

Class: CSc 335

Date: Mar 2, 2023 (Thursday)

AM Quiz problem without the multiplicity restriction

- Let's write $m \leq n$ to abbreviate `(contains? n m)`
- We want, e.g. $333 \leq 3$ to be true - essentially the mapping requirement we imposed on the first solution, but without the one-to-one requirement.
- Again, let m_0 be the rightmost digit in m and n_0 the rightmost digit in n .

What do we learn by comparing m_0 and n_0 ?

- if $m_0 = n_0$, then $m \leq n$ iff `(quotient m 10) ≤ n` OR if $m < 10$
- if $m_0 \neq n_0$, then just as for the first solution we see that no "m-pattern" can begin with n_0 , so $m \leq$ `(quotient n 10)`
 - i.e., there exists a possibly many-to-one mapping from the digits of m to those of n .

This may be all we need to say about divide & conquer, but we still must address the stopping condition(s).

What if $n < 10$, so $n = n_0$, and yet $n_0 \neq m_0$?

- Must in this case return `#f` - $m \leq n$ is not true.

We maintain an open mind as regards the stopping condition(s).

Code:

```
(define (contains? m n)
  (let ((m0 (modulo m 10))
        (n0 (modulo n 10)))
    (cond ((= m0 n0)
           (if (< m 10)
               #t
               (contains? (quotient m 10) n)))
          (else (contains? m (quotient n 10))))))
```

- why not just `(or (< m 10) (contains? (quotient m 10) n))`?
 - should this call be evaluated first, and m is in fact < 10 , could get incorrect result?

Test

- Worry even before testing about termination, because there is no explicit basis case shown in the code

PM Class Quiz 2

Input $n \geq 0$ an integer and a digit d and return the number obtained from n by removing the left-most occurrence of d .

eg. `(remove-leftmost 121314 1) = 21314`

NOTE: $n \neq d$ should be added to pre-condition. Anything else?

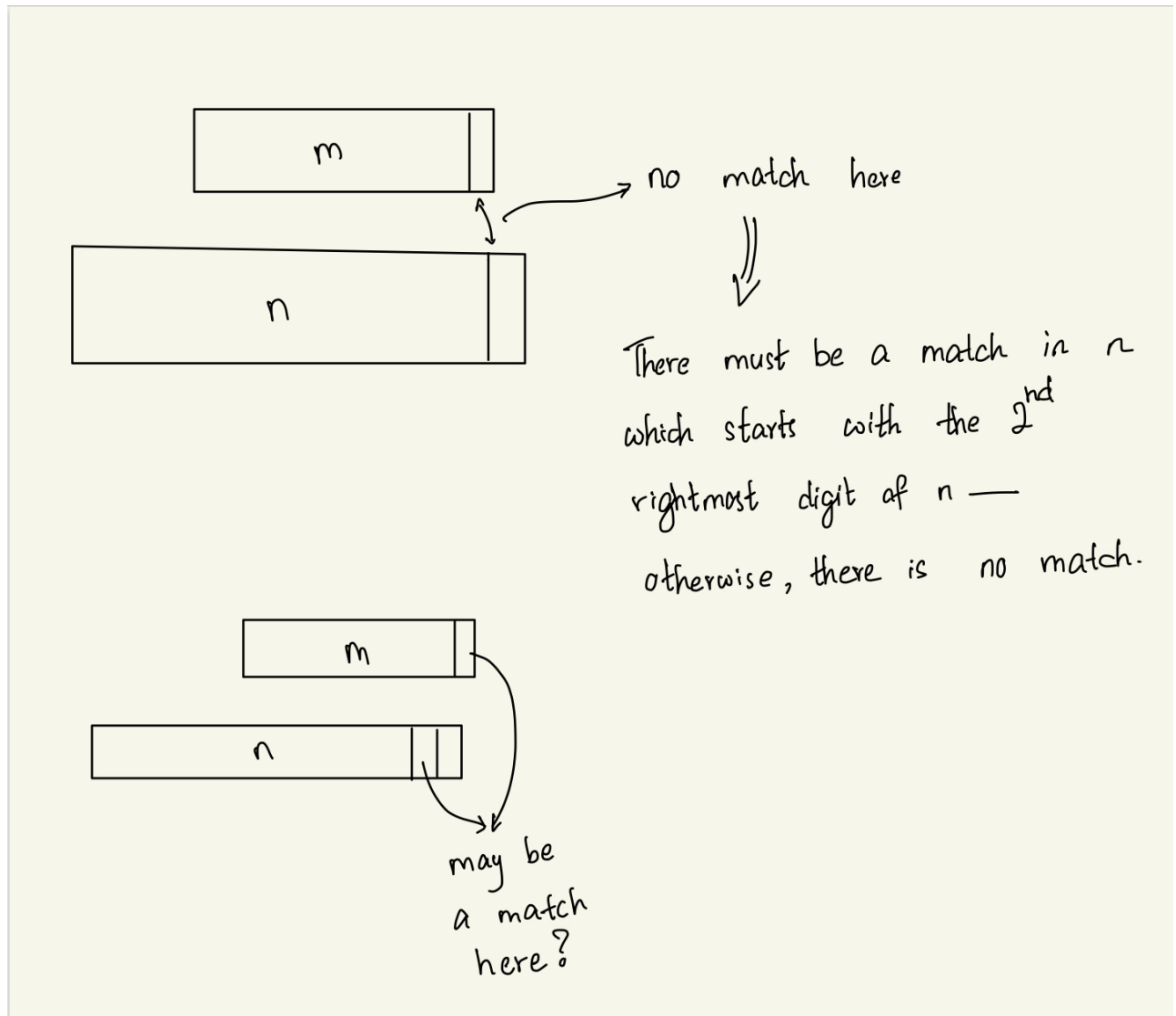
1. One design would make use of a function `reverse-digits` and then to remove the rightmost occurrence of d , and then to reverse again: $121314 \rightarrow 413121 \rightarrow 41312 \rightarrow 21314$.
 - Note: removal of rightmost d is easy.
2. Another design: scan n from the right, making use of a function
 - `(occurs? n d)` which determines whether d occurs in n .

```
(defne (remove-leftmost n d)
  (cond ((= d (modulo n 10))
    (if (not (occurs? (quotient n 10) d))
      (quotient n 10)
      (+ (modulo n 10)
        (* 10 (remove-leftmost (quotient n 10) d)))))
    (else (+ (modulo n 10)
      (* 10 (remove-leftmost (quotient n 10) d)))))
```

- Please use `let` to remove the overuse of `quotient` and `modulo` for repeated tasks.

What about `(occurs? n d)`?

- an iterative design idea: scan n from the right



- GI: d occurs in n iff d occurs in not-yet-processed (nyp)

Design Idea for Reversing n :

$$\left(\text{reverse } \overbrace{\hspace{1.5cm}}^{\text{(quotient } n \text{ } 10)} \mid n_0 \right)$$

some power $\rightarrow (\# \text{ digit in } n) - 1$

$$n_0 * 10 + (\text{reverse (quotient } n \text{ } 10))$$