

## 0.1 Verteilte Systeme/Distributed Systems

### 0.1.1 Orga

VL Di 10-12 (nicht am 23.04.)  
Ue Do 10-12

#### Elektisches

- (kvv)
- Website AG
- Sakai

#### Übungen

- ca. 5 Übungsblätter, 14-tägig
- Vorträge in Gruppen über „verteilte Systeme“

#### Material/Inhalt

##### 1. Hälfte Distributed Systems (Tanenbaum, van Steen)

- Architektur
- Prozesse
- Kommunikation
- Namen
- Synchronisation
- Konsistenz
- Replikation
- Fehlertoleranz

##### 2. Hälfte Distributed Algorithms (Nancy Lynch)

- synchronous network algorithms
- network models (leader election, shortest path, distributed consensus, byzantine agreement)
- asynchronous network algorithms (shared memory, mutual exclusion, resource allocation, consensus)
- timing
- network resource allocation
- failure detectors

## 0.2 Distributed Systems

Def: A distributed System is a collection of independent computers that appears to it's users as a single coherent system.

Characteristics:

- autonomous components
- appears as single system

- communication is hidden
- organisation is hidden  
(could be high-performance mainframe or sensor net)
- heterogenous system offers homogenous look/interface

Objectives:

- provide resources (printer, storage, computing)
  - share in a controlled, efficient way
  - grant access
    - ⇒ connect users and resources

Transparency:

hide the fact that processes and resources are physically distributed.

Types of transparency:

access hide differences in representation and how a resource is accessed

location

migration

relocation

replikation

concurrency

failure

transparency is desirable, but not always perfectly possible

tradeoff between transparency and complexity, maintainability and performance

Open System

- service interfaces specified using Interface Definition Language (IDL)
- service specification as text

Scalability is an important property

- scalable in size (number of nodes)
- scalable in geographic spread
- scalable in administration

Problems

- centralized services
- centralized data
- centralized algorithms

Scaling techniques)

- use only asynchronous communication
- distribution, split components

- replication of components

pitfalls

reliable network

secure network

homogenous network

constant topology

zero latency

infinite bandwidth

zero transport cost

one administrator!

Types of distributed systems

- computing systems
  - cluster computing
  - grid computing(virtual organisation, geographically distributed and heterogenous))
- distributed information systems
  - transaction processing systems (database)  
ACID (atomicity, consistency, isolated, durable)
  - enterprise systems
- Distributed pervasive systems
  - small, wireless, adhoc, no administration
  - home automation, health systems, sensor networks

Why do we need distributed systems?

- performance
- distribution inherent
- reliability
- incremental growth (scalability)
- sharing resources

## 0.3 Architectures of distributed Systems

- how to split software into components  
⇒ Softwarearchitecture
- how to build a system out of the components  
⇒ Systemarchitecture

Middleware can help to create distribution transparency

Architecturestyles:

- Layered architecture  
⇒ network stack, messages or data flow up and down

- control flow between layers
  - requests down
  - reply up
- Object-based architectures
  - interaction between components
  - e.g. remote procedure calls
  - can be client-server system
- data-centered architectures
  - data is key element
  - communication over data, distributed database
  - web-systems mostly data-centric
- event-based architecture
  - publish-subscribe systems
  - processes communicates through events
  - publisher announces events at broker
    - ⇒ loose coupling (publisher and subscriber need not to know each other), decoupled in space
    - ⇒ scalability better than client-server, parallel processing, caching

Event-based and data-based can be combined  
⇒ shared Data space

### 0.3.1 System architectures

centralized architectures  
client - server

- (i) single point of failure
- (ii) performance (server is bottleneck)
  - can request be repeated without harm?
  - ⇒ request is idempotent
- (iii) application layering  
Layers:
  - 1.) User interface
  - 2.) processing
  - 3.) data level
  - ⇒ a lot of waiting
  - ⇒ does not scale

Decentralized architectures

vertical distribution (layering)  
different logic on different machines

horizontal distribution  
replicated client/server operating on different data  
⇒ overlay-underlay hides physical structure by adding logical structure

Structured P2P architectures

- most popular technique is distributed hashtables (DHT)
- randomly 128 bit or 160 bit ke for data and nodes. Two or more keys are very unlikely
- Chord system arranges items in a ring
- data item  $k$  is assigned to node with smallest identifier  $id \geq k$

ie item 1 belongs to node 1  
item 2 belongs to node 2  
for each item  $k_i$   $\text{succ}(k)=id$   
returns the name of the node  $k$  is assigned to  
to find data item  $k$  the function LOOKUP( $k$ ) returns the adress of  $\text{succ}(k)$  in  $O(\log(N))$ (later!)

membership management

join:  
create SHA1 identifier  
 $\text{LOOKUP}(id) = \text{succ}(id)$   
contact  $\text{succ}(id)$  and  $\text{pred}(id)$  to join ring

leave:

node  $id$  informs  $\text{succ}(id)$  and  $\text{pred}(id)$  and assigns it's data to  $\text{succ}(id)$

Content adressable network (CAN)

- d-dimensional cartesian space
- every node draws random number
- space is divided among nodes
- every data draws identifier (coordinates) which assigns a node
- join
  - select random point
  - half the square in which  $id$  falls
  - assign item to centers
- leave
  - one node takes the rectangle
  - ⇒ reassign rectangles periodically

Unstructured P2P Network

- random graph
- each node maintains a list of  $c$  neighbours
- partial view or neighbourhood list with age
- nodes exchange neighbour information
  - active thread
  - select peer

**PUSH**

select  $c/2$  youngest entries+myself  
send to peer

**PULL**

receive peer buffer  
construct new partial view  
increment age

passive thread  
receive buffer from peer

**PULL:**

select  $c/2$   
send to peer  
construct new partial view increment age