

# Eulerian path

Eulerian cycle/path, Chinese postman problem

beCP Training



OLYMPIADE BELGE D'INFORMATIQUE  
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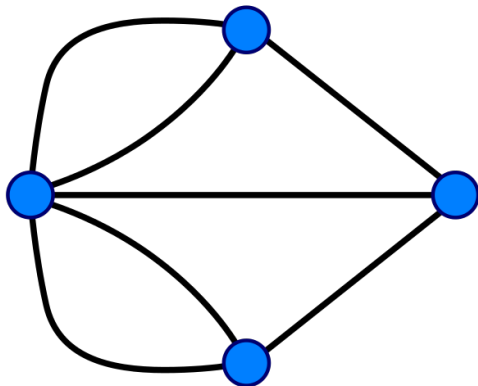
Eulerian cycles and paths

Finding the cycle/path

Chinese postman problem

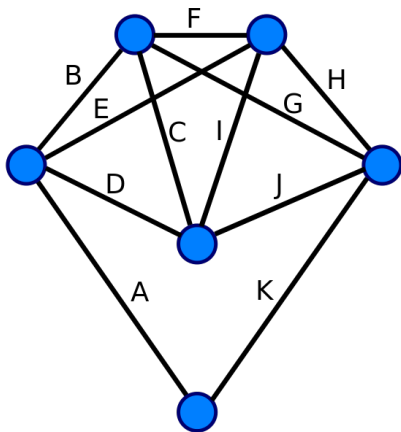
## Euler's problem

Can we make a cycle/path that visits every edge exactly once?



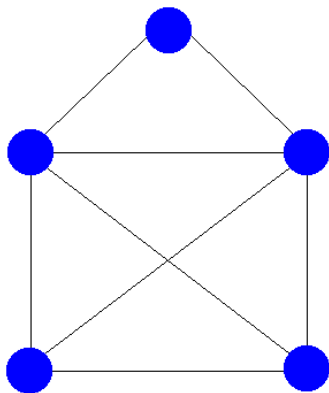
# Eulerian graph

A eulerian graph is a graph that has a Eulerian cycle (must come back to the start).



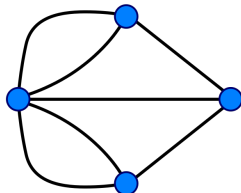
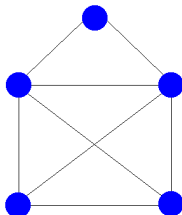
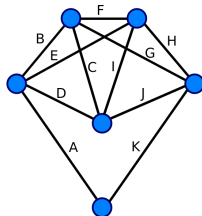
# Semi-eulerian graph

A semi-eulerian graph is a graph that has a Eulerian cycle *or* path (start and end may differ).



# Criteria for undirected graphs

- ▶ The graph must be connected
- ▶ For each node except the start and end, we use two edges: one for entering and one for leaving!
- ▶ So all degrees will be even, except for at most two



# Criteria for directed graphs

Too complicated...

Just check Wikipedia.

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Eulerian cycles and paths

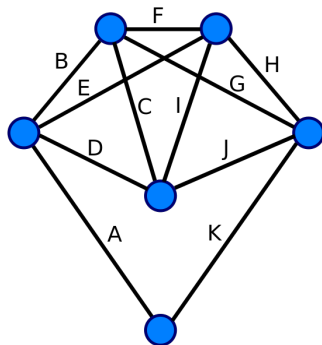
Finding the cycle/path

Chinese postman problem



# Traversal strategy

- ▶ Start at one of the odd-degree nodes (if they exist)
- ▶ Take arbitrary edges and remove them on the way
- ▶ The only possible dead-end is the ending node!
- ▶ If there are edges left, restart in the middle



# Implementation

To erase edges, we need to remember an “edge ID”

```
vector<int> id[MAXV], neigh[MAXV];  
bool visited[MAXE]; // which edges have been taken  
  
// Start on an odd-degree node (if possible)  
void euler(int u, vector<int> &s)  
{  
    for (int i=0; i < (int)neigh[u].size(); i++) {  
        if (!visited[id[u][i]]) {  
            visited[id[u][i]] = true;  
            euler(neigh[u][i], s);  
        }  
    }  
    s.push_back(u);  
}
```

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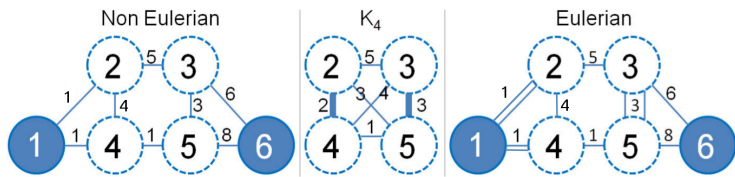
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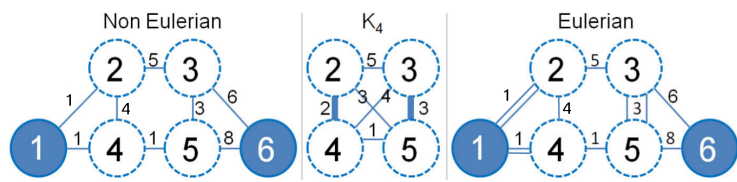
# Chinese postman statement

- ▶ A postman has to make a cycle through every street at least once, while travelling the smallest possible distance.
- ▶ If the graph is Eulerian, then just take a Eulerian cycle
- ▶ Otherwise, there are an odd number of odd-degree nodes
- ▶ Let's add edges to make it Eulerian!



# Choosing the right edges to add

- ▶ We have to add an edge to every odd-degree node
- ▶ But not change the parity of even-degree node
- ▶ So we have to add paths between pairs of odd-degree nodes
- ▶ Compute the best paths and put it in a graph
- ▶ Find the best pairings with complete search



# Sources of figures

- ▶ [https://commons.wikimedia.org/wiki/File:Königsberg\\_graph.svg](https://commons.wikimedia.org/wiki/File:Königsberg_graph.svg)
- ▶ [https://commons.wikimedia.org/wiki/File:Labelled\\_Eulergraph.svg](https://commons.wikimedia.org/wiki/File:Labelled_Eulergraph.svg)
- ▶ <https://en.wikibooks.org/wiki/File:Eulerian3.png>
- ▶ CP3 book