1 Dynamic Programming

1.1 Knapsack

```
vector<vector<ll>> DP;
    vector<ll> Weights;
    vector<ll> Values;
    11 Knapsack(int w, int i)
        if (w == 0 \text{ or } i == -1)
            return 0;
        if (DP[w][i] != -1)
10
            return DP[w][i];
11
12
        if (Weights[i] > w)
            return DP[w][i] = Knapsack(w, i - 1);
13
        return DP[w][i] = max(Values[i] + Knapsack(w - Weights[i], i - 1),
14
             Knapsack(w, i - 1));
15 }
```

1.2 Matrix Chain Multiplication

```
vector<vector<ii>>> DP; //Pair value, op result
                            //Size of DP (i.e. i,j<n)</pre>
4
    ii op(ii a, ii b)
5
        return {a.first + b.first + a.second * b.second, (a.second + b.
             second) % 100}; //Second part MUST be associative, first part
             is cost function
7
    ii MCM(int i, int j)
10
        if (DP[i][j].first != -1)
11
            return DP[i][j];
12
        int ans = 1e9; //INF
13
        int res;
14
        repx(k, i + 1, j)
15
16
            ii temp = op(MCM(i, k), MCM(k, j));
17
            ans = min(ans, temp.first);
18
            res = temp.second;
19
        }
20
        return DP[i][j] = {ans, res};
^{21}
22
23
    void fill()
24
^{25}
        DP.assign(n, vector<ii>(n, {-1, 0}));
26
        rep(i, n - 1) { DP[i][i + 1].first = 1; } // Pair op identity, cost
27
             (cost must be from input)
28 }
```

1.3 Longest Increasing Subsequence

```
1
2
   vi L;
   vi vals;
 4
    int maxl = 1;
    //Bottom up approach O(nlogn)
    int lis(int n)
 9
        L.assign(n, -1);
10
        L[0] = vals[0];
        repx(i, 1, n)
12
13
            auto it = lower_bound(L.begin(), L.begin() + maxl, vals[i]);
14
            if (it == L.begin() + maxl)
15
16
                L[maxl] = vals[i];
17
                maxl++:
18
19
            else
20
21
                *it = vals[i]:
22
23
        return maxl;
24 }
```