Nested Monte Carlo Search for Two-Player Games

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By: Nina Braunmiller

Subject: SE Seminar in Al

Supervisor: Prof. Johannes Fürnkranz

- I. Nested Monte Carlo Tree Search (NMCS)
 - i. definition
 - ii. For one player games
 - iii. Problems for two player games

II. NMCS improvements

- i. Discounting
- ii. Pruning
 - Cut on Win
 - Pruning on Depth
- III. Results
- IV. Conclusions

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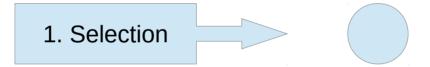
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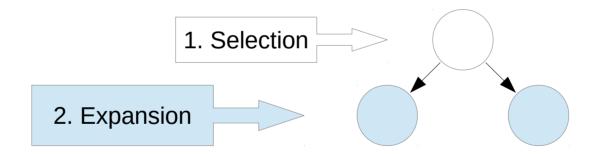
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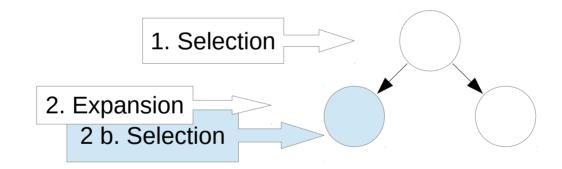
- What is NMCS?
- Monte Carlo Tree search:



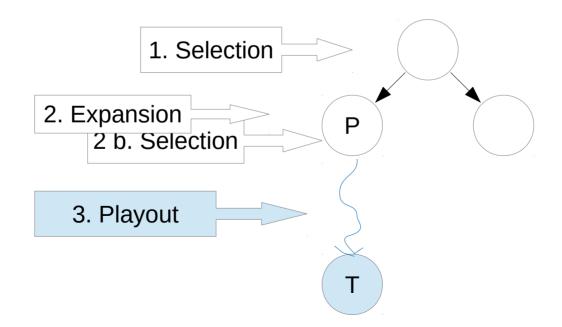
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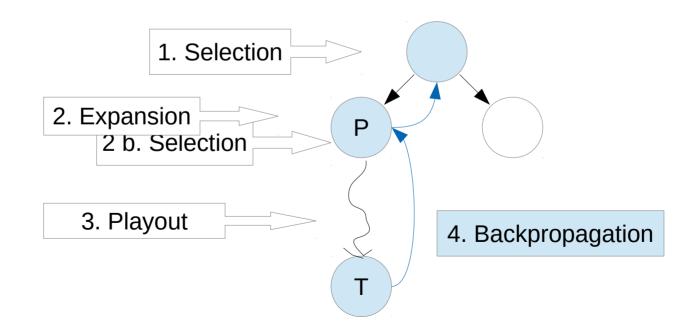
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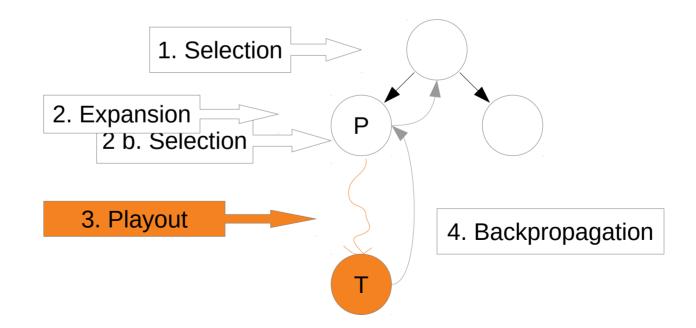
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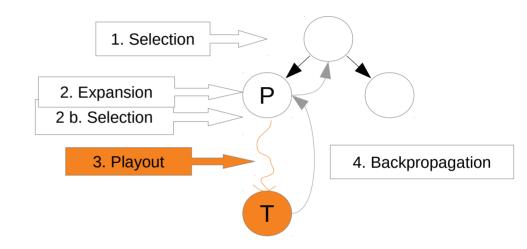


Playout function:

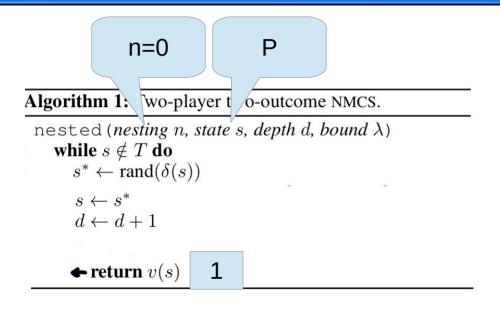
NMC(0)

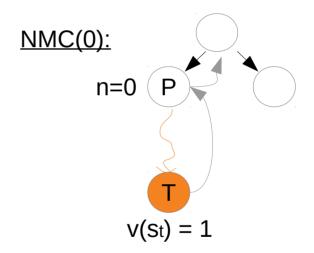
→ Monte Carlo Tree Search

 $NMC(n+1) = P_{\pi(V(NMC(n)))}$



	nesting level
π Pπ V(NMC(n))	policy: tries to find best path from current node playout function: get sequence of states according to π evaluation function: get terminal state value by using NMC(n)





Start with n=1, state s=P

```
Algorithm 1: Two-player two-outcome NMCS.
 1 nested (nesting n, state s, depth d, bound \lambda)
       while s \notin T do
                                       Bad value for a terminal value
           s^* \leftarrow \operatorname{rand}(\delta(s))
          if \tau(s) = \max then l^* \leftarrow \frac{-1}{d} else l^* \leftarrow \frac{1}{d}
          if n > 0 then
              foreach s' in \delta(s) do
                 l \leftarrow \text{nested}(n-1, s', d+1, l^*)
                 if \tau(s)\{l, l^*\} \neq l^* then s^* \leftarrow s'; l^* \leftarrow l
          s \leftarrow s^*
          d \leftarrow d + 1
13
          \leftarrow return v(s)
```

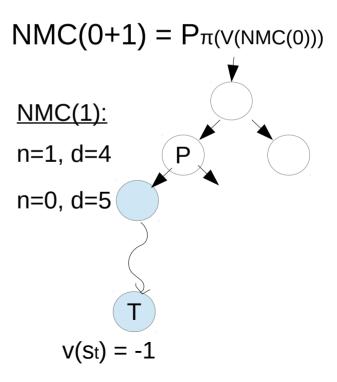
```
NMC(0+1) = P\pi(V(NMC(0)))

NMC(1):

n=1
```

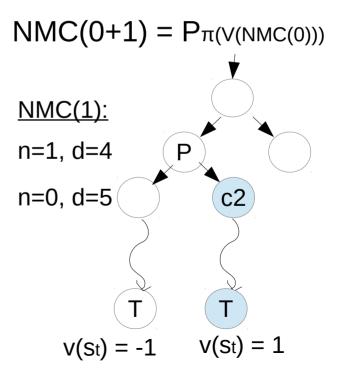
Start with n=1, state s=P, d=4

```
Algorithm 1: Two-player two-outcome NMCS.
 1 nested (nesting n, state s, depth d, bound \lambda)
       while s \notin T do
                                                 Bad v(st), here -1/4
           s^* \leftarrow \operatorname{rand}(\delta(s))
          if \tau(s) = \max then l^* \leftarrow \frac{-1}{d} else l^* \leftarrow \frac{1}{d}
                                                                       Get v(st)
           if n > 0 then
                                                                   for each child
              foreach s' in \delta(s) do
                  l \leftarrow \text{nested}(n-1, s', d+1, l^{\overline{*}})
                  if \tau(s)\{l, l^*\} \neq l^* then s^* \leftarrow s'; l^* \leftarrow l
                       \{-1/4, -1\} = -1/4 \rightarrow Pass
           s \leftarrow \overline{s^*}
           d \leftarrow d + 1
13
          \leftarrow return v(s)
14
```



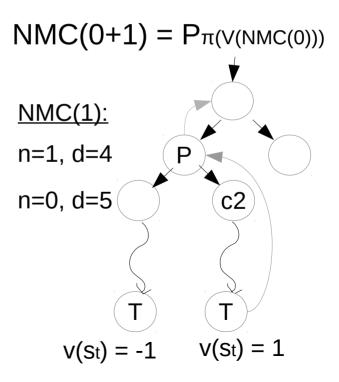
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                 if \tau(s)\{l, l^*\} \neq l^* then s^* \leftarrow s'; l^* \leftarrow l
                  \{-1/4, 1\} = 1 \rightarrow s^* = c2, |* = 1
           s \leftarrow s^*
          d \leftarrow d + 1
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          \leftarrow return v(s)
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Start with n=1, state s=P, d=4

```
Algorithm 1: Two-player two-outcome NMCS.
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                                                                     Get v(st)
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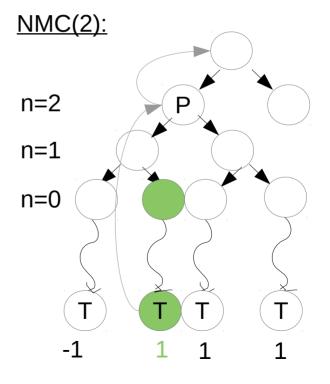


Start with n=2, state s=P

```
Algorithm 1: Two-player two-outcome NMCS.
```

```
1 nested (nesting n, state s, depth d, bound \lambda)
        while s \notin T do
            s^* \leftarrow \operatorname{rand}(\delta(s))
            if \tau(s) = \max then l^* \leftarrow \frac{-1}{d} else l^* \leftarrow \frac{1}{d}
            if n > 0 then
                foreach s' in \delta(s) do
                   l \leftarrow \text{nested}(n-1, s', d+1, l^*)
                   if \tau(s)\{l,l^*\} \neq l^* then s^* \leftarrow s'; l^* \leftarrow l
10
            s \leftarrow s^*
11
            d \leftarrow d + 1
12
13
           \leftarrow return v(s)
14
```

 $NMC(1+1) = P_{\pi(V(NMC(1)))}$



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NMCS: For one player games

- Same author Tristan Cazenave had idea for NMCS for single player games (2009)
- e. g. 16x16 Sudoku:
 - For each cell 16 possible values (domain)
 - Cell gets value → cells of same row/column get updated and value removed from their domain
 - Current cell always cell with smallest domain
 - Score: depth # cells to of which values assigned
 - The deeper the better

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NMCS: Problems for two player games

Single player	Two players
prefering long paths	win as fast as possible
wide range of scores	scores= $\{-1;0;1\} \rightarrow$ which winning path better than the other one?
Save best sequence for faster execution and search improvement	MIN player as information lack for MAX player
Unlimited time	with each nesting level exponential increase of algorithm

Keep perfromance of NMCS without increasing the complexity to much!

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Aims of the two player game:

- Win as fast as possible such that opponent has less options of escaping
- Loosing as slowly as possible such that player has more chances of escaping
- → Express it algoritimically:

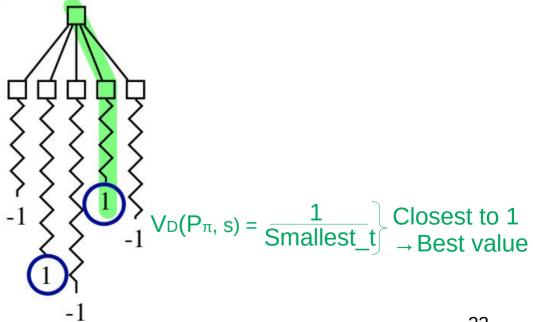
score of terminal state st:

$$V_D(P_{\pi}, s) = V(st)$$
 $t+1$

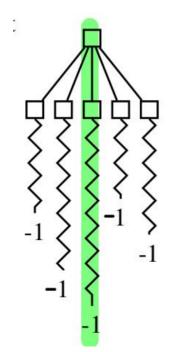
moves from current state s to st

NMC(1): Current state s_0 Future state s_1 Playout $V(P_{\pi}, s) = 1$ Payoff values

NMC_D(1):







$$V_D(P_{\pi}, s) = \frac{-1}{Largest_t}$$
 Closest to 0
 \rightarrow Best value

- Better move selection
- No influence on complexity

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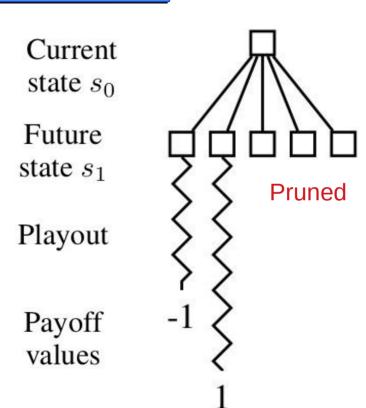
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NMCS improvements: Pruning → Cut on Win (COW)

- $\pi \to \pi$ cow
- Stop with NMCS when terminal state with value 1 found!
 - * Discounting can't be used in combination

- Aim of fast wins ignored
- Complexity decreased



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NMCS improvements: Pruning → Pruning on Depth (POD)

- All playouts skipped which are longer than the playout of the shortest winner path found so far
 - Discounting used

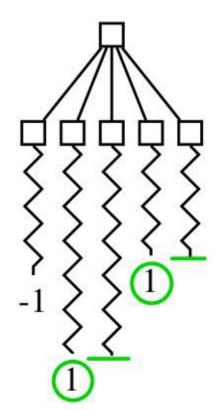
- Shortest wins considered
- Complexity decreased

Current state s_0

Future state s_1

Playout

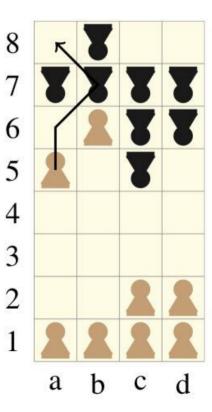
Payoff values



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Rules:

- Move one piece straight/diagnoal forward when there are empty places
- When opponent diagnoally in front of a piece that piece can go to the opponent's one and replace it
- Aim: reach opponent's home row
- Draws are not possible



NMC(3):

Discounting	Dennie	States Wisited(Ir)	Enag(07)
Discounting	Pruning	States Visited(k)	Freq(%)
No	None	$4,459 \pm 27$	11.9 ± 2.2
No	$cow (\leq 1)$	$1,084 \pm 8$	12.3 ± 2.6
No	$cow (\leq 2)$	214 ± 2	10.9 ± 2.0
No	$cow (\leq 3)$	25 ± 1	9.8 ± 2.0
Yes	None	$2,775 \pm 26$	64.1 ± 3.4
Yes	$POD(\leq 1)$	$1,924 \pm 20$	64.7 ± 3.5
Yes	$POD(\leq 2)$	$1,463 \pm 16$	58.6 ± 3.5
Yes	$POD(\leq 3)$	627 ± 19	62.4 ± 3.3
	(a)	(b) (c)	(b) (c)

- (a) Activation in nesting level i or lower
- (b) Average over 900 runs
- (c) Confidence interval

NMC(3):

			— (24)
Discounting	Pruning	States Visited(k)	Freq(%)
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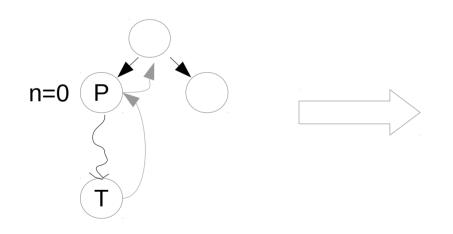
Table 4: Win percentages of NMCS against a standard MCTS player for various settings and thinking times.

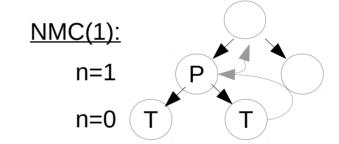
Ga	ame	n COW POD	10ms	20ms	40ms	80ms 1	160ms 3	320ms
Breakthrough		1 1 v 1 v 2 v	3.2 27.6 22.6 4.6	6.0 22.6 25.2 2.0	12.0 16.8 30.4 2.4	11.6 21.6 34.6 1.4	7.8 15.4 35.2 2.4	6.4 20.4 39.6 3.8
Breakt	misère	1 1 1 1 2 2 2	85.4 91.4 95.2 1.0	83.4 95.6 95.2 27.6	70.2 97.0 98.0 43.6	60.8 97.8 99.0 87.0	57.0 98.8 99.8 93.2	56.4 98.8 99.8 95.6

misère: same game with reverse winning condition

Why NMCS better in misère version?

NMCS important in games where the last moves are important.





When nesting level n reaches the depth of winning terminal node with only winning siblings then win.

It is fully included in the MINMAX tree.

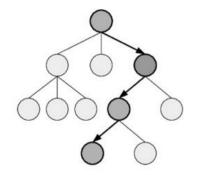
Results: Discounting effect

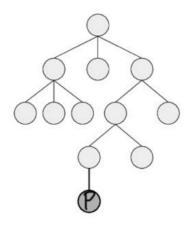
Winrates (%) of NMCS with discounting vs. NMCS without it for nesting levels 0 to 2 and game engine speed.

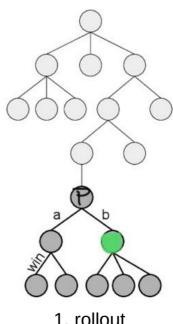
Game	Nesting	Level
	0	1 2
Breakthrough	79.6 < 99.	
misère	42.4 < 80.	8 < 90.0
Knightthrough	78.6 < 100.	0 = 100.0
misère	46.0 < 83.	2 < 85.8
Domineering	71.2 < 77.	0 < 83.8
misère	43.4 < 63 .	2 < 68.4
NoGo	62.8 < 76.	4 < 83.4
misère	53.2 ≺ 65.	6 < 67.2
AtariGo	69.6 < 97.	2 < 100.0

Results: NMCS vs. MCTS-MR

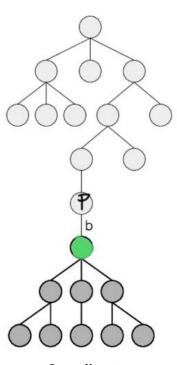
MCTS with Minimax Rollouts:







1. rollout



2. rollout

Results: NMCS vs. MCTS-MR

Win percentages of NMCS and MCTS-MR against standard MCTS with 320ms per move. NMCS uses COW and depth 1 while MCTS-MR uses depth 1 and 2.

NMCS	MCTS MR 1	
	MD 1	
	IVIK I	MR 2
39.6	47.4	50.2
99.8	49.4	48.6
49.6	50.0	50.0
98.6	49.8	45.0
50.0	50.0	49.8
58.6	46.0	44.2
48.0	59.4	50.2
60.8	54.2	67.4
77.2	44.0	47.0
	99.8 49.6 98.6 50.0 58.6 48.0 60.8	99.8 49.4 49.6 50.0 98.6 49.8 50.0 50.0 58.6 46.0 48.0 59.4 60.8 54.2

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Conclusions

- Discounting as tool for better move selection
- COW and POD as policies for reducing the search effort
- NMCS can deal with games in which the last steps are the curcial ones
- NMCS not always outperfoms MCTS

References

Hendrik Baier, Mark H. M. Winands: Monte-Carlo tree search and minimax hybrids. CIG 2013: 1-8

Tristan Cazenave, Abdallah Saffidine, Michael John Schofield, Michael Thielscher: Nested Monte Carlo Search for Two-Player Games. AAAI 2016: 687-693

Tristan Cazenave: Nested Monte-Carlo Search. IJ-CAI 2009: 456-461

Exkursus: Soduko

2		5	3		8	4		9
	7						5	
9		4				6		7
9 5				4				2
			5		7			
6				3				8
4		6				8		1
	2						6	
8		1	2		9	7		4

https://sudoku-club.de/produkt/1000-sudokus-9x9-schwer/