

Sumaiyah Zahid

### ASYNCHRONOUS COUNTERS

In asynchronous counter flip-flops (FF) do not change states at exactly the same time because they do not have a common clock pulse.

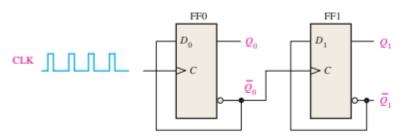
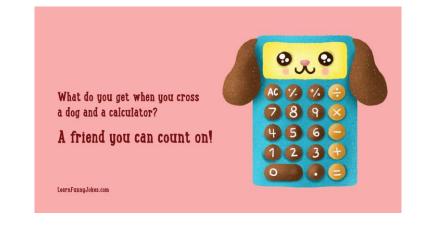


FIGURE 9-4 A 2-bit asynchronous binary counter. Open file F09-04 to verify operation. .

Multisim tutorial is available on the website.

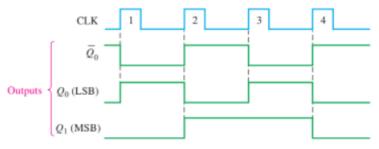


### ASYNCHRONOUS COUNTERS

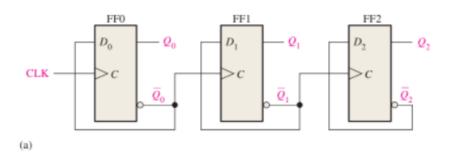
#### TABLE 9-1

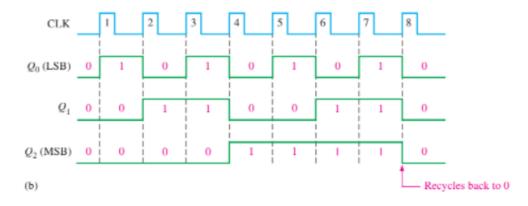
Binary state sequence for the counter in Figure 9-4.

Clock Pulse	$Q_1$	$Q_0$
Initially	0	0
1	0	1
2	1	0
3	1	1
4 (recycles)	0	0



**FIGURE 9–5** Timing diagram for the counter of Figure 9–4. As in previous chapters, output waveforms are shown in green.





## PROPAGATION DELAY

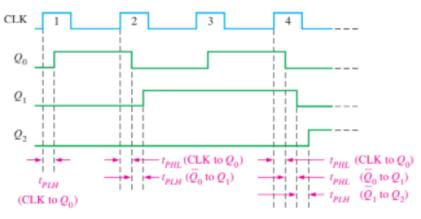
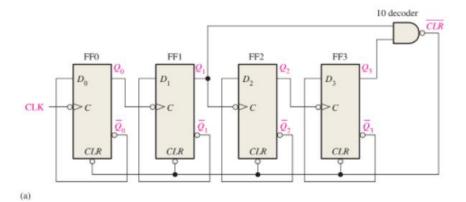
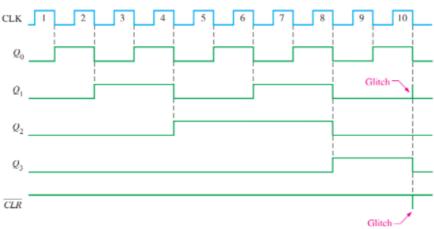


FIGURE 9-7 Propagation delays in a 3-bit asynchronous (ripple-clocked) binary counter.

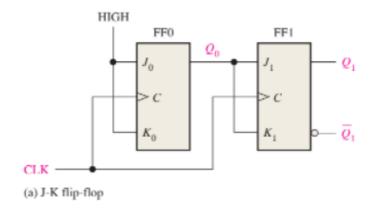
# ASYNCHRONOUS DECADE COUNTERS





### SYNCHRONOUS COUNTERS

In synchronous counter all the flip-flops in the counter are clocked at the same time by a common clock pulse.





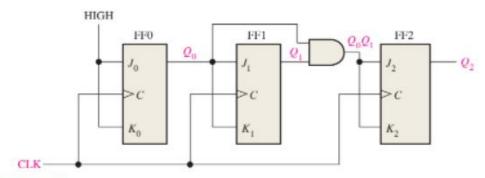


FIGURE 9-15 A 3-bit synchronous binary counter. Open file F09-15 to verify the operation.

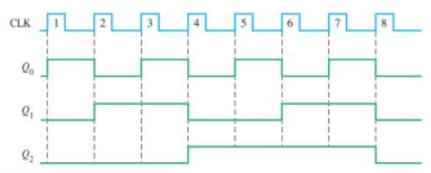
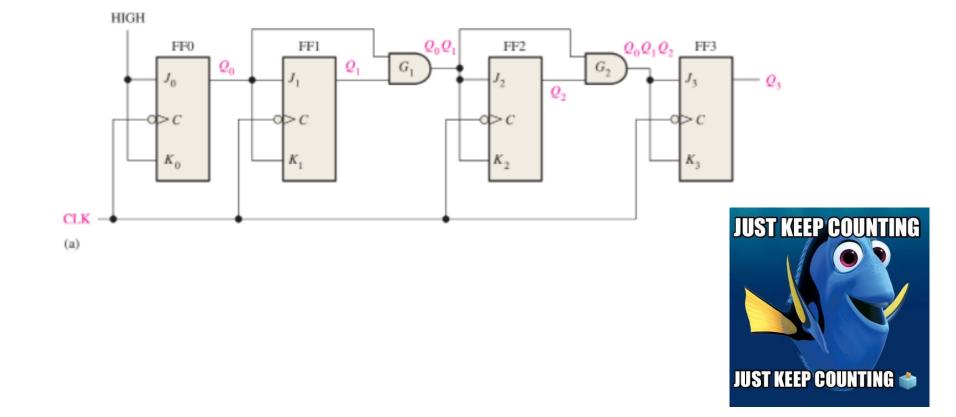
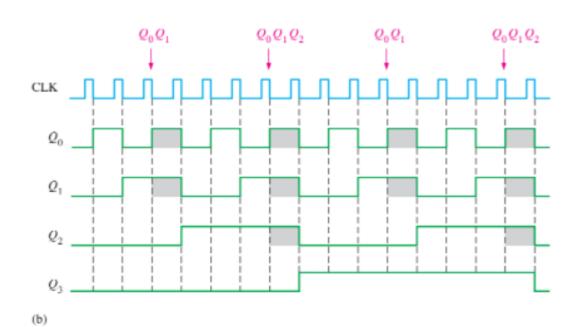


FIGURE 9-16 Timing diagram for the counter of Figure 9-15.





# 4-BIT SYNCHRONOUS DECADE COUNTER

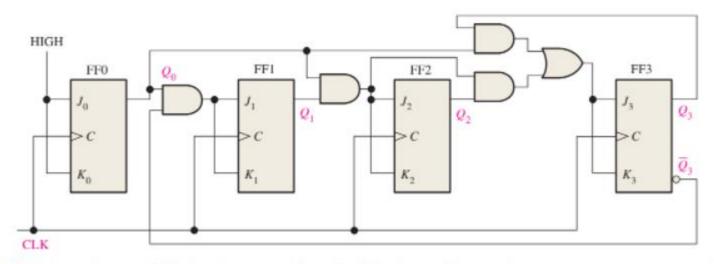
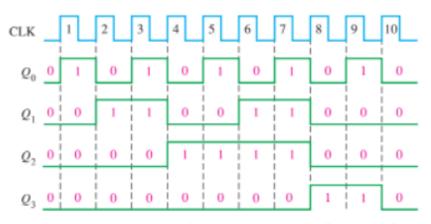


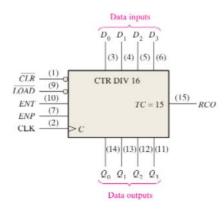
FIGURE 9–18 A synchronous BCD decade counter. Open file F09-18 to verify operation.

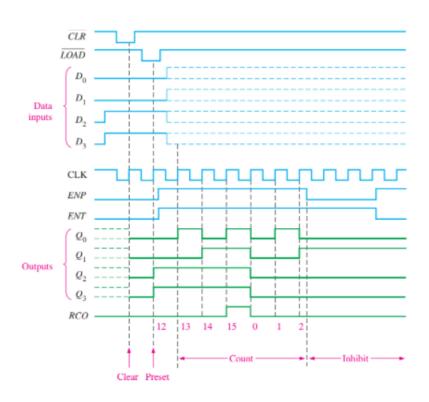


### 4-BIT SYNCHRONOUS DECADE COUNTER

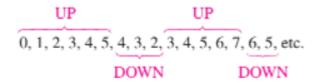


**FIGURE 9–19** Timing diagram for the BCD decade counter ( $Q_0$  is the LSB).

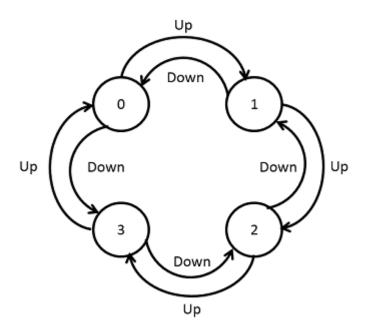




Also called a bidirectional counter.



Clock Pulse	Up	$Q_2$	$Q_1$	$Q_0$	Down
0	40	0	0	0	5 1
1	(	0	0	1	5
2	(	0	1	0	5
3	(	0	1	1	5
4	(	1	0	0	5
5	(	1	0	1	5
6	(	1	1	0	5
7	16	1	1	1	5 ₺



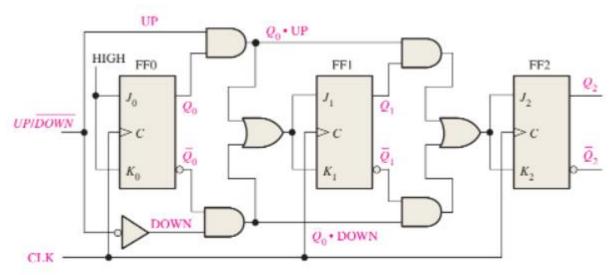


FIGURE 9-22 A basic 3-bit up/down synchronous counter. Open file F09-22 to verify operation.

#### **EXAMPLE 9-3**

Show the timing diagram and determine the sequence of a 4-bit synchronous binary up/down counter if the clock and  $UP/\overline{DOWN}$  control inputs have waveforms as shown in Figure 9–23(a). The counter starts in the all-0s state and is positive edge-triggered.

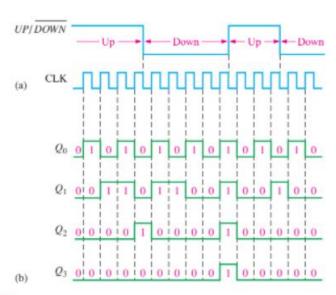
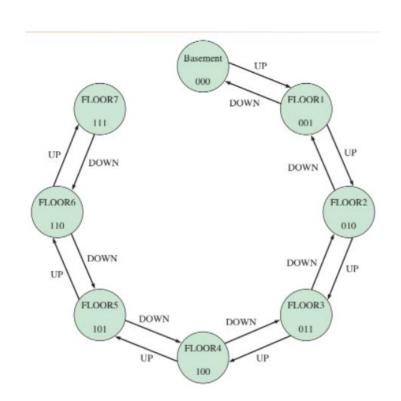


FIGURE 9-23

### **Applications**

Elevators



### Asynchronous Cascading

Total states = 4\*8=32

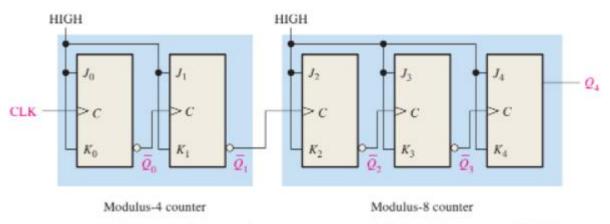
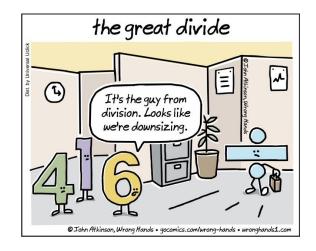


FIGURE 9-35 Two cascaded asynchronous counters (all J and K inputs are HIGH).

### Synchronous Cascading

Total counts = 10\*10=100



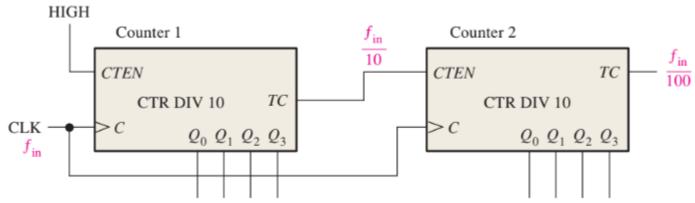
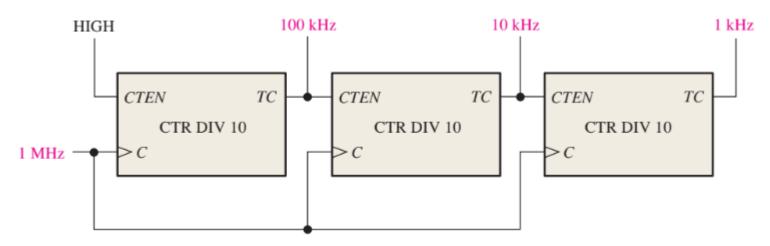


FIGURE 9-37 A modulus-100 counter using two cascaded decade counters.

### Synchronous Cascading

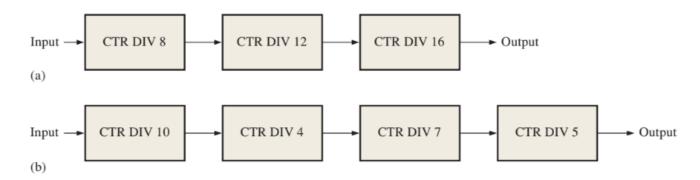
Frequency division by 10\*10\*10= 1000



**FIGURE 9–38** Three cascaded decade counters forming a divide-by-1000 frequency divider with intermediate divide-by-10 and divide-by-100 outputs.

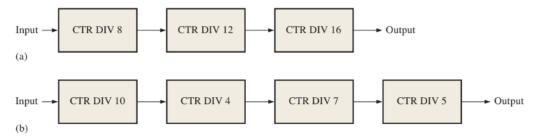
#### **EXAMPLE 9-6**

Determine the overall modulus of the two cascaded counter configurations in Figure 9–39.



#### **EXAMPLE 9-6**

Determine the overall modulus of the two cascaded counter configurations in Figure 9–39.



#### **FIGURE 9-39**

#### Solution

In Figure 9–39(a), the overall modulus for the 3-counter configuration is

$$8 \times 12 \times 16 = 1536$$

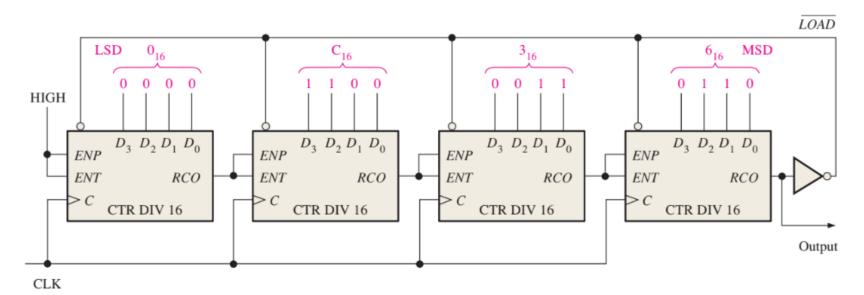
In Figure 9–39(b), the overall modulus for the 4-counter configuration is

$$10 \times 4 \times 7 \times 5 = 1400$$

#### **Related Problem**

How many cascaded decade counters are required to divide a clock frequency by 100,000?

### Cascaded Counters with Truncated Sequences



**FIGURE 9–41** A divide-by-40,000 counter using 74HC161 4-bit binary counters. Note that each of the parallel data inputs is shown in binary order (the right-most bit  $D_0$  is the LSB in each counter).

Cascaded Counters with Truncated Sequences

Also known as full-modulus cascading.

 $2^{16} = 65536$ 

65536 - 40000 = 25536 (63C0 in Hex)

25536 states should be deleted.

#### **SECTION 9-6 CHECKUP**

- 1. How many decade counters are necessary to implement a divide-by-1000 (modulus-1000) counter? A divide-by-10,000?
- 2. Show with general block diagrams how to achieve each of the following, using a flip-flop, a decade counter, and a 4-bit binary counter, or any combination of these:
  - (a) Divide-by-20 counter (b) Divide-by-32 counter
  - (c) Divide-by-160 counter (d) Divide-by-320 counter

Active HIGH Decoding

Active Low Decoding
Will use NAND

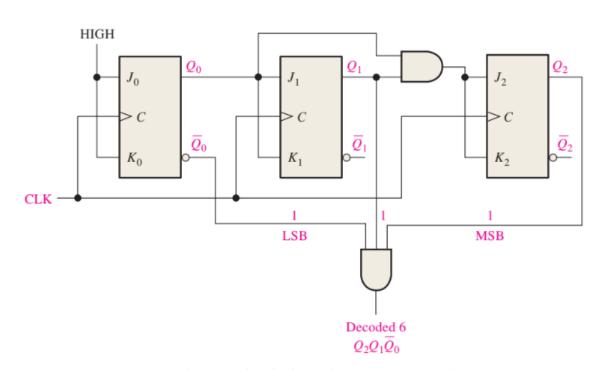
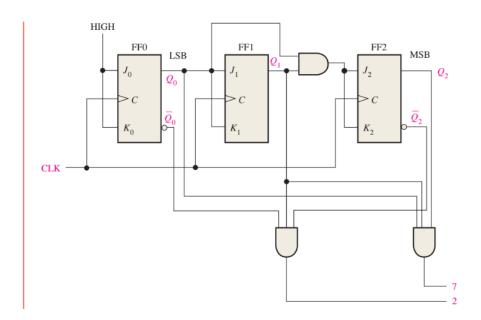


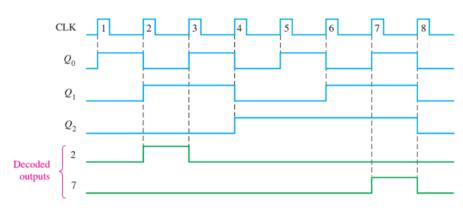
FIGURE 9-42 Decoding of state 6 (110). Open file F09-42 to verify operation.

#### **EXAMPLE 9-8**

Implement the decoding of binary state 2 and binary state 7 of a 3-bit synchronous counter. Show the entire counter timing diagram and the output waveforms of the decoding gates. Binary  $2 = \overline{Q}_2 Q_1 \overline{Q}_0$  and binary  $7 = Q_2 Q_1 Q_0$ .

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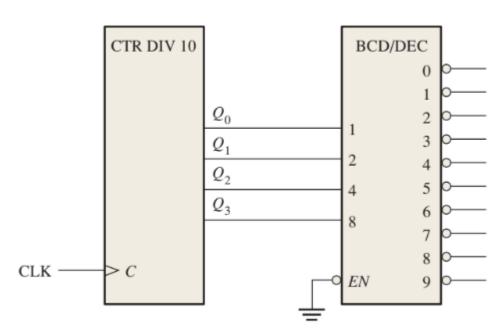
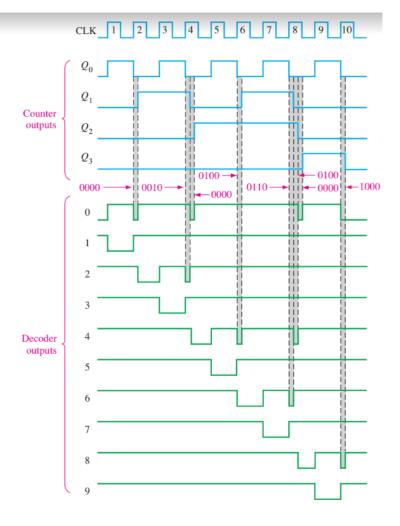


FIGURE 9-44 A basic decade (BCD) counter and decoder.



## COUNTER DECODING ~ STROBING

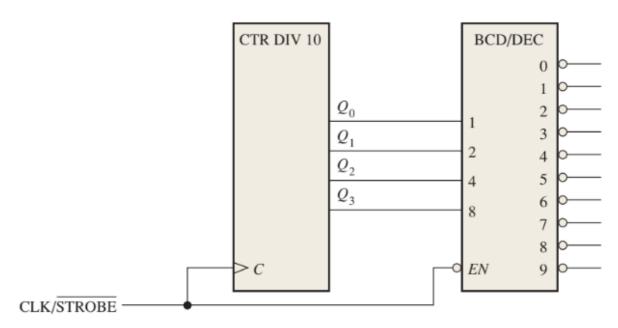


FIGURE 9-46 The basic decade counter and decoder with strobing to eliminate glitches.

# COUNTER DECODING ~ STROBING

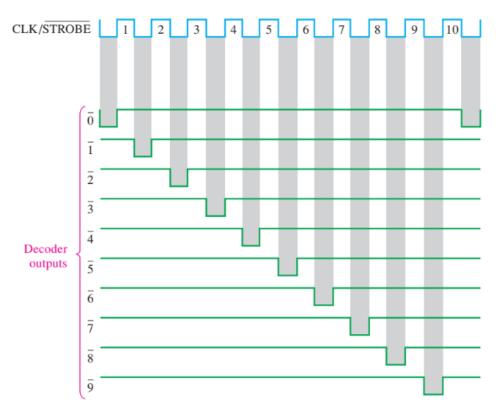
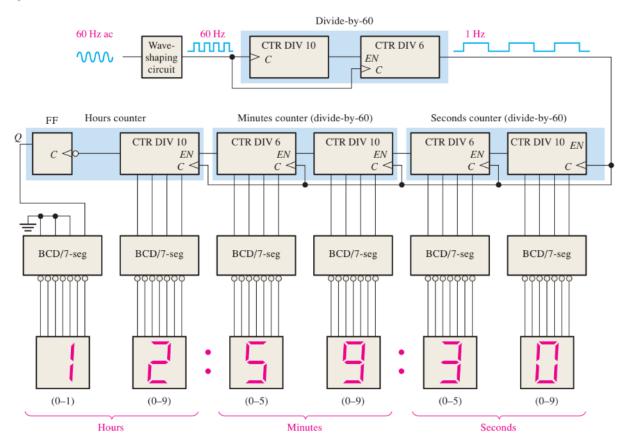


FIGURE 9-47 Strobed decoder outputs for the circuit of Figure 9-46.

### COUNTER APPLICATIONS

A Digital Clock



### FINITE STATE MACHINES

A state machine is a sequential circuit having a limited (finite) number of states occurring in a prescribed order.

- Moore
  - outputs depend only on the internal present state
- Mealy
  - outputs depend on both the internal present state and on the inputs

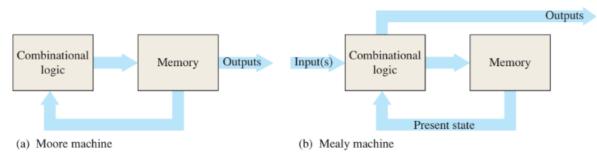
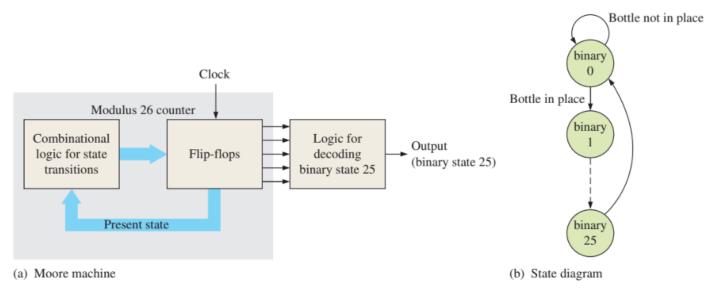


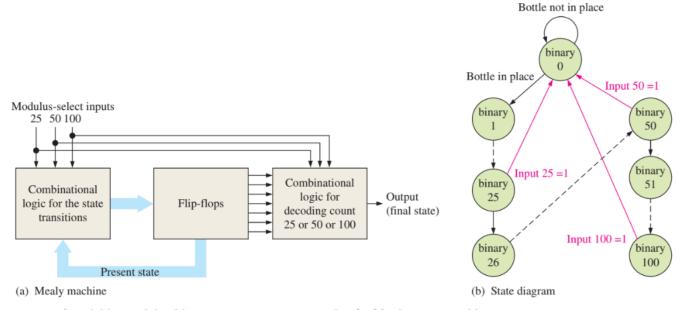
FIGURE 9-1 Two types of sequential logic.

# MOORE MACHINE



**FIGURE 9–2** A fixed-modulus binary counter as an example of a Moore state machine. The dashed line in the state diagram means the states between binary 1 and 25 are not shown for simplicity.

## MEALY MACHINE



**FIGURE 9–3** A variable-modulus binary counter as an example of a Mealy state machine. The red arrows in the state diagram represent the recycle paths that depend on the input number. The black dashed lines mean the interim states are not shown for simplicity.

## DESIGNING A STATE MACHINE

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Step 1: State Diagram
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Step 2: Next-State Table

Step 3: Flip-Flop Transition Table

Step 4: Karnaugh Maps

## STEP 1: STATE DIAGRAM

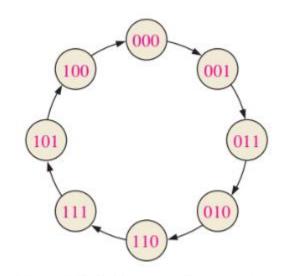


FIGURE 9-26 State diagram for a 3-bit Gray code counter.

## STEP 2: NEXT-STATE TABLE

#### TABLE 9-8

Next-state table for 3-bit Gray code counter.

	Present St	ate		<b>Next State</b>	
$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$
0	0	0	0	0	1
0	0	1	0	1	1
0	1	1	0	1	0
0	1	0	1	1	0
1	1	O	1	1	1
1	1	1	1	0	1
1	0	1	1	0	0
1	0	O	0	0	0

## STEP 3: FLIP-FLOP TRANSITION TABLE

#### **TABLE 9-9**

Transition table for a J-K flip-flop.

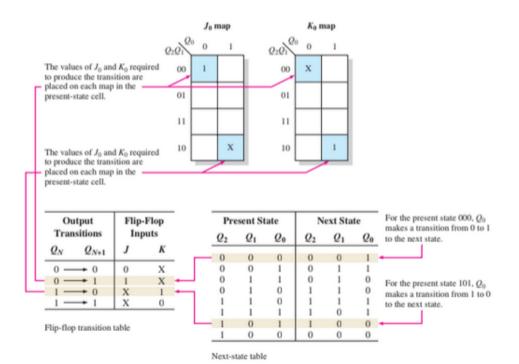
<b>Output Transitions</b>			Flip-Flop Inputs		
$Q_N$		$Q_{N+1}$	J	K	
0	<b>→</b>	0	0	X	
0	$\longrightarrow$	1	1	X	
1	$\longrightarrow$	0	X	1	
1	$\longrightarrow$	1	X	0	

 $Q_N$ : present state

 $Q_{N+1}$ : next state

X: "don't care"

## STEP 4: KARNAUGH MAPS



## STEP 4: KARNAUGH MAPS

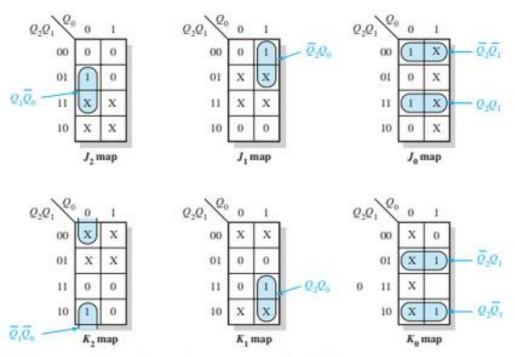


FIGURE 9-28 Karnaugh maps for present-state J and K inputs.

## STEP 5: LOGIC EXPRESSIONS FOR FLIP-FLOP INPUTS

$$J_0 = Q_2Q_1 + \overline{Q}_2\overline{Q}_1 = \overline{Q}_2 \oplus \overline{Q}_1$$
  
 $K_0 = Q_2\overline{Q}_1 + \overline{Q}_2Q_1 = Q_2 \oplus Q_1$   
 $J_1 = \overline{Q}_2Q_0$   
 $K_1 = Q_2Q_0$   
 $J_2 = Q_1\overline{Q}_0$   
 $K_2 = \overline{Q}_1\overline{Q}_0$ 

## STEP 6: COUNTER IMPLEMENTATION

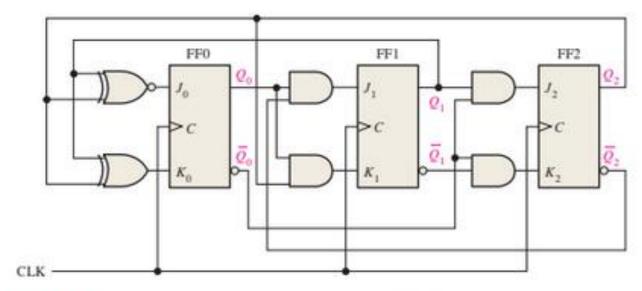


FIGURE 9-29 Three-bit Gray code counter. Open file F09-29 to verify operation.

## SYNCHRONOUS COUNTERS

#### **EXAMPLE 9-4**

Design a counter with the irregular binary count sequence shown in the state diagram of Figure 9-30. Use D flip-flops.

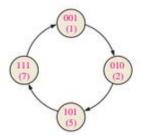


FIGURE 9-30

#### Solution

- Step 1: The state diagram is as shown. Although there are only four states, a 3-bit counter is required to implement this sequence because the maximum binary count is seven. Since the required sequence does not include all the possible binary states, the invalid states (0, 3, 4, and 6) can be treated as "don't cares" in the design. However, if the counter should erroneously get into an invalid state, you must make sure that it goes back to a valid state.
- Step 2: The next-state table is developed from the state diagram and is given in Table 9-10.

Next-state table.	TABLE 9-10		
	Next-state table.		

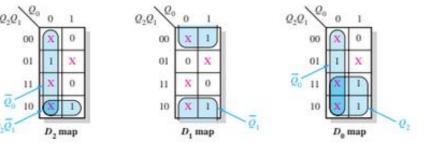
Present State				Next State			
$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$		
0	0	1	0	1	0		
0	1	0	1	0	1		
1	0	1	1	1	1		
1	1	1	0	0	1		

Step 3: The transition table for the D flip-flop is shown in Table 9-11.

TABLE 9-11	
Transition table for a D flip-flop.	

<b>Output Transitions</b>			Flip-Flop Input
$Q_N$		$Q_{N+1}$	D
0	$\rightarrow$	0	0
0	$\longrightarrow$	1	1
1	$x_1 \longrightarrow x_2$	0	0
1	$\longrightarrow$	1	1

Step 4: The D inputs are plotted on the present-state Karnaugh maps in Figure 9-31. Also "don't cares" can be placed in the cells corresponding to the invalid states of 000, 011, 100, and 110, as indicated by the red Xs.



Step 5: Group the 1s, taking advantage of as many of the "don't care" states as possible for maximum simplification, as shown in Figure 9-31. The expression for each D input taken from the maps is as follows:

$$D_0 = \overline{Q}_0 + Q_2$$

$$D_1 = \overline{Q}_1$$

$$D_2 = \overline{Q}_0 + Q_2\overline{Q}_1$$

Step 6: The implementation of the counter is shown in Figure 9-32.

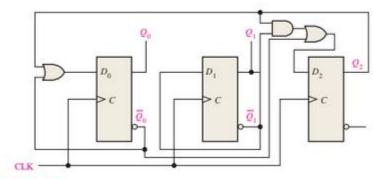


FIGURE 9-32

## SYNCHRONOUS COUNTERS

#### **EXAMPLE 9-5**

Develop a synchronous 3-bit up/down counter with a Gray code sequence using J-K flip-flops. The counter should count up when an UP/DOWN control input is 1 and count down when the control input is 0.

#### Solution

Step 1: The state diagram is shown in Figure 9–33. The 1 or 0 beside each arrow indicates the state of the UP/DOWN control input, Y.

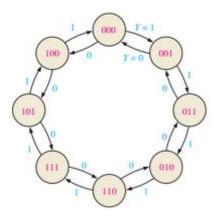


FIGURE 9-33 State diagram for a 3-bit up/down Gray code counter.

## SYNCHRONOUS COUNTERS

Step 2: The next-state table is derived from the state diagram and is shown in Table 9–12. Notice that for each present state there are two possible next states, depending on the UP/DOWN control variable, Y.

#### TABLE 9-12

Next-state table for 3-bit up/down Gray code counter.

			Next State					
Present State		Y	= 0 (DO	WN)	3	Y = 1  (UP)		
$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$	Q2	$Q_1$	$Q_0$
0	0	0	1	0	0	0	0	1
0	0	1	0	0	0	0	1	1
0	1	1	0	0	1	0	1	0
0	1	0	0	1	1	1	1	0
1	1	0	0	1	0	1	1	1
1	1	1	1	1	0	1	0	1
1	0	1	1	1	1	1	0	0
1	0	0	1	0	1	0	0	0

 $Y = UP/\overline{DOWN}$  control input.

**TABLE 9-13** 

Transition table for a J-K flip-flop.

Output Transitions			Flip-Flop Inputs		
$Q_N$		$Q_{N+1}$	J	K	
0		0	0	X	
0		1	1	X	
1	$\longrightarrow$	0	X	1	
1		1	X	0	

Step 4: The Karnaugh maps for the J and K inputs of the flip-flops are shown in Figure 9–34. The UP/DOWN control input, Y, is considered one of the state variables along with Q<sub>0</sub>, Q<sub>1</sub>, and Q<sub>2</sub>. Using the next-state table, the information in the "Flip-Flop Inputs" column of Table 9–13 is transferred onto the maps as indicated for each present state of the counter.

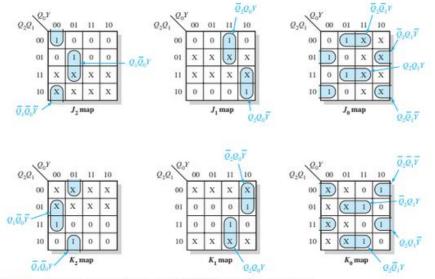


FIGURE 9-34 J and K maps for Table 9-12. The UP/DOWN control input, Y, is treated as a fourth variable

Step 5: The 1s are combined in the largest possible groupings, with "don't cares" (Xs) used where possible. The groups are factored, and the expressions for the J and K inputs are as follows:

$$\begin{array}{lll} J_0 = Q_2Q_1Y + Q_2\overline{Q}_1\overline{Y} + \overline{Q}_2\overline{Q}_1Y + \overline{Q}_2Q_1\overline{Y} & K_0 = \overline{Q}_2\overline{Q}_1\overline{Y} + \overline{Q}_2Q_1Y + Q_2\overline{Q}_1Y + Q_2Q_1\overline{Y} \\ J_1 = \overline{Q}_2Q_0Y + Q_2Q_0\overline{Y} & K_1 = \overline{Q}_2Q_0\overline{Y} + Q_2Q_0Y \\ J_2 = Q_1\overline{Q}_0Y + \overline{Q}_1\overline{Q}_0\overline{Y} & K_2 = Q_1\overline{Q}_0\overline{Y} + \overline{Q}_1\overline{Q}_0Y \end{array}$$

Step 6: The J and K equations are implemented with combinational logic. This step is the Related Problem.

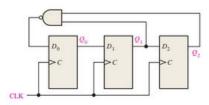
#### Related Problem

Specify the number of flip-flops, gates, and inverters that are required to implement the logic described in Step 5.

## PRACTICE QUESTIONS

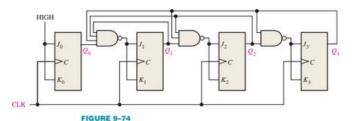
#### Section 9-5 Design of Synchronous Counters

18. Determine the sequence of the counter in Figure 9-73.



#### FIGURE 9-73

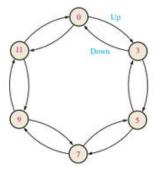
19. Determine the sequence of the counter in Figure 9-74. Begin with the counter cleared.



- Design a counter to produce the following sequence. Use J-K flip-flops.
   00, 10, 01, 11, 00, . . .
- Design a counter to produce the following binary sequence. Use J-K flip-flops.
   1, 4, 3, 5, 7, 6, 2, 1, ...
- 22. Design a counter to produce the following binary sequence. Use J-K flip-flops.

  0, 9, 1, 8, 2, 7, 3, 6, 4, 5, 0, . . .

23. Design a binary counter with the sequence shown in the state diagram of Figure 9-75.



### COUNTER APPLICATIONS

Automobile

Parking Control

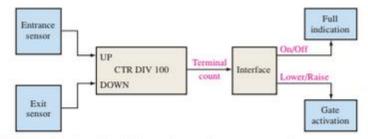


FIGURE 9-51 Functional block diagram for parking garage control.

A logic diagram of the up/down counter is shown in Figure 9–52. It consists of two cascaded up/down decade counters. The operation is described in the following paragraphs.

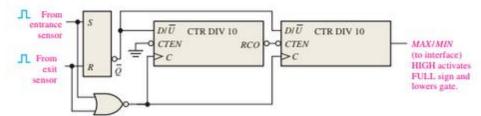


FIGURE 9-52 Logic diagram for modulus-100 up/down counter for automobile parking control.

### COUNTER APPLICATIONS

Parallel-to-Serial

Data Conversion

(Multiplexing)

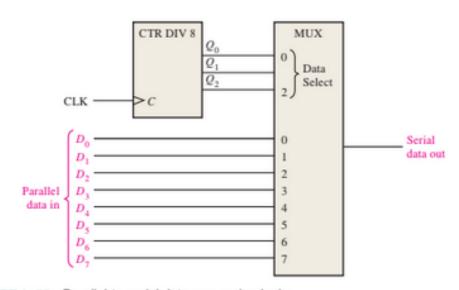


FIGURE 9-53 Parallel-to-serial data conversion logic.

## END IS THE NEW BEGINNING

## BEST OF LUCK FOR YOUR FUTURE

# IT'S BEEN A NICE JOURNEY WITH YOU GUYS