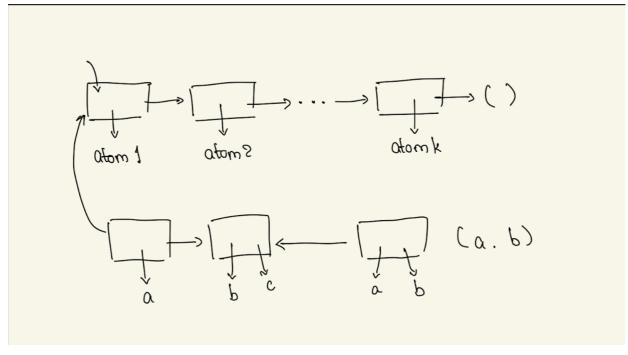
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## Class - CSc 335

## Date - April 18, 2023 (Tuesday)



display as (atom1, atom2,...,atom k) and can be created in a number of ways eg.

```
'(atom1, atom2,...,atom k)
(list atom1, atom2,...,atom k)
```

Data described y (finite list of sexp)

We want a program lat? which checks whether an sexp s is a lat, where

MAIN IDEA: Structure of a program, which processes sexp must mirror the structure of sexps.

lat? is supposed to be a function from sexp to B - the set {#f, #t}

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Several inductive description are possible

- 1. Induction on the length of the sexp, where we agree that the length of an atom is 0.
- 2. We do induction on the number of cons.
- 3. Leverage the BNF definition of sexp as also a definition of component of an sexp, and then use structural induction.

Recall: The IH for a structural induction is to assume the result holds for all proper components of the input.

## **Induction Proof**

• Easily given, say by induction on the length of lat.

## What is GI?

• A possible GI

```
\circ \ a \in LAT \iff a \in lat
```

• Strong enough? We need to check that GI & stopping condition ⇒ post condition.

```
GI && ((null? lat) OR ((car lat) = a)) => returns #f
```

- because GI and (null? lat)  $\Rightarrow a \in LAT$
- Preserved?\_
  - we assume GI true on current call. So the only way  $a \in LAT$  is either

```
1. a = (car \ lat) or 2. a \in (cdr \ lat)
```

- pre: a is an atom, lat is a lat
- post: returns  $a \in lat$

remove an occurence of a from I, leaving I otherwise unchanged.

we are removing the first occurence

As often happens the program is over specified relative to the spec.