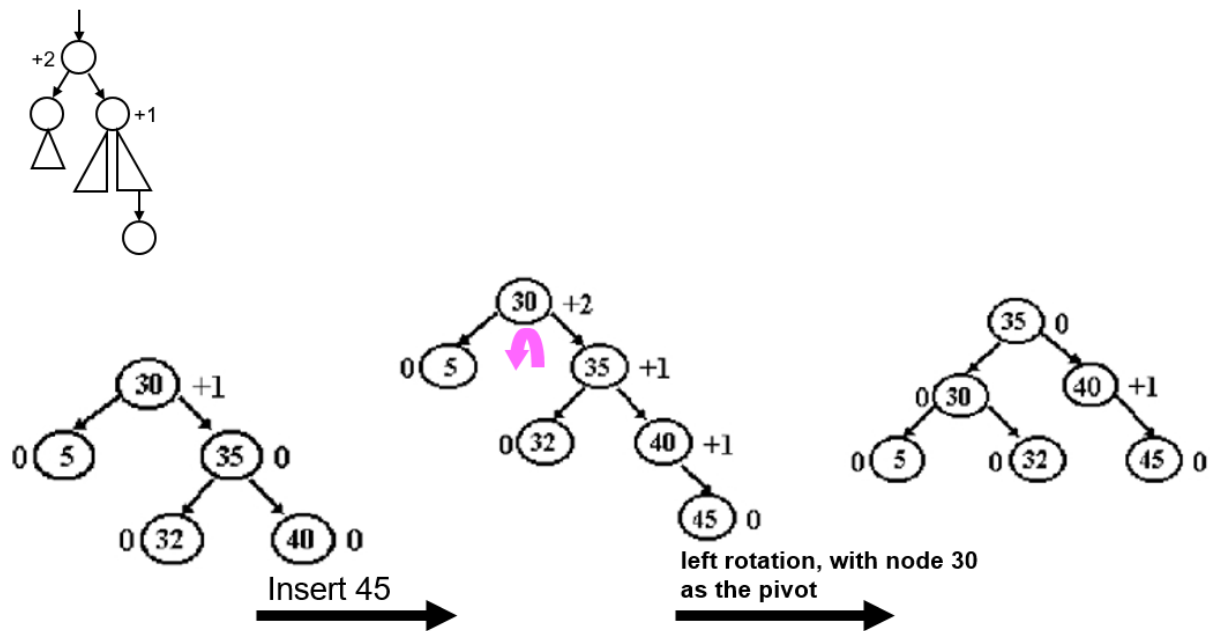


## Insertion Case 2:

- Case 2a: The lowest node (with a balance factor of -2) had a taller left-subtree and the insertion was on the left-subtree of its left child.
- Requires single right rotation to rebalance.

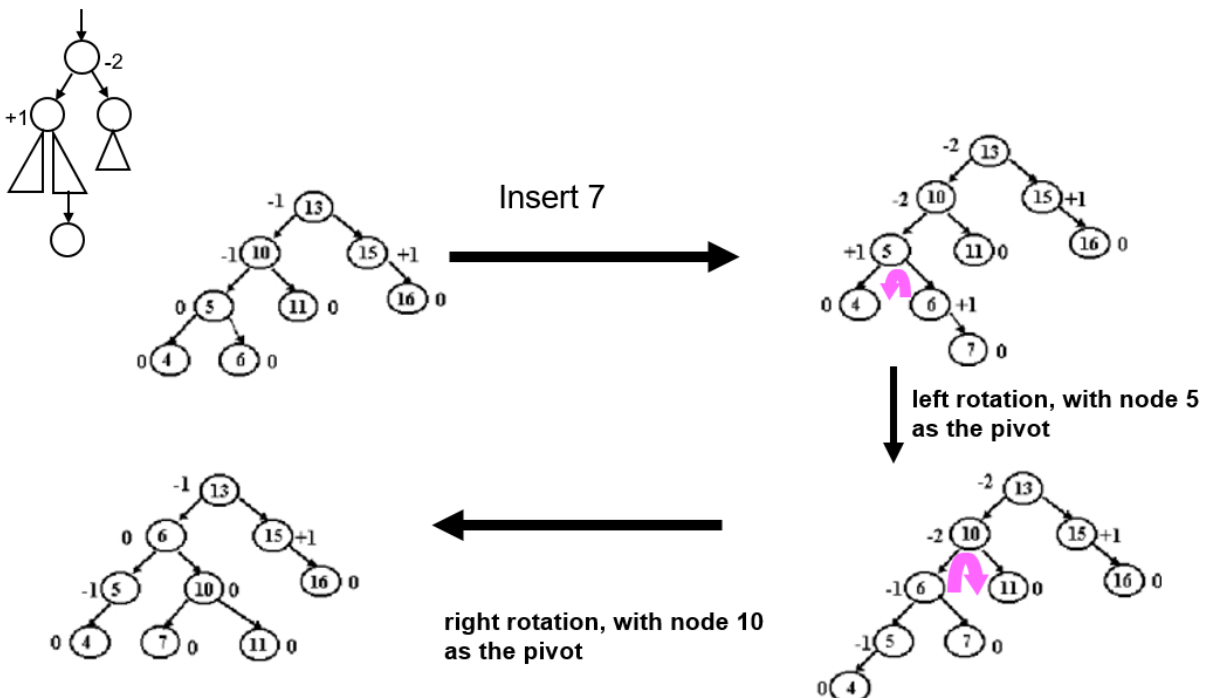


- Case 2b: The lowest node (with a balance factor of +2) had a taller right-subtree and the insertion was on the right-subtree of its right child.
- Requires single left rotation to rebalance.



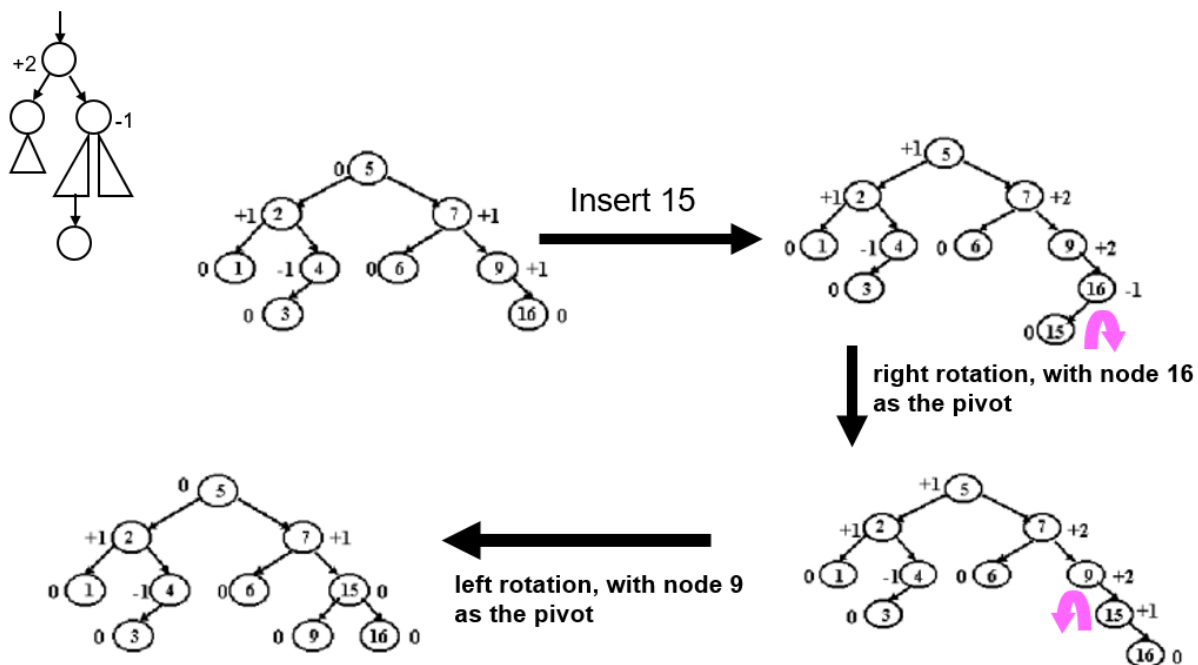
### Insertion Case 3:

- Case 3a: The lowest node (with a balance factor of -2) had a taller left-subtree and the insertion was on the right-subtree of its left child.
- Requires a double left-right rotation to rebalance.





- Case 3b: The lowest node (with a balance factor of +2) had a taller right-subtree and the insertion was on the left-subtree of its right child.
- Requires a double right-left rotation to rebalance.



### Exercise:

#### Question No. 1:

Write a Java program to perform the following operations

- Build BST Tree
- Deletion in BST Tree (All Cases)

#### Question No. 2:

Write a Java program to perform the following operations

- Build AVL Tree
- Write code for AVL Rotations (Single and Double)
- Insertion in AVL Tree (All Cases)

#### Question No. 3:

Write a Java program to perform the following operations

- Build AVL Tree
- Insert Values
- Search for a value in AVL Tree using DFS and BFS
- Calculate time in Second for all operations.

#### Question No. 4:

Write a Java program that uses recursive functions.

- To create a binary search tree.
- Write a Function to find whether the tree is AVL or not
- Write a Function to calculate height of a Tree