Emmy: The Gameboy Color Emulator

6CCS3PRJ Final Project

Classic video game console emulation

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I verify that I am the sole author of this report,
except where explicitly stated to the contrary.

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Abstract

This project aims to create a Gameboy emulator web-application, in other words a program capable of receiving Gameboy game files (commonly refered to as ROMs), and interpreting such ROM to play the game (or execute the program) it contains. The emulator will be usable in browsers, for both desktop computers and mobile devices that may not have access to a physical keyboard. The emulator will also contain debugging capacities, to allow other emulator developers to use it when comparing with their emulator and working on it.

The objective of this project is to create a piece of software that could be used by anyone wanting to emulate retro games, without the need for any technical knowledge on emulators or downloading anything (except the ROMs that need to be obtained separately).

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Chapter 1

Introduction

1.1 Motivation

Emulators are an area of computer science widely used today. Either implemented in hardware or software, they allow replicating the behaviour of one system on another. One of its applications is video game emulation, where a computer simulates a game console (usually a retro console). This allows users to play games that either may not be obtainable in stores anymore, or made for consoles that don't function properly anymore. A wide range of emulators already exist for most consoles. Emulation in general is also widely used in developing new systems, and is an active area of computer science.

This project will seek to create a new emulator for the Gameboy allowing users to play retro games on their computer or mobile device, through the browser. This report will document how the original console works and how the emulator imitates this behaviour to the best accuracy possible, as well as comparing the resulting emulator with other existing ones.

1.2 Scope

The scope of this project is creating a new Gameboy and Gameboy Color emulator, working for browsers. Emulation will be as accurate as achievable with the time available - there may be minor inacuracies in the end product. Extra peripherals and features of the console may be omitted, to allow more focus to be put on the core part of the console.

The emulator will be usable across a range of devices. Debugging tools and additional features may be provided to the user to let them customise their experience to their needs.

1.3 Objectives

The resulting software will allow users to open a game file for the gameboy - also called a ROM - and play it. They may use the emulator on a computer, controlling the console via the keyboard, or on a touch device, using on-screen buttons. The emulated features of the emulator include proper rendering of the screen, simulating the audio of the console, the different buttons, and support for a variety of chip controllers for game cartridges.

The frontend of the emulator may also contain additional quality of life features, such as custom

themes, save states, and debugging options allowing to inspect the state of the Gameboy - a feature vital to emulator developers and retro game developers.

Chapter 2

Background

2.1 Emulation

An emulator is "hardware or software that permits programs written for one computer to be run on another computer" [1]. The imitated computer is the *guest*, and the one that imitates is the *host*. Emulators are nowadays mainly found in the form of software, and have many different uses, from preservation to hardware development.

Emulation was technically born with the first computers: the very first computer, the Colossus made in 1941, was built to imitate the Enigma machine [2]. However emulation was properly studied towards the end of the twentieth century, when computing power started to steadily increase. One of the earliest instances of emulation as an actual feature is with the IBM System/360. This computer supported emulation of previous models, such as the IBM 709, 7090, 7094 and 7094 IIx [3].

Emulation is also vital for preservation: as transistors and motherboards age, old systems become unusable, and with them the software they ran. Emulating these systems is often the only future-proof and sustainable way to keep this software usable.

Finally, another common use for emulation is virtual machines. These programs allow running another Operating System (OS) on a computer, which can be used for instance when developing for other systems, without needing to use the physical device directly, for instance when developing a Windows-compatible app with a Linux computer.

2.2 Video Game Emulation

Video game emulation is the art of emulation applied to video game hardware systems. This allows the host to run games destined for the original console. This usually requires precise understanding of the console's hardware and functioning, as games may rely on specific behaviours and edge cases to function. This task is rendered harder by the fact that often game consoles are poorly documented, as they are proprietary hardware and programmer guides written for them cannot be legally distributed.

Video game emulation started in the 90s when computers were powerful enough to properly simulate console systems. Although precise dates are hard to get, the first console emulators seem to be either from 1990 or 1993 [4], and were able to run some NES games. The first

Gameboy emulators were in the late 90s, with the Virtual GameBoy[1] in 1995 and NO$GMB[2] in 1997 (although it's history page[3] seems to indicate development started in 1993) [5].

The original game files and assembled code for video games is copyrighted material, and is referred to as the Read-Only Memory (ROM) of the game. Although distributing these ROMs is usually illegal, there also exist copyright free ROMs: games created by developers that chose to license them under Creative Commons licenses, for instance. Websites such as Retro Veteran[4] host wide collections of legal ROMs.

2.3 Gameboy, Gameboy Color

The Gameboy (GB) is an 8-bit handheld video game console, released in 1989. It has a small 160×144 pixel screen, and has a Sharp LR35902 as its Central Processing Unit (CPU), clocked at 4.19MHz [6, Specifications]. In 1998, the Gameboy Color (GBC) was then released. Seen as the successor of the GB, it contains a screen of the same resolution, but supporting colour, from a palette of 32768 options (15 bits per color). It contains the same CPU as its predecessor, a Sharp LR35902, with now two modes: a 4.19MHz mode and a 8.38MHz mode (double-speed mode). This allows the GBC to be backwards compatible with most GB games - there are a few exceptions to this, games that used hardware bugs of the original GB that were fixed in the GBC.

From an emulation perspective, the Gameboy Color can thus be seen as an extension of the Game Boy - it has an identical CPU (although with a toggle-able double speed mode), and most of the memory layout is identical. To keep the remaining of this document simple, if not stated, "GB" will refer to both the original Gameboy and the Gameboy Color, as they are very similar. Dot Matrix Game (DMG) refers to the original Gameboy model.

2.4 Existing Literature

2.4.1 Gameboy Documentation

The Gameboy is one of the best documented consoles for emulation, and an array of resources exist to get started writing a new one. Some useful resources explaining it's behaviour are:

- Pandocs[5] is a technical reference of how the GB works. It is extremely complete and covers a wide range of topics, so it is useful to get a global view of a problem. It is one of the most referenced pieces of literature on the console.
- GB CPU Instructions[6] is a table containing all instructions the GB CPU has, as well as information on the amount of cycles taken by the instruction, the bytes of memory used, the flags affected by the operation, and a description of the instruction.
- Gameboy Complete Technical Reference[7] (GBCTR) is an unfinished document that contains very detailed information on the CPU and other components of the GB. Although incomplete, it provides a much lower-level view of the details of the GB (compared to Pandocs), making it useful to emulate very specific behaviour like the cycle-by-cycle timing of

[1] http://fms.komkon.org/VGB/
[2] https://problemkaputt.de/gmb.htm
[3] https://problemkaputt.de/gmbhist.htm
[4] https://www.retroveteran.com/category/nintendo-game-boy-color/
[5] https://gbdev.io/pandocs/
[6] https://meganesu.github.io/generate-gb-opcodes/
[7] https://gekkio.fi/files/gb-docs/gbctr.pdf

the CPU.
- GB dev wiki[8] is a wiki containing additional information on the GB, including guides to making games and explanations on some hardware quirks, and in particular a very precise description of the Audio Processing Unit (APU).

2.4.2 Existing Emulators

A wide range of emulators for the GB and GBC already exist, many of them being open-source. These are useful when developing a new emulator, to see how they work internally. For performance reasons they're usually written in compiled languages, such as C++ and Rust, but some interpreted language alternatives exist. These emulators include:

- Game Boy Crust[9] is a simple GB emulator written in Rust. It is quite incomplete but has a comprehensive structure, so it's a good project to first figure out how emulators work.
- AccurateBoy[10] is a highly accurate emulator, in particular for its Picture Processing Unit (PPU) that has pixel-perfect accuracy.
- oxideboy[11] is another GB emulator written in Rust, that is much more complete and helpful for some edge cases.
- SameBoy[12] is one of the most accurate open source GB and GBC emulators, written in C. It is much more technically complex but still useful to understand edge cases, especially since it is the emulator I use and compare mine with.
- Mooneye GB[13] is a GB research emulator written in Rust. It passes most of the Mooneye test ROMs, making it helpful when encountering issues with these tests.
- GameBoy-Online[14] is a high-accuracy JavaScript emulator, that I used when unsure on how to interface the emulator with the browser (notably for the APU).
- Gameboy.js[15] is another JavaScript emulator. It is fairly simply and inaccurate, but is easily hackable. As such it's the emulator I used when starting the emulator, to compare mine with.
- rboy[16] is an emulator written in Rust, that I used when developing the APU to compare mine with, as it passes some test ROMs I struggled with.

2.4.3 Gameboy Test ROMs

A core set of resources to develop an emulator is the test ROMs for that console. These are valid source code for the console, that instead of playing a game will run a set of tests on the console. These tests are first written to pass on the physical console itself, and are then used to ensure they also pass on the emulator. This means issues in specific components can be easily diagnosed (so long as the rest of the emulator responsible for running the test ROM works itself). These test ROMs also have the advantage of being open source, meaning their source code can be referred to, to understand what they expect of the console.

An other advantage to using test ROM is that they tend to re-use the same framework across a given test suite to report results. This means the result of the test is usually logged somewhere

[8] https://gbdev.gg8.se/wiki/
[9] https://github.com/mattbruv/Gameboy-Crust
[10] https://github.com/Atem2069/accurateboy
[11] https://github.com/samcday/oxideboy
[12] https://github.com/LIJI32/SameBoy
[13] https://github.com/Gekkio/mooneye-gb
[14] https://github.com/taisel/GameBoy-Online/
[15] https://github.com/juchi/gameboy.js/
[16] https://github.com/mvdnes/rboy

in the console, and testing can easily be automated by inspecting specific registers/memory addresses, rather than having to store an "expect result" image for each test.

The test ROMs used for this project are:

- Blaarg test ROMs[17] are some of the most well-known and used GB test ROMs. They include tests for the CPU, the timings of instructions, and some other functionality.
- Mooneye test ROMs[18] is a very complete test suite, that verifies most components of the GB: CPU instructions, memory timings of specific instructions, behaviour of Memory Bank Controllers (MBCs), timing of the Object Attribute Memory (OAM) Direct Memory Access (DMA), PPU timings, timer timings, etc.
- Acid Test (DMG[19], GBC[20]) is a test that verifies the PPU of the GB displays data properly (to line-rendering accuracy), for both Gameboy and Gameboy Color displays.
- SameSuite[21] is a test suite that is valuable for its APU tests: it uses the PCM12 and PCM34 registers exclusive to the GBC to inspect the exact ouput of the APU (whereas other test ROMs tend to inspect the on/off status of the channels, which is much less accurate).

[17] https://github.com/retrio/gb-test-roms/
[18] https://github.com/Gekkio/mooneye-test-suite
[19] https://github.com/mattcurrie/dmg-acid2
[20] https://github.com/mattcurrie/cgb-acid2
[21] https://github.com/LIJI32/SameSuite/

Chapter 3

Requirements and Specification

3.1 Requirements

3.1.1 User Requirements

U1. Run Gameboy games to a satisfiable fidelity, with proper rendering and controls emulation.

U2. Run Gameboy Color games to a satisfiable fidelity, with proper rendering and controls emulation.

U3. Allow the user to run GB and GBC games for both a Gameboy and a Gameboy Color (ie. allow using a GBC for a `.gbc` file, and a GB for a `.gb` file).

U4. Allow the user to save the state of the game, to continue their playthrough later. The state can simply be saved as a downloaded file, and re-uploaded later to continue the game.

U5. Allow the user to change the speed at which the game is played: double speed mode, half speed mode, etc.

U6. Have some debug functionality, to inspect the state of the console at any given time.

U7. Allow users to pause the console, and add breakpoints to stop execution at specific moments.

U8. Allow the user to switch between rendering modes (nearest-neighbour, LCD display, scale2, etc.)

U9. Allow the user to switch the colour palette of the DMG emulation.

3.1.2 System Requirements

F1. The system can receive a ROM file, construct an instance of the emulated console, and run the code inside said ROM.

F2. The system emulates different components of the GB and GBC, with as much precision as possible (M-Cycle precision).

F3. The system renders the output of the emulator to a Web `<canvas />`.

F4. The system creates the required DOM elements for the web-app, and updates them as needed.

F5. The system listens to key presses and releases to emulate controls through the keyboard.

F6. For touch devices, the system may render buttons to simulate the console's controls.

3.1.3 Non-Functional Requirements

N1. The emulator should be accessible on computers through a web browser equipped with a recent version of JavaScript.

N2. The emulator should be accessible on mobile devices through a web browser equipped with a recent version of JavaScript.

N3. The emulator should be accessible on computers through a standalone app.

N4. Maximise the tests passed by the emulator (see Gameboy Test ROMs)

N5. Have the code be well documented, allowing new-comers to the project and to GB emulation to easily understand what is going on - if possible with links to relevant Gameboy emulation resources.

3.2 Specification

Code	Specification	Importance
U1	User can upload a GB ROM file (`.gb`), and the emulator will run the game. The keyboard can be used to control the game, and the output is displayed.	High
U2	User can upload a GBC ROM file (`.gbc`), and the emulator will run the game. The keyboard can be used to control the game, and the output is displayed.	Medium
U3	User can upload a GB ROM file (`.gb`). The user can switch between a DMG and a GBC emulator.	Medium
U4	User can press a button to download a save of their game (or, alternatively, the save can be stored inside the browser with a technology like IndexedDB[1])	Low
U5	User can select the speed of emulation, to dynamically accelerate/decelerate the game.	Medium
U6	User can see debug information of the emulator. This information includes the current tileset, background map, time to draw a frame, and register information.	Low
U7	User can pause the console emulation through a button. They can also input conditions for which the console should break execution.	Low
U8	User can dynamically switch the rendering filter via a dropdown button.	Low
U9	User can dynamically switch the colour palette of the GameBoy via a dropdown button.	Low
F1	A ROM file can be uploaded, is transformed into an `UInt8Array` (because the GB is an 8-bit system), and the appropriate object is created to run the code.	High
F2	Different components exists as different classes, respecting typical OOP principles such as encapsulation and inheritance when relevant.	High

[1] https://developer.mozilla.org/en-US/docs/Web/API/IndexedDB_API

Code	Specification	Importance
F3	A `<canvas />` element is created, and is updated with the output of the emulator after every frame is drawn (ie. at the start of each VBlank mode).	High
F4	The Preact[2] framework is used to handle the UI of the web-app.	High
F5	Listeners are added to the environment's `window` to listen to all key presses and releases. The emulator can then request for a control update, by reading the state of keys.	High
F6	If a touch device is detected, button are added to the UI and are used by the emulator as inputs.	Medium
N1	A deployed version of the web-app is accessible on a desktop browser and provides full functionality, via keyboard and mouse inputs.	High
N2	A deployed version of the web-app is accessible on a mobile device browser and provides full functionality, via touch controls.	Medium
N3	A downloadable version of the web-app can be used on a computer and provides full functionality, via keyboard and mouse inputs.	Low
N4	As many possible tests as possible should be passed, while ensuring previously passing tests don't start failing.	Medium
N5	Main methods and variables must be properly documented, and have links to appropriate online resources to documentation about said element.	High

[2] https://preactjs.com/

Chapter 4

Design

In this chapter we will outline the main components of the Gameboy and of this emulator. For the different GB components, we will briefly go over their role, and how they interact with other components and to what end.

In emulators, the different components are usually split in three parts:

- The CPU, responsible for reading instructions and changing the state of the console. This is what drives the emulation, as other components usually idle unless acted upon.
- Input and output components, such as the APU, the PPU and the joypad. These components interact with the outside user, by either outputing the game state, or reading inputs from the user.
- A memory system, that handles addressing within the console. This part is essential to ensure components are communicating between each other properly, since different addresses may map to different components.

Because the emulator should not rely on the environment it is running it to work, the functionality of the project can be split into two parts: the emulator core, that is responsible for simulating the Gameboy, and the emulator's frontend, that allows interfacing with the user, and may be changed to work for different platforms (see figure 4.1).

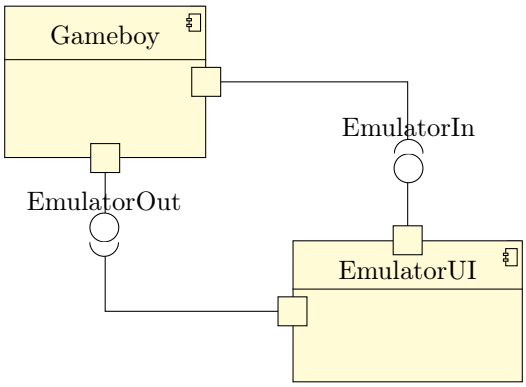

Figure 4.1: Split of emulator core and frontend

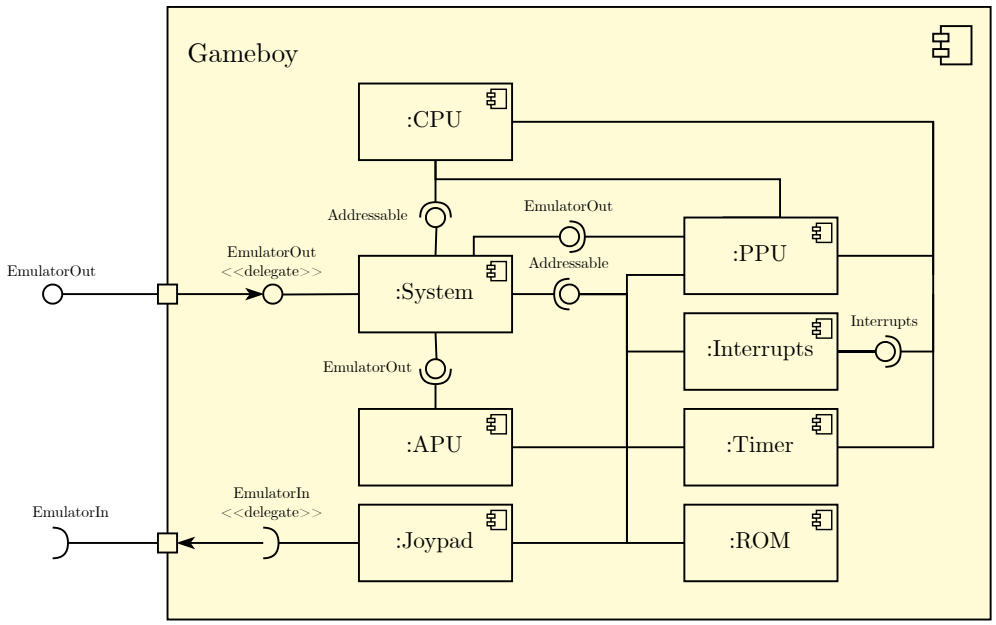

Figure 4.2: Simplified design of the emulator's core
Note utility classes have been excluded for readability purposes

4.1 Emulator Frontend

To allow the user to access the emulator, a frontend needs to be built with it. This frontend is responsible for outputting the graphics and audio data of the emulator, and also for receiving input from the user to pass down to the emulator. Aside from that, extra UI components have been added, to make the usage of the emulator more convenient and add more features: a play/pause button, custom themes, debugging tools, upscaling filters and keybinding options. All of these are kept independent from the emulator, and control transformations of the input/output, or changes in the running of the emulator.

Because this implementation is merely to allow access to the emulator and doesn't have any outstanding functionality, it's description will be brief - most of the focus will be put on the emulator's core.

4.2 CPU

The Central Processing Unit (CPU) is perhaps the most complex part of the GB. It is however also one of best documented parts of it, making it quite easy to emulate properly. The CPU model is a LR35902, which is a hybrid of the Intel 8080 chip and the Zilog Z80 [7].

Most emulator are typically "CPU-driven". This means that the CPU is what drives the emulation. If an instruction takes multiple hardware cycles to execute, the CPU is responsible for asking the rest of the system to tick at the appropriate time. In the design chosen for this project however, the emulator will be "System-driven", with the CPU only being part of the overall running of the emulator, and will not be responsible for ticking the rest of the system. This difference will be made obvious by the interfaces exposed to the CPU; it allows for nicer encapsulation and better separation of concerns.

The CPU is thus responsible for stepping through the program it's given. It has 6 different 16-bit

registers, with different roles. Some registers' bytes can be accessed independently, operating as two 8-bit registers or one 16-bit register as needed [6, CPU Registers and Flags].

- `AF`: the accumulator-flag register. Most arithmetic operations can be done on `A`. `F` holds the CPU's flags, and can only be altered through arithmetic opeartions.
- `PC`: the program counter register. It cannot be accessed directly, and is used to store the current position of the CPU in the program. It can be altered via `CALL`, `RET`, `JP` operations, for instance.
- `SP`: the stack pointer register. It can be modified or incremented, and stores the pointer to the top of the "stack". A typical use of the stack is for handling functions, by pushing the current address to the stack whenever a procedure is called, and popping it back to the stack pointer when it returns.
- `BC`, `DE`: two simple registers that can be used for arithmetic operations.
- `HL`: similar to `BC` and `DE`, it can also be used as an address by some operations - allowing the use of indirect addressing.

Aside from it's registers, the CPU needs to interact with system interrupts, and needs to access memory for reading/writing operations.

4.3 PPU, APU, Joypad

The Picture Processing Unit (PPU) is a complex part of the Gameboy, that handles outputting the game state to a 160×144 screen, that may either be monochrome on the DMG or with color support on the GBC. It may raise interrupts, and needs to be able to access memory, due to it being responsible for the Object Attribute Memory (OAM) - a feature allowing transfer of data from an arbitrary location in memory.

The Audio Processing Unit (APU) is responsible for playing the audio output of the console. It has a few registers to control it, as well as a short memory area called the wave RAM. It is partially controlled by the timer, but is otherwise independent. It's sound output is derived from four separate channels: two square wave channels, a custom wave chanel, and a noise channel [6, Audio]. These channels are here modelled as separate entities, but are purely internal.

The former two components need to output information to the user, either graphic or audio data. To do this in a portable way, the emulator will expose an interface, that can be implemented by the front-end. This allows the emulator core to be platform-agnostic, and be easily portable.

The joypad, on the other hand, is the one input component the GB has. It contains 4 directional arrow buttons, and 4 input buttons: A, B, start and select (see 4.3). To use these inputs, the CPU needs to use two different registers.

Similarly as for the output, the emulator core will expose an interface to read the user's input, ie. the state of all 8 buttons.

4.4 System Bus

Components withing the Gameboy need to communicate. To do this, components typically have a set of registers that control how they behave - for instance, the timer has a `TAC` register at `0xFF07` to control it's frequency. To allow this interaction without having all components directly depend on each other, a "system bus" component is created, responsible for handling all

Figure 4.3: Picture of a Gameboy, with the screen and joypad visible

components and managing memory-addressing among them.

This component, called `System`, is thus responsible for instantiating the other components, and providing them access to eachother, via a read and write method.

4.5 Other Components

Aside from the aforementioned components, the GB has a few components that allow it to run but aren't part of the CPU or the input/output components. These will be briefly outlined here, as they have a lesser impact on the emulator's design.

4.5.1 Timer

The timer, responsible for incrementing a counter (the `DIV`) at a set frequency. It can raise interrupts, and owns a few registers.

4.5.2 Interrupts

The interrupt system, that allows components such as the timer or the PPU to interrupt the CPU to run another process. It can also be enabled and disabled by the CPU, with the `EI` and `DI` instructions.

This sub-system is designed as a separate component to reduce coupling between other components, and instead have a small object that can be passed to components as needed.

4.5.3 MBC

The Memory Bank Controller (MBC) is an extra chip contained within some game cartridges, to allow access to more ROM data (as well as external RAM in some cases) via banking [6, MBCs]. There are multiple different MBCs, and so it is convenient to define a common interface for all of them, which can then be implemented according to specification for each MBC type.

The type of the MBC is in the header of the ROM [6, The Cartridge Header]. To separate this logic from the base system, a `ROM` class is used. It is responsible for reading the cartidge's

header, and creating the appropriate MBC instance.

4.6 Useful Classes

To have similar interfaces over all components, we will also declare specific classes and interfaces to be implemented by them.

4.6.1 Addressable

The `Addressable` interface provides a read and write method. This allows all components to communicate between each other without needing to be aware of what the component they're communicating with is. Aside from the CPU, all components of the emulator implement this method.

4.6.2 Memory and registers

Simple utility classes can be declared to manage memory in a simple way.

The `RAM` class implements `Addressable`. It can be instantiated with a set size, and can be read and written to. It is used in components that have large blocks of writable data.

The `Register` class implements `Addressable`, but ignores the address parameter of the read and write operations, as it contains only one byte. A `DoubleRegister` class may also be implemented, backed by two `Register`s, to provide support for 16-bit registers, used for instance in the CPU.

Chapter 5

Implementation

The project uses Preact[1], a light-weight alternative to the more popular React[2] framework. Preact was chosen because the front-end of the web-app is extremely lightweight, so a smaller framework with less features is enough. This also avoids bloating the app with a heavy framework such as React: its GZipped and minified size is around 31.8Kb, while Preact is only 4Kb (87% less).

The language used for the project is TypeScript[3], a typed version of JavaScript. This is essential for the project, as ensuring the correctness of code can be extremely hard without proper typing constraints, especially as a project grows in size.

The project is divided into two parts:

- The `frontend/` directory contains the UI for the web-app. The main logic to create the emulator and run it is contained in `app.tsx`.
- The `emulator/` directory contains the actual GB emulator. Although most classes and interfaces used are exported, only three elements are needed to properly interact with the emulator:
  - `GameBoyColor.ts` handles the core loop of the system. It contains the `GameBoyColor` class, the emulator. Instantiating the emulator creates all the necessary sub-components, and calling the `drawFrame()` method runs the emulator for one frame (0.16 seconds).
  - `GameBoyOutput.ts` contains a simple interface, with optional methods to receive any output produced by the emulator (see figure 5.1). The two main methods of this are `receiveGraphics` and `receiveSound`, which use the output of the actual console.
  - `GameBoyInput.ts` contains a simple interface with a required `read()` method that returns an object with the current inputs for the console (see figure 5.2).

5.1 Emulator Frontend

The frontend of the emulator is written in Preact, allowing us to create a simple, fast and lightweight UI to control it. It contains the emulator title, the control buttons and the emulator video output (see figure 5.3). Along the left of the screen is a sidebar with more options, allowing

[1] https://preactjs.com/
[2] https://reactjs.org/
[3] https://www.typescriptlang.org/

```
interface GameBoyOutput {
    receiveGraphics?(data: Uint32Array): void;
    receiveSound?(data: Float32Array): void;

    // Debugging methods:
    debugBackground?(data: Uint32Array): void;
    debugTileset?(data: Uint32Array): void;
    serialOut?(data: number): void;
}
```

Figure 5.1: `GameBoyOutput` interface methods

```
type GameBoyInputRead = {
    up: boolean;
    down: boolean;
    left: boolean;
    right: boolean;

    a: boolean;
    b: boolean;
    start: boolean;
    select: boolean;
};

interface GameBoyInput {
    read(): GameBoyInputRead;
}
```

Figure 5.2: `GameBoyInput` interface method

the user to customise the emulator to their needs and debug the state of the console if needed.

5.1.1 Main Controls

The main controls for the emulator are the 6 buttons above the screen. These are, from left to right:

- Pause/play: pauses or resumes the emulator.
- Step by one CPU instruction: this is a feature useful for debugging the emulator, when precise information of the state of the emulator is needed.
- Sound toggle: allows enabling or disabling the sound of the emulator. By default the sound is turned off, because modern browsers don't allow websites to play any sound before the user interacts with the website [8].
- Triple speed toggle: a toggle button that speeds up the emulator to emulate the GB thrice as fast as usual. This is a common feature found in most emulators.
- Save state: saves the current state of the cartridge in the browser's storage, allowing the user to resume playing the game later. Note this doesn't save the full state of the emulator, but that of the cartridge, making it equivalent to a real-life save on the GB where only the battery-backed storage of the cartridge persists through power-offs.

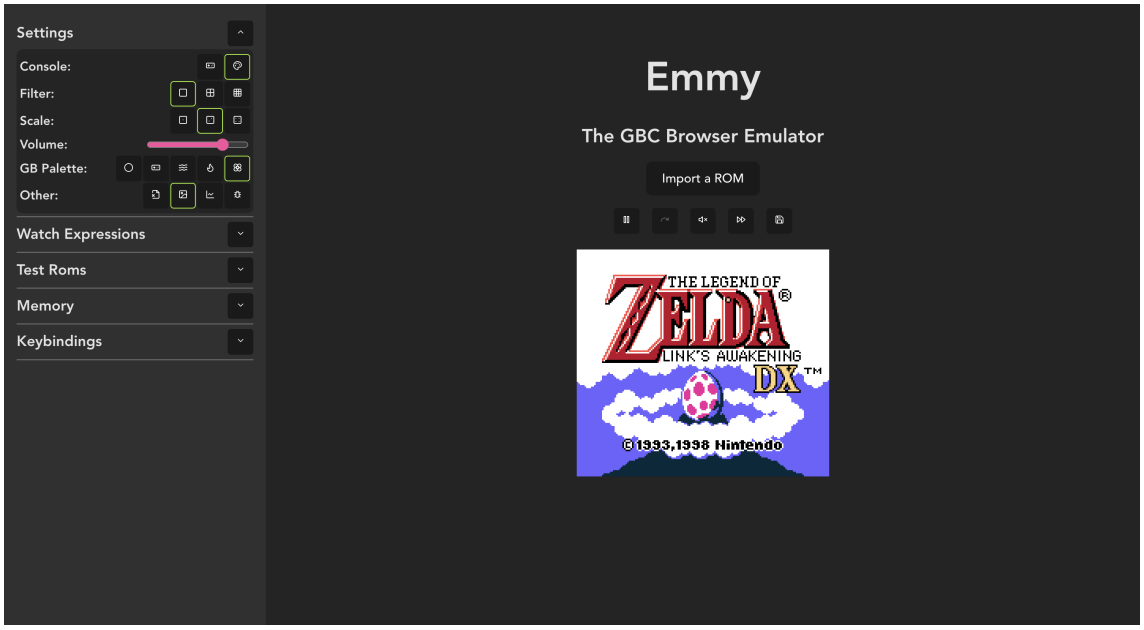

Figure 5.3: Home page of the emulator

5.1.2 Settings

In the side drawer, the first tab is "Settings". It contains general settings for the emulator, such as:

- the console used by the emulator. This may either be the DMG or the GBC.
- an extra filter to be applied to the output. This increases the resolution of the screen, using the Scale2x or Scale4x[4] algorithm (see 5.4).
- the size of the emulator's screen, allowing resizing to two or four times larger.
- a slider to change the volume of the emulator's audio output.
- a palette selector, to change the four hues of the DMG's screen.
- miscellaneous togglable settings, grouped together:
  - an option to play with or without the initial boot ROM of the emulator.
  - a toggle to enable frame-blending, meaning for every new frame the output is mixed with the previous frame. This is a nice addition to have, because certain games made certain objects flicker on screen to make them appear translucent (since the flicker wasn't visible to the eye).
  - a button to show the performance statistics of the emulator. This is mainly useful when developing the emulator to make sure it is still efficient.
  - a toggle to enable the debug view of the emulator, where the currently loaded tileset and background map are displayed (see 5.5).

5.1.3 Watch Expressions

The second drawer allows the user to define custom JavaScript functions to inspect the state of the emulator regularly (see 5.7). This requires knowledge of the inner structure of the emulator, as the field names need to be used, but is quite useful when needing to inspect parts of the console that aren't the memory, like internal counters used for components, or register values.

[4] https://www.scale2x.it/

Figure 5.4: Output of the GBC with, from left to right: no filter, Scale2x and Scale4x

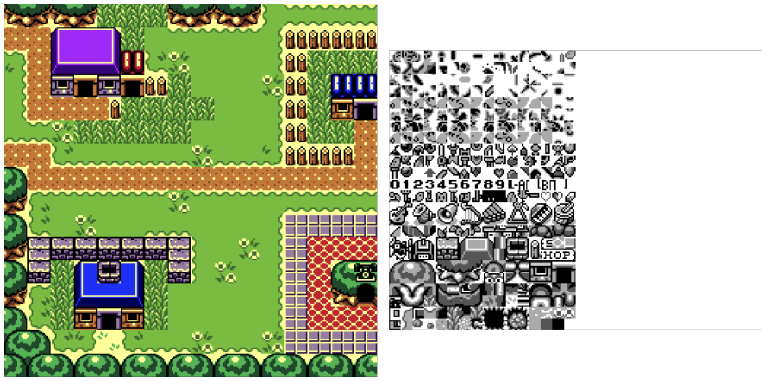

Figure 5.5: Debug view of the emulator

To implement this, the `Function`[5] constructor is used, which takes in a string with the function's code. This allows the user to dynamically change the expression, and the component will simply update the function, without needing to reload the whole application. These expressions are also automatically saved to `localStorage`[6], meaning they will be kept between sessions.

The user-defined function is then repeatedly invoked, and the result output below the expression, allowing for a live-status of the emulator. If an error is thrown by the function (due to a null value, or invalid expression), the error is caught and '`Error`' is displayed.

5.1.4 Test ROMs

To ensure emulators work properly, a variety of test ROMs have been made, that test most aspects of the GB (see 2.4.3). The front-end of the emulator supports running a large number of them in an automated way (see **??**). The user can select the group of tests desired, and they will run internally (without receiving any input our outputting anything directly). The status of each individual test is displayed, and the user can click on the test name to run it on the main emulator, for further debugging or inspection.

To run the tests, an emulator instance is created, with the test ROM as the input data and "spy" objects provided for input and output - these spies store the output of the emulator, to evaluate the state of the test. To stop execution, the emulator's state is inspected regularly, and the outcome of the test is checked (with execution terminating if the test takes too long to end).

[5] https://developer.mozilla.org/en-US/docs/Web/JavaScript/Reference/Global_Objects/Function/Function

[6] https://developer.mozilla.org/en-US/docs/Web/API/Window/localStorage

Figure 5.6: Watch expressions drawer

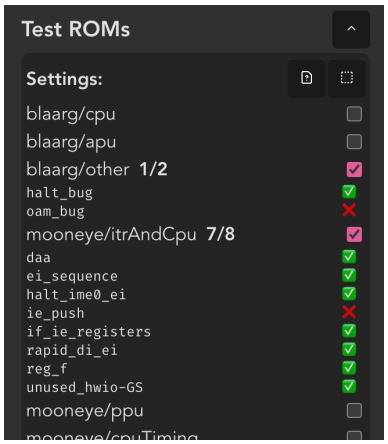

Figure 5.7: Test ROMs drawer

To detect the outcome of each test, these are grouped according to what test suite they belong to. Each test suite is then associated to a *success function*, that can either return `"success"`, `"failure"` or `null` if no outcome has been reached yet. As such, all tests in a test suite have a similar way of reporting success and failure.

The currently automated test suites are:

- The Blaarg test ROMs[7]. The success function of this suite is quite complex, as this suite doesn't have a standard way of outputting the result. As such multiple parts of the emulator are inspected simultaneously:
  - `"Passed"`[8] and `"Failed"`[9] may be output to the console's serial port.
  - If the memory at `0xa001-0xa003` is equal to `0xdeb061`, then the byte stored at `0xa000` is the status of the test[10].
  - For the `halt_bug` test, there doesn't seem to be anywhere where the result is output,

[7] https://github.com/retrio/gb-test-roms
[8] See lines 50-54 of `mem_timing/source/common/testing.s` in https://github.com/retrio/gb-test-roms
[9] See lines 112-139 of `mem_timing/source/common/runtime.s` in https://github.com/retrio/gb-test-roms
[10] See 'Output to memory' in `dmg_sound/readme.txt` of https://github.com/retrio/gb-test-roms

so the graphical output of the emulator is verified.

- The Mooneye test ROMs[11] and SameSuite[12], of which all tests are verified the same way. In case of success, the `B`, `C`, `D`, `E`, `H` and `L` registers hold the values 3, 5, 8, 13, 21 and 34 respectively (Fibonacci's sequence). If they instead all hold the value `0x42`, the test failed[13].
- The Acid test ROMs (DMG[14], GBC[15]). These need to be tested graphically, as their purpose is verifying the actual output of the emulator rather than it's behaviour.

Thanks to this, the emulator's frontend supports a total of 191 automated test ROMs, that verify the behaviour of most of the emulator. Of these 191 tests, 101 pass.

5.1.5 Memory Inspect

When debugging an emulator, being able to inspect it's memory is essential, as some bugs may be caused by the wrong mapping of components, or a fault when writing data. To help debugging this, the frontend comes with a basic memory inspection tool, that can show the entirety of the *addressable* data of the emulator (see 5.8). This means data that is not accessible by the game (for instance, because the appropriate ROM bank is not selected) cannot be inspected here. A simple offset can also be indicated, to restrain the data to a certain area.

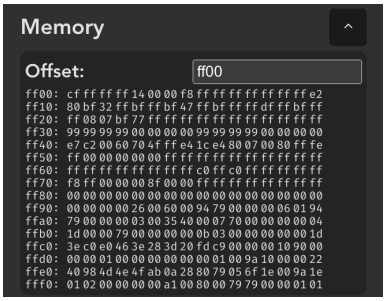

Figure 5.8: The memory inspection drawer

5.1.6 Keybindings

The user can also customise their keybindings for the emulator, by mapping each input to a separate key (see 5.9). This is only relevant on keyboard-equipped devices.

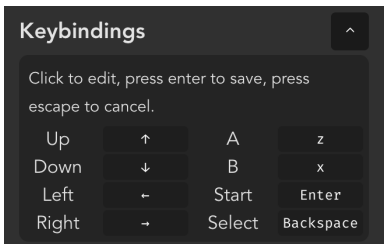

Figure 5.9: The keybindings settings

[11] https://github.com/Gekkio/mooneye-test-suite
[12] https://github.com/LIJI32/SameSuite/
[13] See https://github.com/Gekkio/mooneye-test-suite/#passfail-reporting for the Mooneye test suite, see lines 265-281 of `include/base.inc` in https://github.com/LIJI32/SameSuite/ for the SameSuite
[14] https://github.com/mattcurrie/dmg-acid2
[15] https://github.com/mattcurrie/cgb-acid2

5.2 CPU

The CPU is the core of the emulator, and allows running the code from the ROM by reading the operation code (or opcode) and executing the matching action.

5.2.1 Instruction Set Decoding

Because the Gameboy is an 8-bit system, opcodes are 8 bits (or one byte) long, giving in theory a maximum of $2^8 = 256$ operations. However in the GB the operation `0xCB` gives access to an extended instruction set, meaning that when reading `0xCB` the CPU will read the next byte and use a different logic to execute the operation. This means there are now $2^8 - 1 + 2^8 = 511$ operations. The GB also has 11 'unused' opcodes, that will lock the console when used[6, CPU Comparison with Z80], meaning there are in total $2^8 - 12 + 2^8 = 500$ operations to implement.

Multiple techniques exist to handle this large number of operations:

- Have a large switch-statement for all possible operations. This is the most straight-forward option, but can result in quite large switch-statements, especialy for CPUs with more opcodes. The GB has comparatively few opcodes as it is an 8-bit system; the Gameboy Advance, the console that followed it, had 32-bit opcodes: way too many for a switch statement to be appropriate.
- Decode the operation by reading specific parts of the opcode, and generate the instructions dynamically. This method is what is done for larger instruction sets, where opcodes can be split into separate parts to describe the operation's behaviour [9, ARM CPU Reference]. It however comes with the cost of extra processing for each instruction, as it needs to be decoded. This method is applicable to the GB, as many instructions follow a pattern - see, for instance the `LD` instructions, that all follow the same order of registers: `B`, `C`, `D`, `E`, `H`, `L`, `(HL)` (the byte at address `HL`) and `A`.
- Create a dispatch table a map that associates each opcode to a function to execute [10]. This technique is also only viable for small numbers of opcodes, as the map can result quite large. An advantage of this method is that the map doesn't have to be explicitly written out entirely - generators can be used to populate some chunks of it, making it a hybrid of the two previous solutions: there is very little overhead to execute an instruction, as each opcode is associated to a function, but there is also no need to write out every instruction separately, as the map can be generated in parts or in its entirety (inducing a light setup cost).

Initially, the emulator had a large map, with all the opcodes as keys. The functions associated would then execute the instruction and return the number of cycles taken by the instruction (see figure 5.10). This however proved quite repetitive and prone to errors. To solve this, a *generator function* was created. To use it, two parameters must be given: a map of opcodes to values (of an arbitrary type), and a helper function that for each arbitrary value returns a function that executes the instruction (see figure 5.11). This allows generating repetitive instructions more easily, by only specifying what opcodes and objects are used, and not what the whole body of the instruction is (see figure 5.12). This method is used in other emulators - for instance, Sameboy has a map that has an opcode-function mapping for all opcodes, and uses macros to generate the different functions[16].

[16] https://github.com/LIJI32/SameBoy/blob/master/Core/sm83_cpu.c

```
protected instructionSet: Partial<Record<number, InstructionMethod>> = {
    // NOP
    0x00: () => 1,
    // LD BC/DE/HL/SP, d16
    0x01: (s) => { this.regBC.set(this.nextWord(s)); return 3; },
    0x11: (s) => { this.regDE.set(this.nextWord(s)); return 3; },
    0x21: (s) => { this.regHL.set(this.nextWord(s)); return 3; },
    0x31: (s) => { this.regSP.set(this.nextWord(s)); return 3; },
    ...
}
```

Figure 5.10: Initial instruction set implementation

```
protected generateOperation<K extends number, T>(
    items: Record<K, T>,
    execute: (r: T) => InstructionMethod
): Record<K, InstructionMethod> {
    const obj: Record<K, InstructionMethod> = {};
    for (const [opcode, item] of Object.entries(items)) {
        obj[opcode] = execute(item);
    }
    return obj;
}
```

Figure 5.11: Generator function used for the CPU

5.2.2 A cycle accurate CPU

The Gameboy is a memory-bound system, meaning that it is limited by it's memory accesses. The CPU can only execute either one read or one write per M-Cycle [11, CPU core timing]. The CPU also needs to retrieve the opcode for each instruction, which takes an additional cycle, meaning an instruction performing no memory accesses lasts one cycle, and an instruction performing n memory accesses lasts at least $n + 1$ cycles. Note also that the GB overlaps the last cycle of the execution with the fetching cycle for the next opcode - this will thus be of importance when implementing the fetching of the opcode.

The current implementation is not M-Cycle accurate: the CPU instruction is executed as one monolithic block, rather than in different smaller parts. This becomes crucial when the timer, OAM and PPU are involved, as they run in parallel with the CPU, so memory accesses to these components may return different values depending on the M-Cycle.

In all emulators that were found when researching for this project, M-Cycle accuracy is reached by making the emulator "CPU-driven". What this means is that inside each instruction, between each cycle, the CPU is responsible for ticking the rest of the system - the main loop is then only responsible for continuously running the CPU, and nothing else. This approach is probably the simplest and most straightforward one, as it is quite simple to implement - all one needs to do is call the system tick method when relevant (see figure 5.14).

Some emulator where this method was found:

- In Mooneye GB, this can be seen by the usage of the `CPUContext.tick_cycle()` method

```
protected instructionSet: Partial<Record<number, InstructionObject>> = {
    // NOP
    0x00: () => 1,
    // LD BC/DE/HL/SP, d16
    ...generateOperation(
        {
            0x01: this.regBC,
            0x11: this.regDE,
            0x21: this.regHL,
            0x31: this.regSP,
        },
        (register) => (s) => {
            register.set(this.nextWord(s));
            return 3;
        }
    ),
    ...
}
```

Figure 5.12: Improved instruction set implementation

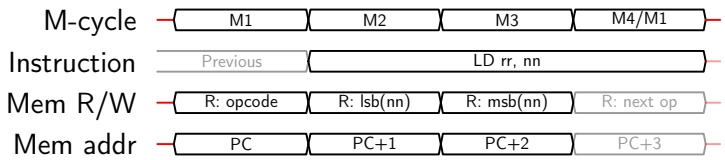

Figure 5.13: Timing of the `LD rr, nn` instruction
Taken from GBCTR [11, Sharp SM83 instruction set]

when appropriate[17].

- In accurateboy, the `Bus.tick()` method is called from within the instructions (or it is done implicitly by methods such as `Bus.read()`)[18].
- In Sameboy, the CPU uses functions like `cycle_read` or `cycle_no_access` to tick the rest of the system[19].

```
protected ld_bc_hl(system: System) { // LD BC, (HL)
    const lower = system.read(this.regPC.inc());
    system.tick();
    const upper = system.read(this.regPC.inc());
    system.tick();
    this.regBC.set(upper << 8 | lower);
}
```

Figure 5.14: CPU-driven `LD BC, (HL)`

This solution however comes with the cost of coupling the CPU to the system, or at least to a way of ticking said system. A particularity of this emulator is that the CPU is almost autonomous,

[17] See https://github.com/Gekkio/mooneye-gb/blob/master/core/src/cpu.rs
[18] See https://github.com/Atem2069/accurateboy/blob/master/accurateboy/CPU.cpp
[19] See https://github.com/LIJI32/SameBoy/blob/master/Core/sm83_cpu.c

in that it doesn't interact with any other functionality of the system, except the interrupts. The only other interface it requires being an `Addressable`, to allow access to memory. This also means the emulator is no longer "CPU-driven", since the CPU *cannot* tick the rest of the system. It is instead up to the root system to tick all of the components (including the CPU). This is thus a code quality improvement, that isn't present in other emulators.

This was done by splitting all the instructions into their respective steps, and have each step return the next step (or `null` if it's the last step). The CPU now must simply store whatever the instruction returns: if it's `null` then it needs to fetch an instruction to prepare for the next cycle, otherwise it's a function and must be executed (and it's result stored for the next step). See figures 5.15 and 5.16 for an example of this.

In the final version of the CPU's tick method, a method `loadNextOp` is implemented. It not only handles getting an instruction associated with the desired opcode, but it also is responsible for checking if an interrupt was raised (in which case the interrupt handling procedure occurs [6, Interrupts]), and for ensuring the CPU isn't halted before running an instruction.

```
protected ld_bc_hl(system: Addressable) { // LD BC, d16
    const lower = system.read(this.regPC.inc());
    return () => {
        const upper = system.read(this.regPC.inc());
        return () => {
            this.regBC.set(upper << 8 | lower);
            return null;
        }
    }
}
```

Figure 5.15: System-driven `LD BC, d16`

```
step(system: Addressable, interrupts: Interrupts) {
    if (this.nextStep === null) {
        // looks up opcode in instruction table
        const nextStep = this.loadNextOp(system, interrupts, verbose);
        if (nextStep === "halted") return true;
        this.nextStep = nextStep;
    }
    this.nextStep = this.nextStep(system, interrupts);
    if (this.nextStep === null) {
        // opcode is fetched on the last cycle of execution
        this.currentOpcode = system.read(this.regPC.inc());
    }
}
```

Figure 5.16: System-driven step of the CPU. Note how the opcode is fetched directly when the next instruction is over, to emulate the overlap between the fetch and the execute steps.

5.3 System

The `System` class implements the system bus, as well as the ticking of all other components. It handles the ticking of the PPU, timer, APU and interrupt logic. Furthermore, it is the

component that links all of the data together: whenever a component is ticked (any of the above or the CPU), the `System` instance is passed, so that the components can read and write to the rest of the system. It implements `Addressable` (see **??**), and internally has a `getAdress` method that returns the `Addressable` at this specific address, to be accessed in the `read` and `write` methods. This ensures that the logic used to determine what component is accessed depending on the address isn't duplicated.

5.3.1 `getAddress` optimisation

Because `System` is a higly used component and is accessed for almost every read and write, the `getAddress` method is under a lot of pressure: for a game like 'The Legend of Zelda: Link's Awakening DX', the method is called around 650,000 times per second. For a more complex and resource-intensive game such as 'Alone in the Dark: The New Nightmare', around 1,900,000 calls are made. This intensity in usage can easily be explained, as the GB is a memory-bound system: almost all interactions between components occur through memory reads and write. `getAddress` must thus be optimised as much as possible, as it is one of the main bottle-necks of the emulator.

An initial implementation of `getAddress` used the combination of a list of if-statements for different ranges of the Memory Map, as well as a map where all register addresses where mapped to the component responsible for them. The code would first check if the key exists in this map,,and if not it would then go through a series of if-conditions (see figure 5.17). It would not return a tuple, containing both `Addressable` to use and the address within in - this was needed as some components had a particular mapping to memory. This is the case for the Work RAM (WRAM), who's memory starts at `0xC000` and ends at `0xFDFF` (spanning over 15.5KB bytes) despite it only being 8KB long, because the last 7.5KB of its address range map back to the beginning of it (this is called the Echo RAM [12]).

```
protected getAddress(pos: number): AddressData {
    const register = {
        0xff00: this.joypad,
        ...
        0xffff: this.intEnable,
    }[pos];
    if (register !== undefined) return [register, pos];

    if (0x0000 <= pos && pos <= 0x7fff) return [this.rom, pos]; // ROM Bank
    ...
    if (0xfe00 <= pos && pos <= 0xfe9f) return [this.oam, pos]; // OAM

    // Unmapped area, return 'fake' register
    return [{ read: () => 0xff, write: () => {} }, 0];
}
```

Figure 5.17: Initial implementation of `getAddress`

This proved quite costly:

- This code creates a new map every time it is called, when checking for the registers' addresses.
- Having to return both an `Addressable` and an address is quite costly, as a new array with

two items must be created every time. It also adds unecessary complexity to the method, as the system bus shouldn't be responsible for handling the details of address mapping, and instead simply delegate the task to the appropriate component.

- The if-conditions for the ranges of the biggest areas of memory (everything between `0x0000` and `0xEFFF`) happened after the register checks, which delayed response for these reads and writes. This is all the more important that this range represents $\frac{\texttt{0xEFFF}}{\texttt{0xFFFF}+1} \approx 93\%$ of the total address space.
- Chaining if-conditions is unneficient, as the JS engine must step through all conditions and check the values each time. Furthermore, although having both the lower and upper bound of the memory section indicated in the condition (e.g. `0x0000 <= pos && pos <= 0x7FFF` for [`0x0000`; `0x7FFF`]) makes the translation from Memory Map to code easier, it is slower, since only the upper bound of the range is needed for the condition if all the ranges below have already returned.

First, we may move the fine-grained addressing logic to sub-components like the WRAM. This removes the need to return an address from the method: only an `Addressable` is enough, as that component will then be responsible for decoding the address further.

The last two points may then be addressed, by removing the use of if-conditions, and moving this part of the code to the beginning of the function.

With the memory map of the console (see figure **??**) we can notice how the largest chunks of memory all end with the last three nibbles of the address as `0xFFF`. This means that the component mapped to an address can be determined by simply looking at the first nibble of said address. This is probably done to simplify the circuitry responsible for addressing, as only the most significant 4 bits need to be compared. The emulator can take advantage of this.

Figure 5.18: Memory map for the largest chunks of memory [12]

Start	End	Description
0000	3FFF	16 KB ROM bank 00
4000	7FFF	16 KB ROM Bank 01∼NN
8000	9FFF	8 KB Video RAM (VRAM)
A000	BFFF	8 KB External RAM
C000	CFFF	4 KB Work RAM (WRAM)
D000	DFFF	4 KB Work RAM (WRAM)
E000	FDFF	Mirror of 0xC000∼0xDDFF

For this area of memory (`0x000-0xEFFF`) the system bus may simply isolate the most significant nibble, and then use a map to associate each address block to a component. Only if the address corresponds to `0xF---` do we try matching it to a register address. This removes the need for the if-conditions, and is also faster as it is evaluated much earlier on (see figure 5.19). The map can be created on instantiation and kept for later use, to avoid unnecessary memory allocations.

To ensure placing the "main block" first was the best choice, measurements have been taken of four different GB games. The emulator would log all memory accesses by groups of $100\,000\,000 = 10^8$, and categorise them by address "blocks": the main block, the OAM block, the illegal block (called like this because this area of memory is restricted by Nintendo, and only returns `0x00`) and the register block. This separation is justified by the fact these are the five chunks of memory that must be checked separately, due to their irregular boundaries. The *second* set of 10^8 accesses was then used to gather the statistics of what blocks are used the most. The first 10^8 accesses aren't used, as they may include setup operations that only happen when the game

```
protected getAddress(pos: number): Addressable {
    // Checking last nibble
    let addressable = this.addressesLastNibble[pos >> 12];
    if (addressable) return addressable;

    // Registers
    if((address & 0xff00) === 0xff00) {
        addressable = this.addressesRegisters[pos & 0xff];
        if (addressable) return addressable;
    }

    if (pos <= 0xfdff) return this.wram;  // Echo RAM
    if (pos <= 0xfe9f) return this.ppu;   // OAM
    if (pos <= 0xfeff) return Register00; // Illegal Area

    return RegisterFF; // fake register
}
```

Figure 5.19: Optimised implementation of `getAddress`

loads, and as such don't represent what the average execution will look like. see figure 5.20 for the results.

Figure 5.20: Rounded access rate per memory blocks, for the second set of 10,000,000 accesses of 4 different games. Games are, respectively, "Alone in the Dark: The New Nightmare", "The Legend of Zelda: Link's Awakening DX", "Tetris" and "Pokémon Silver".

Name	Memory Area	Game 1	Game 2	Game 3	Game 4
Main Block	0x0000-0xEFFF	88.869%	95.679%	84.289%	95.740%
Echo RAM	0xF000-0xFDFF	0.000%	0.000%	0.000%	0.000%
OAM Block	0xFE00-0xFE9F	0.001%	0.000%	0.000%	0.000%
Illegal Block	0xFEA0-0xFEFF	0.001%	0.000%	0.000%	0.000%
Register Block	0xFF00-0xFFFF	11.129%	4.321%	15.711%	4.260%

As we can see, the great majority of memory accesses go to the main part of memory, with almost all of the rest going to the "register area", the last 256 bytes of the address space. This can easily be explained by the fact all registers that allow interaction with other components (the PPU, the APU, the timer, etc.) are in this narrow range, so it is bound to have a high usage. It is also quite interesting to note that neither the Echo RAM or the OAM block are used at all, except very rarely for one of the tested games. It is thus safe to assume that the conditions responsible for mapping these areas can be left at the end of the function, as they will rarely match an address.

5.3.2 Evaluating the `getAddress` optimisation

A simple experiment was then run, to verify the performance improvement. The first 25 million instructions of the `cpu_instrs`[20] test ROM were run. This sample was chosen because it is considerably large, because the test itself requires around 25 million instructions to complete. For the measurement, `window.performance.now()`[21] was used before and after each drawn

[20] https://github.com/retrio/gb-test-roms/tree/master/cpu_instrs
[21] https://developer.mozilla.org/en-US/docs/Web/API/Performance/now

frame, and the values were then summed.

The result was the following: 33 955.9ms before the change, and 20 039.1 after the change. The relative difference is thus $\frac{20\,039.1-33\,955.9}{33\,955.9} = -0.4098$, thus reducing time taken by 40.98%. By measuring the time spent within `getAddress` for these 25 million instructions, we get a total spent time in the method of 0.000 151ms on average, showing that this method is no longer a bottleneck for the system, as its impact on performance is minimal.

5.4 PPU

The Picture Processing Unit (PPU) is the component responsible for rendering the game onto the Gameboy's screen. It is one of the most complex components of the GB, with intricate timings, and a behaviour that changes between the DMG and the GBC, due to the addition of colour-support, as well as Video RAM (VRAM) banking.

5.4.1 Presentation

The screen rendering is divided intro three layers, drawn on top of each other. From bottom to top, these are:

- The background, a 256×256 image loaded into memory that support scrolling on both axis.
- The window, similar to the background, is a 256×256 image that can be moved around the screen, and is toggleable. It however doesn't support scrolling.
- The objects, that are drawn at the very top, are smaller 8×8 tiles. These can be moved freely, and support some transformations, such as horizontal or vertical flipping. Their data is stored in the Object Attribute Memory (OAM).

It renders onto the screen line by line, meaning it will draw a 1-pixel wide line (the scanline) from left to right, and then proceed to the line below. Drawing an entire line takes 114 M-cycles, and there is an 1140 M-cycle delay when the bottom of the screen is reached. The advantage of having per-line rendering is that by modifying the position of the background between each drawn line, complex visual effects can be obtained quite easily [13], see figure 5.21 for an example.

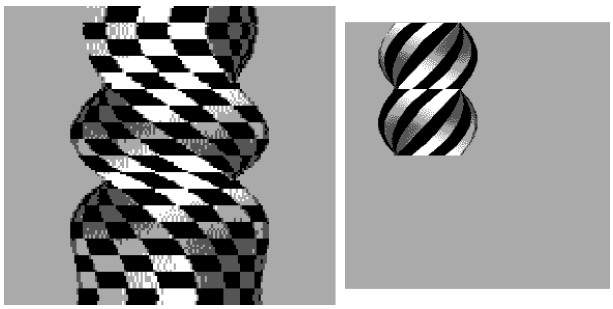

Figure 5.21: Visual effects made by scrolling the background
Gameboy output (right), and loaded background data (right)

One of the main characteristics of the PPU is that it operates in modes. During each mode, it perform a different set of operations, and can only be interacted with in certain way [6, LCD Status Registers]. These four modes are:

- Mode 2: the PPU looks through the OAM, to look for all the sprites to draw. During this

time, the OAM is not accessible.
- Mode 3: the PPU reads both the Video RAM (VRAM) and OAM, and draws the line. None of the video data is accessible here.
- Mode 0: the PPU does nothing for a short length of time after each line. An interrupt is raised at the start of this period, allowing the CPU to be notified now is the time to update what's rendered if necessary. This is called the horizontal blank, or HBlank.
- Mode 1: the PPU does nothing for a long length of time after the bottom of the screen is loaded. This is usually where the bulk of the graphics update go, as it lasts a considerable amount of time (1140 M-Cycles). An interrupt is also raised here, to notify the CPU. This is called the vertical blank, or VBlank.

Rendering of the line is thus done during mode 3, when no graphics data is accessible, using a FIFO queue (First In First Out) of pixels that can be pushed in and out as the data is read. Some of the registers are still writable, but most games don't use them during this mode because they would require very precise timing. An example of a game relying on this is "Prehistorik Man", that changes the colour palette registers during the scanline in its intro-scene [14, Tricky-to-emulate games].

Since this is only the case for a minority of games, we can go past this inaccuracy. We'll thus take advantage of this to simplify greatly the rendering logic and make the renderer "scanline based", drawing an entire line all at once. This shortcut is commonly used by emulators that aren't too accuracy-oriented.

To handle each mode separately, we have four distinct "mode" objects of type `PPUMode`, that hold some basic information on the operation of the mode: it's length, it's flag for the `STAT` register and the name of the method in `PPU` responsible for ticking said mode (see figure 5.22). Note `KeyForType<T, V>` is the union of all keys `k` of `T` such that `T[k]` is of type `V`.

```
type PPUMode = {
    doTick: KeyForType<PPU, (interrupts: Interrupts) => void>;
    flag: number;
    cycles: number;
};
```

Figure 5.22: Type definition of `PPUMode`

5.4.2 Rendering Logic

The first step needed to render the game is to select and order all the sprites to be shown on screen. One of the limitations of the GB is that despite having enough space in memory for 40 sprites, only 10 may be rendered at a time in a scanline [6, OAM]. The sprites are selected during mode 2 - as such, during the last cycle of mode 2, the emulator will look through the available sprites in the OAM, and select the appropriate ones. This can be done quite elegantly in a functional manner, using an array (see figure 5.23).

First, we retrieve the sprites from the OAM. Internally, sprite data is cached between scanlines and invalidated when written to - `getSprites()` updates the dirty tiles in the cache and returns them. This avoid decoding the sprite every line if it doesn't receive changes, while also ensuring excess decoding isn't done if the sprite is modified twice between scanlines. We then select the first 10 sprites to be part of the scanline - this selection is done based on index: the PPU looks through the OAM sequentially to find matching sprites [6, OAM]. Finally, a re-ordering of the

```
this.readSprites = this.oam
    .getSprites()
    .filter((sprite) => sprite.y <= y && y < sprite.y + objHeight)
    .slice(0, 10)
    .map((sprite, index) => [sprite, index])
    .sort(objPrioritySort)
    .map(([sprite]) => sprite);
```

Figure 5.23: Retrieve the selected sprites for the scanline

sprites needs to be done, to determine which sprite goes above which - this may depend on the sprite's position in the OAM, or on it's X coordinate. This is handled in `objPrioritySort`. To allow sorting based on both attributes, the sprites must be briefly remapped to a tuple with the sprite and it's index (as JavaScript doesn't expose the indices of elements when ordering them). These sprites are then stored in the PPU object, to be used when rendering at the end of mode 3.

The PPU uses a form of indirect addressing to handle data. It has an area in VRAM that stores image data (tiles), 8×8 images. Whenever one of these tiles needs to be used, for the background, or for an object, the identifier of the tile needs to be used - this identifier being derived from the tile's address. This allows for the very simple re-use of tiles, which is particularly relevant for backgrounds where they may be repeated a lot. This background data is stored in a *tile map*, a 256×256 map of 1-byte indices to the actual tile data. The GB has two such maps, and both the background and window can display either of them - this is controlled via a flag in the `LCDC` register.

To render the scanline, the layers need to be drawn on top of each other: background, then window, then objects. Due to the similarities between background and window, a `drawLayer` method can be shared to handle the bulk of the work.

This method will loop over all tiles it needs to draw in the current scanline. It first needs to determine the tile's index. This index can be used to the retrieve the tile's address in VRAM. It can also be used to fetch the attributes of the tile; this is a GBC-exclusive feature, that allows transforming tiles (for instance, flipping them, or selecting a different palette for the tile). These attributes are stored in the second bank of VRAM, at the same address, in a single byte of data. The DMG/GBC distinction isn't needed however, because the attributes for a `0x00` value match the attributes used by the DMG. This means that instead of having two different branches to split the console versions, we can setup the DMG emulator to always return `0x00` for the second bank of VRAM, and keep the GBC handling of attributes.

Once the tile index has been acquired, the PPU may retrieve the tile data address from VRAM. This one-byte address needs to be converted to a valid address in the tile data range of VRAM, as it contains two partially overlapping tile data areas - a flag in the `LCDC` register controls this. Once the address where the tile data (ie. it's texture) is obtained, the PPU may finally retrieve said data, and draw it in the scanline.

This tile data, an 8×8 images where each pixel can be one of four shades is encoded in 16 bytes. Because tile data is rarely changed and the procedure to extract the image data from these 16 bytes is quite complex and requires some bit manipulation, tiles are cached.

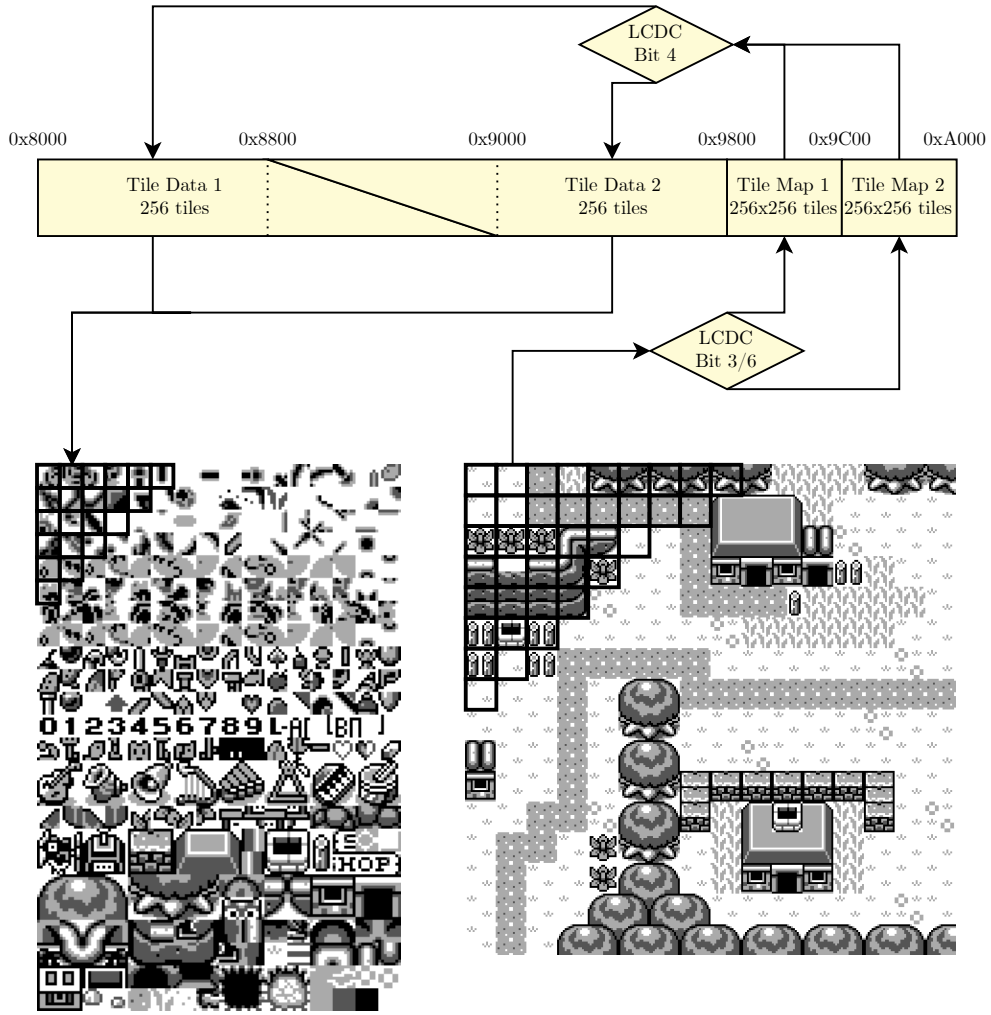

Figure 5.24: PPU logic to obtain tile data
This applies to the DMG - the CGB has two VRAM banks, thus having twice as many tile textures, and an additional tile attribute map.

5.4.3 Subcomponents

Although it is present to the system as a unit, the PPU is implemented as multiple smaller-components, that handle different parts of the screen logic. This is all the more important that the PPU is extremely complex, so care is needed to ensure the code remains maintainable.

The Object Attribute Memory (OAM) is the memory area in the PPU that stores object data. This small 160-byte long area has a feature called OAM Direct Memory Access (DMA), allowing for direct transfers from ROM or RAM to the OAM. Because this transfer runs in parallel to the other components, the OAM itself is its own component, contained inside the PPU. This allows for a better separation of concerns between the two - the PPU handles rendering, the OAM handles OAM DMA and sprite storing and decoding.

The colour management of the PPU also requires some extra logic that can be extracted into its own class. Although on the DMG this is limited to splitting a byte into 4 shades of 2 bits (white, light gray, dark gray and black), the GBC has a 16 palettes of 4 colours, 8 palettes for the background and window, and 8 for sprites [6, Palettes]. As such, a `ColorController` class was created - it implements `Addressable`, and has a method to retrieve the palette for

the background, and a method to retrieve the palette of a sprite. This class is extended by a `DMGColorController`, and a `CGBColorController`. The `PPU` picks one of the two on instantiation, based on the emulated console, and can then use both, without needing to know if the screen is monochrome or coloured.

Similarly, the VRAM behaves slightly different between the DMG and the GBC: the latter's VRAM has two banks it can switch between, via an additional register, and also supports VRAM DMA transfers, similar to what is possible with the OAM. A `VRAMController` class is created - it also implements `Addressable`, and has a few extra methods, to allow getting tiles from it, and reading each bank individually (see figure 5.25).

```
abstract class VRAMController implements Addressable {
    read(pos: number): number;
    write(address: number, value: number): void;

    tick(system: Addressable, isInHblank: boolean, isLcdOn: boolean): boolean;

    getTile(tileAddress: number, bankId: 0 | 1): Int2[][];
    readBank0(pos: number): number;
    readBank1(pos: number): number;
}
```

Figure 5.25: Interface of `VRAMController`

5.5 APU

The Audio Processing Unit (APU) is the component responsible for producing the sounds and music of the GB. It's output is the combination of four different channels, each with their own properties: two square channels, a wave channel and a noise channel [6, Audio]. These channels run independently, can be turned on or off individually, and are merged to form the output.

Most features of the channels can be derived to a timer, that ticks down from a value and has an effect when reaching 0. And although they differ in output, all channels have some attributes in common. For example, they all have a *length timer*, that turns off the channel when it reaches 0. Channels 1, 2 and 4 also have an *envelope*, that allows updating the volume of the channel at a set frequency. To have these behaviours work across all channels, an abstract class `SoundChannel` was created. It holds the registers all channels have in common, as well as methods like `tickLengthTimer` to handle common mechanisms across all channels.

To return their output, all channels have a `getSample` method that returns the current output value of the channel, a value between 0 and 15. This output can then be combined, and output to the frontend to handle. It is of interest to note that the GB is always outputting sound, with a resolution of 1 M-Cycle, ie. 2^{20}Hz $\approx$ 1.05MHz. This is significantly more than the sample rate computers usually use, 44.1Hz. To fit to this standard, the emulator's APU has an internal counter that ticks down and only produces a sample every $\frac{2^20}{44.1*10^3} = 23.7$ cycles. This isn't the most accurate way of producing the audio data, but is far simpler and faster than downsampling the audio. An example of very accurate emulator that does downsampling is SameBoy[22].

[22]See Accuracy, https://sameboy.github.io/features/

5.6 MBCs and ROMs

A majority of GB cartridges came shipped with a Memory Bank Controller (MBC). This was done to circumvent the memory limit of the GB: due to it's 16-bit addresses, only memory in `0x0000-0x7FFF` is mapped to the ROM, limiting it's size to 32KB. MBCs provided a way of extending this memory limit, by doing *bank switching* [6, MBCs]. The cartridge holds more data than it can address, and the CPU can modify which part of the memory it's accessing by writing to a register in the MBC. Because ROMs are read-only, the write-operations can be safely re-used to instead write to these built-in registers.

On top of more memory space, MBCs would sometimes come with additional features, such as external RAM, a battery (to be able to keep the state of the RAM between sessions, which allows saving the game), a rumble motor, or a battery-backed real time clock, for example. The most common MBCs found were the MBC1, with space for a maximum of 2MB [6, MBC1], and the MBC5 - the only MBC to officially support the GBC's double speed mode [15]. see figure 5.26 for statistics of the usage of different MBCs.

Figure 5.26: Statistics of different MBCs [16]

Name	Count	Percentage
No MBC	2150	23.0%
MBC1	4010	43.0%
MBC2	227	2.4%
MBC3	367	3.9%
MBC5	2532	27.1%
Others	53	0.6%

To ensure the emulator can run as many games as possible, the main MBCs have been implemented: MBC1, MBC2, MBC3 and MBC5 (although support for the MBC3's real time clock hasn't been added). Since the GB interacts with the cartridge in the exact same way whether there is an MBC or not, the choice of making a `MBC` abstract class came quite naturally: it's interface is that of `Addressable`, and it also supports a `save` and `load` method, for save files.

To decide which MBC to use, we have a `GameCartridge` class, that wraps around `MBC`. It's role is decoding the cartridge's header, contained in `0x0100-0x014F` [6, The Cartridge Header], to decide on which MBC to use, as well as some of the properties of the cartridge: does it have RAM, is it battery backed (in which case it supports saving), what's the game's title and identifier.

The `read` and `write` methods received an optimisation similar to what was done with `System.getAddress` - because the internal addressing logic of the MBCs only relies on the most significant nibble, a simple switch statement can be made (see figure 5.27).

Internally, the `MBC`s have a `ROM` instance to store the cartridge data (this is obtained from the ROM uploaded by the user), as well as an optional RAM. This RAM is what the game can edit. If it is battery backed, memory is kept when the GB is turned off, allowing data like scoreboards and progress to be saved. To replicate this saving behaviour, the emulator exposes a `save` and `load` method, that respectively return and set the data in the RAM. The emulator's core is thus not responsible for handling these saves, and the frontend can decide freely how to store them.

In the implemented frontend, this is done by saving the RAM data in the browser's available

```
read(pos: number): number {
    switch (pos >> 12) {
        case 0x0: // ROM bank 00
        case 0x1:
        case 0x2:
        case 0x3:
            return this.data[pos & addressMask];
        case 0x4: // ROM bank 01-ff
        case 0x5:
        case 0x6:
        case 0x7: {
            const address =
                (pos & ((1 << 14) - 1)) |
                (this.romBankLower8.get() << 14) |
                (this.romBankUpper1.get() << 22);
            return this.data[address & addressMask];
        }
        case 0xa: // ERAM
        case 0xb: {
            if (this.ramEnable.get() !== RAM_ENABLED) return 0xff;
            const address = this.resolveERAMAddress(pos);
            return this.ram.read(address);
        }
    }
    throw new Error(`Invalid address`);
}
```

Figure 5.27: `read` method of `MBC5` [6, MBC5]

storage, using the "localForage" library[23]. When a ROM is loaded, the frontend checks if a save for this game exists, by using the ROM's identifier decoded in the `GameCartridge` class. If it does, the save is then loaded. Similarly, when changing ROMs, closing the window, or pressing the "Save" button, the frontend retrieves the RAM data from the emulator and saves it, by setting the key of the entry in the local storage to the identifier of the ROM.

5.7 Timer

The timer is a component in the GB that ticks regularly. It allows the control of two independent mechanisms [6, Timer and Divider Registers].

First is the *divider counter*, accessed and controlled via the `DIV` register. This counter is internally 16-bits, although only the upper 8-bits can be accessed. It is incremented every machine cycle (ie. increased by 4 every M-Cycle), can be read, and writing to the divider resets the counter (see figure 5.28).

The second, more complex part of the timer is a customisable timer. It is made of three registers:

- `TIMA`: the timer counter. Every time a falling edge is detected on one of the bits of

[23] https://github.com/localForage/localForage

```
protected divider = new DoubleRegister(0xab00);
protected addresses: Record<number, Register> = {
    0xff04: this.divider.h, // we only ever read the upper 8 bits
    ...
};

tick(interrupts: Interrupts): void {
    const newDivider = wrap16(this.divider.get() + 4);
    this.divider.set(newDivider);
    ...
}

write(pos: number, data: number): void {
    if (pos === 0xff04) { // Writing anything to DIV clears it.
        this.divider.set(0);
        return;
    }
    ...
}
```

Figure 5.28: Implementation of divider counter

`DIV`, this register is incremented. What bit is inspected can be controlled via the `TAC` register, effectively changing the frequency of the timer. When this register overflows (is incremented when equal to `0xFF`), an interrupt is raised.

- `TAC`: the timer control register. It allows enabling or disabling the timer, and changing the inspected bit of `DIV`.
- `TMA`: the timer modulo. It defines what value `TIMA` is reset to when overflowing.

Although it is a seemingly simple system made of 3 registers, the behaviour of this system is quite complex, as it has multiple edge cases [6, Timer obscure behaviour]. The code that handles the increase of `TIMA` is thus quite long, and required some re-factoring to be easily readable (see figure 5.29). As can be seen from this implementation, the `TIMA` register can actual be incremented for multiple reasons (lines 19–22). A falling edge on the inspected bit of `DIV` is one of the reasons (when `bitStateBefore==0` and `bitStateAfter!=0`, and the timer is enabled), but if the timer is disabled while the bit was in a high position then the register is also incremented. It is also of interest to note that `TIMA` is not set to the value of `TMA` when overflowing - this is actually done the following cycle (lines 2–9). This logic requires some extra fields to be present, recording the values of the registers on the previous tick. Some additional logic is required in the `write` method, because writes to `TIMA` are ignored when the timer overflows.

5.8 Helpful Components

To avoid repeating logic throughout the main components, small classes have been written. One of these is the `Register` class - it holds a single integer, that can be read or written. It also comes with a `flag` method, that returns whether a given flag is set in the register, as this is a frequent operation.

This class is extended by `MaskRegister`, a class to support registers of which all bits aren't used. This is the case for instance for `TAC`: bits 0–2 control the timer, and bits 3–7 are hardwired to 1,

```
tick(interrupts: Interrupts): void {
    this.previousTimerOverflowed = false;
    if (this.timerOverflowed) {
        const modulo = this.timerModulo.get();
        this.timerCounter.set(modulo);
        interrupts.requestInterrupt(IFLAG_TIMER);
        this.timerOverflowed = false;
        this.previousTimerOverflowed = true;
    }

    const timerControl = this.timerControl.get();
    const speedMode = (timerControl & 0b11);
    const checkedBit = TIMER_CONTROLS[speedMode];

    const bitStateBefore = (this.previousDivider >> checkedBit) & 1;
    const bitStateAfter = (newDivider >> checkedBit) & 1;
    const timerIsEnabled = timerControl & TIMER_ENABLE_FLAG;

    if (
        bitStateBefore &&
        (timerIsEnabled ? !bitStateAfter : this.timerWasEnabled)
    ) {
        const result = (this.timerCounter.get() + 1) & 0xff;
        this.timerCounter.set(result);
        if (result === 0) this.timerOverflowed = true;
    }

    this.timerWasEnabled = timerIsEnabled;
    this.previousDivider = newDivider;
}
```

Figure 5.29: Code to manage the `TIMA` increments

and cannot be reset. A `MaskRegister` allows defining this behaviour directly in the register, avoiding needing additional logical in the class using the register (here, `Timer`). On instantiation, a *mask* is passed as an argument, and is applied whenever the value of the register needs to be changed (see figure 5.30).

To accomodate the 16-bit registers of the CPU, a `DoubleRegister` class was implemented. It is made of two `Register` instances, and provides methods to get and set its value as if it held a 16-bit number (despite being implemented as two 8-bit numbers). This provides flexibility to the CPU, enabling both 16-bit arithmetic and 8-bit arithmetic, without having to manage the way the register is implemented.

Classes to manage memory were also created, to avoid having high-level components access data structures directly. For instance, a `ROM` class was made. It is backed by a `UInt8Array`[24], and implements `Addressable`. It is extended by `RAM`, to allow `write` operations.

To simplify the operation of certain components, a `CircularRAM` class was also created. It

[24] https://developer.mozilla.org/en-US/docs/Web/JavaScript/Reference/Global_Objects/Uint8Array

```
class MaskRegister extends Register {
    protected mask: number;

    constructor(mask: number, value: number = 0) {
        super(value | mask);
        this.mask = mask;
    }

    override set(value: number): void {
        super.set(value | this.mask);
    }
}
```

Figure 5.30: Implementation of `MaskRegister`

extends `RAM`, and must be provided an *offset* on creation. When reading or writing to it, the offset is substracted from the address, and the modulo of the RAM's length is then applied to it (see figure 5.31). This allows this part of memory to handle addressing in an independent way: the high-level component simply has to provide it with it's address on creation.

```
class CircularRAM extends RAM {
    protected offset: number;

    constructor(size: number, offset: number, data?: Uint8Array) {
        super(size, data);
        this.offset = offset;
    }
    override read(pos: number): number {
        return super.read((pos - this.offset) % this.size);
    }
    override write(pos: number, data: number): void {
        super.write((pos - this.offset) % this.size, data);
    }
}
```

Figure 5.31: Implementation of `CircularRAM`

This can be used to implement the WRAM. It can be addressed from `0xC000` to `0xFDFF`, but `0xE000`–`0xFDFF` maps back to `0xC000`–`DDFF`. As such, a `CircularRAM` with an offset of `0xC000` and a length of `0x2000` will properly wrap address in the `0xE000`–`0xFDFF` range back to the beginning of its memory. This avoids handling this logic in high-level components, and provides a generic solution to this kind of memory component. Other places where this class is used include the wave RAM of the APU, and the VRAM of the PPU.

Chapter 6

Evaluation

Acronyms

APU Audio Processing Unit.

CPU Central Processing Unit.

DMA Direct Memory Access.

DMG Dot Matrix Game.

GB Gameboy.

GBC Gameboy Color.

MBC Memory Bank Controller.

OAM Object Attribute Memory.

OS Operating System.

PPU Picture Processing Unit.

ROM Read-Only Memory.

VRAM Video RAM.

WRAM Work RAM.

Glossary

Dot Matrix Game The model name of the Gameboy. It is often used to refer to the "base" Gameboy (in contrast with GBC for the Gameboy Color).

M-Cycle An M-Cycle (or machine cycle) is the smallest step the CPU of the Gameboy can do. Because all instructions of the CPU are multiples of 4, instruction lengths and timings are usually referred to in M-cycles (e.g. `LD A, B` takes 4 T-cycles, thus 1 M-cycle).

Memory Map The memory map is what determines where each address leads to - it can be seen as a list of non-overlapping ranges.

Picture Processing Unit The part of the Gameboy that is responsible for rendering the game.

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