

APPS@UCU.

Linux course

Linux File systems

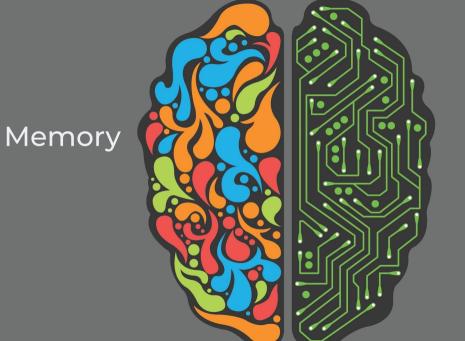
Morhunenko Mykola

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Intro

- This is not an overview of some hardware memory staff
- Neither a presentation with deep File systems implementation details
- More about that you should learn at the Operating systems course
- This is just an overview of file systems that system administrators use in their everyday life
- If you think that you are not a system administrator think one more time, because you administrate your own system every day



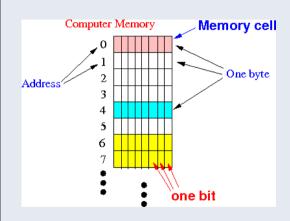
Drives

- All data stored on some physical devices
- It has different storage approaches on each device (HDD, SSD, CD, DVD, Flash, RAM, DDR memory modules)
- But now we are going to overview the memory from user point of view
- How to manage files and file systems, how to choose the most suitable



Memory storage

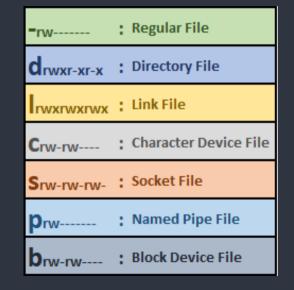
- Memory as abstraction looks like an array, where bites are stored one by one in a row
- File system a method of data structure that the operating system uses to control how data is stored and retrieved
- A file is an ordered collection of data blocks
- In Linux system, everything is a file, and if it is not a file, it is a process
- So File systems are very important for this OS





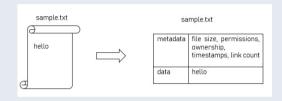
File types

- There are a lot of file types, but the most important for us are:
- Regular Files some files with data stored inside
- Directories files, that allowed to group other files and keep tree filesystem structure
- Character files for simulating character devices as terminals, keyboard, network etc
- Block files for modelling block devices as disks, flash drives
- Links entry points to other files
- There are pipes, sockets



File metadata

- File also save a metadata about itself, as:
- Protection, password
- Creator, owner
- Flags (r w x)
- Size
- Creation time, last update time (timestamp)

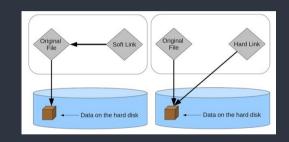


Inode

- The inode is data structure, that describes files on Unix-like OS's
- Each inode stores disk block location, some atributes, file's metadata
- Directory just a file with list of inodes
- File's inode number can be found with Is -i command
- From the inode number, the kernel's file system driver can access the inode contents, including the location of the file, thereby allowing access to the file
- More about inodes in the Operating systems course

Links

- There are two types of links: symbolic (soft) and hard
- They are totally different types of file
- Maybe first few years you will not use
- But with experience it becomes more and more useful
- Here we will make only a brief overview and comparison



Hard links

- Exact replica of a file
- Share same inode with other hard links
- Can not be made across filesystems
- Changes in hl will reflect in other files
- Deleting of a hardlink wil not affect other files
- Can links to files only

Soft links

- Alias to a file
- Has another inode
- Can be established outside filesystem
- Link becomes inaccessible without original file
- Can links to both files and directories

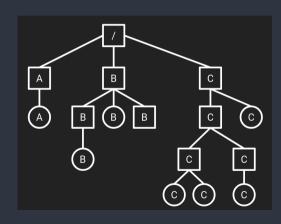


Fyle systems types overview

- There are several file systems types. Just for your information. the most important will be in orange colour
- Disk file systems for simple disks, a.e. FAT16/32, NTFS, ext2-4, brtfs etc
- Flash file systems consider speciality of flesh memory devices
- Database file systems another concept for file management
- Transactional file systems
- Network file systems acts as a client for a remote file access protocol, providing access to files on a server, a.e. FTP
- Shared disk file systems a number of machines (usually servers) all have access to the same external disk subsystem
- Flat file systems no subdirectories, directory entries for all files are stored in a single directory

Fyle system abstraction

- We used to see a filesystem as a tree. It is really the most comfortable structure as for now
- There is a CLI tool to see your filesystem structure called tree
- Using such abstraction programmer works with files and directories, not with memory cells or some low level staff, but with files, directories and subdirectories





Review of previous topics

- It's part of presentation abot shell, but let's make a brief overview
- Every proces has it's own working directory.
- pwdx \$(pgrep process_name) show working directories for a process_name
- pwd print working (current) directory
- Is list what is instde the working directory
- cd change directory
- ./ spesial, current directory
- ../ spesial, parrent directory (in a tree struct)
- ~/ \$HOME directory for current user
- cp <from....> <to> copy
- mv <from....> <to> rename (inside one fs, or move from one to another fs)
- mkdir make directory
- touch <filename> update the last acces date (if no such file create)
- rm <filename> remove
- cat show the file content

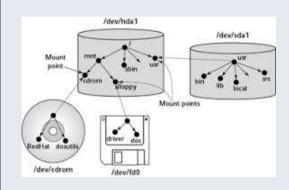
Devices

- Everything is a file, devices are not an exception
- All devices are in /dev folder
- Devices could be either secondary storages or mous, keyboard, terminals, cpu, gpu etc
- Devices can be either block or character devices
- Easy to remember:
- Block devices store or hold data
- Character devices transmit or transfer data



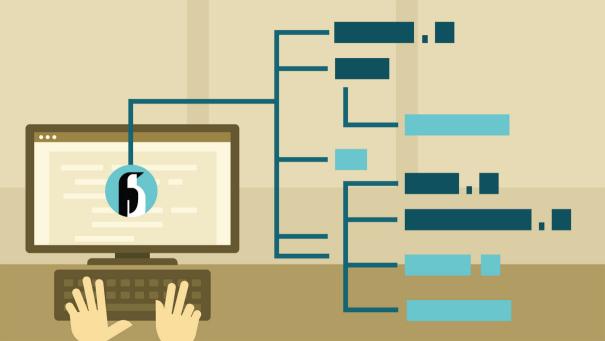
Mounting

- Mounting attaching some additional fs to already mounted
- By default, user use only one filesystem, and it's mountpoint is /
- Then /boot is also another fs, that is not used by user directly (more about that bootloader topic)
- On one physical device there could be few filesystems (more about that partition tables topic)



A smorgasbord of important commands and

- mkfs make a file system on a device
- mkfs.filesystemtype /dev/X create a filesystem on existing logical device
- fdisk -l to see all devices available for mounting
- mount /dev/X /mnt to mount device. mnt is used as a convension, and it is important
- /etc/fstab file with all "default mountings" during startup. Usually has only / and boot fs's
- UUID Universally unique identifier 128bit label. The probability that a UUID will be duplicated is close enough to zero to be negligible. HERE - unique device identifier



Linux Filesystems Hierarchy (LFS)

This topic worth a separate lecture, but lets make it short

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Sources

Sources

- UCU Linux Club
- File systems Wiki
- Linux file systems
- Differences between hard and soft links on Unix systems
- Mounting and unmounting on Linux
- UUID Wiki