

MSc Project Notes

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1 Memory model constructor cheatsheet

In italic, “primitives”, ie. types that aren’t memory models.

Note $X^? \stackrel{\text{def}}{=} X \uplus \perp$, $X^\emptyset \stackrel{\text{def}}{=} X \uplus \emptyset$

Language	Memory Model
WISL	$\text{PMap}(\text{Loc}, \text{OneShot}(\text{List}(\text{Exc}(\text{Val}))))$
JSIL	$\text{PMap}(\text{Loc}, \text{PMap}(\text{Str}, \text{Exc}(\text{Val}^\emptyset)) \times \text{PMap}(\text{Loc}, \text{Ag}(\text{Val})))$

Name	Purpose	Type	Actions	Predicates
Exc	Exclusive ownership of a specific resources	$\tau^?$	load , store ^{*1}	PointsTo
Ag	Multiple parties agree on the same value for a resource	τ		Agree
Frac	Allow partial (readonly) ownership of an object	$\tau \times (0, 1]$		Frac
List	Ensure continuous memory allocation	$(\mathbb{N} \xrightarrow{\text{fin}} \tau) \times \mathbb{N}^? \text{ } ^{*2}$		
OneShot	The program only has one go at something (eg. freeing memory)	$\text{Exc}(\tau) + \text{Ag}(\{\emptyset\})$	free	
PMap	Define memory as a map of address (a sort I) to value	$(\mathbb{I} \xrightarrow{\text{fin}} \tau) \times \mathcal{P}(\text{I})^? \text{ } ^{*3}$		lift with index in-param
Product (\times)	Two simultaneous states, each being updated separately (eg. List)	$\tau_1 \times \tau_2$	lift with A1 , A2	
Sum ($+$)	Either of two states existing (eg. OneShot)	$\tau_1 \uplus \tau_2$	lift with A1 , A2	

^{*1} Would we define **load** and **store** at this level, or at a more primitive “Value” memory model level?

^{*2} Full definition: $\{(b, n^?) \in (\mathbb{N} \xrightarrow{\text{fin}} \tau) \times \mathbb{N}^? \mid \text{dom}(b) \subseteq [0, n^?]\}$

^{*3} Full definition: $\{(h, d) \mid h \in (\mathbb{I} \xrightarrow{\text{fin}} \tau) \wedge d \in \mathcal{P}(\text{I})^? \wedge \text{dom}(h)^? \subseteq d\}$