

Introduction to SQL

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Overview of The SQL

1. SQL Classification

SQL can be categorised into 2 major classes:

i. Data-definition Language (DDL)

- Provides commands for
 - Create table
 - Drop table
 - Modify table

ii. Data-manipulation Language (DML)

- Provides ability to query database
 - Update
 - Insert
 - Delete

2. Data Definition Language

It allows the specification of information about relations including:

- The **schema** for each relation.
- The **type** of values associated with each **attribute**.
- The **Integrity constraints** (such as primary constraint and others)
- The set of **indices** to be maintained for each relation.
- **Security and authorization** information for each relation.
- The **physical storage** structure of each relation on disk.

3. Domain Types in SQL

- char(n)
- varchar(n)
- int
- smallint
- numeric(p,d)
- real, double precision
- float(n)

Oracle Implementation

- varchar2(n)
- number(p,s)
- Date

4. Create Table Construct

```
CREATE TABLE table_name
(
    attribute1 datatype [ NULL | NOT NULL ],
    attribute2 datatype [ NULL | NOT NULL ],
    ...,
    [CONSTRAINT constraint_name] PRIMARY KEY (primary_attribute1, ...),
    [CONSTRAINT constraint_name] FOREIGN KEY (foreign_attribute1, ...)
    REFERENCES reference_table_name [ON DELETE CASCADE | SET NULL | SET DEFAULT]
    [ON UPDATE CASCADE | SET NULL | SET DEFAULT]
);
```

- **r** → name of relation
- **A_i** → attribute
- **D_i** → Data type
- Integrity constraints:
 - primary key (A₁, ..., A_n)
 - foreign key (A_m, ..., A_n) references r

- not null

Example:

```
create table instructor(
    ID char(5),
    name varchar(20) not null,
    dept_name varchar(20),
    salary numeric(8,2),

    constraint pk_instructor primary key (ID),
    constraint fk_instructor_dept foreign key (dept_name) references department
);
```

5. Dropping Table

1. To delete the schema, we use: (structure and information)

```
DROP TABLE table_name ;
```

2. To delete the table with constraints:

```
DROP TABLE table_name CASCADE CONSTRAINTS;
```

Example:

```
DROP TABLE instructor ;
```

6. Alter Table

```
ALTER TABLE table_name ADD attribute_name datatype;
```

Example:

```
alter table instructor add address varchar2( 40 ) ;
```

Basic Structure of SQL Queries

1. Single Relation

- Three clauses of Basic SQL Structure:
 - SELECT,

- FROM
- WHERE

```
select id, name, salary
from instructors
where dept_name='CSE' or dept_name='EEE';
```

- Select all can be expressed by the `*` sign as follows

```
SELECT * FROM STUDENTS;
```

- Keyword **DISTINCT** is used to eliminate duplicate records (not duplicate column)

```
SELECT DISTINCT DEPT, NAME FROM STUDENTS;
```

2. Multiple Relations

Queries on a Cartesian Product. Remember it is all possible combinations between 2 relations.

```
create table dept(
    dept_name varchar2(10),
    location varchar2(10),
    year_est number(4,0),
    yearly_budget number(10,2),
    constraint pk_dept primary key (dept)
);
create table students(
    name varchar2(20),
    prog varchar2(20),
    dob date,
    cgpa number,
    dept varchar(20),
    constraint fk_students foreign key (dept) references dept(dept)
);

SELECT * FROM dept, students;
```

The Rename Operation

- Renaming in SQL is very useful in a number of cases:
 - attribute name,
 - expression,
 - relation name
 - entire query
 - compare tuples in the same relation (Self-reference)

Renaming Attribute:

```
select name as instructor_name,
course_id
from instructor, teaches
where instructor.ID = teaches.ID
```

Renaming Expression

```
select ID, name, (salary/12) MSalary
from emp;
```

Renaming Relation

Renaming Relation is called **Correlation Name** in the SQL standard, but it is also commonly referred to as a **Table Alias**

```
select name as instructor_name,
course_id
from instructor I, teaches T
where I.ID = T.ID
```

Renaming Query

```
select A.N as Name, A.CID as Course
from
(select name as N, CourseID as CID
from instructor I, teaches T
where I.ID = T.ID) A;
```

Rename for Self-reference

```
select E1.ID, E1.Name, E1.BID  
from Emp E1, Emp E2  
where E1.BID = E2.ID;
```

String Operations

- Two Special characters:
 1. percent(%) → matches **any substring**
 2. underscore (_) → matches **any character**
- equal (=) → **exact** match.
 - Single quote (' ') is used for string matching
 - backslash (\) as escape character
- Pattern matching examples:
 - 'Intro%' matches any string beginning with "Intro".
 - '%Comp%' matches any string containing "Comp" as a substring.
 - '___' matches any string of exactly three characters.
 - '%m%' matches any string containing letter m anywhere.
- SQL supports a variety of string operations such as:
 - concatenation (using || sign)
 - Many built-ins are normally used (i.e. UPPER, LOWER, SUBSTR..)

Set Operation

- Used to extract data from two relations having no references
- Useful to merge data from different sources where conditions vary in each case.
- Alternate way is using **AND, OR, NOT** operators

Example:

```
(select course_id from section where sem = 'Fall' and year = 2017)  
union
```

```
(select course_id from section where sem = 'Spring' and year = 2018)
```

```
(select course_id from section where sem = 'Fall' and year = 2017)
intersect
(select course_id from section where sem = 'Spring' and year = 2018)
```

Null Values

A special problem in relational operations, including arithmetic operations, comparison operations, and set operations.

Solution: use `nvl()` built-in

Example:

```
select Name, (Salary + nvl(bonus,0)) Total
from emp;
```

Now null values will be treated as 0

Aggregate Functions

Take a collection of values and return a single value

- Average: `avg` (must be numbers)
- Minimum: `min`
- Maximum: `max`
- Total: `sum` (must be numbers)
- Count: `count`

Examples:

```
select avg(salary)
from instructor
where dept_name = 'Comp. Sci.';
```

```
select count(distinct ID)
from teaches
```

```
where semester = 'spring' and year = 2018;
```

```
select count(*) from course
```

Aggregation with Grouping

Tuples with the same value on all attributes in the **group by** clause are placed in one group

```
select department, avg(salary) as avg_salary  
from instructor  
group by department;
```

Basic Principle

All selected attributes must be present either in

1. aggregate functions
2. group by clause

Having Clause

- **where** clause can only be applied on the **values already stored**
- **having** clause the **values are not stored** rather will be **computed on-the-fly**

HAVING is used to filter groups after they have been formed

```
select dname, max(budget), count (*) totalstudent  
from depts, students  
where depts.dname = students.dept  
group by dept  
having count (*) >= 200;
```

predicates in the having clause are applied after the formation of groups

predicates in the where clause are applied before forming groups

Nested Sub-query

A subquery is a select-from-where expression that is nested within another query

It can happen in 2 places:

1. In `WHERE` clause (nested subquery)
2. In `FROM` clause (inline view)

Subquery in WHERE clause

```
select id, name, salary
from emp
where salary > (select avg(salary) from emp);
```

- Here subquery in the where part is executed first and the value is locked.

Set Membership

`IN` and `NOT IN` are normally used in the where clause as a mechanism of set-membership

```
select distinct course_id
from section
where semester = 'Fall' and year = 2017 and
course_id in (select course_id
from section
where semester = 'Spring' and year = 2018);
```

Subquery in FORM Clause

```
select ROWNUM as RANK, lastn_name salary
from
(
    select last_name, salary from employees order by salary desc
```

```
)  
where ROWNUM ,= 5;
```

DML Statements

There are 3 DML Statements

1. **INSERT**: new record entry
2. **UPDATE**: existing record modification
3. **DELETE**: Existing record remove

1. Insert

General Insert

- Add a new tuple to course:

```
Method-1:  
INSERT INTO COURSE VALUES('CS-437', 'Database Systems', 'Comp. Sci.', 4);  
  
Method-2:  
INSERT INTO COURSE (course_id, title, dept_name, credits)  
VALUES('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

- Add a new tuple to student with tot_creds set to null

```
INSERT INTO STUDENT  
VALUES ('3003', 'Green', 'Finance', null);
```

Using Query

- Make each student in the Music department who has earned more than 144 credit hours an instructor in the Music department with a salary of \$18, 000 (fixed values can be passed as well).

```
INSERT INTO INSTRUCTOR  
SELECT ID, name, dept_name, 18000
```

```
FROM STUDENT  
WHERE dept_name = 'Music' and total_cred > 144;
```

2. Update Statement

Simple Update

- Give a 5% salary raise to those instructors who earn less than 70000

```
update instructor set salary = salary * 1.05 where salary < 70000;
```

- Give a 5% salary raise to instructors whose salary is less than average

```
update instructor set salary = salary * 1.05  
where salary < (select avg(salary) from instructor);
```

Compound Update

- Increase salaries of instructors whose salary is over \$100, 000 by 3%, and all others by a 5%

```
update instructor set salary = salary * 1.03 where salary > 100000;  
update instructor set salary = salary * 1.05 where salary <= 100000;
```

- Using case

```
update instructor set salary = case  
when salary <= 100000 then salary * 1.05  
else salary * 1.03 end;
```

3. Delete Statement

- Delete all instructors:

```
Delete [from] instructors;
```

- Delete all tuples in the instructor relation for those instructors associated with a department located in the Watson building

```
delete from instructor  
where dept_name in (select dept_name from department where building = 'Watson');
```

- Delete all instructors whose salary is less than the average salary of instructors

```
delete from instructor where salary < (select avg(salary) from instructor);
```