操作系统lab5: RV64 缺页异常处理与fork机 制

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一、实验目的

- 通过 vm area struct 数据结构实现对进程多区域虚拟内存的管理
- 在 Lab4 实现用户态程序的基础上,添加缺页异常处理 page fault handler
- 为进程加入 fork 机制, 能够支持通过 fork 创建新的用户态进程

二、实验过程

2.1 准备工程

从仓库同步 user/main.c 文件并删除原来的 getpid.c

修改 user/Makefile:

补充SYS_CLONE的宏,不然会一直显示报错的,虽然在fork阶段才用到

```
arch > riscv > include > C syscall.h > ...

1  #ifndef __SYSCALL_H__
2  #define __SYSCALL_H__
3  #include "stdint.h"
4  #include "stddef.h"
5  #include "defs.h"
6
7  #define SYS_WRITE 64
8  #define SYS_GETPID 172
9  #define SYS_CLONE 220

10
11  int sys_write(unsigned int fd, const char *buf, size_t count);
12  int sys_getpid();
13  uint64_t do_fork(struct pt_regs *regs);
14  #endif
```

2.2 缺页异常处理

2.2.1 实现虚拟内存管理功能

先添加vma的定义与数据结构

```
arch > riscv > include > C defs.h >  VM_READ

1  #ifndef __DEFS_H__
2  #define __DEFS_H__
3
4  #include "stdint.h"
5
6  // lab5 vma
7  #define VM_ANON 0x1
8  #define VM_READ 0x2
9  #define VM_WRITE 0x4
10  #define VM_EXEC 0x8
11
```

```
/* 线程状态段数据结构 */
    struct thread struct
                                         // 32
        uint64 t ra;
        uint64 t sp;
42
        uint64 t s[12];
        uint64 t sepc, sstatus, sscratch; // 144 152 160
    };
    /* 线程数据结构 */
    struct task struct
48
        uint64 t state; // 线程状态 0
        uint64 t counter; // 运行剩余时间 8
        uint64_t priority; // 运行优先级 1 最低 10 最高, 16
        uint64 t pid;
                         // 线程 id 24
        struct thread struct thread;
        uint64 t *pgd;
                            // 用户态页表 168
        struct mm struct mm; // 176
    };
```

每一个 vm_area_struct 都对应于 task 地址空间的唯一连续区间。

为了支持 demand paging,我们需要支持对 vm_area_struct 的添加和查找:

find_vma 函数: 实现对vm_area_struct的查找

- 根据传入的地址 addr, 遍历链表 mm 包含的 VMA 链表, 找到该地址所在的 vm_area_struct
- 如果链表中所有的 vm_area_struct 都不包含该地址, 则返回 NULL

```
// lab 5 function
struct vm_area_struct *find_vma(struct mm_struct *mm, uint64_t addr)
{
    // 遍历查找
    struct vm_area_struct *vma = mm->mmap;
    while (vma)
    {
        if (vma->vm_start <= addr && addr < vma->vm_end)
            return vma;
        vma = vma->vm_next;
    }
    return NULL;
}
```

do_mmap 函数: 实现 vm_area_struct 的添加

• 新建 vm_area_struct 结构体,根据传入的参数对结构体赋值,并添加到 mm 指向的 VMA 链表中

2.2.2 修改task_init

本次注释掉之前实验对用户栈,代码load segment的映射操作 (alloc和create_mapping).

调用 do_mmap 函数,建立用户 task 的虚拟地址空间信息,在本次实验中仅包括两个区域:

- 代码和数据区域: 该区域从 ELF 给出的 Segment 起始用户态虚拟地址 phdr->p_vaddr 开始,对应文件中偏移量为 phdr->p_offset 开始的部分
- 用户栈: 范围为 [USER_END PGSIZE, USER_END) , 权限为 VM_READ | VM_WRITE , 并且是匿名的区域 (VM_ANON)

```
void load_program(struct task_struct *task)
{
    INFO("\ntask init - load program");
    Elf64_Ehdr *ehdr = (Elf64_Ehdr *)_sramdisk;
    Elf64_Phdr *phdrs = (Elf64_Phdr *)(_sramdisk + ehdr->e_phoff);

    // uint64_t uapp_size = _eramdisk - _sramdisk;
    // int page_num = ((uapp_size + PGSIZE + 1) / PGSIZE);
    // DEBUG("uapp page size: %llx page_num: %llx", uapp_size, page_num);
    // char *new_uapp = (char *)alloc_pages(page_num);
    printk("enter load_program, e_phnum = %d\n", ehdr->e_phnum);
    for (int i = 0; i < ehdr->e_phnum; ++i)
```

```
Elf64_Phdr *phdr = phdrs + i;
        printk("i = %d, type = %d\n",i, phdr->p_type);
        if (phdr->p_type == PT_LOAD)
        {
            // alloc space and copy content
            uint64_t offset = PGOFFSET(phdr->p_vaddr);
            uint64_t size = offset + phdr->p_memsz;
            uint64_t perm = 0x0; // 初始化
            perm = (phdr->p_flags & PF_X) ? perm | PTE_X : perm;
            perm = (phdr->p_flags & PF_W) ? perm | PTE_W : perm;
            perm = (phdr->p_flags & PF_R) ? perm | PTE_R : perm;
            DEBUG("perm: %11x", perm);
            // char *va = (char *)alloc_pages((size + PGSIZE - 1) / PGSIZE);
            DEBUG("size : %11x", phdr->p_memsz + offset);
           DEBUG("page number : %11x , offset number : %11x", (size + PGSIZE -
1) / PGSIZE, offset);
            // segment 使用的内存就是 [p_vaddr, p_vaddr + p_memsz) 这一连续区间,然后
将 segment 的内容从 ELF 文件中读入到这一内存区间
           // 为了实现对齐, 实际上 0 - offset之间的为无效内容
           // for (int i = 0; i < phdr->p_filesz; i++)
           // va[offset + i] = ((char *)ehdr)[phdr->p_offset + i];
           // DEBUG("va : %11x", va);
           // 将 [p_vaddr + p_filesz, p_vaddr + p_memsz) 对应的 *物理区间* 清零
           // memset(va + offset + phdr->p_filesz, 0, phdr->p_memsz - phdr-
>p_filesz);
           // DEBUG("va: %11x, pa: %11x, size: %11x, perm: %11x", va,
(uint64_t)va - PA2VA_OFFSET, phdr->p_memsz + offset, perm | PTE_U | PTE_V);
            // create_mapping(task->pgd, phdr->p_vaddr, (uint64_t)va -
PA2VA_OFFSET, size, perm | PTE_U | PTE_V);
            // 1ab5
            \label{local_mmap} $$ do_mmap(\&task->mm, phdr->p\_vaddr, phdr->p\_memsz, phdr->p\_offset, $$
phdr->p_filesz, VM_READ | VM_WRITE | VM_EXEC);
        }
   }
   // DEBUG("user_start: %llx", USER_START);
    // uint64_t user_stack = (uint64_t)alloc_page();
   // DEBUG("user_stack: %llx", user_stack);
   // create_mapping(task->pgd, USER_END - PGSIZE, user_stack - PA2VA_OFFSET,
PGSIZE, PTE_U | PTE_V | PTE_R | PTE_W);
    do_mmap(&task->mm, USER_END - PGSIZE, PGSIZE, 0, 0, VM_ANON | VM_READ |
VM_WRITE);
    task->thread.sepc = ehdr->e_entry;
}
```

task init函数也需要修改, 修改如下:

2.3 实现page fault handler

修改 trap.c, 为 trap_handler 添加捕获 page fault 的逻辑, 分别需要捕获 12, 13, 15 号异常。

实现缺页异常的函数 do_page_fault, 具体逻辑如下:

- 1. 通过 stval 获得访问出错的虚拟内存地址 (Bad Address)
- 2. 通过 find_vma() 查找 bad address 是否在某个vma 中
 - 。 如果不在,则出现非预期错误,可以通过 Err 宏输出错误信息
 - o 如果在,则根据vma 的 flags 权限判断当前page fault 是否合法
 - 如果非法(比如触发的是 instruction page fault 但 vma 权限不允许执行),则 Err 输 出错误信息
 - 其他情况合法,需要我们按接下来的流程创建映射
- 3. 分配一个页,接下来要将这个页映射到对应的用户地址空间
- 4. 通过(vma->vm_flags & VM_ANON) 获得当前的 VMA 是否是匿名空间
 - 如果是匿名空间,则直接映射即可
 - o 如果不是,则需要根据 vma->vm_pgoff 等信息从 ELF 中读取数据,填充后映射到用户空间

2.4 测试缺页处理

使用 make run TEST=PFH1 在task_init中并未进行映射,直到page fault触发用户态进程的拷贝和映射,且只有第一次触发

使用 make run TEST=PFH2 类似

```
...budy int done!
...mm init done!
...mm init done!
...mm init done!
INFO] setup wn final starting
IVm.c,125,create mapping] root : ffffffe00020c000, [80204000, 80206000] -> [fffffe000204000, ffffffe000204000], perm: b
IVm.c,125,create mapping] root : ffffffe00020c000, [80206000, 80206000] -> [fffffe000206000, fffffe000206000], perm: 3
IVm.c,125,create mapping] root : fffffe00020c000, [80206000, 80000000] -> [fffffe000206000, fffffe000206000], perm: 7
IVm.c,125,create mapping] root : fffffe00020c000, [80206000, 80000000] -> [fffffe000206000, fffffe000206000], perm: 7
IUFO] setup, vm final done
[INFO] setup, vm final done
[INFO] setup, vm final done
[INFO] task init - load program
enter load program, e.phnum = 3
1 = 0, type = 1879048195
1 = 1, type = 1
IOEBUG) perm: e
IOEBUG) per uniber : 4, offset number : e8
I = 2, type = 1065362481
INFO = 1 priority = 7
INFO
```

3.1 实现fork系统调用 准备工作

修改proc.c相关代码,使得其只初始化一个进程,其他进程保留为 NULL 等待 fork 创建

- 定义nr_tasks记录当前进程数并作为栈顶指针
- task_init: 只初始化一个进程,将nr_tasks置为2指向栈顶
- Schedule:将NR_TASKS替换为nr_tasks

```
for (int i = 1; i < nr_tasks; i++)
{
    task[i] = (struct task_struct *)kalloc();
    task[i]->state = TASK_RUNNING;
    task[i]->counter = 0;
    task[i]->priority = PRIORITY_MIN + rand() % (PRIORITY_MAX - PRIORITY_MIN + 1);
```

```
// lab 5 fork
// 记录当前进程数,并作为 tasks 的栈顶指针来使用
uint64_t nr_tasks = 2;
```

添加系统调用处理

在syscall.h加入#define SYS_CLONE 220(之前已经加入了)

在trap_handler调用do_fork, 处理regs->a7 == SYS_CLONE的情况

```
case 8:

// 处理用户态系统调用
int sys_call_num = regs->x[17]; // 获取系统调用编号
switch (sys_call_num)
{
    case SYS_WRITE:
        sys_write((unsigned int)regs->x[10], (const char *)regs->x[11], (size_t)regs->x[12]); // 处理write系统调用
        break;
case SYS_GETPID:
        regs->x[10] = sys_getpid(); // 处理getpid系统调用
        break;
case SYS_CLONE:
        regs->x[10] = do_fork(regs); // 处理fork系统调用
        break;
default:
        break;
}
regs->sepc += 4; // 更新程序计数器,跳过系统调用指令
break;
```

3.2 do_fork实现

fork的实现:

- 创建一个新进程:
 - 。 拷贝内核栈 (包括了 task_struct 等信息)
 - 。 创建一个新的页表
 - 拷贝内核页表 swapper_pg_dir
 - 遍历父进程 vma, 并遍历父进程页表
 - 将这个 vma 也添加到新进程的 vma 链表中
 - 如果该 vma 项有对应的页表项存在(说明已经创建了映射),则需要深拷贝一整页的内容并映射到新页表中

- 将新进程加入调度队列
- 处理父子进程的返回值
 - 父进程通过 do_fork 函数直接返回子进程的 pid, 并回到自身运行
 - 子进程通过被调度器调度后(跳到 thread.ra), 开始执行并返回 0

```
// 遍历父进程的虚拟内存区域 (VMA)
struct vm_area_struct *pyma = current->mm.mmap;
for (; pyma != NULL; pyma = pyma->vm_next)

// 根据父进程的VMA信息,映射相应的内存区域到新进程
do_mmap(6_task->mm, pyma->vm_start, pyma->vm_start, pyma->vm_goff, pyma->vm_filesz, pyma->vm_flags);

// 遍历VMA中的每一页,复制父进程的物理页到新进程
for (uint64_t va = PGROUNDDOWN(pyma->vm_start); va < pyma->vm_end; va += PGSIZE)

// int64_t pa = PTE2PA(pte_res); // 获取物理地址
if (pte_res == 0)
continue;
uint64_t *page = (uint64_t *)alloc_page(); // 为新页分配内存
memcpy((void *)page, (void *)(pa + PAZVA_OFFSET), PGSIZE); // 复制父进程的页到新进程

Log(*copy page: %lX -> %lX*, (uint64_t)(pa), (uint64_t)page);
create_mapping(_task->pgd, va, (uint64_t)page - PAZVA_OFFSET, PGSIZE, PTE_U | PTE_V | pyma->vm_flags); // 创建新进程的页表映射
}

// 更新新进程的寄存器状态。设置正确的返回值和SEPC
struct pt_regs *new_regs = (struct pt_regs *)((PGOFFSET((uint64_t)regs) + (uint64_t)_task));
memcpy((void *)new_regs, (void *)regs, sizeof(struct pt_regs));
new_regs->x[l0] = 0; // 返園新进程的PID

// 更新SEPC, 魏过fork指令
return_task->pid; // 返回新进程的PID
```

3.3 处理进程返回逻辑

修改entry.S如下:

```
| Solite, 272(sp) | Solite, 2
```

此部分还有设置do_fork的父进程、子进程的用户栈和内核栈指针部分,在前面已经展示

3.4 测试fork

make run TEST=FORK1

fork时为已映射页拷贝并映射,fork出的子进程与父进程global_variable相互独立

fork在PID=2 forked from PID=1后调用完cause = 1f的 create_mapping 后完成

make run TEST=FORK2

fork出PID=2的子进程实现深拷贝,正确输出父进程改变的字符串及变量值,且后续与父进程独立.

```
[syscall.c,41,do_fork] current: thread sp : ffffffe0002cf000, thread sscratch : 4000000000 [syscall.c,42,do_fork] sp : fffffffe0002ce08 [syscall.c,46,do_fork] cur_sscratch : 3fffffffc0 [INFO] [PID = 2] [syscall.c,84,do_fork] cur_sscratch : 3fffffffc0 [INFO] [PID = 2] [syscall.c,87,do_fork] copy page : 802d5000 -> ffffffe0002cd000 [sozdf000, 802c0000] -> [3fffffff000, 400000000], perm: 17 [syscall.c,87,do_fork] copy page : 802d5000 -> ffffffe0002cs5000 [vm.c,125,create_mapping] root : ffffffe0002d5000, [802c3000, 802c4000] -> [10000, 11000], perm: 1f [syscall.c,87,do_fork] copy page : 802d900 -> ffffffe0002c5000 [vm.c,125,create_mapping] root : ffffffe0002d5000, [802c8000, 802c7000] -> [11000, 12000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c5000 [sozc8000] -> [12000, 13000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c5000 [vm.c,125,create_mapping] root : ffffffe0002dd5000, [802c7000, 802c8000] -> [12000, 13000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c8000 [sozc8000] -> [13000, 14000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c8000 [sozc8000] -> [13000, 14000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c8000 [sozc8000] -> [14000, 15000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c8000 [sozc8000] -> [14000, 15000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c8000 [sozc8000] -> [14000, 15000], perm: 1f [syscall.c,87,do_fork] copy page : 802d800 -> ffffffe0002c8000 [sozc8000] -> [14000, 15000], perm: 1f [syscall.c,87,do_fork] copy page : 802d800 -> ffffffe0002c8000 [sozc8000] -> [14000, 15000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c8000 [sozc8000] -> [14000, 15000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c8000 [sozc8000] -> [14000, 15000], perm: 1f [syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002c8000 [sozc800] copy page : 802d800 [sozc800], sozc800] copy page : 802d800
```

make run TEST=FORK3

进程能够正确实现多层fork(),并且子进程正确复制父进程global variable值,后续父子进程相互独立

```
... butdy init done!
...mm init done!
..
```

```
| vm.c,125,create_mapping| root: ffffffe0002db000, [802e1000, 802e2000) -> [10000, 11000), perm: 1f | syscall.c,87,do fork| copy page: 802d9000 -> ffffffe0002c4000 | vm.c,125,create_mapping| root: ffffffe0002db00, [802e3000] -> [11000, 12000), perm: 1f | syscall.c,87,do fork| copy page: 802d8000 -> ffffffe0002c5000 | vm.c,125,create_mapping| root: ffffffe0002c5000 | vm.c,125,create_mapping| root: ffffffe0002c5000 | vm.c,125,create_mapping| root: fffffffe0002c5000 | vm.c,125,create_mapping| root: fffffff60002c5000 | vm.c,125,create_mapping| root: fffffff60002c5000 | vm.c,125,create_mapping| root: fffffff60002c5000 | vm.c,125,create_mapping| root: ffffff60002c5000 | vm.c,125,create_mapping| root: ffffff60002c5000 | vm.c,125,create_mapping| root: fffff60002c5000 | vm.c,125,create_mapping| root: ffffff60002c5000 | vm.c,125,create_mapping| root: f
```

三、讨论和心得

本次写代码的过程还是比较困难的,战线拉得比较长,细节问题也比较多,很难处理,之前的代码也有很多要注释掉的部分,我对注释掉哪个部分也不算特别清晰,因而踩了一些坑。

还有会有莫名其妙的问题,解决起来更是痛苦,只能通过大改代码来解决。

我在这里设置新进程的内核栈位置错误,没有-35*8,会在切换进程的时候卡死,因为栈指针的位置不对。子线程跳回的地址在_traps的call trap_handler之后,所以要手动调整-35 * 8.

```
uint64_t do_fork(struct pt_regs *regs)

uint64_t do_fork(struct pt_regs *regs)

uint64_t cul_sscratch : %lx", cur_sscratch);

// g制父进程的内存映射结构

memcpy((void *)_task, (void *)current, PGSIZE);

_task->pid = nr_tasks; // 设置新进程的PID

nr_tasks++; // 增加全局任务计数

INFO("[PID = %d] forked from [PID = %d]", _task->pid, current->pid); // 记录新进程创建信息

if (_task->pid >= NR_TASKS)

Err("task number exceed"); // 如果任务数量超过上限,报错

task[_task->pid] = _task; // 将新进程添加到任务数组中

_task->mm.mmap = NULL; // 清空新进程的内存映射结构

// 设置新进程的返回地址和栈指针
_task->thread.ra = (uint64_t)__ret_from_fork;

task->thread.sp = (uint64_t)_task + PGSIZE; // 设置新进程的内核栈位置

_task->thread.sscratch = cur_sscratch; // 设置新进程的sscratch寄存器
```

```
[INFO] [PID = 4] forked from [PID = 1]
[syscall.c,87,do_fork] copy page : 802d5000 -> ffffffe0002f5000
[vm.c,124,create_mapping] root : ffffffe0002f3000, [802f5000, 802f6000) -> [3ffffff000, 400000000), perm: 17
[syscall.c,87,do_fork] copy page : 802d2000 -> ffffffe0002f9000
[vm.c,124,create_mapping] root : ffffffe0002f3000, [802f9000, 802fa000) -> [10000, 11000), perm: 1f
[syscall.c,87,do_fork] copy page : 802d9000 -> ffffffe0002f3000
[vm.c,124,create_mapping] root : ffffffe0002f3000, [802fd000, 802fd000) -> [11000, 12000), perm: 1f
[syscall.c,87,do_fork] copy page : 802d8000 -> ffffffe0002fd000
[vm.c,124,create_mapping] root : ffffffe0002f3000, [802fd000, 802fe000) -> [12000, 13000), perm: 1f
[U] pid: 1 is running! global_variable: 2
[U] pid: 1 is running! global_variable: 3
SCHEDULE(): INFO OF ALL PROCESSES
pid = 0, counter = 0, priority = 0
pid = 1, counter = 0, priority = 7
pid = 2, counter = 6, priority = 7
pid = 2, counter = 6, priority = 7
pid = 3, counter = 6, priority = 7
switch from 1 to [PID = 2 PRIORITY = 7 COUNTER = 6]
```

我在fork的最后一开始没有给new regs的sepc+4,这样会一直进行fork指令,直到超过nr_tasks数目。

```
[U] pid: 4 is running! global variable: 2
[U] pid: 4 is running! global variable: 3
SCHEDULE(): INFO OF ALL PROCESSES
pid = 0, counter = 0, priority = 0
pid = 1, counter = 0, priority = 7
pid = 2, counter = 0, priority = 7
pid = 2, counter = 0, priority = 7
pid = 3, counter = 0, priority = 7
pid = 4, counter = 0, priority = 7
pid = 5, counter = 0, priority = 7
pid = 5, counter = 6, priority = 7
pid = 6, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 2, counter = 6, priority = 7
pid = 3, counter = 6, priority = 7
pid = 3, counter = 6, priority = 7
pid = 4, counter = 6, priority = 7
pid = 5, counter = 6, priority = 7
pid = 6, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 8, counter = 6, priority = 7
pid = 9, counter = 6, priority = 7
pid = 9, counter = 6, priority = 7
pid = 9, counter = 6, priority = 7
pid = 9, counter = 6, priority = 7
pid = 9, counter = 6, priority = 7
pid
```

四、思考题

1. 呈现出你在 page fault 的时候拷贝 ELF 程序内容的逻辑。

通过vma获取bad_addr所在segment和page的起始地址

 126
 uint64_t bad_seg_start = vma->vm_pgoff + (uint64_t)_sramdisk; // 文件段起始位置

 127
 uint64_t bad_seg_now = bad_seg_start + PGROUNDDOWN(bad_addr) - vma->vm_start; // 当前映射段的起始位置

拷贝起始地址:

- 如果bad_seg_now和bad_addr在同一页,从seg_start开始拷贝
- bad_seg_now处于一个整页page,直接拷贝所在页的内容
- 如果bad_seg_now和和file size的末尾在同一个page,只对在file size范围内部分进行拷贝,但建立映射的范围是整个page
- bad_seg_now和vm_end在同一个page,从vm_end所在页的起始部分开始拷贝, size=PGOFFSET(vm_end)

2. 回答4.3.5 中的问题:

· 在 do_fork 中,父进程的内核栈和用户栈指针分别是什么?

内核栈指针为thread_struct.sp, 用户栈指针为thread_struct.sscratch

• 在 do_fork 中,子进程的内核栈和用户栈指针的值应该是什么?

内核栈指针的值为子进程的pt_regs地址,用户栈指针的值为当前sscratch寄存器的值,即父进程用户态 栈指针

• 在 do_fork 中,子进程的内核栈和用户栈指针分别应该赋值给 谁?

用户态进程的thread_struct.sp和thread_struct.sscratch

3. 为什么要为子进程 pt_regs 的 sepc 手动加四?

子进程应该和父进程一样回到用户程序的对应位置。父进程在do_fork返回用户程序经过trap_handler时,因为发生的是系统调用异常,需要手动+4,避免重复系统调用语句从而死循环。而子进程返回用户程序不经过trap_handler,为达到相同效果需要在此时为其手动+4,使得_traps通过sret返回时,回到下一条语句。

4. 对于 Fork main #2 (即 FORK2), 在运行时, ZJU OS Lab5 位于内存的什么位置? 是否在读取的时候产生了page fault? 请给出必要的截图以说明。

位于[p_vaddr + p_filesz, p_vaddr + p_memsz)

没有发生page fault,在fork进程时父进程已经映射其所在的内存段,子进程也完成复制以及映射在[12000, 13000)这段里面

```
| Switch from 0 to [PID = 1 PRIORITY = 7 COUNTER = 7|
| Itrap.c,82.do_page_fault| [PID = 1 PC = 100e8] valid page fault at 100e8 with cause 12
| Ivm.c,125,create_mapping| root : ffffffe0002cf000, [802d2000, 802d3000] -> [10000, 11000), perm: 1f
| Itrap.c,82.do_page_fault| [PID = 1 PC = 101e3] valid page fault at 12000 with cause 15
| Ivm.c,125,create_mapping| root : ffffffe0002cf000, [802d5000, 802d5000] -> [12000, 13000], perm: 17
| Itrap.c,82.do_page_fault| [PID = 1 PC = 101e4] valid page fault at 12000 with cause 13
| Ivm.c,125,create_mapping| root : ffffffe0002cf000, [802d5000, 802d5000] -> [12000, 13000], perm: 1f
| Itrap.c,82.do_page_fault| [PID = 1 PC = 112e4] valid page fault at 112e4 with cause 12
| Ivm.c,125,create_mapping| root : ffffffe0002cf000, [802d5000, 802d5000] -> [11000, 12000], perm: 1f
| Itrap.c,82.do_page_fault| [PID = 1 PC = 1042e4] valid page fault at 14068 with cause 13
| Ivm.c,125,create_mapping| root : ffffffe0002cf000, [802d5000, 802d5000] -> [14000, 143f8), perm: 1f
| Ivm.c,125,create_mapping| root : ffffffe0002cf000, [802d5000] -> [14000, 143f8), perm: 1f
| Ivm.c,125,create_mapping| root : ffffffe0002cf000, [802d5000] -> [13000, 14000], perm: 1f
| Ivm.c,125,create_mapping| root : ffffffe0002cf000, [802d5000] -> [13000, 14000], perm: 1f
| Ivs.call.c,41,do_fork] sp: ffffffe0002cf000, [802d5000], sp: 13000] -> [13000, 14000], perm: 1f
| Ivs.call.c,42,do_fork] sp: ffffffe0002cf000, [802d5000] -> [3ffffff000, 4000000000], perm: 17
| Ivs.call.c,67,do_fork] copy page : 802d5000 -> ffffffe0002df000
| Ivm.c,125,create_mapping| root : ffffffe0002df000, 802e0000] -> [3ffffff000, 4000000000], perm: 1f
| Ivs.call.c,67,do_fork] copy page : 802d5000 -> ffffffe0002df000, 802e0000] -> [10000, 12000], perm: 1f
| Ivs.call.c,67,do_fork] copy page : 802d5000 -> ffffffe0002df000, 802e0000] -> [10000, 12000], perm: 1f
| Ivs.call.c,67,do_fork] copy page : 802d5000 -> ffffffe0002df000, 802e0000] -> [10000, 12000], perm: 1f
| Ivs.call.c,67,do_fork] copy page : 802d5000 -> ffffffe0002df000, 802e0000] -> [100
```

5. 画图分析 make run TEST=FORK3 的进程 fork 过程,并呈现出各个进程的 global_variable 应该从几开始输出,再与你的输出进行对比验证

开始的全局变量值如下:

```
pid 1 - 0
pid 2 - 1
pid 5 - 1
pid 8 - 2
pid 6 - 2
pid 3 - 1
pid 7 - 2
pid 4 - 2
```

和输出的结果相同:

```
[trap.c,82,do_page_fault] [PID = 1 PC = 100e8] valid page fault at `100e8` with cause 12
[vm.c,125,create_mapping] root : ffffffe0002cf000, [802d2000, 802d3000) -> [10000, 11000), perm: 1f
[trap.c,82,do_page_fault] [PID = 1 PC = 101ac] valid page fault at `3ffffffff8` with cause 15
[vm.c,125,create_mapping] root : ffffffe0002cf000, [802d5000, 802d6000) -> [3ffffff000, 4000000000), perm: 17
[trap.c,82,do_page_fault] [PID = 1 PC = 101c8] valid page fault at `12000` with cause 13
[vm.c,125,create_mapping] root : ffffffe0002cf000, [802d8000, 802d83f0) -> [12000, 123f0), perm: 1f
[trap.c,82,do_page_fault] [PID = 1 PC = 11134] valid page fault at `11134` with cause 12
[vm.c,125,create_mapping] root : ffffffe0002cf000, [802d9000, 802da000) -> [11000, 12000), perm: 1f
[U] pid: 1 is running! global variable: 0
[syscall.c,41,do_fork] current: thread sp : ffffffe0002cf000, thread sscratch : 4000000000
[syscall.c,42,do_fork] sp : ffffffe0002ceee8
[syscall.c,42,do_fork] cur_sscratch : 3ffffffd0
[INFO] [PID = 2] forked from [PID = 1]
      [INFO] [PID = 5] forked from [PID = 2]
[syscall.c,87,do_fork] copy page : 802e1000 -> ffffffe000301000
[vm.c,125,create_mapping] root : ffffffe0002ff000, [80301000, 80302000) -> [10000, 11000), perm: 1f
[syscall.c,87,do_fork] copy page : 802e4000 -> ffffffe000305000) -> [11000, 12000), perm: 1f
[syscall.c,87,do_fork] copy page : 802e5000 -> ffffffe000305000
[vm.c,125,create_mapping] root : ffffffe0002ff000, [80305000, 80305000) -> [12000, 13000), perm: 1f
[syscall.c,87,do_fork] copy page : 802dd000 -> ffffffe000307000
[vm.c,125,create_mapping] root : ffffffe0002ff000, [80307000, 80306000) -> [3ffffff000, 400000000), perm: 17
[U] pid: 2 is running! global_variable: 1
[syscall.c,41,do_fork] current: thread sp : ffffffe0002daee8, thread sscratch : 3ffffffd0
[syscall.c,42,do_fork] sp : ffffffe0002ceee8
[syscall.c,46,do_fork] cur_sscratch : 3ffffffd0
[INFO] [PID = 6] forked from [PID = 2]
       [INFO] [PID = 6] forked from [PID = 2]
[syscall.c,87,do_fork] copy page : 802e1000 -> ffffffe00030d000
[vm.c,125,create_mapping] root : ffffffe00030b000, [8030d000, 8030e000) -> [10000, 11000), perm: 1f
[syscall.c,87,do_fork] copy page : 802e4000 -> ffffffe000310000
[vm.c,125,create_mapping] root : ffffffe00030b000, [80310000, 80311000) -> [11000, 12000), perm: 1f
[syscall.c,87,do_fork] copy page : 802e5000 -> ffffffe000311000
[vm.c,125,create_mapping] root : ffffffe00030b000, [80311000, 80312000) -> [12000, 13000), perm: 1f
[syscall.c,87,do_fork] copy page : 802dd000 -> ffffffe000313000
[vm.c,125,create_mapping] root : ffffffe00030b000, [80313000, 80314000) -> [3ffffff000, 400000000), perm: 17
     [U] pid: 3 is running! global variable: 1
[syscall.c,41,do_fork] current: thread sp : ffffffe0002e6ee8, thread sscratch : 3fffffffd0
[syscall.c,42,do_fork] sp : ffffffe0002ceee8
[syscall.c,46,do_fork] cur_sscratch : 3fffffffd0
    [INFO] [PID = 7] forked from [PID = 3]
[syscall.c,87,do_fork] copy page : 802ed000 -> ffffffe000319000
[vm.c,125,create mapping] root : ffffffe000317000, [80319000, 8031a000) -> [10000, 11000), perm: 1f
[syscall.c,87,do_fork] copy page : 802f0000 -> ffffffe00031c000
[vm.c,125,create mapping] root : ffffffe000317000, [8031c000, 8031d000) -> [11000, 12000), perm: 1f
[syscall.c,87,do_fork] copy page : 802f1000 -> ffffffe00031d000
[vm.c,125,create mapping] root : ffffffe000317000, [8031d000, 8031e000) -> [12000, 13000), perm: 1f
[syscall.c,87,do_fork] copy page : 802e9000 -> ffffffe00031f000
[vm.c,125,create mapping] root : ffffffe000317000, [8031f000, 80320000) -> [3ffffff000, 4000000000), perm: 17
                                          from 3 to [PID = 4 PRIORITY = 7 COUNTER = 6]
   [U] pid: 4 is running! global variable: 2
   [U] pid: 4 is running! global_variable: 3
 pid = 1, counter = 0, priority = 7
pid = 2, counter = 0, priority = 7
 pid = 6, counter = 6, priority = 7
pid = 7, counter = 6, priority = 7
switch from 4 to [PID = 5 PRIORITY = 7 COUNTER = 6]
 [U] pid: 5 is running! global_variable: 1
[syscall.c,41,do_fork] current: thread sp : ffffffe0002feee8, thread sscratch : 3fffffffd0
[syscall.c,42,do_fork] sp : ffffffe0002ceee8
[syscall.c,46,do_fork] cur_sscratch : 3fffffffd0
[INFO] [PID = 8] forked from [PID = 5]
```

```
switch from 5 to [PID = 6 PRIORITY = 7 COUNTER = 6]
[U] pid: 6 is running! global variable: 2
[U] pid: 6 is running! global variable: 3
SCHEDULE(): INFO OF ALL PROCESSES
pid = 6, counter = 0, priority = 7
pid = 8, counter = 6, priority = 7
switch from 6 to [PID = 7 PRIORITY = 7 COUNTER = 6]
[U] pid: 7 is running! global variable: 2
[U] pid: 7 is running! global variable: 3
SCHEDULE(): INFO OF ALL PROCESSES
pid = 3, counter = 0, priority = 7
pid = 4, counter = 0, priority = 7
pid = 5, counter = 0, priority = 7
pid = 8, counter = 6, priority = 7
switch from 7 to [PID = 8 PRIORITY = 7 COUNTER = 6]
[U] pid: 8 is running! global variable: 2
[U] pid: 8 is running! global variable: 3
SCHEDULE(): INFO OF ALL PROCESSES
```