

Ant Algorithms: Part 1

Ant Algorithms

Ant's are social insects, living in colonies of 30 to millions of individuals.

- Current total ant population is estimated at 10^{16} individuals
- The complex behaviours that emerge from colonies of ants have intrigued humans, and there have been many studies of ant colonies aimed at a better understanding of these collective behaviours
- Some interesting and well studied behaviours are:
 - ▶ foraging behaviour
 - ▶ division of labour
 - ▶ cemetery organization and brood care
 - ▶ construction of nests

Ant Algorithms

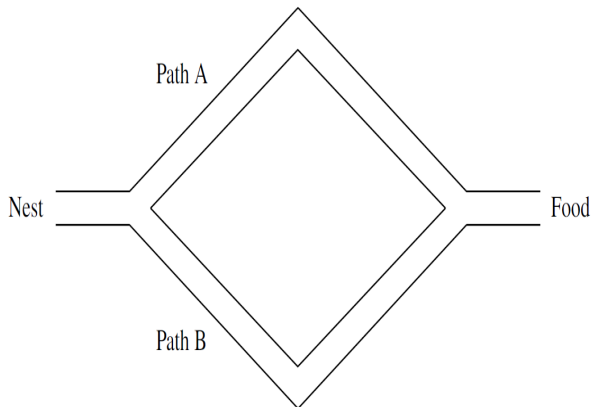
Most (though not all) ants communicate using some form of indirect communication, which was termed as **stigmergy** by the biologist Grassé.

- The indirect communication is often achieved via an interaction with the environment
- The most classic example of **stigmergy** is that of **pheromone** trails.
 - ▶ The first algorithmic models of foraging behaviour used pheromones
 - ▶ There are two clear advantages of this type of communication
 - ★ One insect need not know the location of the others with whom it is communicating.
 - ★ The communication will potentially outlast the sender. (say the ant died)

- In the context of collective behaviour, social insects are basically **stimulus-response agents**.
 - ▶ Based on information perceived from the local environment, an individual performs a simple, basic action. Often with a random subcomponent.
 - ▶ Despite this simplicity in individual behavior, social insects collectively form a highly structured social organism.
 - ▶ Emergent behavior can arise from a very simple set of rules.

Ant Algorithms

The origins: Pasteels tried to mathematically model the behavior of the ant species: *Iridomyrmex humilis* in the famous binary bridge experiment



Ant Algorithms

With the knowledge that the ants drop a pheromone trails, Pasteels wanted to model which path the ants would eventually favor. Empirically it was derived that the guiding equation was simply:

$$P_A(t+1) = \frac{(c + n_A(t))^\alpha}{(c + n_A(t))^\alpha + (c + n_B(t))^\alpha} = 1 - P_B(t+1) \quad (1)$$

where

- $n_A(t)$ and $n_B(t)$ ants on paths A and B respectively at time step t .
- c quantifies the degree of attraction of an unexplored branch
- α is the bias to using pheromone deposits in the decision process.

For example: $n_A(t) = 5, n_B(t) = 100, c = 20$, and $\alpha = 2$ then

$$P_A(t+1) = 0.04 \quad (2)$$

which is to be expected as path A has very few ants on it.

Ant Algorithms

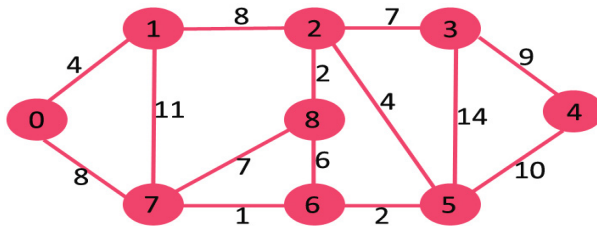
What is more interesting in-terms of the application of AA to optimization is the extended experiment by Goss *et al* where one of the two bridges is longer than the other.

- It was found that given enough time the ants favored the shorted path, with a substantially larger percentage of ants choosing to follow the shortest path.
- Selection is biased towards the shortest path, since ants that follow the shortest path return to the nest earlier than ants on the longer path.
 - ▶ The pheromone on the shorter path is therefore reinforced sooner than that on the longer path.
- Goss et al. found that the probability of selecting the shorter path increases with the **length ratio** between the two paths. Referred to as the differential path length effect

Ant Algorithms

We will now go over a simplified version of the original ant colony optimization-meta-heuristic (ACO-MH), called simple ACO (SACO). The intention is to get you used to the ACO concepts in a simple context.

- SACO and ACO-MH are predominantly used for solving types of combinatorial problems.
- Let us consider the problem of finding the shortest path between two nodes in graph.



Ant Algorithms

The following are the main components of SACO.

- Begin by assigning a small amount of pheromone to all links, $\tau_{i,j}$.
- n_k ants are placed at the source node.
- Each ant must then attempt to construct a path from the source to the destination but using the following equation to guide it's transition probability

$$P_{i,j} = \begin{cases} \frac{\tau_{i,j}^\alpha(t)}{\sum_{e \in \mathcal{N}_i^k} \tau_{i,e}^\alpha} & \text{if } j \in \mathcal{N}_i^k \\ 0 & \text{if } j \notin \mathcal{N}_i^k \end{cases} \quad (3)$$

where \mathcal{N}_i^k is set of feasible nodes connected to node i , with respect to ant k . α is a positive constant used to amplify the influence of pheromone concentrations.

Ant Algorithms

- If for any node i and ant k , if $\mathcal{N}_i^k = \emptyset$, then the predecessor to node i is included in \mathcal{N}_i^k
 - ▶ Note that this may cause loops to occur within constructed paths.
 - ▶ These loops are removed once the destination node has been reached.
- Once all ants have constructed a complete path from the origin node to the destination node, and all loops have been removed, each ant retraces its path to the source node deterministically, and deposits a pheromone amount,
 - ▶ The amount dropped should be inversely proportional to the lengths of the found path.

$$\Delta\tau_{i,j}^k(t) \propto \frac{1}{L^k(t)} \quad (4)$$

- ▶ Shorter paths \implies more pheromone dropped.

Ant Algorithms

- This implies that after all ants have dropped their pheromone trails we have

$$\tau_{i,j}(t+1) = \tau_{i,j}(t) + \Delta\tau_{i,j}(t) \quad (5)$$

where

$$\Delta\tau_{i,j}(t) = \sum_{k=1}^{n_k} \Delta\tau_{i,j}^k(t) \quad (6)$$

- If the pheromone trail lasted forever it is likely that paths will become over saturated, and force the ants to over exploit found solutions.
 - ▶ To avoid this an evaporation rate is added ρ .
 - ▶ **Before** the ants lay new pheromone trails the current pheromone level is updated with

$$\tau_{i,j}(t) = (1 - \rho)\tau_{i,j}(t) \quad (7)$$

with $\rho \in [0, 1]$

Ant System

The first ant algorithm was developed by Marco Dorigo, who has over 87000 citations, and is referred to as the **ant system**.



Ant System

The SACO algorithm can be greatly improved with just a few changes. The most important change proposed in ant system (AS) are

- The inclusion of heuristic information into the transition probability.
- specifically

$$P_{i,j} = \begin{cases} \frac{\tau_{i,j}^{\alpha}(t) \eta_{i,j}^{\beta}(t)}{\sum_{e \in \mathcal{N}_i^k} \tau_{i,e}^{\alpha} \eta_{i,e}^{\beta}(t)} & \text{if } j \in \mathcal{N}_i^k \\ 0 & \text{if } j \notin \mathcal{N}_i^k \end{cases} \quad (8)$$

where $\eta_{i,j}^{\beta}(t)$ indicates the desirability of the moves, which is computed based on heuristic information.

- The transition probability used by AS is a balance between pheromone intensity and heuristic information. The balance is controlled by α and β

The other change is the inclusion of a memory.

- Specifically, the inclusion of a tabu list was proposed by Dorigo.
- A list of all visited nodes are stored in a tabu list. Specifically the sets \mathcal{N}_i^k will exclude any node on the tabu list.

Ant System

There where originally three ways of calculating $\Delta\tau_{i,j}^k(t)$

- Ant-cycle AS:

$$\Delta\tau_{i,j}^k(t) = \begin{cases} \frac{Q}{f(x^k(t))} & \text{if link } (i,j) \text{ occurs in path } x^k(t) \\ 0 & \text{otherwise} \end{cases} \quad (9)$$

Q is a positive constant, and $f(x^k(t))$ is the quality of the complete path constructed by the ant. (if we are minimizing)

- Ant-density AS:

$$\Delta\tau_{i,j}^k(t) = \begin{cases} Q & \text{if link } (i,j) \text{ occurs in path } x^k(t) \\ 0 & \text{otherwise} \end{cases} \quad (10)$$

- Ant-quantity AS:

$$\Delta\tau_{i,j}^k(t) = \begin{cases} \frac{Q}{d_{i,j}} & \text{if link } (i,j) \text{ occurs in path } x^k(t) \\ 0 & \text{otherwise} \end{cases} \quad (11)$$

$d_{i,j}$ is the length of the link.

Other addition to the AS is to add the concept of elitism

- Specifically, the pheromone update equation changes to

$$\tau_{i,j}(t+1) = \tau_{i,j}(t) + \Delta\tau_{i,j}(t) + n_e\Delta\tau'_{i,j}(t) \quad (12)$$

where

$$\Delta\tau'_{i,j}(t) = \begin{cases} \frac{Q}{f(\hat{x}(t))} & \text{if } (i,j) \in \hat{x}(t) \\ 0 & \text{otherwise} \end{cases} \quad (13)$$

an n_e is positive constant controls the strength of the elitist force.

- The elitist strategy has as its objective directing the search of all ants to construct a solution to contain links of the current best route.

Ant Colony System

The ant colony system (ACS) was developed by Gambardella and Dorigo to improve the performance of AS. ACS differs from AS in four aspects

- a different transition rule is used
- a different pheromone update rule is defined
- local pheromone updates are introduced
- candidate lists are used to favor specific nodes (added in scalability)

Ant Colony System

The ACS transition rule, also referred to as a pseudo-random-proportional action rule is:

- Ant k , currently located at node i , selects the next node j to move to using the rule,

$$j = \begin{cases} \operatorname{argmax}_{u \in \mathcal{N}_i^k(t)} \tau_{i,u}(t) \eta_{i,u}^\beta(t) & \text{if } r \leq r_0 \\ J & \text{if } r > r_0 \end{cases} \quad (14)$$

where $r \sim U(0, 1)$, and $r_0 \in [0, 1]$ is user-specified parameter

- $J \in \mathcal{N}_i^k(t)$ is a node randomly selected using according to probability

$$p_{i,J}^k = \frac{\tau_{i,J}^\alpha(t) \eta_{i,J}^\beta(t)}{\sum_{u \in \mathcal{N}_i^k} \tau_{i,u}^\alpha(t) \eta_{i,u}^\beta(t)} \quad \text{if } j \in \mathcal{N}_i^k \quad (15)$$

\mathcal{N}_i^k is a set of valid nodes to visit from i .

Ant Colony System

The pheromones updates rule changes such that

- Only the “best” ant lays down new pheromones.
- let $\hat{x}(t)$ represent the best path, then the global pheromone update rule becomes

$$\tau_{i,j}(t+1) = (1 - \rho_1)\tau_{i,j}(t) + \rho_1\Delta\tau_{i,j}(t) \quad (16)$$

where

$$\Delta\tau_{i,j}(t) = \begin{cases} \frac{Q}{f(\hat{x}(t))} & \text{if link } (i,j) \text{ occurs in path } \hat{x}(t) \\ 0 & \text{otherwise} \end{cases} \quad (17)$$

- The ACS global update rule causes the search to be more directed, by encouraging ants to search in the vicinity of the best solution found thus far.

Ant Colony System

Gambardella and Dorigo implemented two methods of selecting the “best” path:

- **iteration-best**, where $\hat{x}(t)$ represents the best path found during the last completed iteration, t
- **global-best**, where $\hat{x}(t)$ represents the best path found from the first iteration of the algorithm.

An additional local update rule is used when an ant traverses a link, specifically

$$\tau_{i,j}(t) = (1 - \rho_2)\tau_{i,j}(t) + \rho_2\tau_0. \quad (18)$$

where both ρ_2 and τ_0 are positive control parameters. It was found that $\tau_0 = \frac{1}{n_g L}$ was effective for TSP problems. (n_g =number of nodes, L is the shortest tour produced by the nearest neighbor heuristic). Both are however problem dependent

Ant Colony System

The last change present in ACS is the use of a candidate list.

- This was a TSP specific addition.
 - ▶ when trying to solve big TSP problems it is common practice to use a data structure known as a candidate list.
 - ▶ A candidate list is a list of preferred cities to be visited;
 - ★ it is a static data structure which contains, for a given city , the closest cities, ordered by increasing distances;
 - ▶ IN ACS's introduction a size of 15 was used.

This was a way of inserting heuristic (bias) to speed up the process that was specific to TSP. Full detail can be found at <http://people.idsia.ch/~luca/acs-ec97.pdf>

Max-Min Ant System

A variant to AS was proposed by Stützle and Hoos. The primary changes in max-min ant system (MMAS) from AS are

- Pheromone intensities are restricted within given intervals.
- Only the best ant may reinforce pheromones (a mixed approach was used in subsequent versions)
- Initial pheromones are set to the max allowed value
- a pheromone smoothing mechanism is used.

Max-Min Ant System: Pheromone Interval

- Set a minimal amount of pheromone τ_{min} allowed on all links.
 - ▶ This ensures that all links have a non-zero chance of being selected.
 - ▶ Forces some minimum amount of exploration.
- Set a maximum allowable amount of pheromone τ_{max}
 - ▶ This prevents a path becoming over saturated, which would cause all ants to follow the same path.
 - ▶ Prevents over exploitations

Max-Min Ant System: Trail smoothing

While MMAS was less susceptible to stagnation than AS, it still occurred. As a result Stützle and Hoos added the following mechanism

- If **stagnation** was detected then all arcs (i,j) undergo “trail smoothing” by adding

$$\Delta\tau_{i,j} \propto \tau_{max} - \tau_{i,j} \quad (19)$$

namely, all pheromone concentrations are increased proportional to the difference with the maximum bound.

- Using this smoothing strategy, stronger pheromone concentrations are proportionally less reinforced than weaker concentrations.

Max-Min Ant System: Stagnation

The stagnation detection method was simple: if the mean λ -branching factor ($\bar{\lambda}$) dropped below a very small predefined constant the search was in a state of stagnation: where

$$\bar{\lambda} = \frac{\sum_{i \in V} \lambda_i}{N_G} \quad (20)$$

where λ_i is defined as the number of links leaving node i with τ_{ij} values greater than $\lambda(\tau_{i,max} - \tau_{i,min}) + \tau_{i,min}$, where

- Let $\tau_{i,max}$ and $\tau_{i,min}$ be the max and min pheromone intensity of links leaving node i .

Fast Ant System

The fast ant system (FANT) was designed to solve the quadratic assignment problem (QAP). A QAP is a placement problem. There are many ways to understand the QAP but the simplest is in graph form.

- There are n facilities that need to be assigned to n locations.
- There exists a link between each location pair, which contains a distance (could be some other cost)
- There exists a link between each facility pair, which contains a flow (how much product needs to be moved between the facilities)
- Placement of facilities must be made to minimize the sum of products between the distance and flow on each arc.
- *In the above description a 0 weight is how the lack of an arc is modeled

Fast Ant System

The two main difference between FANT and AS are:

- FANT uses only one ant.
- A different pheromone update rule is applied which does not make use of any evaporation.

The pheromone update rule is:

- Let $l(t)$, and $g(t)$ be iteration best and global best paths respectively
- Then

$$\tau_{i,j}(t+1) = \tau_{i,j}(t) + w_1 \Delta_{i,j}^{(l)} + w_2 \Delta_{i,j}^{(g)} \quad (21)$$

where

$$\Delta_{i,j}^{(l)} = \begin{cases} 1 & \text{if } (i,j) \in l(t) \\ 0 & \text{otherwise} \end{cases} \quad (22)$$

$$\Delta_{i,j}^{(g)} = \begin{cases} 1 & \text{if } (i,j) \in g(t) \\ 0 & \text{otherwise} \end{cases} \quad (23)$$

Fast Ant System

- $\tau_{i,j}(0)$ is initialized to 1
- If a new global best path is found $\tau_{i,j}$ is reinitialized.
- If a iteration yields no improved global best path, w_1 is increased by 1 to promote exploration.