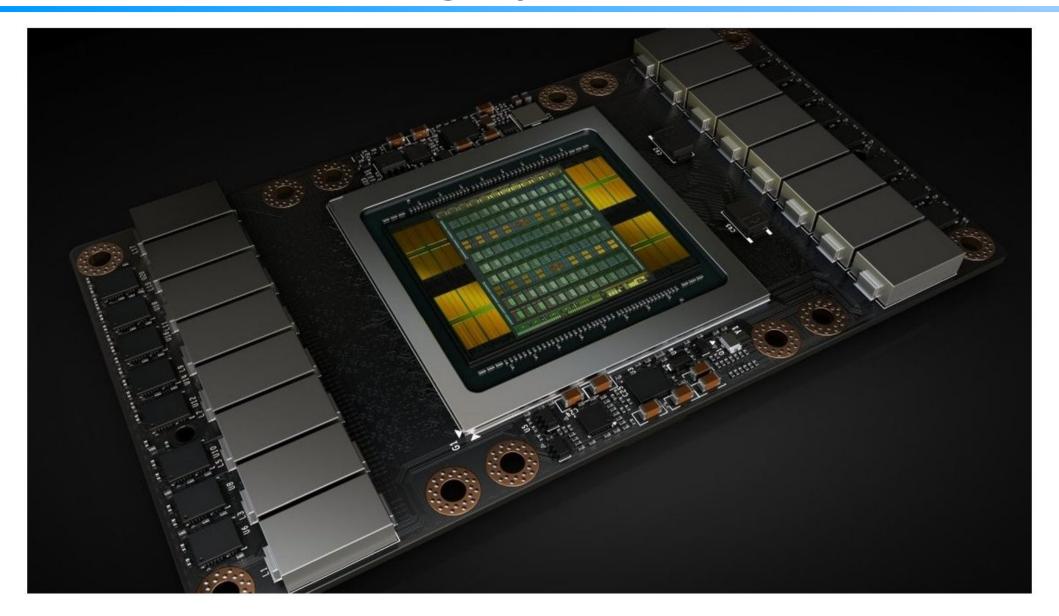
Optimization Techniques for Parallel Code 2. Introduction to GPU architecture

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http://www.irisa.fr/alf/collange/ sylvain.collange@inria.fr

OPT - 2017

What is inside a graphics card?



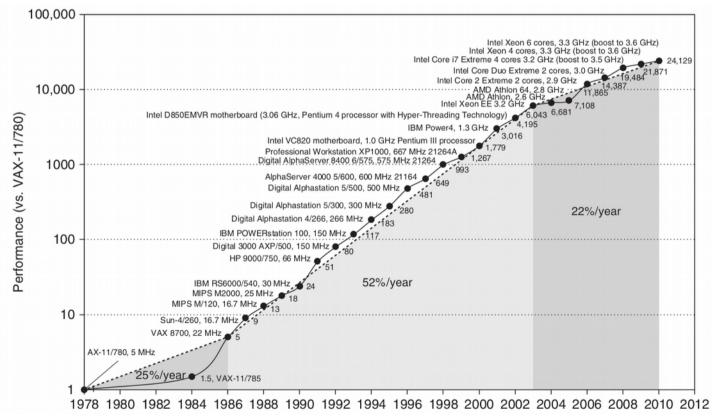
NVIDIA Volta V100 GPU. Artist rendering!

Outline

- GPU, many-core: why, what for?
 - Technological trends and constraints
 - From graphics to general purpose
 - Hardware trends
- Forms of parallelism, how to exploit them
 - Why we need (so much) parallelism: latency and throughput
 - Sources of parallelism: ILP, TLP, DLP
 - Uses of parallelism: horizontal, vertical
- Let's design a GPU!
 - Ingredients: Sequential core, Multi-core, Multi-threaded core, SIMD
 - Putting it all together
 - Architecture of current GPUs: cores, memory
- High-level performance modeling

The free lunch era... was yesterday

- 1980's to 2002: Moore's law, Dennard scaling, micro-architecture improvements
 - Exponential performance increase
 - Software compatibility preserved



Hennessy, Patterson. Computer Architecture, a quantitative approach. 5th Ed. 2010

Do not rewrite software, buy a new machine!

Technology evolution

Memory wall

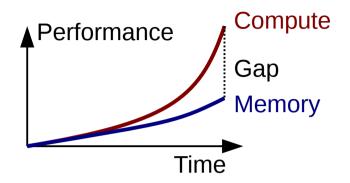
- Memory speed does not increase as fast as computing speed
- Harder to hide memory latency

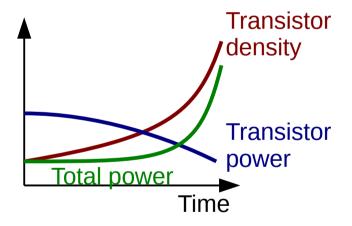
Power wall

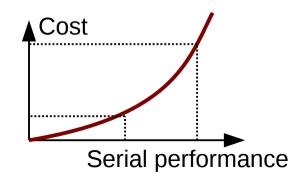
- Power consumption of transistors does not decrease as fast as density increases
- Performance is now limited by power consumption

ILP wall

- Law of diminishing returns on Instruction-Level Parallelism
- Pollack rule: cost ≃ performance²







Usage changes

- New applications demand parallel processing
 - Computer games : 3D graphics
 - Search engines, social networks...
 "big data" processing
- New computing devices are power-constrained
 - Laptops, cell phones, tablets...
 - Small, light, battery-powered
 - Datacenters
 - High power supply and cooling costs





Latency vs. throughput

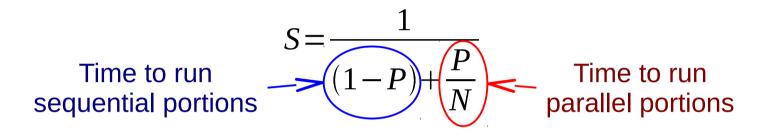
- Latency: time to solution
 - Minimize time, at the expense of power
 - Metric: time e.g. seconds
- Throughput: quantity of tasks processed per unit of time
 - Assumes unlimited parallelism
 - Minimize energy per operation
 - Metric: operations / time e.g. Gflops / s
- CPU: optimized for latency
- GPU: optimized for throughput



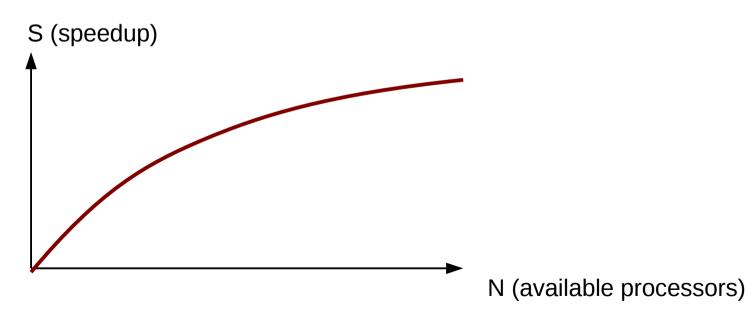


Amdahl's law

Bounds speedup attainable on a parallel machine



- S Speedup
- P Ratio of parallel portions
- N Number of processors

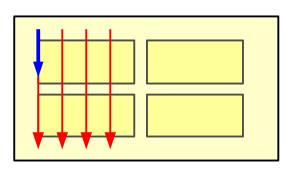


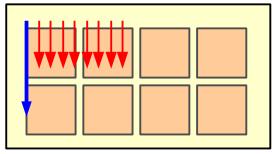
G. Amdahl. Validity of the Single Processor Approach to Achieving Large-Scale Computing Capabilities. AFIPS 1967.

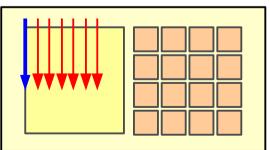
Why heterogeneous architectures?

Time to run sequential portions $S = \frac{1}{(1-P) + (P)}$ Time to run parallel portions

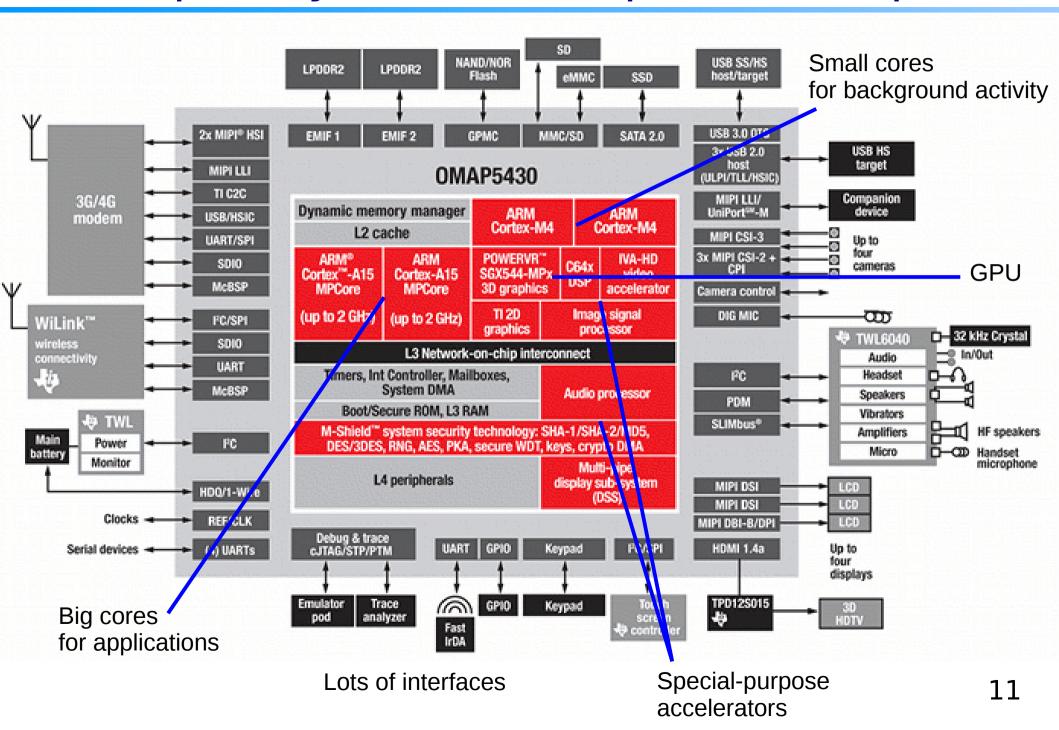
- Latency-optimized multi-core (CPU)
 - Low efficiency on parallel portions: spends too much resources
- Throughput-optimized multi-core (GPU)
 - Low performance on sequential portions
- Heterogeneous multi-core (CPU+GPU)
 - Use the right tool for the right job
 - Allows aggressive optimization for latency or for throughput







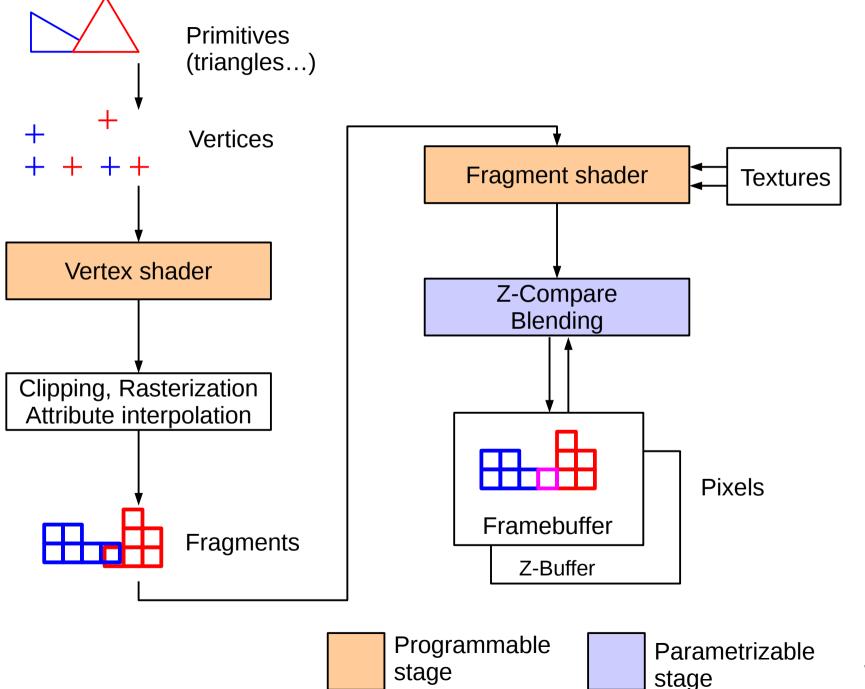
Example: System on Chip for smartphone



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The (simplest) graphics rendering pipeline



How much performance do we need

... to run 3DMark 11 at 50 frames/second?

Element	Per frame	Per second
Vertices	12.0M	600M
Primitives	12.6M	630M
Fragments	180M	9.0G
Instructions	14.4G	720G



- Intel Core i7 2700K: 56 Ginsn/s peak
 - We need to go 13x faster
 - Make a special-purpose accelerator

Beginnings of GPGPU

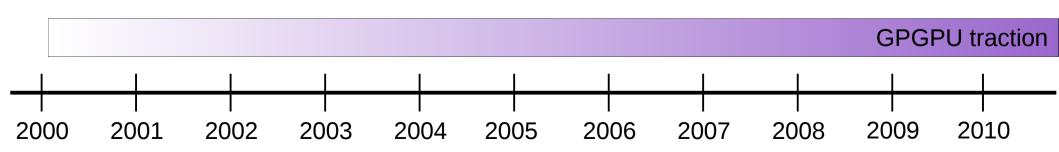
Microsoft DirectX

7.x	8.0	8.1	9.0 a	9.0b	9.0c	10.0	10.1	11
Unified shaders								

NVIDIA

NV10	NV20	NV30	NV40	G70	G80-G90	GT200	GF100
FP 16	Programmable shaders	FP 32	Dynamic control flow	SIMT	CUDA		





Today: what do we need GPUs for?

1. 3D graphics rendering for games

Complex texture mapping, lighting computations...

2. Computer Aided Design workstations

Complex geometry

3. High-performance computing

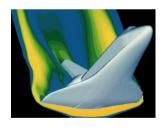
 Complex synchronization, off-chip data movement, high precision

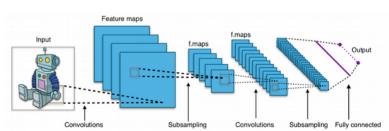
4. Convolutional neural networks

- Complex scheduling of low-precision linear algebra
- One chip to rule them all
 - Find the common denominator





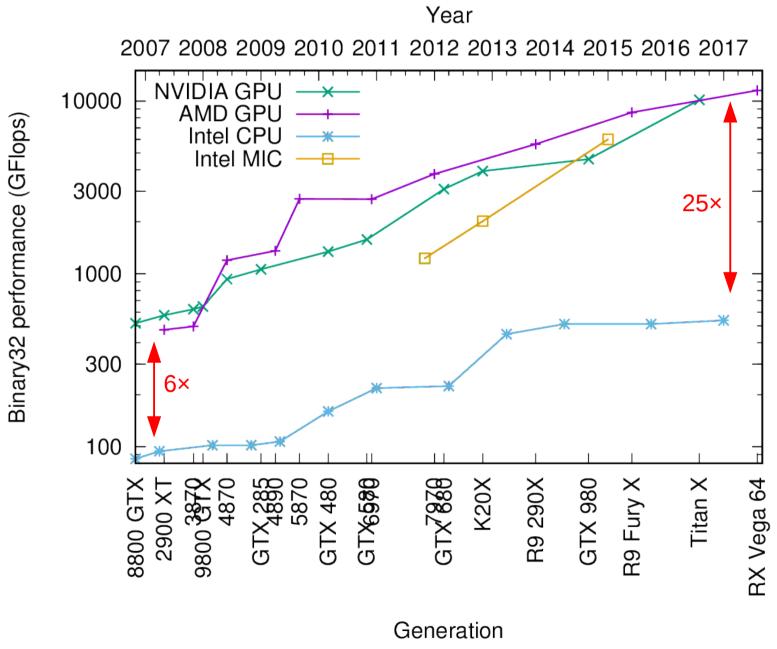




Outline

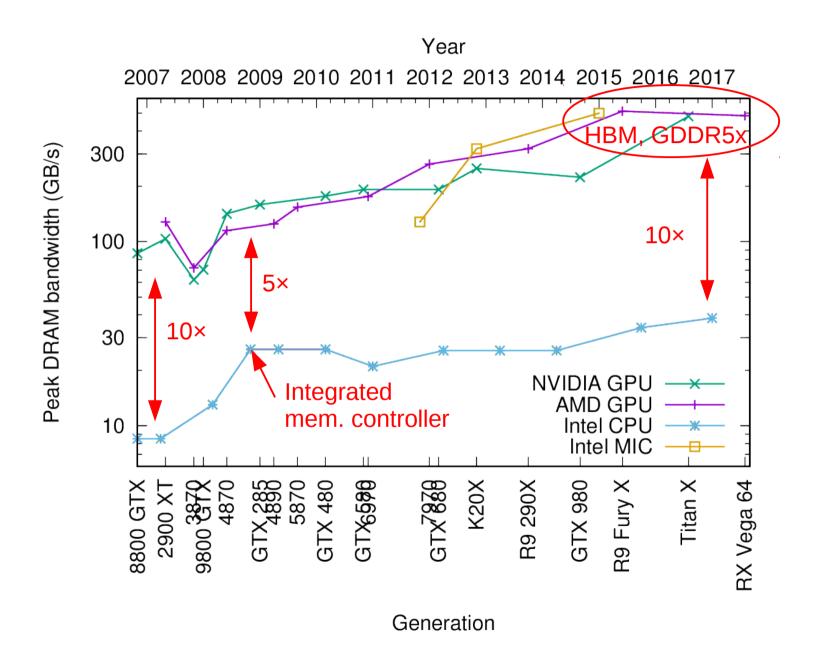
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Trends: compute performance

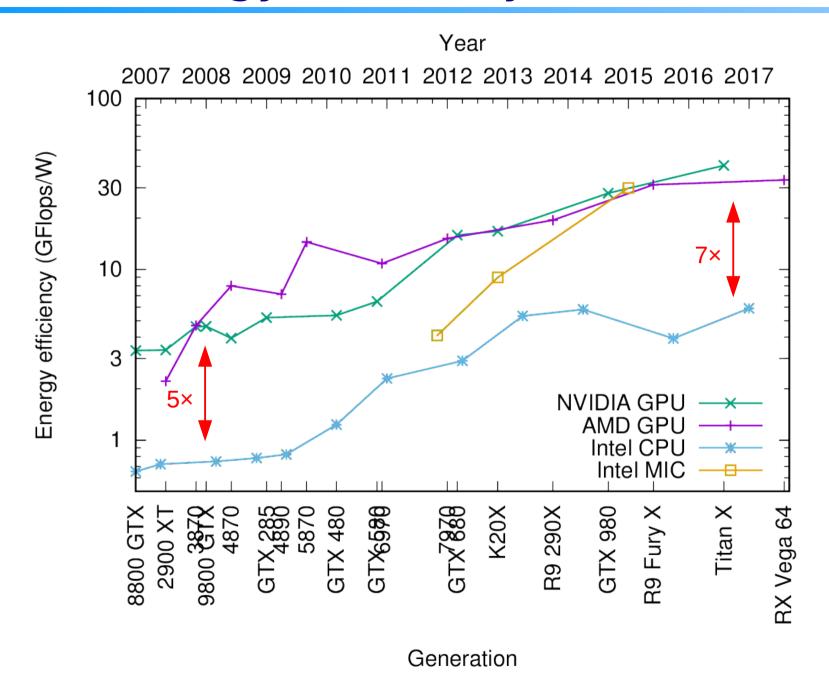


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Trends: memory bandwidth



Trends: energy efficiency



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What is parallelism?

Parallelism: independent operations which execution can be overlapped

Operations: memory accesses or computations

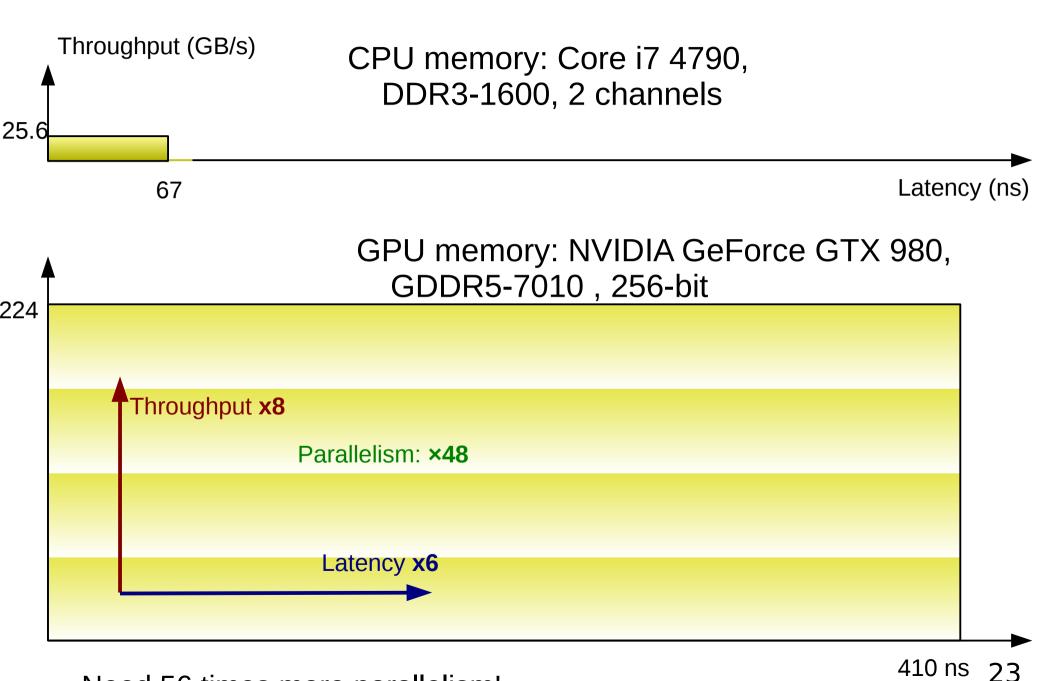
How much parallelism do I need?

- Little's law in queuing theory
 - Average customer arrival rate λ ← throughput

 - Average number of customers
 L = λ×W
- ← Parallelism = throughput × latency

- Units
 - For memory: B = GB/s × ns
 - For arithmetic: flops = Gflops/s × ns

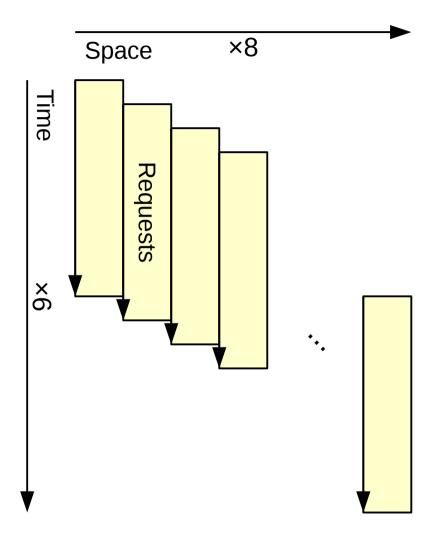
Throughput and latency: CPU vs. GPU



→ Need 56 times more parallelism!

Consequence: more parallelism

- GPU vs. CPU
 - 8× more parallelism to feed more units (throughput)
 - 6× more parallelism to hide longer latency
 - 48× more total parallelism
- How to find this parallelism?



Sources of parallelism

- ILP: Instruction-Level Parallelism
 - Between independent instructions in sequential program

add
$$r3 \leftarrow r1$$
, $r2$
mul $r0 \leftarrow r0$, $r1$
sub $r1 \leftarrow r3$, $r0$

- TLP: Thread-Level Parallelism
 - Between independent execution contexts: threads

- DLP: Data-Level Parallelism
 - Between elements of a vector: same operation on several elements

vadd r
$$\leftarrow$$
 a, b $a_1 \ a_2 \ a_3 \ b_1 \ b_2 \ b_3 \ \hline r_1 \ r_2 \ r_3$

Example: $X \leftarrow a \times X$

In-place scalar-vector product: X ← a×X

Sequential (ILP) For
$$i = 0$$
 to $n-1$ do:
 $X[i] \leftarrow a * X[i]$

Threads (TLP)

Launch n threads:

X[tid] ← a * X[tid]

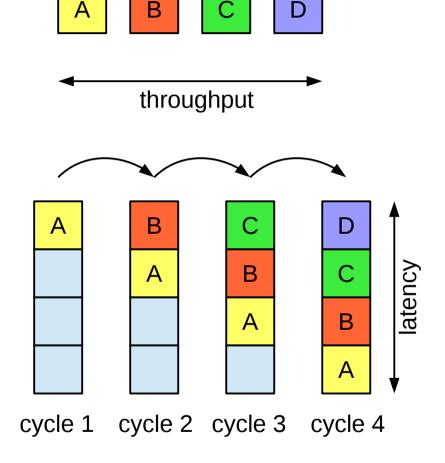
Vector (DLP) X ← a * X

Or any combination of the above

Uses of parallelism

- "Horizontal" parallelism for throughput
 - More units working in parallel

- "Vertical" parallelism for latency hiding
 - Pipelining: keep units busy when waiting for dependencies, memory



How to extract parallelism?

	Horizontal	Vertical
ILP	Superscalar	Pipelined
TLP	Multi-core SMT	Interleaved / switch-on-event multithreading
DLP	SIMD / SIMT	Vector / temporal SIMT

- We have seen the first row: ILP
- We will now review techniques for the next rows: TLP, DLP

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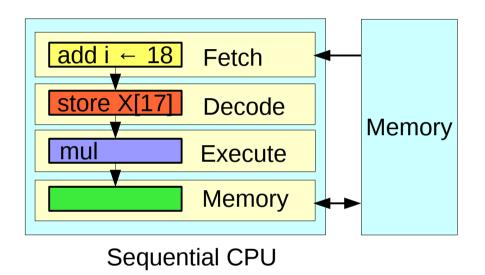
Sequential processor

```
for i = 0 to n-1
  X[i] ← a * X[i]

Source code

move i ← 0
loop:
  load t ← X[i]
  mul t ← a×t
  store X[i] ← t
  add i ← i+1
  branch i<n? loop

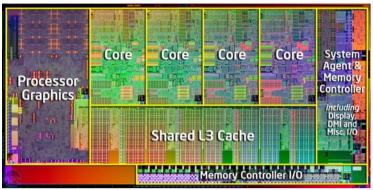
Machine code</pre>
```



- Focuses on instruction-level parallelism
 - Exploits ILP: vertically (pipelining) and horizontally (superscalar)

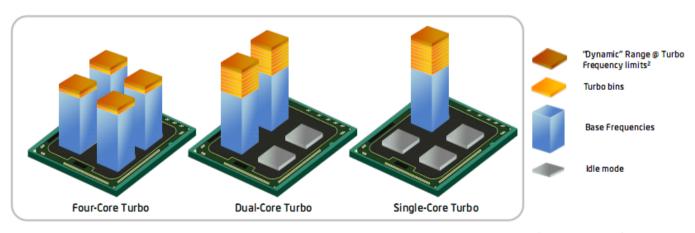
The incremental approach: multi-core

Several processors
 on a single chip
 sharing one memory space



Intel Sandy Bridge

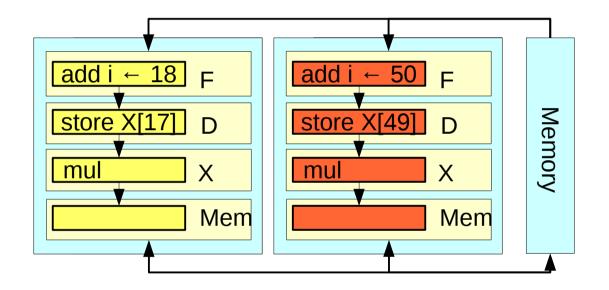
- Area: benefits from Moore's law
- Power: extra cores consume little when not in use
 - e.g. Intel Turbo Boost



Source: Intel

Homogeneous multi-core

Horizontal use of thread-level parallelism

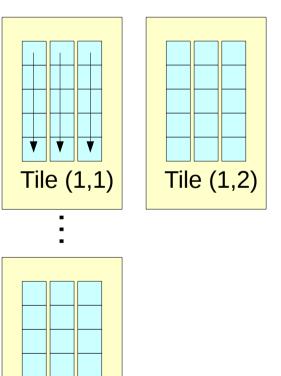


Threads: T0 T1

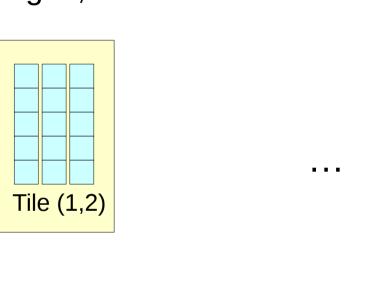
Improves peak throughput

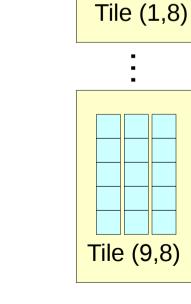
Example: Tilera Tile-GX

- Grid of (up to) 72 tiles
- Each tile: 3-way VLIW processor,
 5 pipeline stages, 1.2 GHz



Tile (9,1)

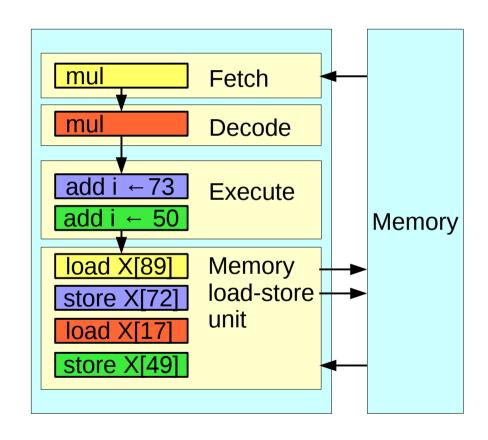






Interleaved multi-threading

Vertical use of thread-level parallelism

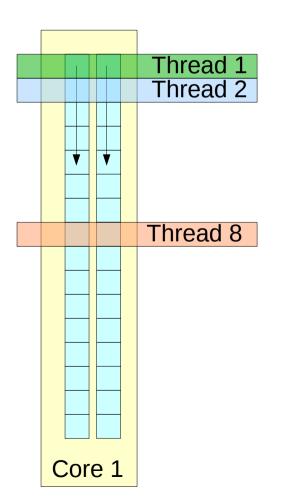


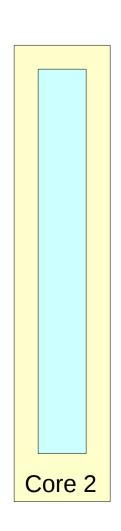
Threads: T0 T1 T2 T3

 Hides latency thanks to explicit parallelism improves achieved throughput

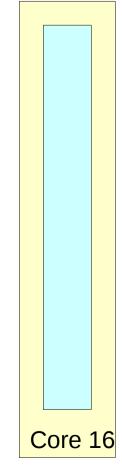
Example: Oracle Sparc T5

- 16 cores / chip
- Core: out-of-order superscalar, 8 threads
- 15 pipeline stages, 3.6 GHz



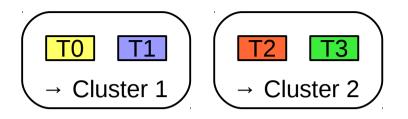


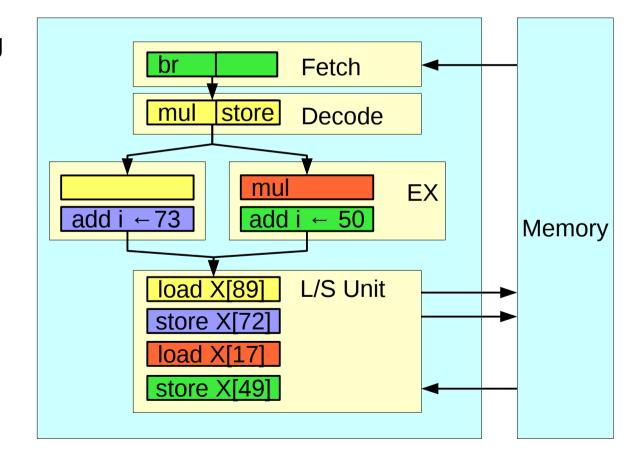




Clustered multi-core

- For each individual unit, select between
 - Horizontal replication
 - Vertical time-multiplexing
- Examples
 - Sun UltraSparc T2, T3
 - AMD Bulldozer
 - IBM Power 7, 8, 9

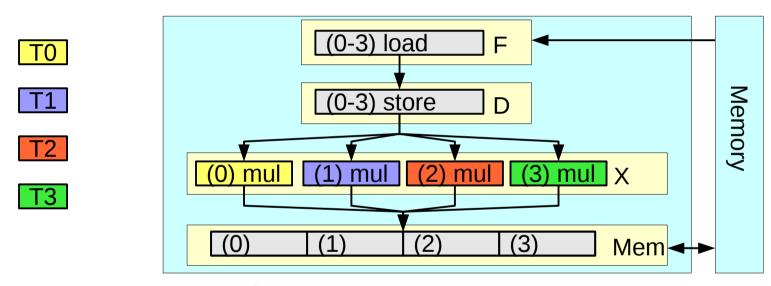




- Area-efficient tradeoff
- Blurs boundaries between cores

Implicit SIMD

- Factorization of fetch/decode, load-store units
 - Fetch 1 instruction on behalf of several threads
 - Read 1 memory location and broadcast to several registers

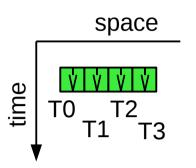


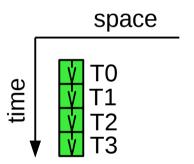
- In NVIDIA-speak
 - SIMT: Single Instruction, Multiple Threads
 - Convoy of synchronized threads: warp
- Extracts DLP from multi-thread applications

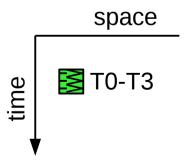
How to exploit common operations?

Multi-threading implementation options:

- Horizontal: replication
 - Different resources, same time
 - Chip Multi-Processing (CMP)
- Vertical: time-multiplexing
 - Same resource, different times
 - Multi-Threading (MT)
- Factorization
 - If we have common operations between threads
 - Same resource, same time
 - Single-Instruction Multi-Threading (SIMT)





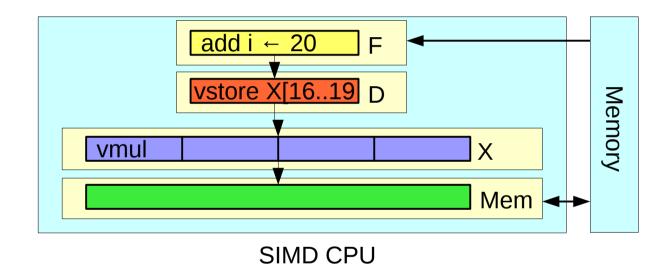


Explicit SIMD

- Single Instruction Multiple Data
- Horizontal use of data level parallelism

```
loop:
    vload T ← X[i]
    vmul T ← a×T
    vstore X[i] ← T
    add i ← i+4
    branch i<n? loop

Machine code
```



- Examples
 - Intel MIC (16-wide)
 - AMD GCN GPU (16-wide×4-deep)
 - Most general purpose CPUs (4-wide to 16-wide)

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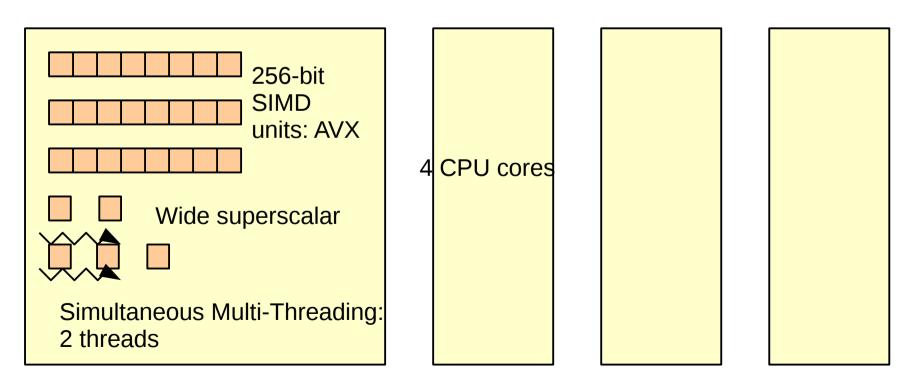
Hierarchical combination

There is no silver bullet!

- All of these techniques face the law of diminishing returns
 - More cores → complex interconnect, hard to maintain cache coherence
 - More threads/core → more register and cache pressure
 - Wider vectors → more performance lost to irregular control/data flow
- Both CPUs and GPUs combine techniques
 - Multiple cores
 - Multiple threads/core
 - SIMD units

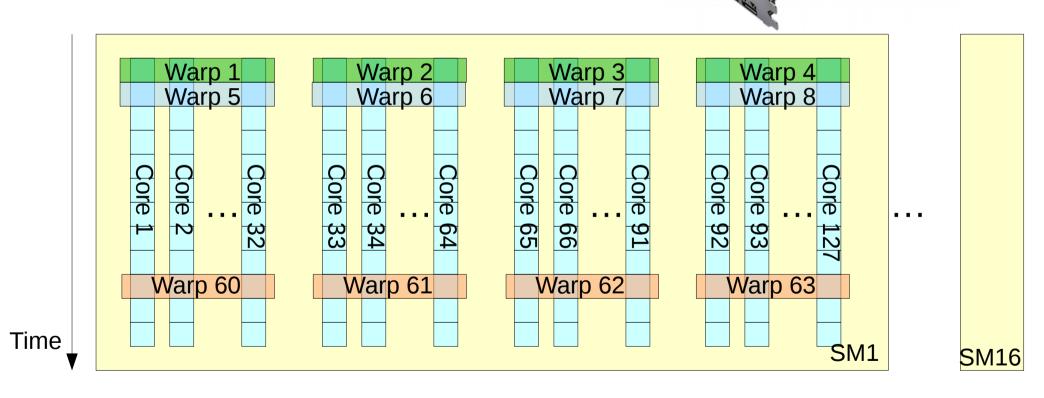
Example CPU: Intel Core i7

- Is a wide superscalar, but has also
 - Multicore
 - Multi-thread / core
 - SIMD units
- Up to 117 operations/cycle from 8 threads



Example GPU: NVIDIA GeForce GTX 980

- SIMT: warps of 32 threads
- 16 SMs / chip
- 4×32 cores / SM, 64 warps / SM

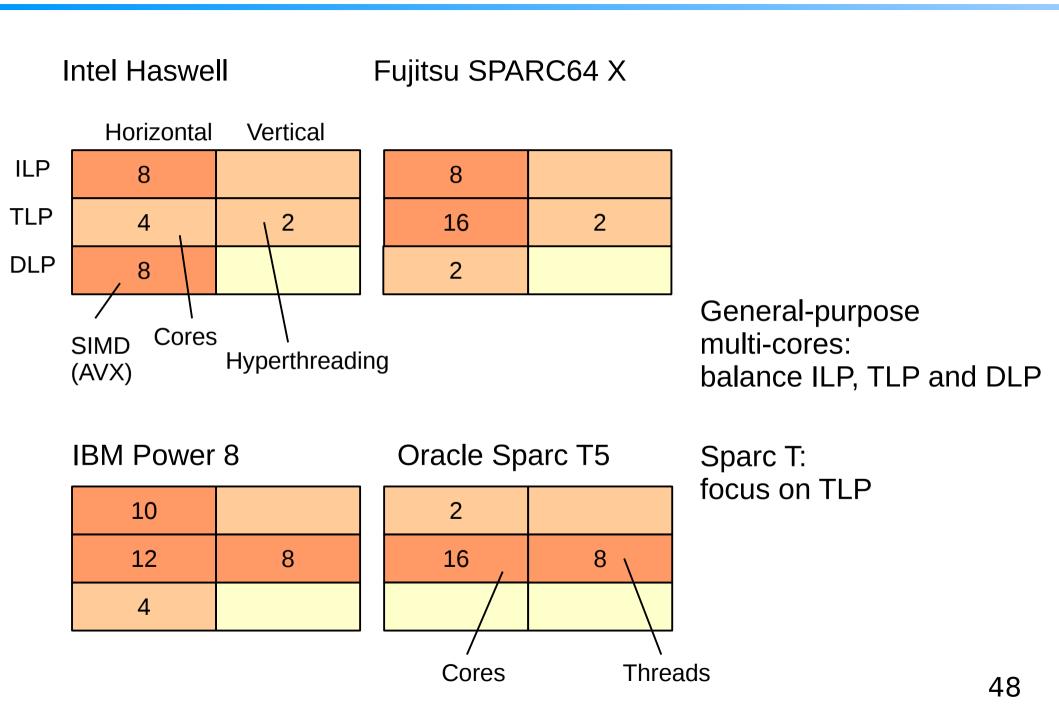


- 4612 Gflop/s
- Up to 32768 threads in flight

Taxonomy of parallel architectures

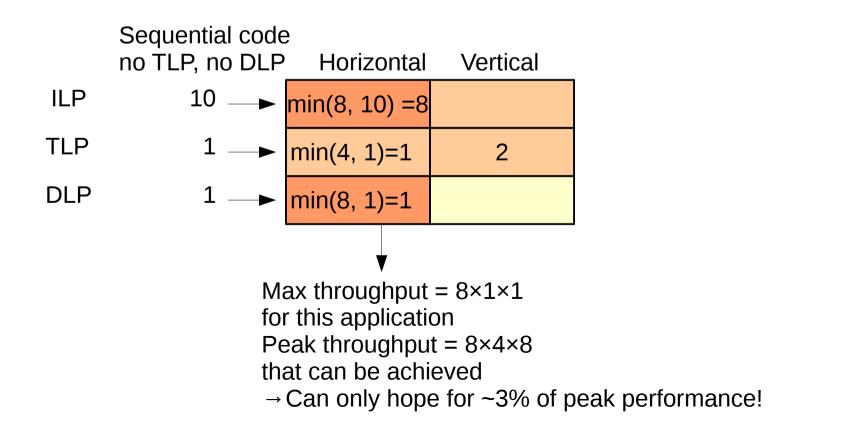
	Horizontal	Vertical
ILP	Superscalar / VLIW	Pipelined
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DLP	SIMD / SIMT	Vector / temporal SIMT

Classification: multi-core

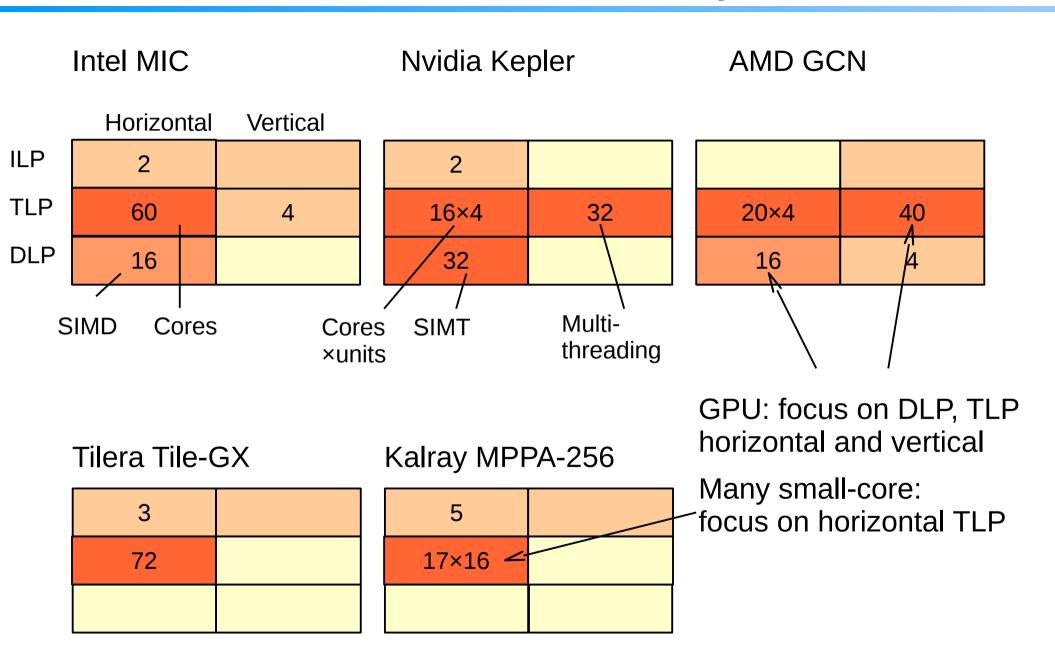


How to read the table

- Given an application with known ILP, TLP, DLP how much throughput / latency hiding can I expect?
 - For each cell, take minimum of existing parallelism and hardware capability
 - The column-wise product gives throughput / latency hiding



Classification: GPU and many small-core



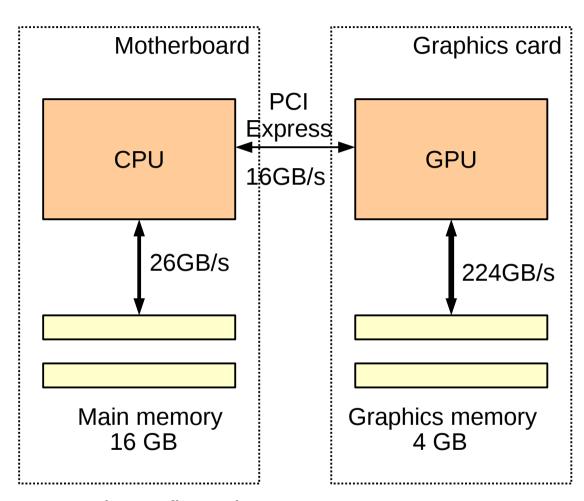
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External memory: discrete GPU

Classical CPU-GPU model

- Split memory spaces
- Need to transfer data explicitly
- Highest bandwidth from GPU memory
- Transfers to main memory are slower

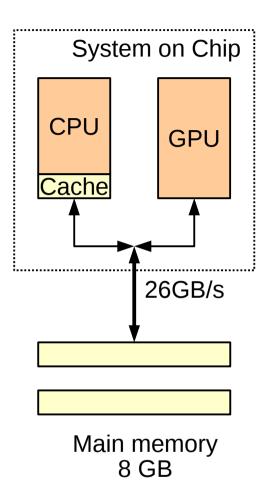


Example configuration: Intel Core i7 4790, Nvidia GeForce GTX 980

External memory: embedded GPU

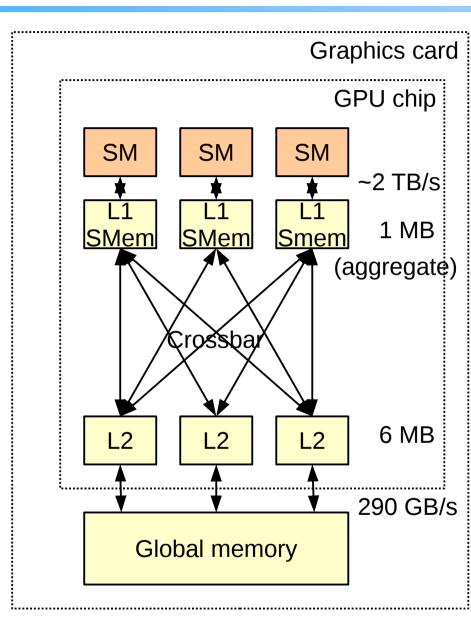
Most GPUs today are integrated

- Same physical memory
- May support memory coherence
 - GPU can read directly from CPU caches
- More contention on external memory



GPU high-level organization

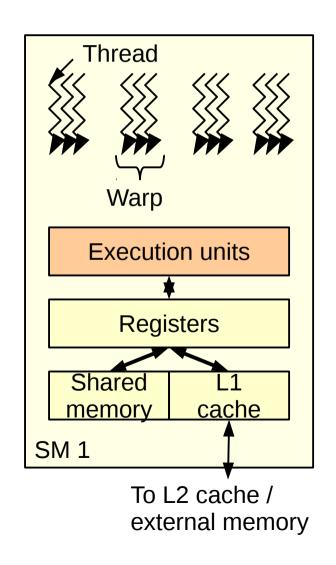
- Processing units
 - Streaming Multiprocessors (SM) in Nvidia jargon
 - Compute Unit (CU) in AMD's
 - Closest equivalent to a CPU core
 - Today: from 1 to 20 SMs in a GPU
- Memory system: caches
 - Keep frequently-accessed data
 - Reduce throughput demand on main memory
 - Managed by hardware (L1, L2) or software (Shared Memory)

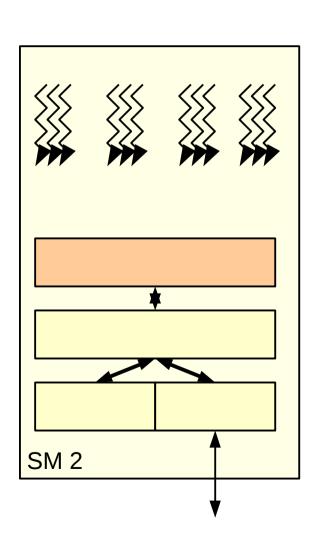


GPU processing unit organization

Each SM is a highly-multithreaded processor

- Today: 24 to 48 warps of 32 threads each
 - → ~1K threads on each SM, ~10K threads on a GPU





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First-order performance model

Questions you should ask yourself, before starting to code or optimize

- Will my code run faster on the GPU?
- Is my existing code running as fast as it should?
- Is performance limited by computations or memory bandwidth?

Pen-and-pencil calculations can (often) answer such questions

Performance: metrics and definitions

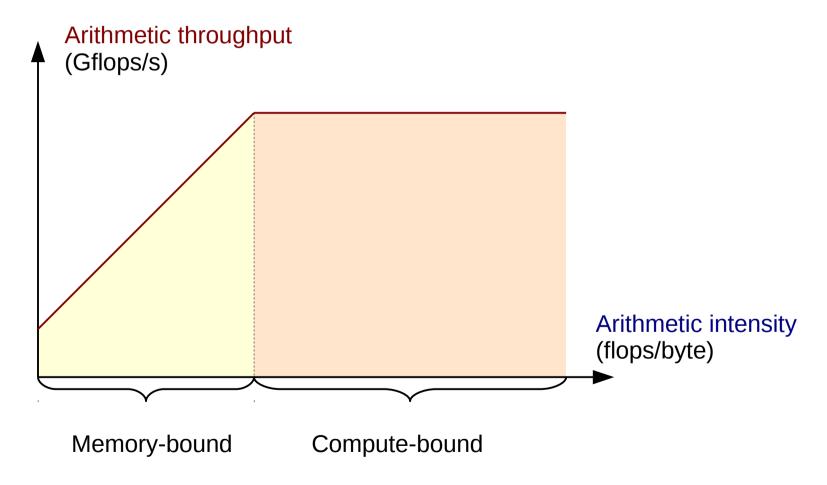
Optimistic evaluation: upper bound on performance

Assume perfect overlap of computations and memory accesses

- Memory accesses: bytes
 - Only external memory, not caches or registers
- Computations: flops
 - Only "useful" computations (usually floating-point) not address calculations, loop iterators..
- Arithmetic intensity: flops / bytes
 = computations / memory accesses
 - Property of the code
- Arithmetic throughput: flops / s
 - Property of code + architecture

The roofline model

- How much performance can I get for a given arithmetic intensity?
 - Upper bound on arithmetic throughput, as a function of arithmetic intensity
 - Property of the architecture

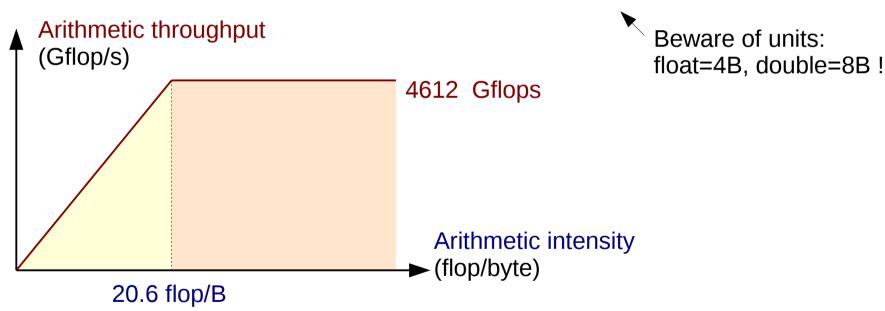


S. Williams, A. Waterman, D. Patterson. *Roofline: an insightful visual performance model* 64 *for multicore architectures.* Communications of the ACM, 2009

Building the machine model

- Compute or measure:
 - Peak memory throughput GTX 980: 224 GB/s
 - Ideal arithmetic intensity = peak compute throughput / mem throughput

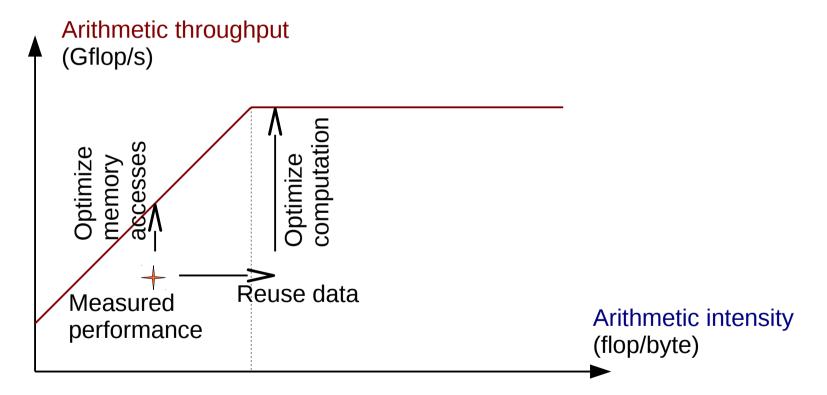
```
GTX 980: 4612 (Gflop/s) / 224 (GB/s) = 20.6 flop/B \times 4 (B/flop) = 82 (dimensionless)
```



- Achievable peaks may be lower than theoretical peaks
 - Lower curves when adding realistic constraints

Using the model

- Compute arithmetic intensity, measure performance of program
- Identify bottleneck: memory or computation
- Take optimization decision



Example: dot product

```
for i = 1 to n
r += a[i] * b[i]
```

- How many computations?
- How many memory accesses?
- Arithmetic intensity?
- Compute-bound or memory-bound?
- How many Gflop/s on a GTX 980 GPU?
 - With data in GPU memory?
 - With data in CPU memory?
- How many Gflop/s on an i7 4790 CPU?

GTX 980: 4612 Gflop/s, 224 GB/s i7 4790: 460 Gflop/s, 25.6 GB/s

PCIe link: 16 GB/s

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1790 CPU?

→ 2 n flops

 \rightarrow 2 n words

 \rightarrow 1 flop/word = 0.25 flop/B

→ Highly memory-bound

25.6 GB/s × 0.25 flop/B → 6.4 Gflop/s Conclusion: don't bother porting to GPU!

224 GB/s \times 0.25 flop/B \rightarrow 56 Gflop/s

16 GB/s \times 0.25 flop/B \rightarrow 4 Gflop/s

GTX 980: 4612 Gflop/s, 224 GB/s i7 4790: 460 Gflop/s, 25.6 GB/s

PCIe link: 16 GB/s

Takeaway

- Result of many tradeoffs
 - Between locality and parallelism
 - Between core complexity and interconnect complexity
- GPU optimized for throughput
 - Exploits primarily DLP, TLP
 - Energy-efficient on parallel applications with regular behavior
- CPU optimized for latency
 - Exploits primarily ILP
 - Can use TLP and DLP when available
- Performance models
 - Back-of-the-envelope calculations and common sense can save time
- Next time: GPU programming in CUDA