Question

**Editorial Solution** 

# Question

Given an array of integers, return indices of the two numbers such that they add up to a specific target.

You may assume that each input would have exactly one solution.

#### Example:

```
Given nums = [2, 7, 11, 15], target = 9,

Because nums[0] + nums[1] = 2 + 7 = 9,

return [0, 1].
```

### UPDATE (2016/2/13):

The return format had been changed to zero-based indices. Please read the above updated description carefully.

### **Quick Navigation**

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# Solution

# Approach #1 (Brute Force) [Accepted]

The brute force approach is simple. Loop through each element x and find if there is another value that equals to target-x.

```
public int[] twoSum(int[] nums, int target) {
    for (int i = 0; i < nums.length; i++) {
        for (int j = i + 1; j < nums.length; j++) {
            if (nums[j] == target - nums[i]) {
                return new int[] { i, j };
            }
        }
    }
    throw new IllegalArgumentException("No two sum solution");
}</pre>
```

## **Complexity Analysis**

- Time complexity:  $O(n^2)$ . For each element, we try to find its complement by looping through the rest of array which takes O(n) time. Therefore, the time complexity is  $O(n^2)$ .
- Space complexity : O(1).

## Approach #2 (Two-pass Hash Table) [Accepted]

To improve our run time complexity, we need a more efficient way to check if the complement exists in the array. If the complement exists, we need to look up its index. What is the best way to maintain a mapping of each element in the array to its index? A hash table.

We reduce the look up time from O(n) to O(1) by trading space for speed. A hash table is built exactly for this purpose, it supports fast look up in near constant time. I say "near" because if a collision occurred, a look up could degenerate to O(n) time. But look up in hash table should be amortized O(1) time as long as the hash function was chosen carefully.

A simple implementation uses two iterations. In the first iteration, we add each element's value and its index to the table. Then, in the second iteration we check if each element's complement (target-nums[i]) exists in the table. Beware that the complement must not be nums[i] itself!

```
public int[] twoSum(int[] nums, int target) {
   Map<Integer, Integer> map = new HashMap<>();
   for (int i = 0; i < nums.length; i++) {
        map.put(nums[i], i);
   }
   for (int i = 0; i < nums.length; i++) {</pre>
```

```
int complement = target - nums[i];
  if (map.containsKey(complement) && map.get(complement) != i) {
     return new int[] { i, map.get(complement) };
  }
}
throw new IllegalArgumentException("No two sum solution");
}
```

### Complexity Analysis:

- Time complexity : O(n). We traverse the list containing n elements exactly twice. Since the hash table reduces the look up time to O(1), the time complexity is O(n).
- Space complexity: O(n). The extra space required depends on the number of items stored in the hash table, which stores exactly n elements.

## Approach #3 (One-pass Hash Table) [Accepted]

It turns out we can do it in one-pass. While we iterate and inserting elements into the table, we also look back to check if current element's complement already exists in the table. If it exists, we have found a solution and return immediately.

```
public int[] twoSum(int[] nums, int target) {
    Map<Integer, Integer> map = new HashMap<>();
    for (int i = 0; i < nums.length; i++) {
        int complement = target - nums[i];
        if (map.containsKey(complement)) {
            return new int[] { map.get(complement), i };
        }
        map.put(nums[i], i);
    }
    throw new IllegalArgumentException("No two sum solution");
}</pre>
```

### **Complexity Analysis:**

- Time complexity : O(n). We traverse the list containing n elements only once. Each look up in the table costs only O(1) time.
- Space complexity: O(n). The extra space required depends on the number of items stored in the hash table, which stores at most n elements.

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Strill commented 5 days ago

@miziaiba (https://discuss.leetcode.com/uid/62080) You'll get [0,3].

miziaiba commented last week

For approach 2 and 3, if there is two repeated value in the array(e.g. [0, 3, 4, **ஜ்** அரசு ஊத்த இடிக்கும் இருக்கு விடுகள்ளின்ற விடுகள்ளின் விடுகள்ளின்ற விடுகள்ளின் விடுகள்ளின் விடுகள்ளின்ற விடுகள்ளின் விடுகள் விடுகள்ளின் விடுகள்ளின்

brianliu7 commented last week

@ayuanx (https://discuss.leetcode.com/uid/58648) I believe you are technically correct about the O(1) lookup vs the O(logN). I uss, leetcode.com/user/brianliu? Think for learning purposes, e.g. this problem is meant to illustrate the advantages of hashing vs. array as a data structure and ignores whats going on in reality in the background (using a tree).

reading more about it it seems that it performs closer to O(1), but given a really really unlucky set of collisions worst case running scenario can be O(n), but extremely rare / only with edge cases

 $\label{lem:http://stackoverflow.com/questions/1055243/is-a-java-hashmap-really-o1 (http://stackoverflow.com/questions/1055243/is-a-java-hashmap-really-o1)} \\$ 

brianliu7 commented last week

@chenxiaokai (https://discuss.leetcode.com/uid/58972) I just ran your code (c code) and in the comments it mentions you must ss.leetcode.com/user/brianliur?)
use malloc() to allocate memory and the caller will use free on the object; so the code did not work for me

brianliu7 commented last week

@chenxiaokai (https://discuss.leetcode.com/uid/58972) you are returning the INDICES of the array, which are correct (0 and 2), uselectede.com/uip/brightlift/leset.ves\_ (0 and 3)