

1 Dynamic Programming

1.1 Knapsack

```

1 vector<vector<ll>> DP;
2 vector<ll> Weights;
3 vector<ll> Values;
4
5 ll Knapsack(int w, int i) {
6     if (w == 0 or i == -1)
7         return 0;
8     if (DP[w][i] != -1)
9         return DP[w][i];
10    if (Weights[i] > w)
11        return DP[w][i] = Knapsack(w, i - 1);
12    return DP[w][i] = max(Values[i] + Knapsack(w - Weights[i], i - 1),
13                          Knapsack(w, i - 1));
14 }
```

1.2 Matrix Chain Multiplication

```

1 vector<vector<ii>> DP; // Pair value, op result
2 int n; // Size of DP (i.e. i,j<n)
3 ii op(ii a, ii b) {
4     return {
5         a.first + b.first + a.second * b.second,
6         (a.second + b.second) %
7         100}; // Second part MUST be associative, first part is cost
8         function
9     }
10
11 ii MCM(int i, int j) {
12     if (DP[i][j].first != -1)
13         return DP[i][j];
14     int ans = 1e9; // INF
15     int res;
16     repx(k, i + 1, j) {
17         ii temp = op(MCM(i, k), MCM(k, j));
18         ans = min(ans, temp.first);
19         res = temp.second;
20     }
21     return DP[i][j] = {ans, res};
22 }
23
24 void fill() {
25     DP.assign(n, vector<ii>(n, {-1, 0}));
26     rep(i, n - 1) {
27         DP[i][i + 1].first = 1;
28     } // Pair op identity, cost (cost must be from input)
29 }
```

1.3 Longest Increasing Subsequence

```

1 vi L;
```

```

2 vi vals;
3 int maxl = 1;
4
5 // Bottom up approach O(nlogn)
6 int lis(int n) {
7     L.assign(n, -1);
8     L[0] = vals[0];
9     repx(i, 1, n) {
10         auto it = lower_bound(L.begin(), L.begin() + maxl, vals[i]);
11         if (it == L.begin() + maxl) {
12             L[maxl] = vals[i];
13             maxl++;
14         } else
15             *it = vals[i];
16     }
17     return maxl;
18 }
```