MPI: Microprocessors and Interfacing Academic Year: 2022 - 23

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Department of Computer Science and Engineering Ecole Centrale School of Engineering



Lab-6 Due Date: November 24, 2023

- 1. Write an assembly language program to perform the addition of two $N \times N$ matrices.
- 2. Write an assembly language program to perform the multiplication of two N \times N matrices.

Note: Use NASM and consider the value of N during run time.

Lab-5 Due Date: November 05, 2023

Take any five problems that we discussed in the first four lab assignments (i.e., Lab-1 to Lab4) and implement them using the MASM assembler.

Lab4- Due Date: October 20, 2023

Assume that you have a processor (P) and it supports **ADD** (Addition), **SUB** (Subtraction), **INC** (Increment), **DEC**, **CMP** (Compare), **Jcc** (Jump if a condition is met), **JMP**, **XOR**, and **MOV**operations. Assume that A and B are positive integers and A > B.

- 1. Write an assembly language program to compute the $(A^2 B^2)$ on P.
- 2. Write an assembly language program to compute the quotient of $\frac{(A^3-B^3)}{(A+B)}$ on P.
- 3. Compute A choose B (i.e., $\binom{A}{B}$) on P

Lab3- Due Date: September 30, 2023

You are given with n positive Integers. Write functions to compute the following using **assembly language instructions**.

- 1. To find the number of even numbers
- 2. To find the number of odd numbers
- 3. To find the number of prime numbers
- 4. To find the GCD of *n* numbers
- 5. To find the LCM of *n* numbers

Lab2- Due Date: September 15, 2023

Write assembly language programs to compute the following (use NASM assembler). Assume that all elements are integers.

- 1. Find the **maximum** of *n* numbers.
- 2. Find the **minimum** of *n* numbers.
- 3. Assume that you are given with a list of *n* elements then find the **mode**.
- 4. Assume that you are given with a list of *n* elements then find the **median**.
- 5. Assume that you are given with a list of *n* elements and **key**, search the **key** using **linear search**. If the required **key** is found, return its position. Otherwise, return -1.

Lab1 - Due date: August 31, 2023.

- 1. Give the configuration of your laptop.
- 2. Give the configuration of a system in IT-Lab.
- 3. Develop C-Programs for the following problem statements:
 - 3.1. Print the internal representation of data stored in primary data types: **int, float, and double**.
 - 3.2. Perform addition and multiplication of two 32-bit numbers (Please remember that the input is a 32-bit binary number). The numbers can either be integers or real numbers.

Submission Guide Lines

Max. team size is 6.

Mail-ID: cs3106.mpi@gmail.com

Sub:TEAM_NUM_LAB_NUM

Attach.Name and Type: (Sub.).zip

Late Submission:50%.

Write a readme file to understand your solutions.

Number Systems

Representation of Integer Numbers
Signed Magnitude Representation

1's Complement Representation

2's Complement Representation

Representation of Real Numbers Fixed Point Representation

Floating Point Representation

Resolution is difference between two successive numbers.

Representation of Integer Numbers

Let $A = a_{n-1}a_{n-2}a_{n-3}...a_0$ is an n-bit binary number if A is an **unsigned integer**, then value of A is : $\sum_{i=0}^{n-1} (2^i \times a_i)$. if A is a **signed integer**:

Signed Magnitude Representation:

$$A = \sum_{i=0}^{n-2} (2^i \times a_i)$$
, if $a_{n-1} = 0$
 $A = -\sum_{i=0}^{n-2} (2^i \times a_i)$, if $a_{n-1} = 1$

1's Complement Rep.:
$$A = -(2^{n-1} - 1) \times a_{n-1} + \sum_{i=0}^{n-2} (2^i \times a_i)$$

2's Complement Rep.:
$$A = -2^{n-1} \times a_{n-1} + \sum_{i=0}^{n-2} (2^i \times a_i)$$

Resolution: 1

Range of Numbers

Let $A = a_{n-1}a_{n-2}a_{n-3}...a_0$ is an n bit binary number if A is an **unsigned integer**, then range of A is : 0 to $(2^n - 1)$. if A is a **signed integer**:

Signed Magnitude Rep., range of A is : $-(2^{n-1}-1)$ to $(2^{n-1}-1)$.

1's Complement Rep., range of A is : $-(2^{n-1}-1)$ to $(2^{n-1}-1)$.

2's Complement Rep., range of A is : -2^{n-1} to $(2^{n-1}-1)$.

```
Add additional bit positions to the left and fill in with value of the sign bit. Let A=1\ 0\ 1\ 0 is a 4-bit binary number, Representation of A using 8-bits (i.e. B): 1\ 1\ 1\ 1\ 1\ 0\ 1\ 0. is A=B?

In 2's Complement Rep.: Yes.
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In 2's Complement Rep.: Yes.

In 1's Complement Rep.: Yes.
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In Signed Magnitude Rep.: No.

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Add additional bit positions to the left and fill in with value of the sign bit. Let A = 1 \ 0 \ 1 \ 0 is a 4-bit binary number, Representation of A using 8-bits (i.e. B): 1 \ 1 \ 1 \ 1 \ 0 \ 1 \ 0. is A = B?

In 2's Complement Rep.: Yes.

In 1's Complement Rep.: Yes.
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12 / 18

In Signed Magnitude Rep.: No.

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In 2's Complement Rep.: Yes.

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- 2. $(8.25)_{10} = (1000.01)_2$
- 3. $(16.125)_{10} = (10000.001)_2$
- 4. $(0.875)_{10} = (0.111)_2$
- 5. $(4.5)_{10} = (1.001)_2 \times 2^2$
- 6. $(8.25)_{10} = (1.00001)_2 \times 2^3$
- 7. $(16.125)_{10} = (1.0000001)_2 \times 2^4$
- 8. $(0.875)_{10} = (1.11)_2 \times 2^{-1}$

Representations 5 to 8 are called Normalized Representations.

Normalized Rep.: $(\pm 1.xxxxx)_2 \times 2^E$, Where 'E' is a **True Exponent**, 'xxxxxx' is a **Fraction/Mantissa**.

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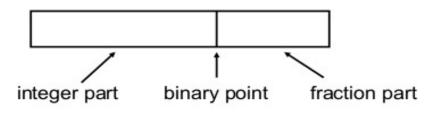
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Fixed Point (FP) Representation



Find the decimal equivalent of the following binary numbers. Assume that the binary numbers are represented using FP representation (6,2), i.e., 6 bits for integer part and 2 bits for fractional part.

$$(00101011)_2 =$$
? $(11111011)_2 =$?

Smallest +ve number that can be represented using FP (6,2) rep.: ? Biggest +ve number that can be represented using FP (6,2) rep.: ?

Fixed Point Arithmetic

Consider a 16-bit binary representation, in which least significant 8 bits are used for precision then give the range of values that can be represented using the 16-bit representation.

$$+ve Values: 2^{-8} to 2^7 - 2^{-8}$$

-ve Values: -2^{7} to -2^{-8}

Zero

Resolution is ?

Advantages of FP Arithmetic:

Easy to implement and occupies less space.

If performance is important than precision.

Once can choose a trade off between range and precision.

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IEEE 754 format for Real Numbers.

	Sign	Biased Exponent	Mantissa/Fraction]
Single Precision N=32	1 bit	8 bits	23 bits E	Bias Value:+127
Double Precision N=64	1 bit	11 bits	52 bits Bi	as Value :+1023

Biased Exponent=True Exponent + Bias Value

Rep. of
$$(4.5)_{10}$$
 using Single Precision $(4.5)_{10} = (100.1)_2 = (1.001)_2 \times 2^2$

Normalized Rep.: $(\pm 1.xxxxxx)_2 \times 2^E$, Where 'E' is a **True Exponent**, 'xxxxxx' is a Fraction/Mantissa.

Biased Exponent =
$$2 + 127 = 129 = 1000 \ 0001$$

$$\mathsf{Mantissa} = \mathbf{001} = \mathbf{0010000} \ \mathbf{0000} \ \mathbf{0000} \ \mathbf{0000} \ \mathbf{0000}$$

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IEEE 754

Biased Exponent=True Exponent + Bias Value, where $1 \le$ Biased Exponent $\le (2^{Length \ of \ Biased \ Exponent} - 2)$.

Single Precision (N=32), $1 \le$ Biased Exponent ≤ 254 .

Biased Exponent = 0,

Mantissa = ± 0 , then Value is ± 0 .

Mantissa \neq 0, then Value is **not a normalized number**.

Biased Exponent = 255,

Mantissa = ± 0 , then Value is $\pm \infty$.

Mantissa $\neq 0$, then Value is **NAN**.

Range of positive values: $[1.0 \times 2^{-126}, (2-2^{-23}) \times 2^{127}]$

Range of negative values: $[-(2-2^{-23}) \times 2^{127}, -1.0 \times 2^{-126}]$

Single Precision Number Resolution: $2^{-23} \times 2^{TrueExponent}$

BCD Representation

BCD: Binary Coded Decimal

It uses a 4-bit binary number to represent each decimal digit.

Decimal Number	BCD Rep.
0	0000
1	0001
2	0010
3	0011
9	1001
10	0001 0000
25	0010 0101
99	1001 1001

Table 1: BCD equivalent of a decimal number.