

ASSIGNMENT 4

- ▶ Overwrite →next in another entry by overflowing the buffer.
- ▶ You don't need to deal with heap *metadata* yet, but heap *layout* matters.
- ▶ Then, proceed as in phonebook.

- You can free arbitrary addresses.
- Build a fake chunk in one of your allocations that you then free (\Rightarrow *House of Spirit*).
- Make the fake chunk large enough, so you get a really big allocation on the stack when you allocate again.
- Write your ROP chain as usual.

- checksec: **No RELRO, no PIE**
- Goal: Overwrite GOT entry with win function
- Straightforward tcache exploit

- checksec: **everything** enabled (except `_FORTIFY_SOURCE`)
- Trivial stack and heap pointer leaks (via `%p`)
- The `strings_p` entry is not zeroed after free → **use-after-free**
- `change` clears the buffer first, so we can't use it on `strings_p`
- Make `malloc/realloc` return the same chunk twice by overwriting → `fd`
- Allocate a string on top of `strings_p`
- Overwrite the pointers to point to the stack (during `alloc`)
- Leak the `libc` from the stack, then write a ROP chain

- magic to do "magic".
- This interprets the scores in sorted order as shellcode and runs it.
- We told you sorted shellcode was going to happen 😊

- ▶ Used to select optimized versions of functions based on available hardware at runtime (remember magic8ball's GLIBC_TUNABLES?)
- ▶ This means entries vary by hardware
- ▶ Some of the unoptimized implementations call other functions, when the optimized version does not
- ▶ *Some* entry will usually work, but it may not be the same as on your machine — try different offsets
- ▶ Typical trick: most print functions (puts, printf, ...) will call strlen or strnlen.