Class: CSc 335

Date: Feb 2, 2023 (Thursday)

### **QUIZ NEXT THURSDAY**

• Enbironment diagram

### **Lambda Function**

- creating a closure
  - closure has 3 parts → reserved word, formal and body

```
(lambda (x) (* x x))
> #<procedure>
```

• to use lambda

```
((lambda (x) (* x x)) 8)
> 64
```

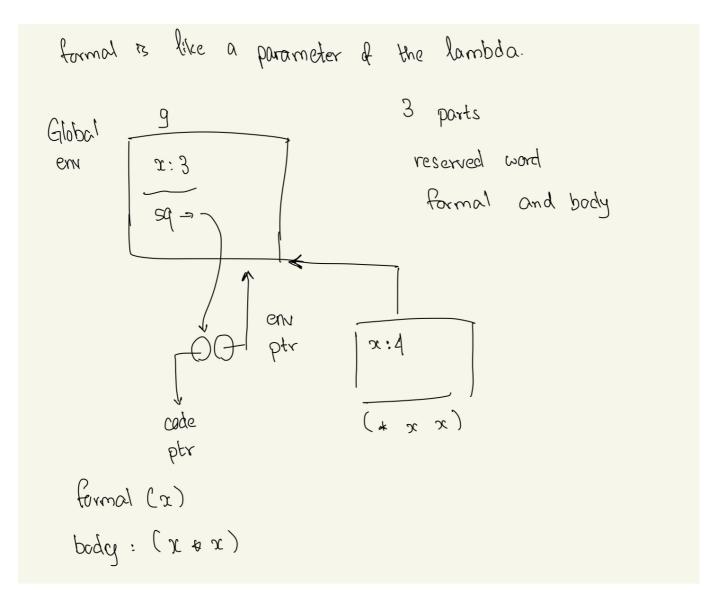
• with two parameters

```
((lambda (x y) (+ x y)) 3 4) > 7
```

• When you use define, you're installing variables in global env (knows all about scheme, but not about x or sq.)

Global Env	Define
x: 3	(define x 3)
sq → ?	

(define sq (lambda (x) (\* x x)))  $\rightarrow$  creates a closure



• A&S use a sugared form of function definition

```
eg. (define(sq x) (* x x))
```

### QUESTION: What happens when (sq 4) is evaluated, given g as shown?

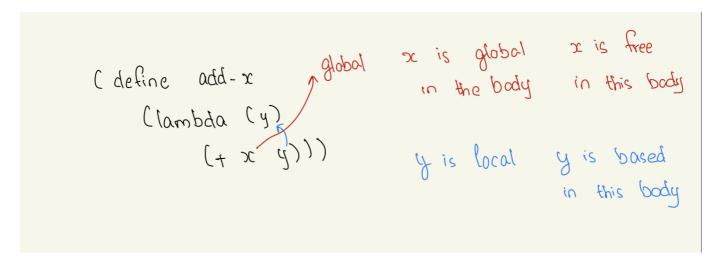
### **ANSWER:**

- 1. Create a new frame with the same enclosing frame as sq
- 2. In the new frame, bind the formal parameter x to the value of the actual parameter 4.
- 3. Evaluate the body of sq in the frame. This evaluation entails same lookups
  - \* is not defined in the new frame so the system looks in the containing frame and finds the primitive x.
  - Take the first value of x found in any of the enclosing frames → in this case, 4. Note that x=3 in g is ignored for this evaluation → one says that it is shadowed. (This is the essence of the implementation of local variables)
- 4. The value (\* 4 4) = 16 is returned.
- 5. Next Garbage collector kicks in → it stores the newly allocated frame (memory) to the keys because there's no directed path from g to the new frame. The contents of the new frame are

immediately inaccessible or disappeared.

# **Functions**

- functions in scheme are First Class
  - o can be returned by a function call
  - $\circ$  can occur in data structure  $\rightarrow$  we can, for example, have lists of functions



### **QUIZ ON THURSDAY FEB 9**

## • Environemnt diagram for sure

