

CS-200

Computer Architecture

Part I: Instruction Set Architecture

Arrays and Data Structures

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Arrays in High-Level Languages

- Java:

```
short[] myData = {10407, -16533, -22715, 29796, 18956,...};
```

- Python:

```
myData = [10407, -16533, -22715, 29796, 18956,...]
```

or

```
import array
```

```
myData = array.array('h', [10407, -16533, -22715, 29796, 18956,...])
```

- Scala:

```
val myData: Array[Short] = Array(10407, -16533, -22715, 29796, 18956,...)
```

Different Ways to Store Arrays

A

myData →

10407
-16533
-22715
29796
18956
-882
-27277
16066
-30868
0

16-bit

B

myData →

9
10407
-16533
-22715
29796
18956
-882
-27277
16066
-30868

16-bit

C

myData →

10407
-16533
-22715
29796
18956
-882
-27277
16066
-30868

16-bit

Adding Positive Elements

- Add all positive elements in an array of signed 16-bit integers
 - At call time, `a0` points to the array
 - At return time, `a0` contains the result
- Write it three times, for arrays of type **A**, **B**, and **C**

Type C

```
short sum = 0;
int i;
for (i = 0; i < N; i++) {
    if (myData[i] > 0) {
        sum += myData[i];
    }
}
```

A

myData →

10407
-16533
-22715
29796
18956
-882
-27277
16066
-30868
0

16-bit

B

myData →

9
10407
-16533
-22715
29796
18956
-882
-27277
16066
-30868

16-bit

Adding Positive Elements (Variation on C)

- Add all positive elements in an array of signed 16-bit integers
 - At call time, `a0` points to the array and `a1` is the length of the array
 - At return time, `a0` contains the result
- Write it by incrementing the index of the array

```
int i = 0;
while (i < N) {
    if (myData[i] > 0) {
        ...
    }
    i++;
}
```

=

```
int i;
for (i = 0; i < N; i++) {
    if (myData[i] > 0) {
        ...
    }
}
```

C

myData →

10407
-16533
-22715
29796
18956
-882
-27277
16066
-30868

16-bit

Which is Better?

Pointer to memory

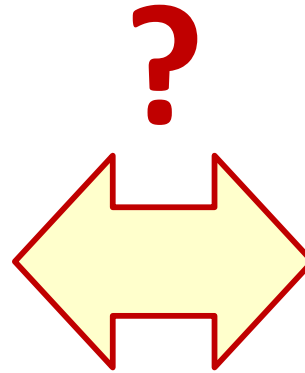
```
add_positive:
    li    t0, 0
    mv    t1, a1

next_short:
    beq    t1, zero, end
    lh     t2, 0(a0)
    bltz   t2, negative
    add    t0, t0, t2

negative:
    addi   a0, a0, 2
    addi   t1, t1, -1
    j      next_short

end:
    mv     a0, t0
    ret
```

Less instructions → Faster



Index in array

```
add_positive:
    li    t0, 0
    li    t1, 0

next_index:
    bge    t1, a1, end
    slli   t2, t1, 1
    add    t2, a0, t2
    lh     t3, 0(t2)
    bltz   t3, negative
    add    t0, t0, t3

negative:
    addi   t1, t1, 1
    j      next_index

end:
    mv     a0, t0
    ret
```

Which is Better?

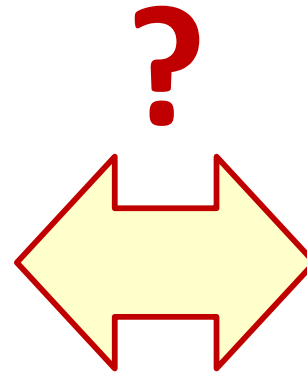
Pointer to memory

```
short sum = 0;
short *ptr = myData;
short *end = myData + N;

while (ptr < end) {
    if (*ptr > 0) {
        sum += *ptr;
    }
    ptr++;
}
```

A pointer to **short** integers
= address of **myData**

Advances the **pointer** instead of recomputing the element address



Index in array

```
short sum = 0;
int i;
for (i = 0; i < N; i++) {
    if (myData[i] > 0) {
        sum += myData[i];
    }
}
```

Better code

We Need Good Compilers!

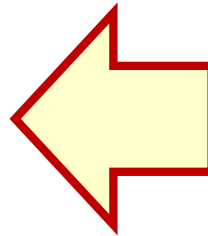
Pointer to memory

```
add_positive:
    li    t0, 0
    mv    t1, a1

next_short:
    beq    t1, zero, end
    lh     t2, 0(a0)
    bltz   t2, negative
    add    t0, t0, t2

negative:
    addi   a0, a0, 2
    addi   t1, t1, -1
    j      next_short

end:
    mv     a0, t0
    ret
```

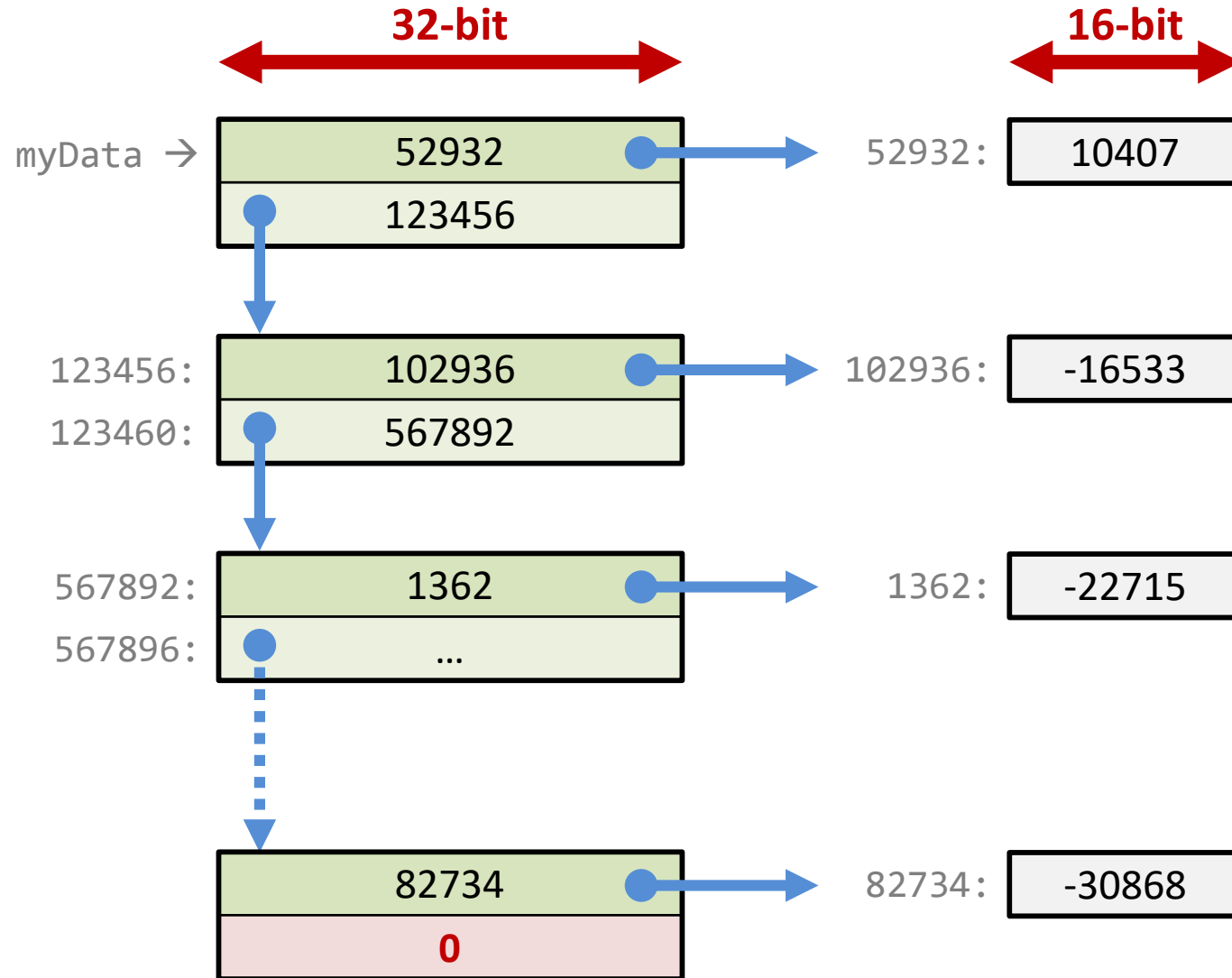


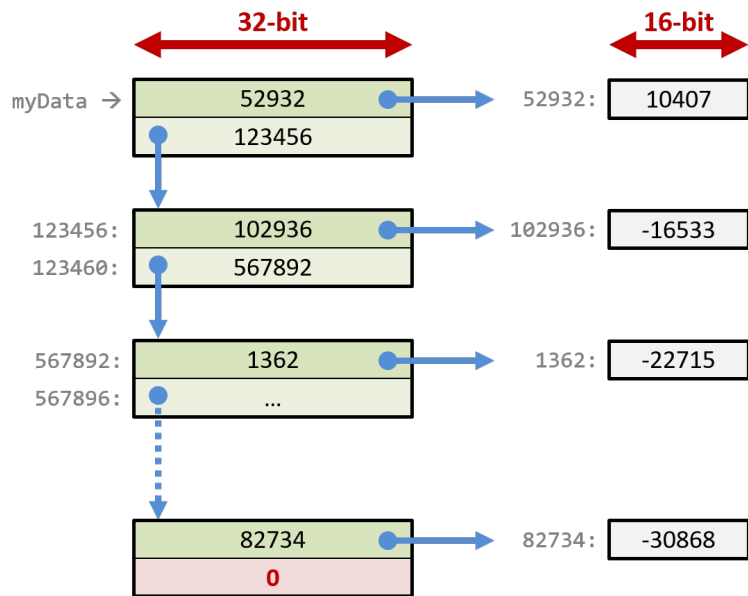
Index in array

```
short sum = 0;
int i;
for (i = 0; i < N; i++) {
    if (myData[i] > 0) {
        sum += myData[i];
    }
}
```

Write **better code** but get the **best performance**!

One More Way to Store Arrays





Not Much More Complex...

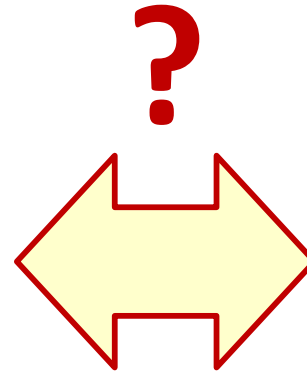
Pointer to memory

```
add_positive:
    li    t0, 0
    mv    t1, a1

next_short:
    beq    t1, zero, end
    lh     t2, 0(a0)
    bltz   t2, negative
    add    t0, t0, t2

negative:
    addi   a0, a0, 2
    addi   t1, t1, -1
    j      next_short

end:
    mv     a0, t0
    ret
```



Linked list

```
add_positive:
    li    t0, 0

next_element:
    beqz   a0, end
    lw     t1, 0(a0)
    lh     t1, 0(t1)
    bltz   t1, negative
    add    t0, t0, t1

negative:
    lw     a0, 4(a0)
    j      next_index

end:
    mv     a0, t0
    ret
```

Multiple noncontiguous loads

References

- Patterson & Hennessy, COD – RISC-V Edition
 - **Chapter 2** and, in particular, **Sections 2.9, 2.13, and 2.14**