## CS341 #15

Condition Variables. Implement a Mutex Lock. The Critical Section Problem

- 1. How do I block a thread (= send it to 'sleep')?
- 2. How do I wake up threads that are blocked on a condition variable?

Example: Fix the following methods using a condition variable and mutex lock to ensure the cake integer is never negative.

```
pthread mutex t m = PTHREAD=MUTEX INITIALIZER;
    pthread cond t cv = PTHREAD COND INITIALIZER;
04
   int cake = 0;
    void decrement() { // Will block if zero
      while(cake == 0) {
10
          sleep(1)
11
12
      cake --;
14
15
16
17
    void increment() {
18
        cake ++;
19
```

3. How does pthread cond wait really work?

4. Challenge. A fixed size stack that blocks

```
01 pthread mutex t m = PTHREAD MUTEX INITIALIZER;
02 pthread cond t cv = PTHREAD COND INITIALIZER;
03 double array[10];
04 \text{ int } n = 0;
06 // blocks while full (n ==10)
07 void push(double v) {
09
10
11
12
13
14
15
16 }
17 //  blocks while empty (n == 0)
18 double pop() {
19
20
21
22
23
24
25
26 }
27 // Test with 2+ threads that add values...
28 void* generator(void*){
     for(int i =0; i < 100000; i++)
         push(i);
31
    Return NULL;
32 }
33 // And one thread that remove values
34 void * consumer(void*result) {
    double sum = 0, i=0;
    while (i=pop() != -1) sum += i;
    printf("%.0f", sum);
   Return NULL;
39 }
40
```

## How can you implement a reliable mutex lock?

5. Let's try writing a simple implementation...

```
01    pthread_mutex_init(int * m) { *m= 0; }
02
03    pthread_mutex_lock(int* m) {
04         while(*m ==1) {
05             pthread_yield(); /*sleeps for a short time */
06         }
07         *m = 1;
08    }
09    pthread_mutex_unlock(int* m) {?______}
```

Problems?

## 6. CPU support: Use an atomic CPU instruction.

Imagine a special *Atomic\_Exchange* instruction 'exch' that swaps the values at two addresses as an *indivisible*, *uninterruptable* operation

```
01    pthread_mutex_init(int* m) { *m= 0; }
02
03    pthread_mutex_lock(int* m) {
        for(int q = 1; q;) { ______} }
04
05    pthread_mutex_unlock(int* m) { _______} }
```

## 7. The Critical Section Problem

while(running) {

- 1. Wait to enter the critical section if another thread is in the CS.
- 2. Critical Section Code here. Only one thread in here at a time!
- 3. Leave critical section. Allow another waiting thread to enter.
- 4. // do other stuff most of the time

~~ Welcome to the **Critical Section Problem** game show! ~~ Today's prizes: mutual exclusion and progress

Candidate #1. Use a single, boolean "flag"

boolean flag

| Thread A                   | Thread B                   |
|----------------------------|----------------------------|
| wait while the flag is up  | wait while the flag is up  |
| raise the flag!            | raise the flag!            |
| Critical Section code here | Critical Section code here |
| lower the flag!            | lower the flag!            |
|                            |                            |

// Then each thread does other work but will repeat this again sometime in the future Problems?

Candidate #2. Give each thread its own a flag.

boolean flagA, flagB

| wait while B's flag is up  | wait while A's flag is up  |
|----------------------------|----------------------------|
| raise A flag               | raise B flag               |
| Critical Section code here | Critical Section code here |
| lower A flag               | lower B flag               |
|                            |                            |

Problems?

Candidate #3. Change the sequence order

| raise A flag               | raise B flag               |
|----------------------------|----------------------------|
| wait until B flag is down  | wait until A flag is down  |
| Critical Section code here | Critical Section code here |
| lower A flag               | lower B flag               |

Problems?