



CS215 DISCRETE MATH

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k -Element Permutations/Combinations of a Set

- *k -element permutation of N* : a **list** of k distinct elements from $\{1, 2, \dots, n\}$. How many are there?



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- *k -element combination of N* : a **set** of k distinct elements from $\{1, 2, \dots, n\}$. How many are there?

$$\binom{n}{k} = C(n, k) = \frac{P(n, k)}{k!} = \frac{n!}{k!(n-k)!}.$$



Some Properties of Binomial Coefficients

- $\binom{n}{k} = \frac{n!}{k!(n-k)!}$ is the number of k -element subsets of an n -element set.

$$\binom{n}{0} = 1 \text{ only one set of size } 0.$$

$$\binom{n}{n} = 1 \text{ only one set of size } n.$$

$\binom{n}{k} = \binom{n}{n-k}$ Obvious from equation. Can you think of a simple bijection that explains this?

Some Properties of Binomial Coefficients (cont.)



$$\sum_{i=0}^n \binom{n}{i} = 2^n$$



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Use Sum Rule

Let P = set of all subsets of $\{1, 2, \dots, n\}$

S_i = set of all i -subsets of $\{1, 2, \dots, n\}$

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$$\Rightarrow |P| = \sum_{i=0}^n |S_i| = \sum_{i=0}^n \binom{n}{i}$$

Some Properties of Binomial Coefficients (cont.)

■ Let $L = L_1 L_2 \dots L_n$ be a list of size n from $\{0, 1\}$

If \mathcal{L} = set of all such lists $\Rightarrow |\mathcal{L}| = 2^n$

There is a *bijection* between \mathcal{L} and P so
 $|P| = 2^n$ and we are done.

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Define the following function $f : \mathcal{L} \rightarrow P$

If $L \in \mathcal{L}$ then $f(L)$ is the set $S \subseteq \{1, 2, \dots, n\}$ defined by

$$i \in S \Leftrightarrow L_i = 1$$

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Ex: $n = 5$

$$f(10101) = \{1, 3, 5\}, \quad f(11101) = \{1, 2, 3, 5\}, \quad f(00000) = \emptyset$$

Binomial Coefficients

$n \backslash k$	0	1	2	3	4	5	6
0	1						
1	1	1					
2	1	2	1				
3	1	3	3	1			
4	1	4	6	4	1		
5	1	5	10	10	5	1	
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because $\binom{n}{0} = 1$

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Second half of each row is the reverse of the first half.

Sum of items on n -th row is 2^n

Pascal's Triangle

Take the table

$n \backslash k$	0	1	2	3	4	5	6
0	1						
1	1	1					
2	1	2	1				
3	1	3	3	1			
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5	1	5	10	10	5	1	
6	1	6	15	20	15	6	1

and shift each row slightly
so that middle element is
in middle

--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--

Pascal's Triangle

				1				
			1		1			
		1		2		1		
	1		3		3		1	
	1	4		6		4	1	
1		5	10		10	5		1
1	6	15	20	15	6		1	

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				1			
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	1	4		6		4	1
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1	6	15	20	15	6	1	

What is the next row in the table?

Pascal's Triangle

				1				
			1		1			
		1		2		1		
	1		3		3		1	
	1	4		6		4		1
	1	5	10		10	5		1
	1	6	15	20		15	6	1
1	7	21	35	35	21	7		1

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				1				
			1		1			
		1		2		1		
	1		3		3		1	
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1	6	15		20		15	6	1
1	7	21	35		35	21	7	1

Pascal identity

Each (non-1) entry in Pascal's Triangle is the sum of the two entries directly above it (to left and to right).



Pascal's Triangle

				1				
			1		1			
		1		2		1		
	1		3		3		1	
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1	5	10		10		5	1	
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We will use a *combinatorial proof*.

A Combinatorial Proof

- $\binom{n}{k}$ is the number of k -element subsets of an n -element set.



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$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$$

Therefore, each term (left and right) represents the number of subsets of a particular size chosen from an appropriately sized set.

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Number of $(k-1)$ -subsets of an $(n-1)$ -element set.

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Try to use sum principle to explain relationship among these three terms.

Example: $n = 5, k = 2$

$$\binom{5}{2} = \binom{4}{1} + \binom{4}{2}.$$

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Consider $S = \{A, B, C, D, E\}$.



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Set S_1 of 2-subsets of S

$$S_1 = \{\{A, B\}, \{A, C\}, \{A, D\}, \{A, E\}, \{B, C\}, \\ \{B, D\}, \{B, E\}, \{C, D\}, \{C, E\}, \{D, E\}\}.$$



A Combinatorial Proof

$$\binom{5}{2} = \binom{4}{1} + \binom{4}{2}.$$

Consider $S = \{A, B, C, D, E\}$.

Set S_1 of 2-subsets of S can be partitioned into 2 disjoint parts.

S_2 : the 2-subsets that contain E and

S_3 : the set of 2-subsets that do not contain E .

$$S_1 = \{\{A, B\}, \{A, C\}, \{A, D\}, \{A, E\}, \{B, C\}, \\ \{B, D\}, \{B, E\}, \{C, D\}, \{C, E\}, \{D, E\}\}.$$

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A Combinatorial Proof

- If n and k are integers satisfying $0 < k < n$, then

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}.$$



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To apply **sum rule**, partition S_1 into S_2 and S_3 .

Let S_2 be set of k -element subsets that **contain** x_n .

Let S_3 be set of k -element subsets that **don't contain** x_n .



Blaise Pascal

Born 1623; Died 1662

French Mathematician

A Founder of Probability Theory

Inventor of one of the first mechanical
calculating machines

Pascal Programming Language named for him



The Binomial Theorem

$$(x + y) = \binom{1}{0}x + \binom{1}{1}y$$



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Proof?



Application of the Binomial Theorem

- We may use the Binomial Theorem to prove

$$\sum_{i=0}^n \binom{n}{i} = 2^n$$

Labelling and Trinomial Coefficients

- Suppose we have k labels of one kind, e.g., **red** and $n - k$ labels of another, e.g., **blue**. In how many different ways can we apply these labels to n objects?



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Show that if we have k_1 labels of one kind, e.g., red, k_2 labels of a second kind, e.g., blue, and $k_3 = n - k_1 - k_2$ labels of a third kind, then there are $\frac{n!}{k_1!k_2!k_3!}$ ways to apply these labels to n objects



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What is the coefficient of $x^{k_1}y^{k_2}z^{k_3}$ in $(x + y + z)^n$?

Labelling and Trinomial Coefficients

- There are $\binom{n}{k_1}$ ways to choose the red items. There are then $\binom{n-k_1}{k_2}$ ways to choose the blue items from the remaining $n - k_1$. The remaining k_3 items get labelled a third color.

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Using the *product rule* the total number of labellings is

$$\begin{aligned}\binom{n}{k_1} \binom{n-k_1}{k_2} &= \frac{n!}{k_1!(n-k_1)!} \frac{(n-k_1)!}{(k_2)!(n-k_1-k_2)!} \\ &= \frac{n!}{k_1!k_2!(n-k_1-k_2)!} = \frac{n!}{k_1!k_2!k_3!}\end{aligned}$$

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a *trinomial coefficient* and denote it as

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This will be very similar to the analysis of hashing n keys into a table of size 365.



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$$\#A_n + \#B_n = 365^n$$



The Birthday Paradox

n	A_n	B_n	n	A_n	B_n
1	0.00000000	1.00000000	16	0.28360400	0.71639599
2	0.00273972	0.99726027	17	0.31500766	0.68499233
3	0.00820416	0.99179583	18	0.34691141	0.65308858
4	0.01635591	0.98364408	19	0.37911852	0.62088147
5	0.02713557	0.97286442	20	0.41143838	0.58856161
6	0.04046248	0.95953751	21	0.44368833	0.55631166
7	0.05623570	0.94376429	22	0.47569530	0.52430469
8	0.07433529	0.92566470	23	0.50729723	0.49270276
9	0.09462383	0.90537616	24	0.53834425	0.46165574
10	0.11694817	0.88305182	25	0.56869970	0.43130029
11	0.14114137	0.85885862	26	0.59824082	0.40175917
12	0.16702478	0.83297521	27	0.62685928	0.37314071
13	0.19441027	0.80558972	28	0.65446147	0.34553852
14	0.22310251	0.77689748	29	0.68096853	0.31903146
15	0.25290131	0.74709868	30	0.70631624	0.29368375

“Birthday” attacks

- Event A : **at least** two people in the room have the same birthday
- Event B : **no** two people in the room have the same birthday

$$\Pr[A] = 1 - \Pr[B]$$

$$\begin{aligned}\Pr[B] &= \left(1 - \frac{1}{365}\right) \cdot \left(1 - \frac{2}{365}\right) \cdots \left(1 - \frac{n-1}{365}\right) \\ &= \prod_{i=1}^{n-1} \left(1 - \frac{i}{365}\right).\end{aligned}$$

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$$p(n; H) := 1 - \prod_{i=1}^{n-1} \left(1 - \frac{i}{H}\right)$$



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- Since $e^x = 1 + x + \frac{x^2}{2!} + \dots$, for $|x| \ll 1$, $e^x \approx 1 + x$



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Recall that $p(n; H) := 1 - \prod_{i=1}^{n-1} (1 - \frac{i}{H})$

This probability can be approximated as

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Let $n(p; H)$ be the **smallest** number of values we have to choose, such that the probability for finding a collision is **at least** p . By inverting the expression above, we have

$$n(p; H) \approx \sqrt{2H \ln \frac{1}{1-p}}.$$

Euclidean Algorithm

- The Euclidean algorithm in pseudocode

ALGORITHM 1 The Euclidean Algorithm.

```
procedure gcd(a, b: positive integers)
  x := a
  y := b
  while y ≠ 0
    r := x mod y
    x := y
    y := r
  return x{gcd(a, b) is x}
```

The number of **divisions** required to find $\text{gcd}(a, b)$ is $O(\log b)$, where $a \geq b$. (this will be proved later.)



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procedure gcd( $a, b$ : positive integers)
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The number of **divisions** required to find $\text{gcd}(a, b)$ is $O(\log b)$, where $a \geq b$. (this will be proved later.)

Why ?

Euclidean Algorithm

- Key steps in the Euclidean algorithm

$$r_0 = r_1 q_1 + r_2 \quad 0 \leq r_2 < r_1,$$

$$r_1 = r_2 q_2 + r_3 \quad 0 \leq r_3 < r_2,$$

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$$r_{i+2} = r_i \bmod r_{i+1}$$

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Case (i): $r_{i+1} \leq \frac{1}{2} r_i$: $r_{i+2} < r_{i+1} \leq \frac{1}{2} r_i$.

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See [Theorem 1 p. 347].

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Solving Linear Recurrence Relations

- **Definition** A *linear homogeneous relation of degree k* with *constant coefficients* is a recurrence relation of the form

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \cdots + c_k a_{n-k},$$

where c_1, c_2, \dots, c_k are real numbers, and $c_k \neq 0$.



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By induction, such a recurrence relation is **uniquely** determined by this recurrence relation, and **k initial conditions** a_0, a_1, \dots, a_{k-1} .



Examples of Linear Recurrence Relations

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Examples

$$P_n = (1.11)P_{n-1}$$

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Examples

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$$f_n = f_{n-1} + f_{n-2} \quad \text{linear homogeneous recurrence relation of degree 2}$$

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$a_n = a_{n-1} + a_{n-2}^2$ NOT linear

$H_n = 2H_{n-1} + 1$ NOT homogeneous

$B_n = nB_{n-1}$ coefficients are not constants



Solving Linear Recurrence Relations

- **Example** Consider the recurrence relation

$$a_n = 2a_{n-1} - a_{n-2},$$

Which of the following are solutions?

◇ $a_n = 3n$:

◇ $a_n = 2^n$:

◇ $a_n = 5$:



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So, try to find any solution of the form $a_n = r^n$ that satisfies the recurrence.



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- ◇ Bring $a_n = r^n$ back to the recurrence relation:

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- ◇ The solutions to the *characteristic equation* can yield an explicit formula for the sequence.

$$(r^k - c_1 r^{k-1} - \cdots - c_k) = 0$$



Recall: Problem IV

■ Fibonacci number

$$F_0 = 0, F_1 = 1, F_n = F_{n-1} + F_{n-2} \text{ for } n \geq 2$$



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◇ What is the closed-form expression of F_n ?



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◇ What is the closed-form expression of F_n ?

Consider $x^n = x^{n-1} + x^{n-2}$, with $x \neq 0$. There are two different roots

$$\phi = \frac{1 + \sqrt{5}}{2}, \quad \psi = \frac{1 - \sqrt{5}}{2}$$

Then F_n can be the form of $a\phi^n + b\psi^n$. By $F_0 = 0$ and $F_1 = 1$, we have $a + b = 0$ and $\phi a + \psi b = 1$, leading to $a = \frac{1}{\sqrt{5}}$, $b = -a$. Therefore,

$$F_n = \frac{\phi^n - \psi^n}{\sqrt{5}}$$

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$$a_n = c_1 a_{n-1} + c_2 a_{n-2}.$$



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See [Theorem 1 p. 515].



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Example $a_n = 2a_{n-1} + 5a_{n-2} - 6a_{n-3}$



The Case of Degenerate Roots

- **Theorem** If the CE $r^2 - c_1r - c_2 = 0$ has **only 1** root r_0 , then

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Exercise.



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The Case of Degenerate Roots in General

- **Theorem** [Theorem 4, p.519] Suppose that there are t roots r_1, \dots, r_t with **multiplicities** m_1, \dots, m_t . Then

$$a_n = \sum_{i=1}^t \left(\sum_{j=0}^{m_i-1} \alpha_{i,j} n^j \right) r_i^n,$$

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Example

$$a_n = -3a_{n-1} - 3a_{n-2} - a_{n-3} \text{ with } a_0 = 1, a_1 = -2, \\ a_2 = -1$$



Linear Nonhomogeneous Recurrence Relations

- **Definition** A *linear nonhomogeneous relation* with constant coefficients may contain some terms $F(n)$ that depend only on n

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \cdots + c_k a_{n-k} + F(n).$$

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Fact: Assume that the sequence b_n satisfies the recurrence. Then another sequence a_n satisfies the *non-homogeneous* recurrence *if and only if* $h_n = a_n - b_n$ is a sequence that satisfies the *associated homogeneous* recurrence.



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Idea: We already know how to find h_n . For many common $f(n)$, a solution b_n to the non-homogeneous recurrence is similar to $f(n)$. We then need find $a_n = b_n + h_n$ to the non-homogeneous recurrence that satisfies both recurrence and initial conditions.



Linear Nonhomogeneous Recurrence Relations

- **Theorem** If $a_n = p(n)$ is any particular solution to the linear nonhomogeneous relation with constant coefficients,

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \cdots + c_k a_{n-k} + F(n),$$

Then all its solutions are of the form

$$a_n = p(n) + h(n),$$

where $a_n = h(n)$ is any solution to the associated homogeneous recurrence relation

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- **Example** $a_n = 3a_{n-1} + 2n$. Which solution has $a_1 = 3$?



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The *characteristic equation* of the associated linear homogeneous recurrence relation is $r^2 - 3r = 0$. Thus, the solution to the original problem are all of the form

$$a_n = \alpha 3^n + p(n) .$$



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Let $p(n) = cn + d$, then

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We get $c = -1$ and $d = -3/2$. Thus,

$$p(n) = -n - 3/2$$



Solving Linear Nonhomogeneous Recurrence Relations



"Science is driven by simulation," says Dongarra. "It's that match between the hardware capability, and the necessity of the simulations to use that hardware, where my software fits in."

"I'm a mathematician, to me, everything is linear algebra, but the world is seeing that as well," he said. "It's a fabric on which we build other things." Most problems in machine learning and AI, he said, go back to an "eternal computational component" in linear algebra.

Next Lecture

- generating function, relation ...

