

Class: CSc 335

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Strong induction

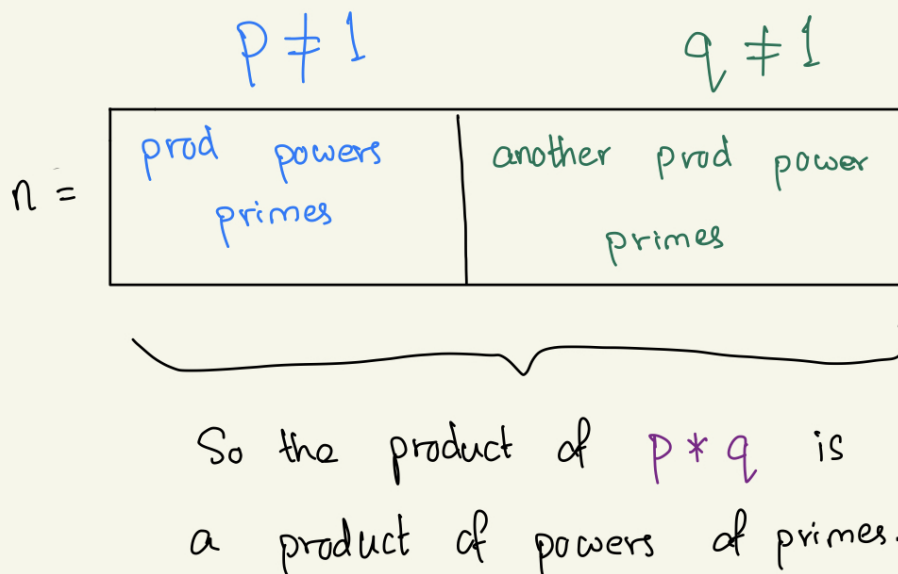
Motivation: We will need strong induction when we begin discussing data structures.

Prove: every integer $n \geq 2$ can be written as a product of powers of primes (proof by strong induction)

Try inducting on n , with divide & conquer: if $n = p * q$, with neither p nor q equal to 1. Then assume p is a product of powers of primes and also that q is a product of powers of primes.

Now that we have the divide & conquer strategy, here is the full argument

- Consider $n \geq 2$
 - If $n = 2$, there's nothing to show, since $2 = 2^1$ is a product of powers of primes.
- Consider $n > 2$
 - Either n is prime or not n i.e. composite
 - If n is prime, we're done.
 - Otherwise, $n = p * q$ for some integers $p, q > 1$ [def. of composite number]
 - Now the argument is the one we just gave.



This is a strong induction because

- We know neither p nor q is $n-1$
- We don't know in advance the values of the factors p and q - So we actually need all of the hypotheses.

Tree Recursion

```
;; Fib-sequence

(define (fib n)
  (cond ((zero? n) 0)
        ((one? n) 1)
        (else (+ (fib (- n 1)) (fib (- n 2))))))
```

- a cool technique is **memoization** → to transform the results of a function into something to remember
 - example: creating a dictionary to remember the value of **fib 2** or **fib 3** in computing **fib 5**

What about an iterative solution?

A design idea flows from the def of $\text{fib}(n+1)$ as $\text{fib}(n) + \text{fib}(n-1)$ when $n \geq 2$. where

```
fib(n+1) = fib(n) + fib(n-1)
  |         |         |
  next      = curr  + prev
```

If we introduce **curr** and **prev** as variables with design roles:

```
curr = (fib n)
prev = (fib (- n 1))
```

with the idea of maintaining these with as count increases from 0 up to n .

i.e.: Take the GI as the logical AND of the design roles:

- GI:

```
next = curr + prev
where next = fib(count)
```