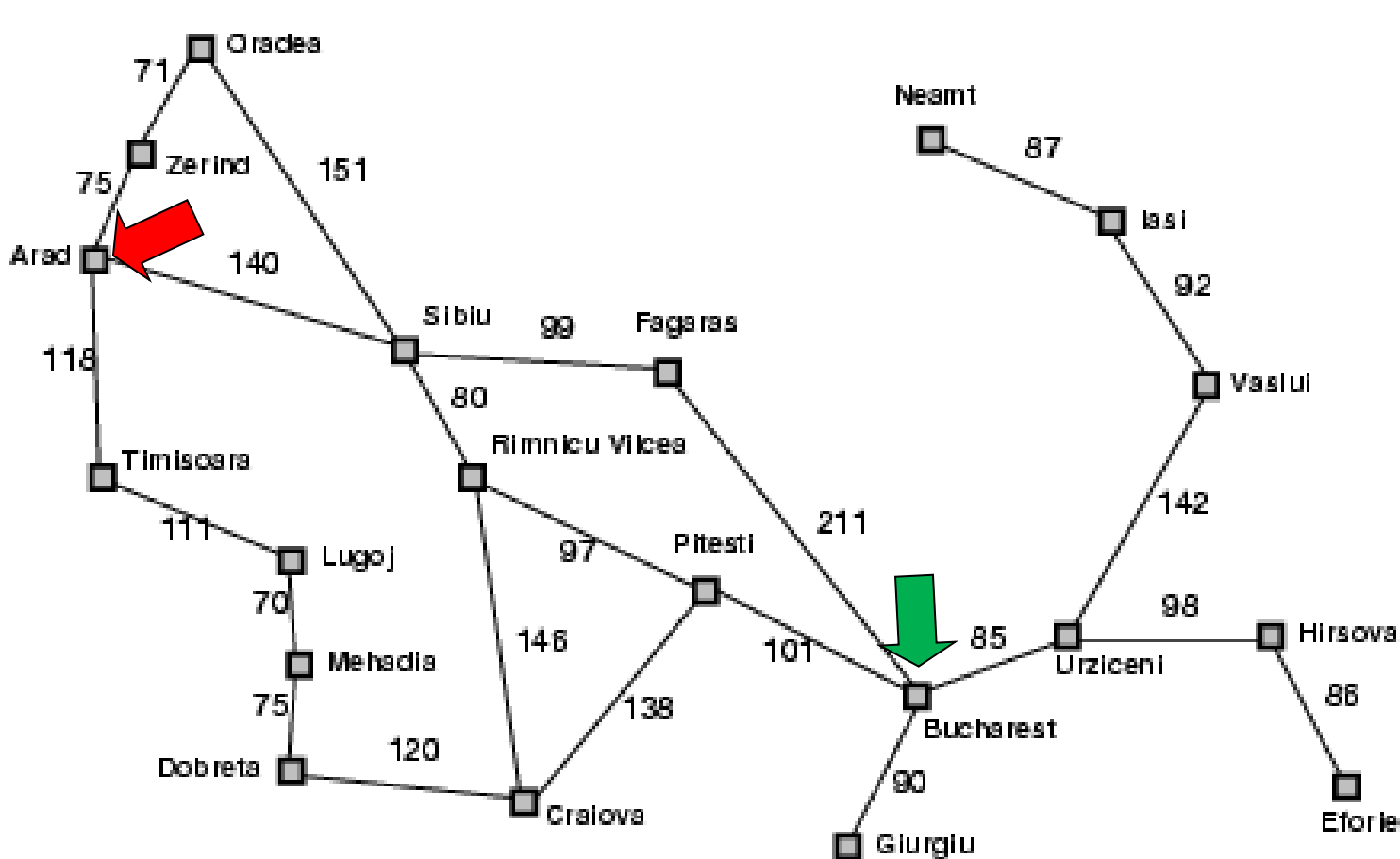


A* Search

 Start Node

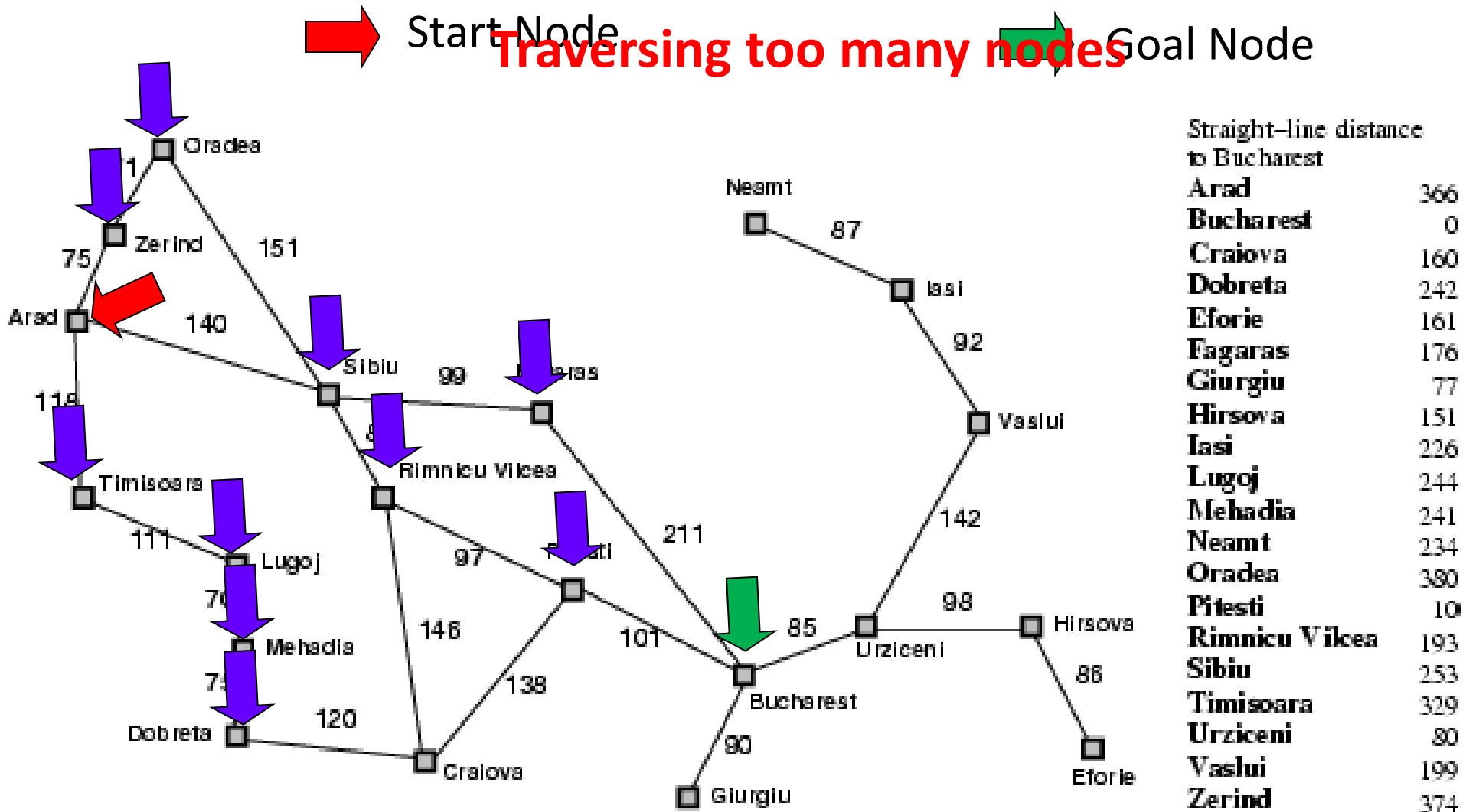
 Goal Node



Straight-line distance
to Bucharest

Arad	366
Bucharest	0
Craiova	160
Dobreta	242
Eforie	161
Fagaras	176
Giurgiu	77
Hirsova	151
Iasi	226
Lugoj	244
Mehadia	241
Neamt	234
Oradea	380
Pitesti	10
Rimnicu Vilcea	193
Sibiu	253
Timisoara	329
Urziceni	80
Vaslui	199
Zerind	374

Dijkstra's Algorithm



Dijkstra's Algorithm

- Dijkstra's algorithm has one cost function, which is real cost value from source to each node:

$$f(n)=g(n).$$

- Use **heuristics** to guide the search.

Heuristic: estimation of how to search for a solution

- Evaluation function $f(n) = g(n) + h(n)$ where

$g(n)$ = cost so far to reach n

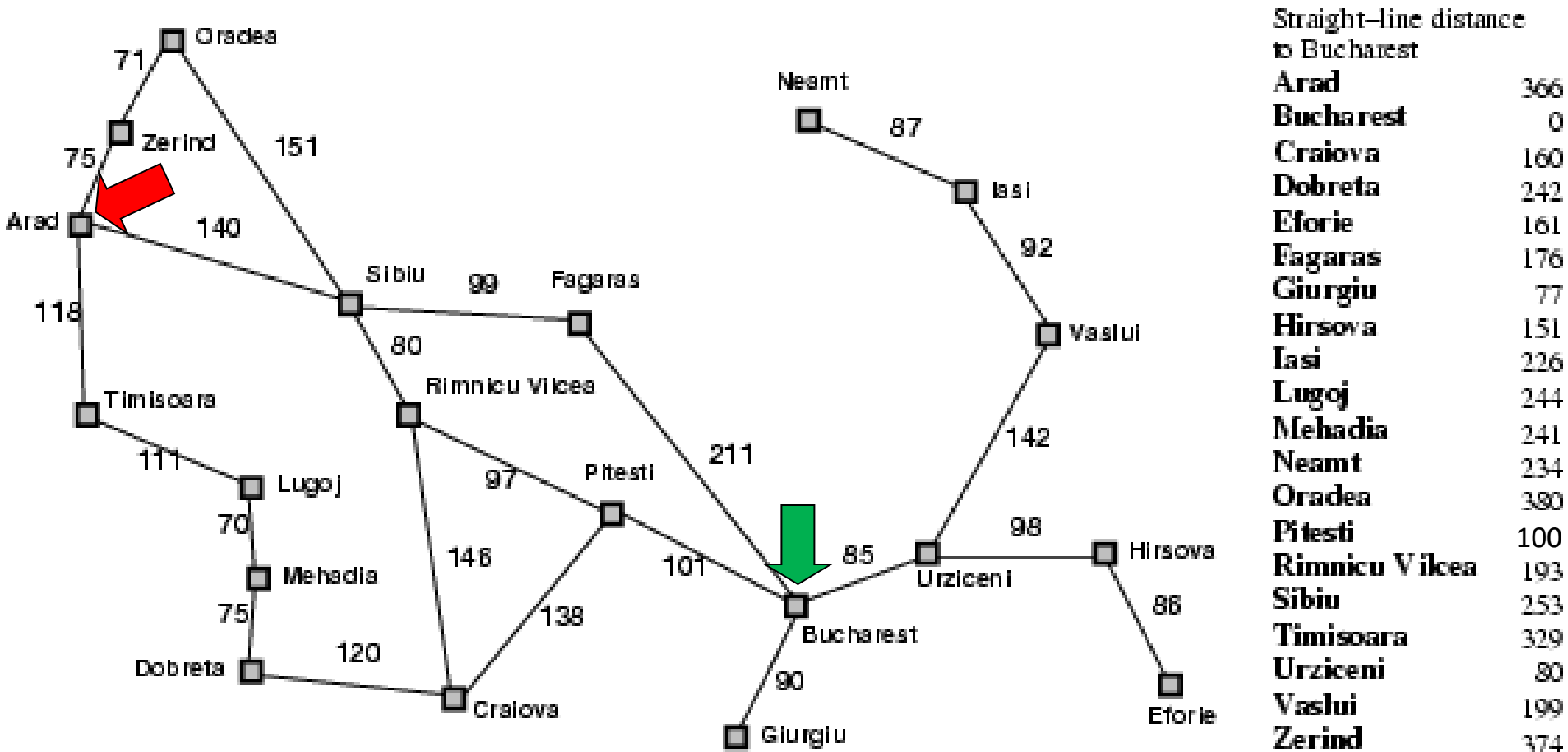
$h(n)$ = estimated cost from n to goal

$f(n)$ = estimated total cost of path through n to goal

A* Search

 Start Node

 Goal Node

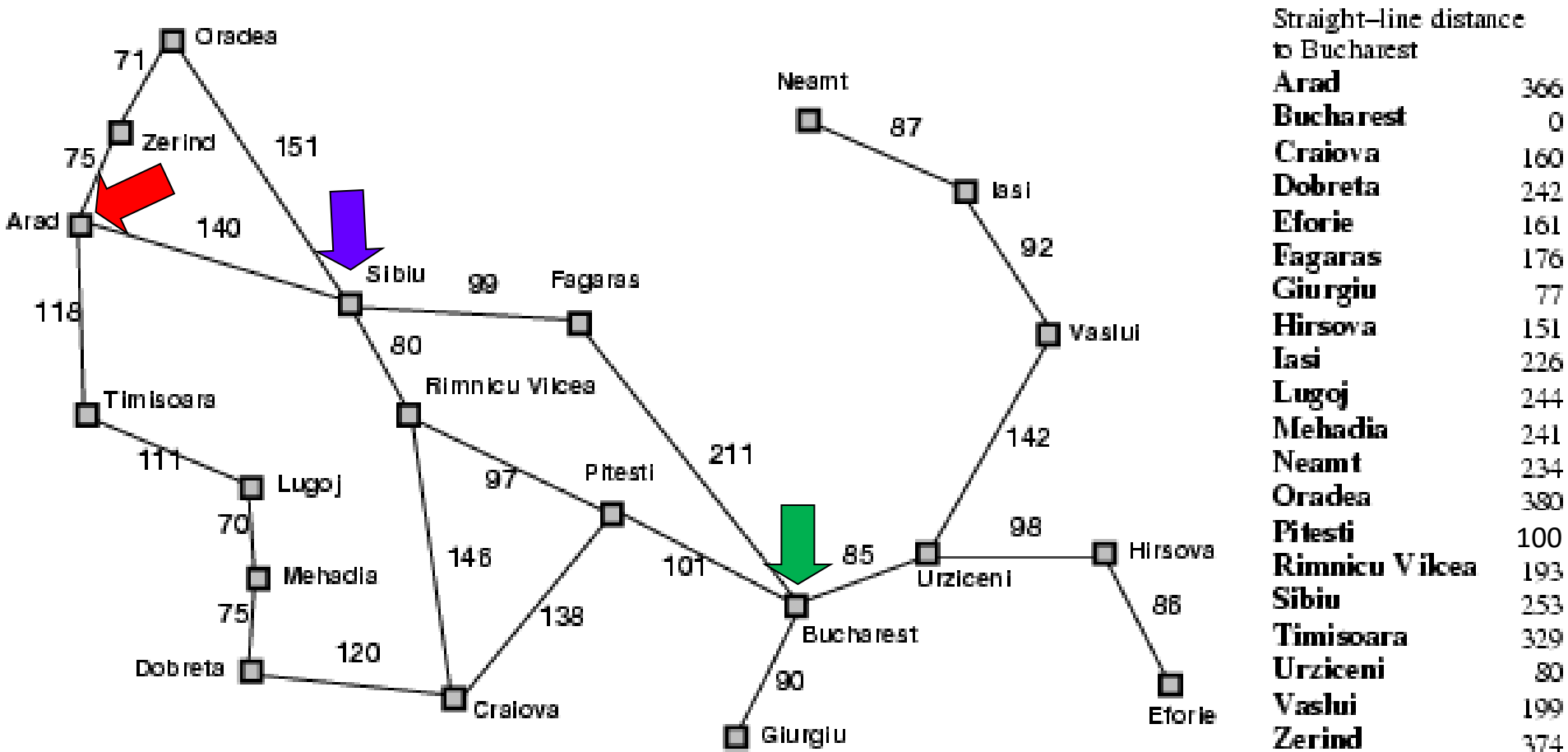


A* Search

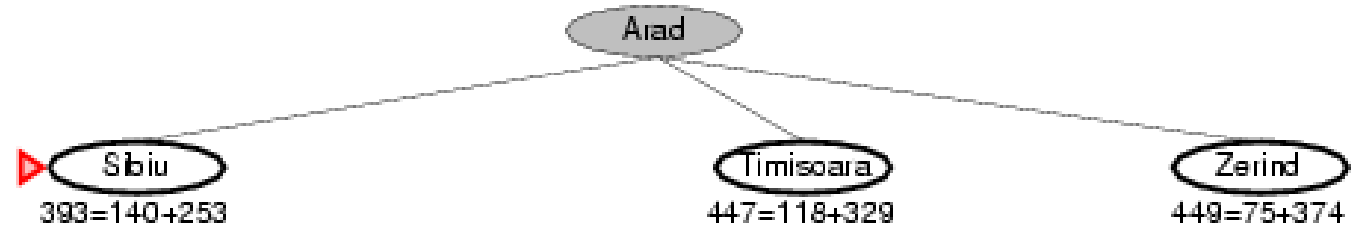
Zerind: $f(n) = g(n) + h(n) = 75 + 374 = 449$

Timisoara: $f(n) = g(n) + h(n) = 118 + 329 = 447$

Sibiu: $f(n) = g(n) + h(n) = 140 + 253 = 393$



A* Search

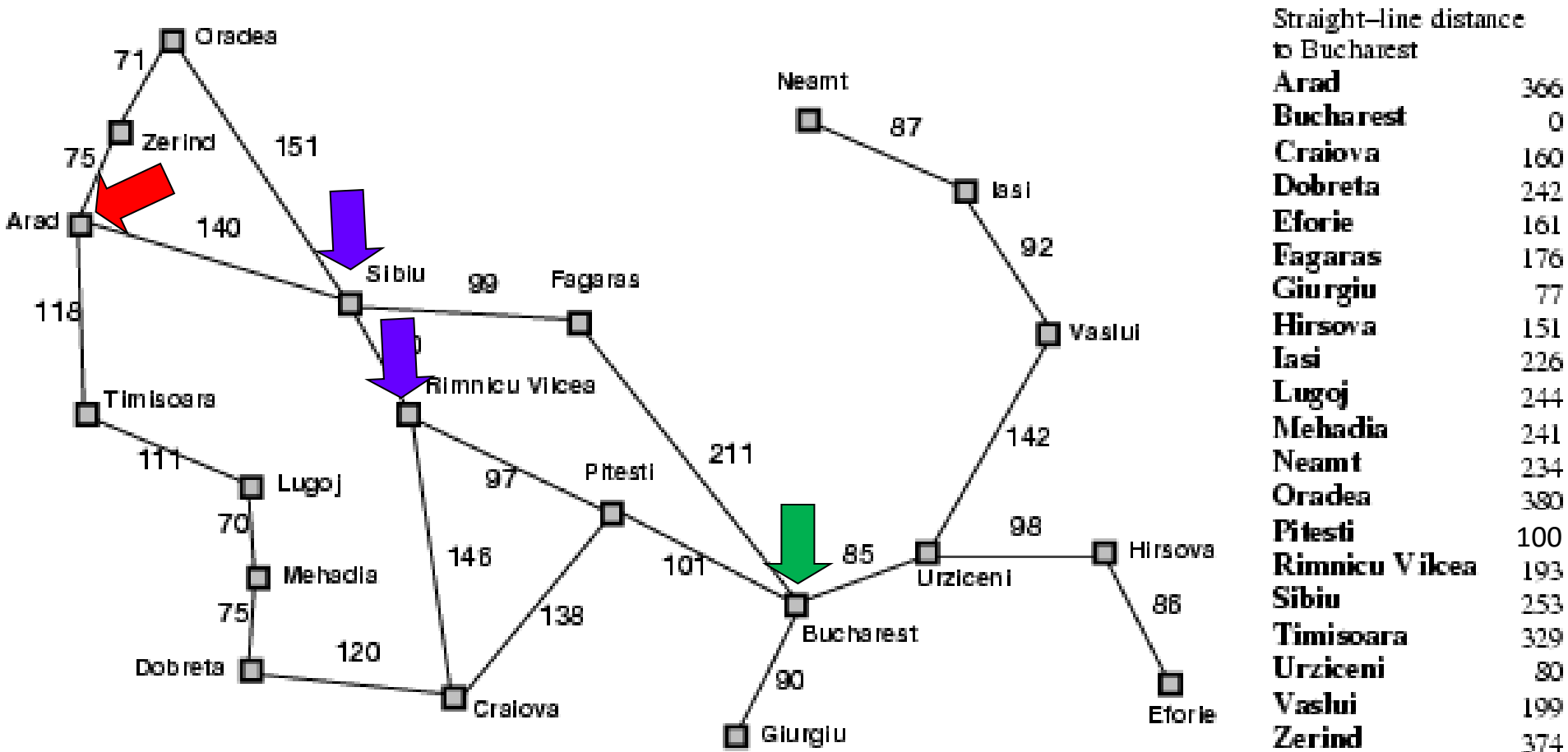


A* Search

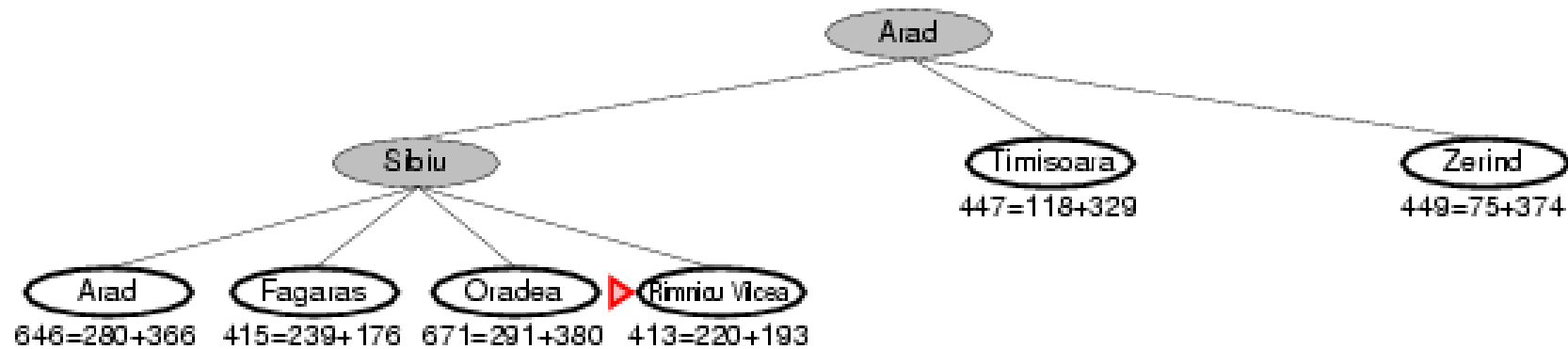
Fagaras : $f(n) = g(n) + h(n) = (140+99) + 176 = 415$

Rimnicu Vilcea : $f(n) = g(n) + h(n) = (140+80) + 193 = 413$

Oradea : $f(n) = g(n) + h(n) = (140+151) + 380 = 671$



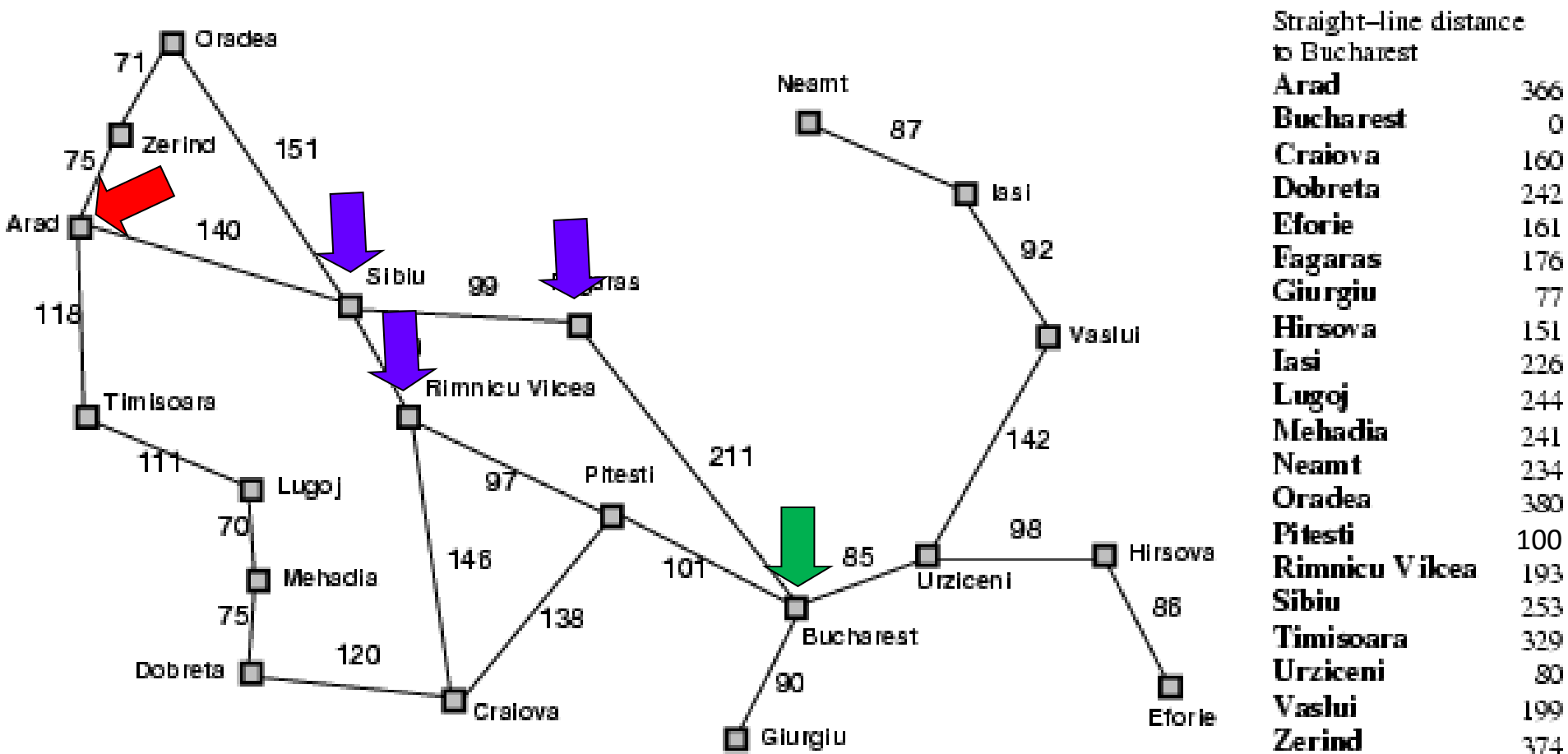
A* Search



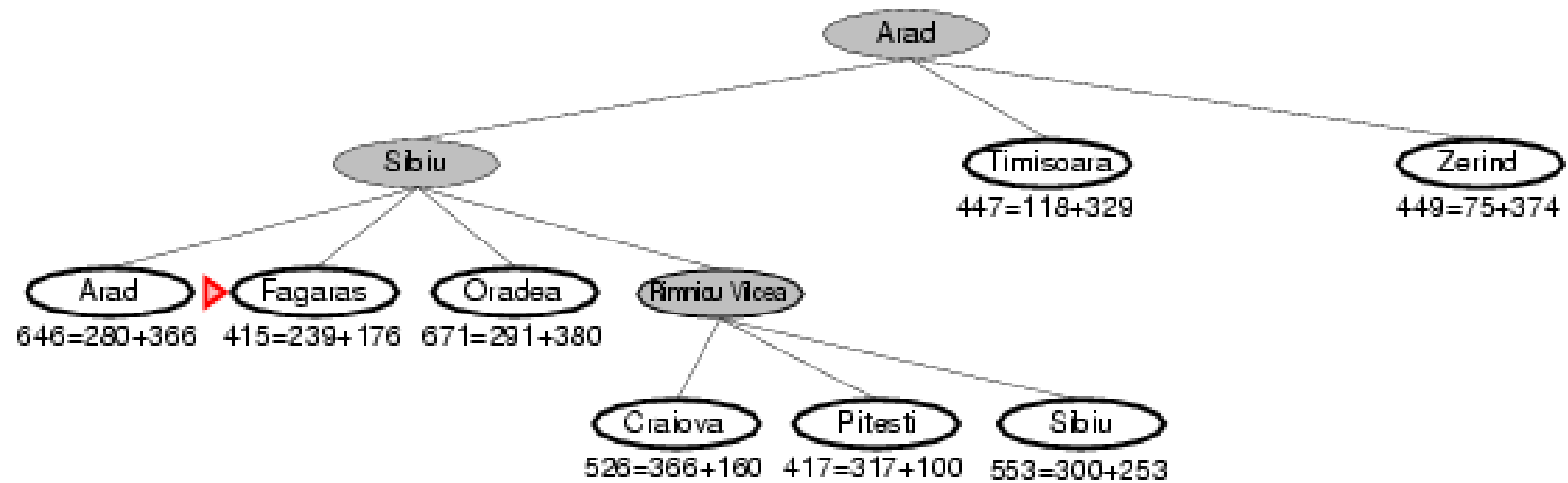
A* Search

Fagaras : $f(n) = g(n) + h(n) = (140+99) + 176 = 415$

Pitesti : $f(n) = g(n) + h(n) = (140+80+97) + 100 = 417$



A* Search

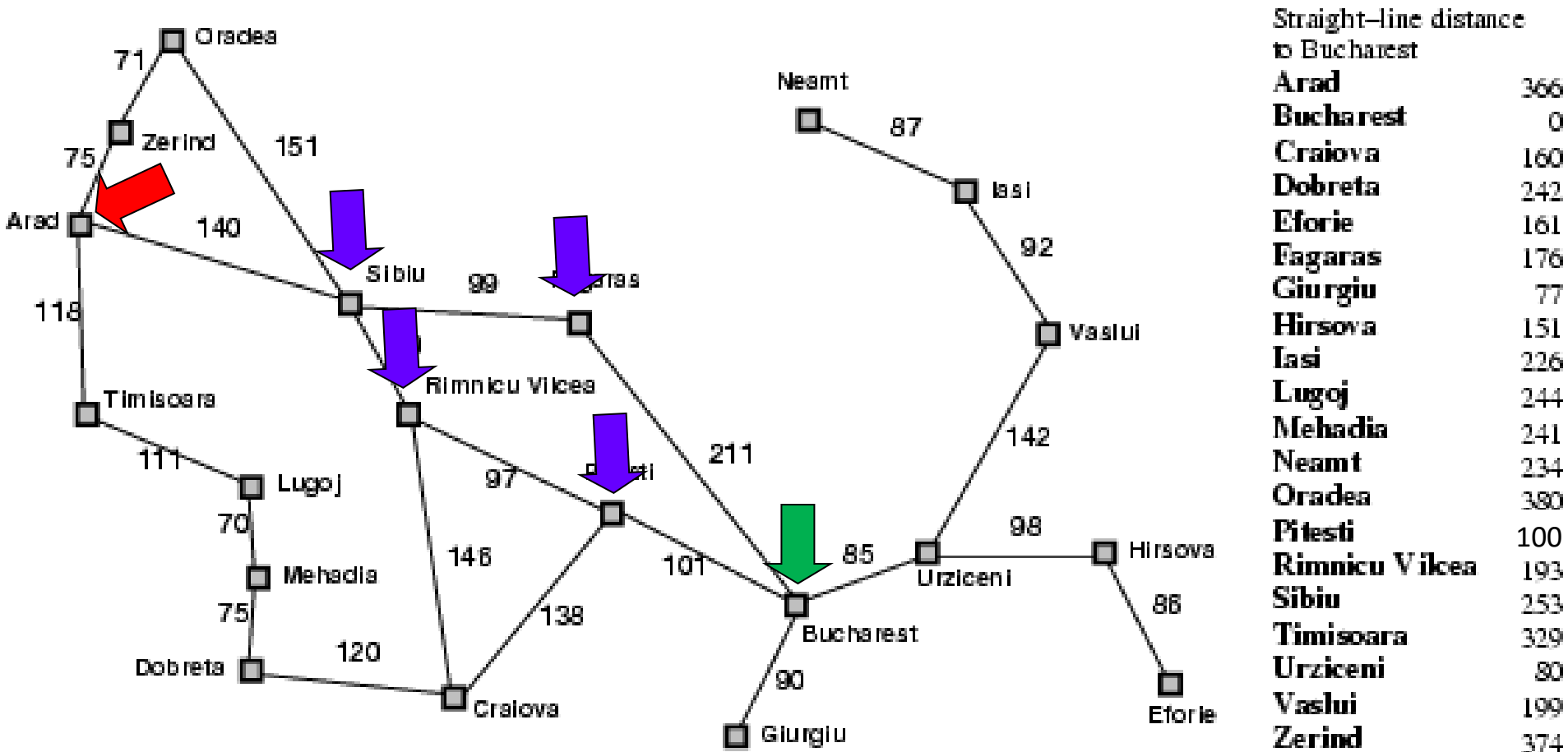


A* Search

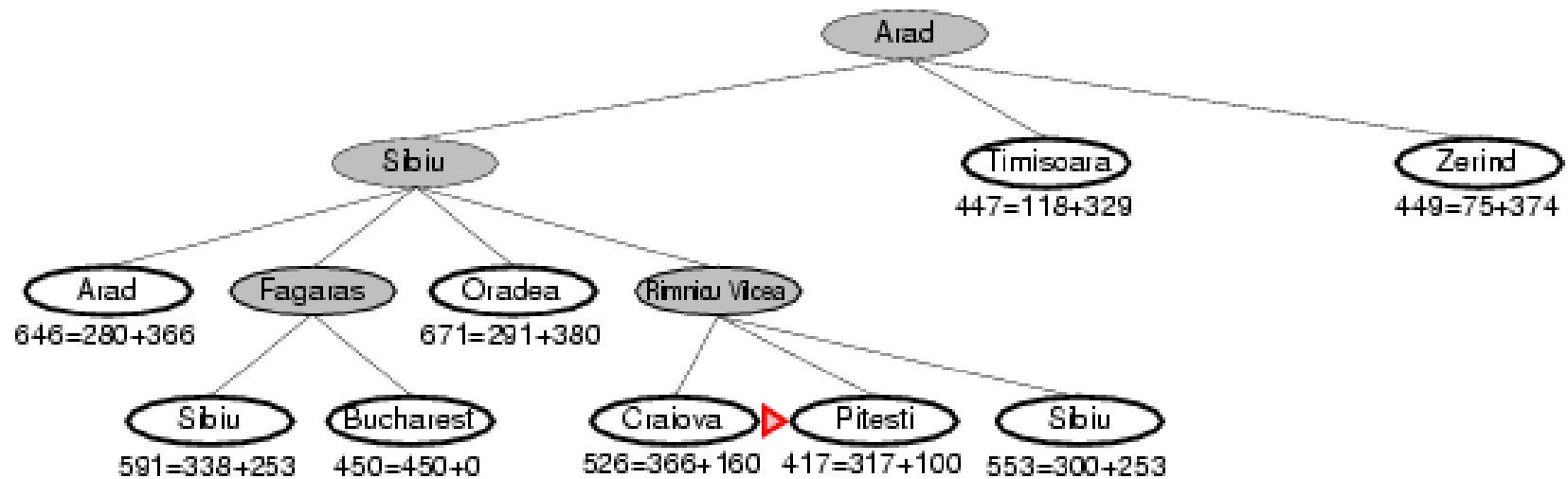
Fagaras : $f(n) = g(n) + h(n) = (140+99) + 176 = 415$

Pitesti : $f(n) = g(n) + h(n) = (140+80+97) + 100 = 417$

Bucharest (via fagaras) : $f(n) = g(n) + h(n) = (140+99+211) + 0 = 450$



A* Search



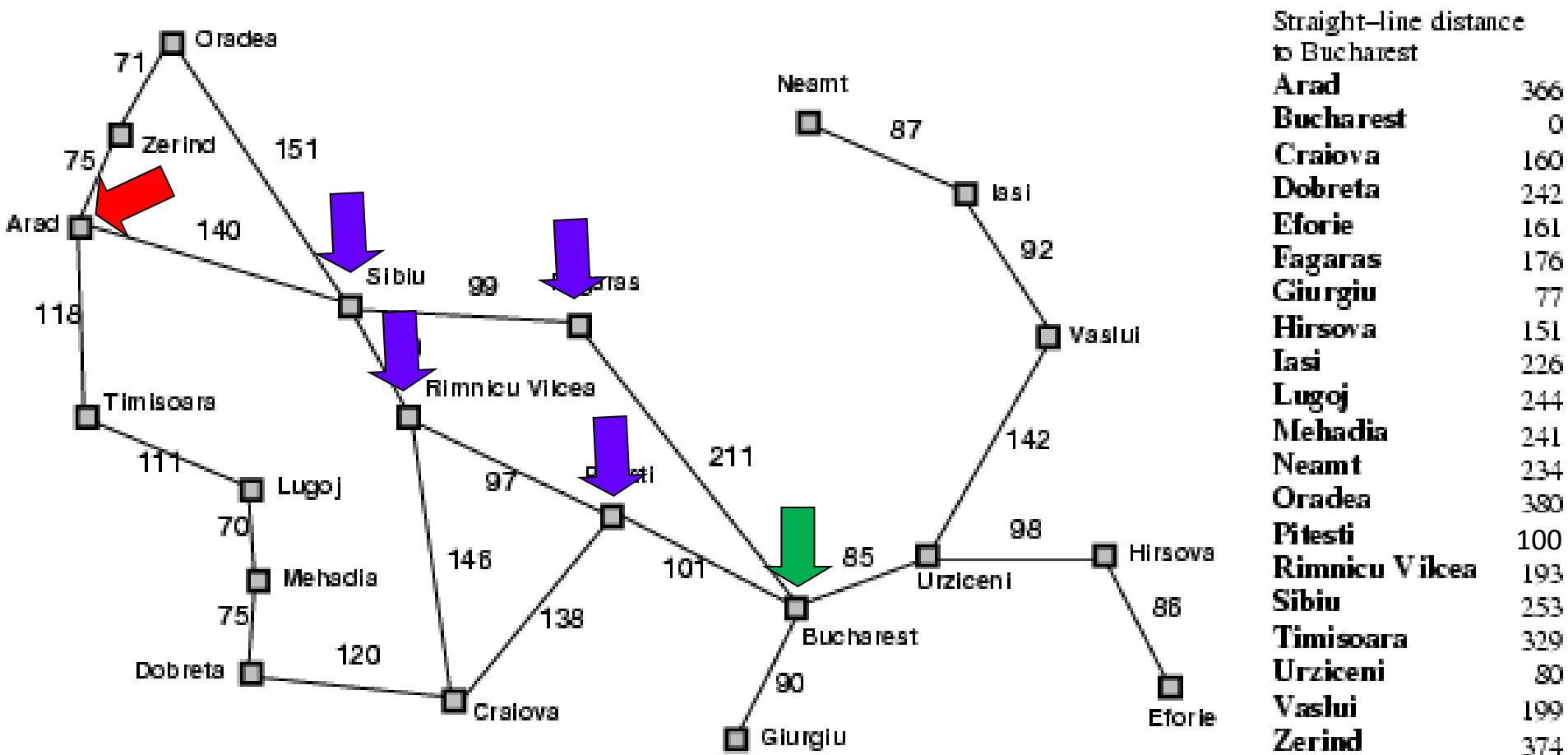
A* Search

Fagaras : $f(n) = g(n) + h(n) = (140+99) + 176 = 415$

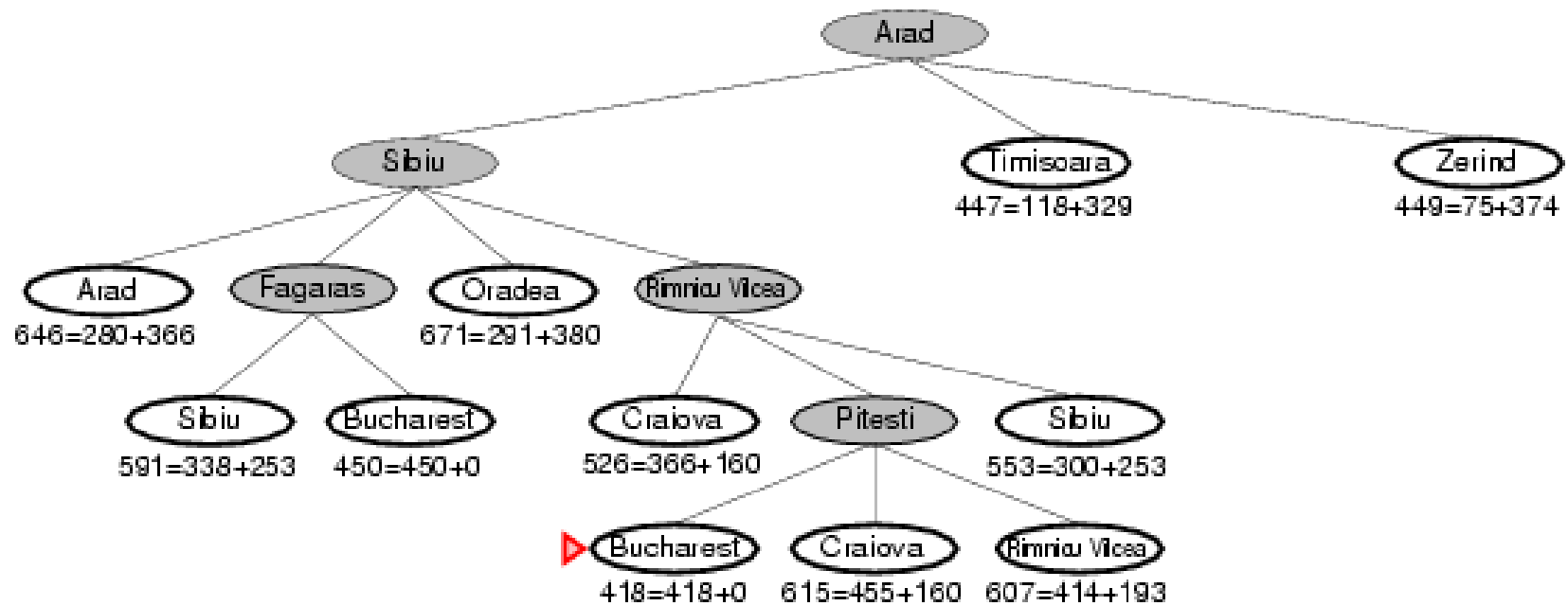
Pitesti : $f(n) = g(n) + h(n) = (140+80+97) + 100 = 417$

Bucharest (via fagaras) : $f(n) = g(n) + h(n) = (140+99+211) + 0 = 450$

Bucharest (via Pitesti) : $f(n) = g(n) + h(n) = (140+80+97+101) + 0 = 418$



A* Search



A* Search

```
function A*(start,goal)
    closedset := the empty set    // The set of nodes already evaluated.
    openset := {start}           // The set of tentative nodes to be evaluated, initially containing the start node
    came_from := the empty map    // The map of navigated nodes.

    g_score[start] := 0           // Cost from start along best known path.
    f_score[start] := g_score[start] + heuristic_cost_estimate(start, goal) // Estimated total cost

    while openset is not empty
        current := the node in openset having the lowest f_score[] value
        if current = goal
            return reconstruct_path(came_from, goal)

        remove current from openset
        add current to closedset
        for each neighbor in neighbor_nodes(current)
            if neighbor in closedset
                continue
            tentative_g_score := g_score[current] + dist_between(current,neighbor)

            if neighbor not in openset or tentative_g_score < g_score[neighbor]
                came_from[neighbor] := current
                g_score[neighbor] := tentative_g_score
                f_score[neighbor] := g_score[neighbor] + heuristic_cost_estimate(neighbor, goal)
                if neighbor not in openset
                    add neighbor to openset
    return failure

function reconstruct_path(came_from,current)
    total_path := [current]
    while current in came_from:
        current := came_from[current]
        total_path.append(current)
    return total_path
```

A* Search

- When $h(n) = \text{actual cost to goal}$
 - Only nodes in the correct path are expanded
 - Optimal solution is found
- When $h(n) < \text{actual cost to goal}$
 - Additional nodes are expanded
 - Optimal solution is found
- When $h(n) > \text{actual cost to goal}$
 - Optimal solution can be overlooked

Admissible heuristics

- A heuristic $h(n)$ is **admissible** if for every node n , $h(n) \leq h^*(n)$, where $h^*(n)$ is the **true cost** to reach the goal state from n .
- An admissible heuristic **never overestimates** the cost to reach the goal, i.e., it is **optimistic**.
- If $h(n)$ is admissible, A^* using **TREE-SEARCH** is optimal

Dominance

➤ If $h_2(n) \geq h_1(n)$ for all n (both admissible)

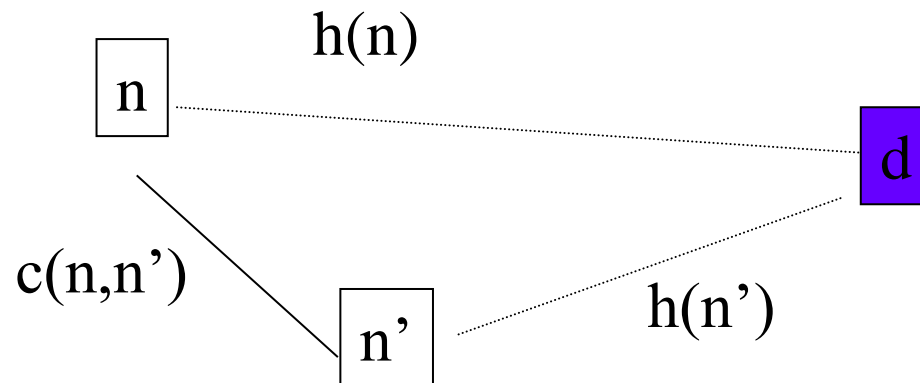
then h_2 dominates h_1

➤ h_2 is better for search: it is guaranteed to expand less or equal number of nodes.

Consistent heuristics

- $h(n)$ is **consistent** if for every node n and for every successor node n' of n :

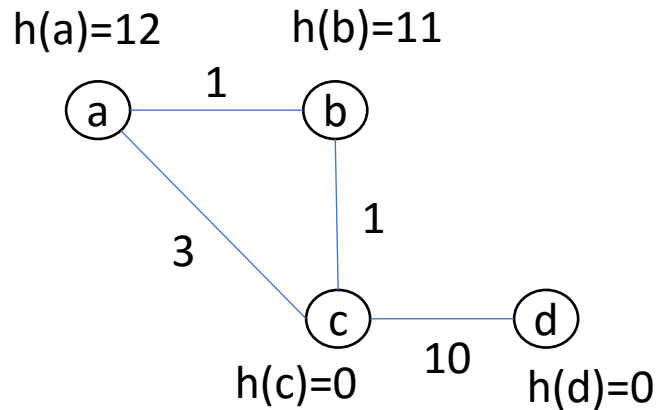
$$h(n) \leq c(n, n') + h(n')$$



Consistent heuristics

- If $h(n)$ is consistent then $h(n)$ is admissible
- Frequently when $h(n)$ is admissible, it is also consistent.
- If $h(n)$ is a consistent heuristic, A^* using graph search is optimal.

Consistent heuristics



- H is admissible but not consistent heuristic
- Graph search will find path $a-c-d$.
 - c will be expanded after a , and there will be no option to expand c again after expanding b (Since c already in closed set)



THANK YOU!!!