

# 10장. I/O Models

# I/O blocking: Why ?

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- CPU operates much faster than the disk or the network does
  - A very fast disk has 5 ms seek time
  - On a 500 MHz Pentium III machine, a task can execute about 1,250,000 assembler instructions during one seek time

# Computer System들의 전형적인 수행 시간

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항목	시간	사람의 시간으로 환산
Processor Cycle	0.5 ns(2GHz)	1초
Cache Access	1ns(1GHz)	2초
Memory Access	15ns	30초
Context Switch	5,000ns(5us)	167분
Disk Access	7,000,000ns(7ms)	162일
Time Quantum	100,000,000ns(100ms)	6.3년

주)  $1\text{ns} = 10^{-9}\text{초}$ ,  $1\text{us} = 10^{-6}\text{초}$ ,  $1\text{ms} = 10^{-3}\text{초}$

# When to Block ?

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- When to Block
  - Reading
    - No data has arrived yet
  - Writing
    - Internal buffers are full and waiting for transmission and your task requests more data to be sent
  - Connecting
    - `accept( )` and `connect( )` system calls find no pending connections in the listening queue

# Alternatives to I/O Blocking

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- While waiting for a system request to finish, the task could
  - Test the integrity of its data
  - Start and track other requests
  - Wait for several socket connections
  - Process some CPU-intensive calculations

# I/O Models

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- Blocking I/O
- Non-blocking I/O
- I/O multiplexing
- Signal Driven I/O Model
- Asynchronous I/O

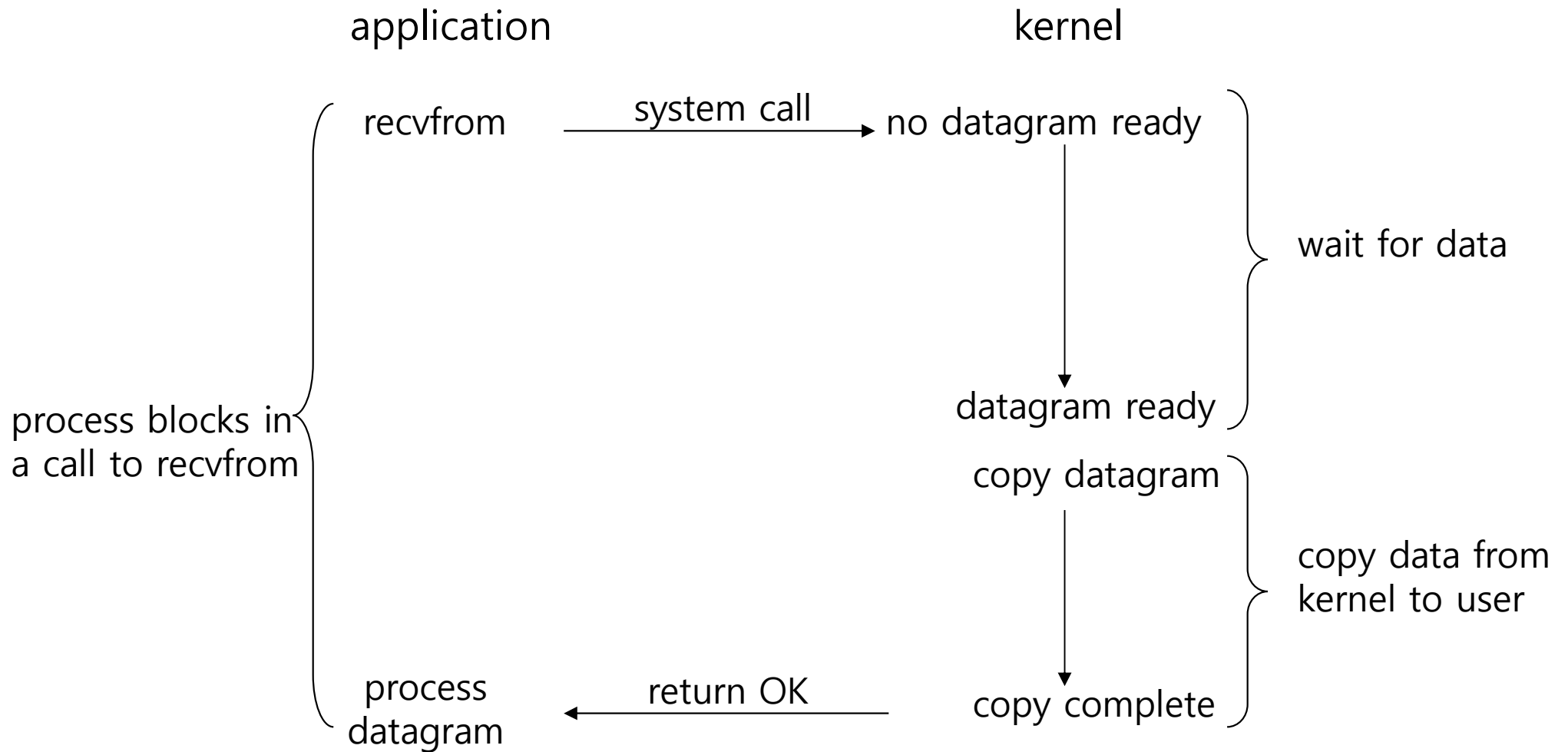
# I/O Operation

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- 입력 동작에는 2 단계가 존재
  - 데이터가 준비될 때까지 기다리는 단계
  - 준비된 데이터를 커널에서 사용자 프로세스로 복사하는 단계
- 네트워크 소켓의 경우
  - 패킷이 네트워크로부터 도착, 커널 버퍼로 복사 됨
  - 커널 버퍼로부터 응용 프로그램의 버퍼로 복사

# Blocking I/O Model

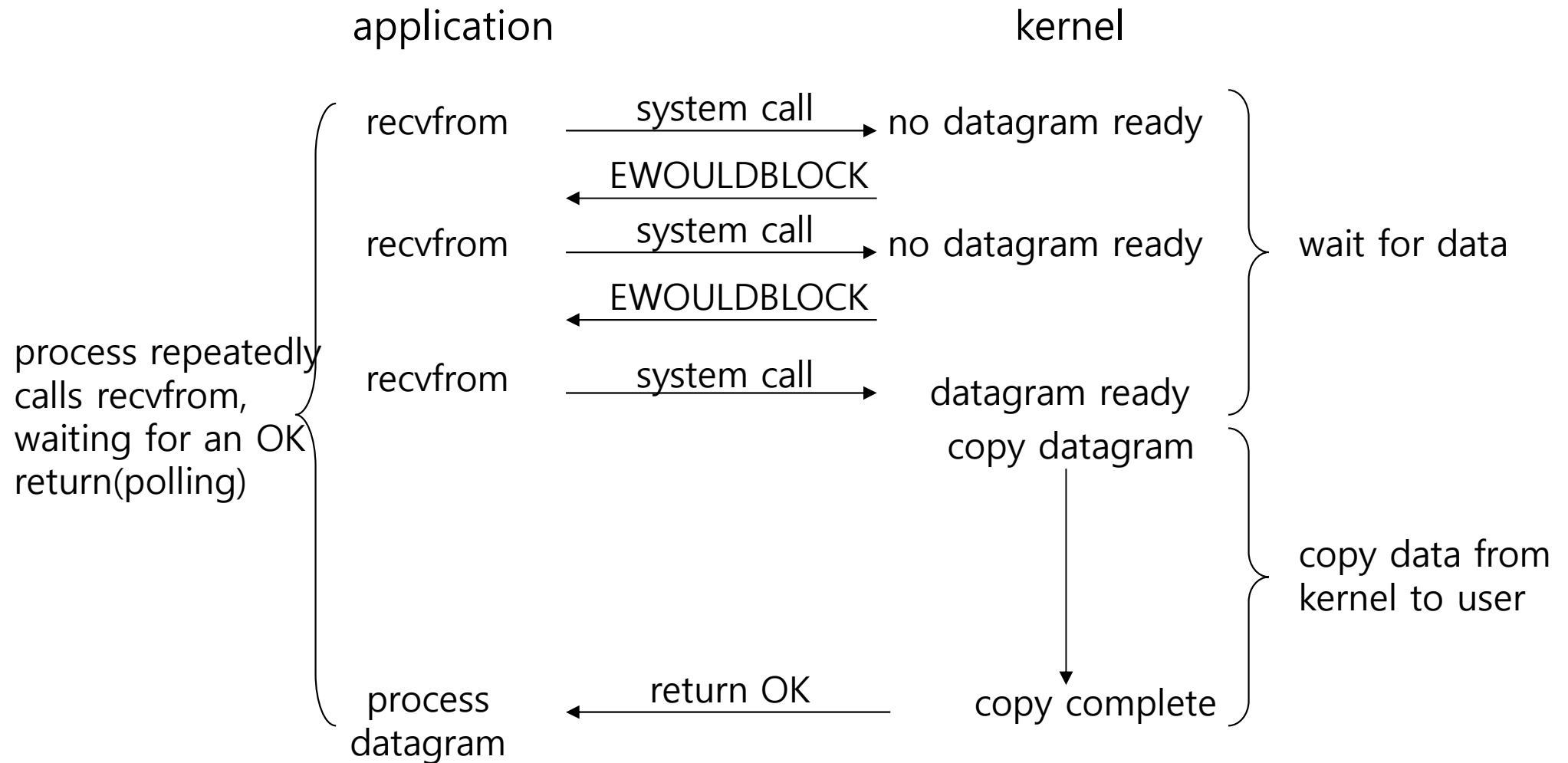
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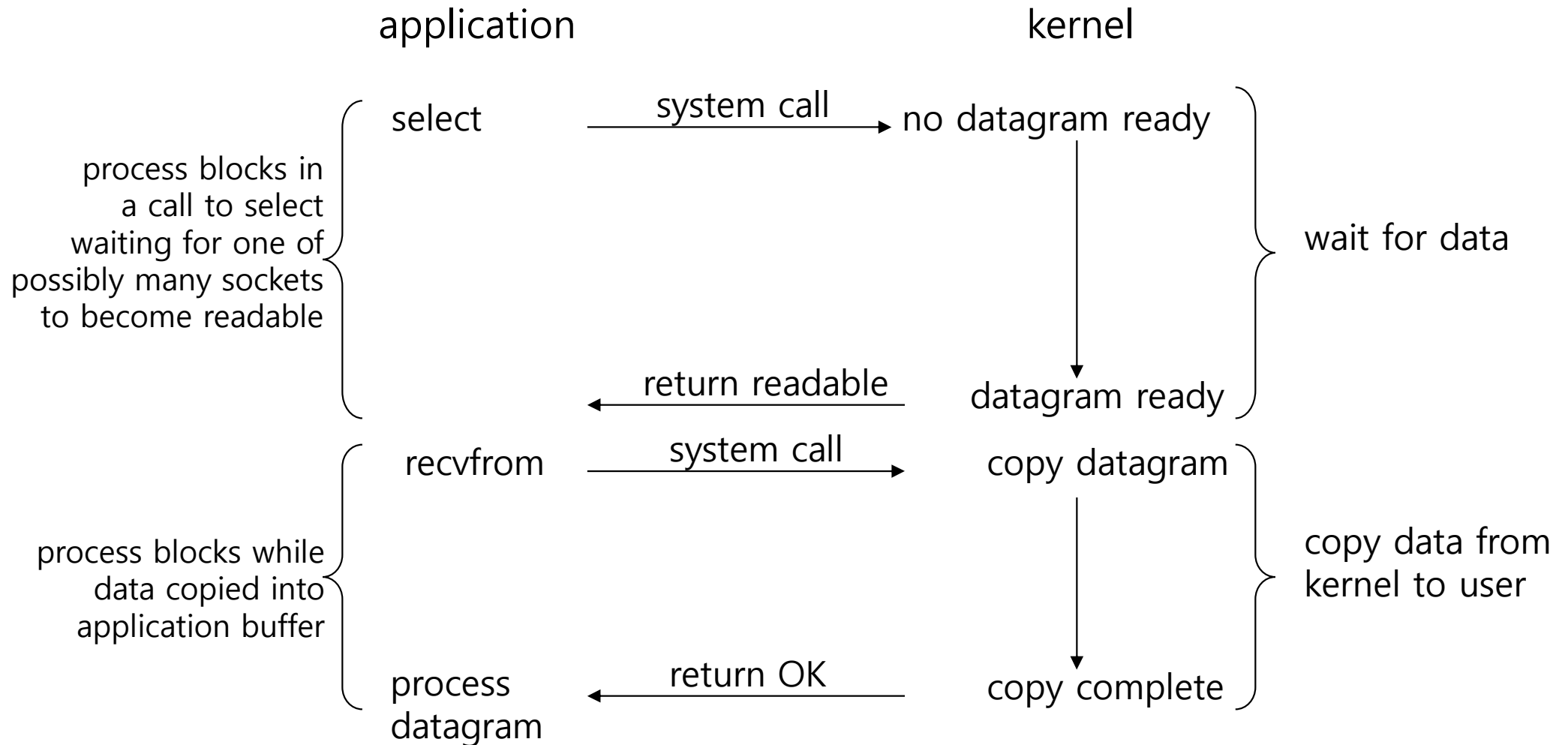
# Non-blocking I/O Model

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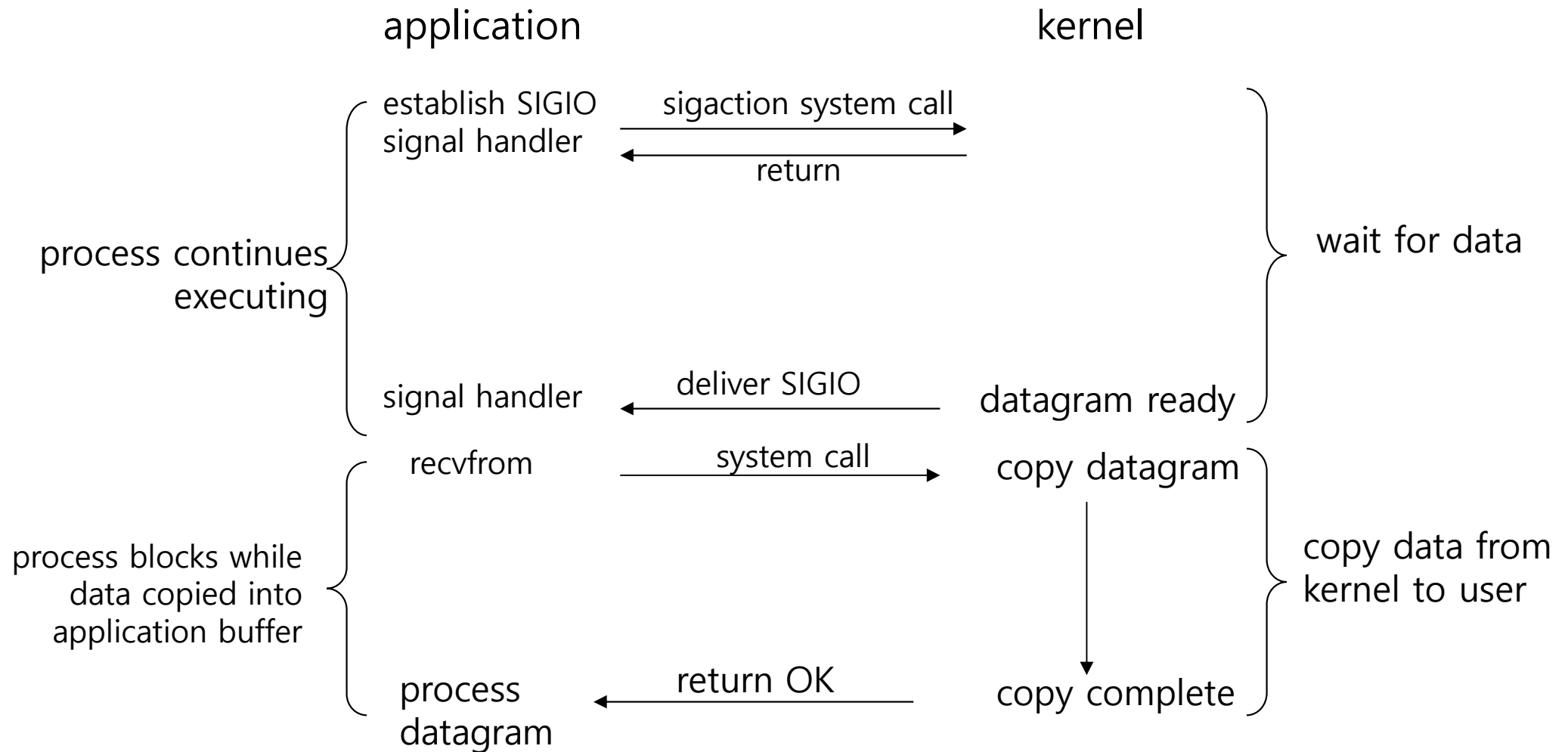
# I/O Multiplexing Model

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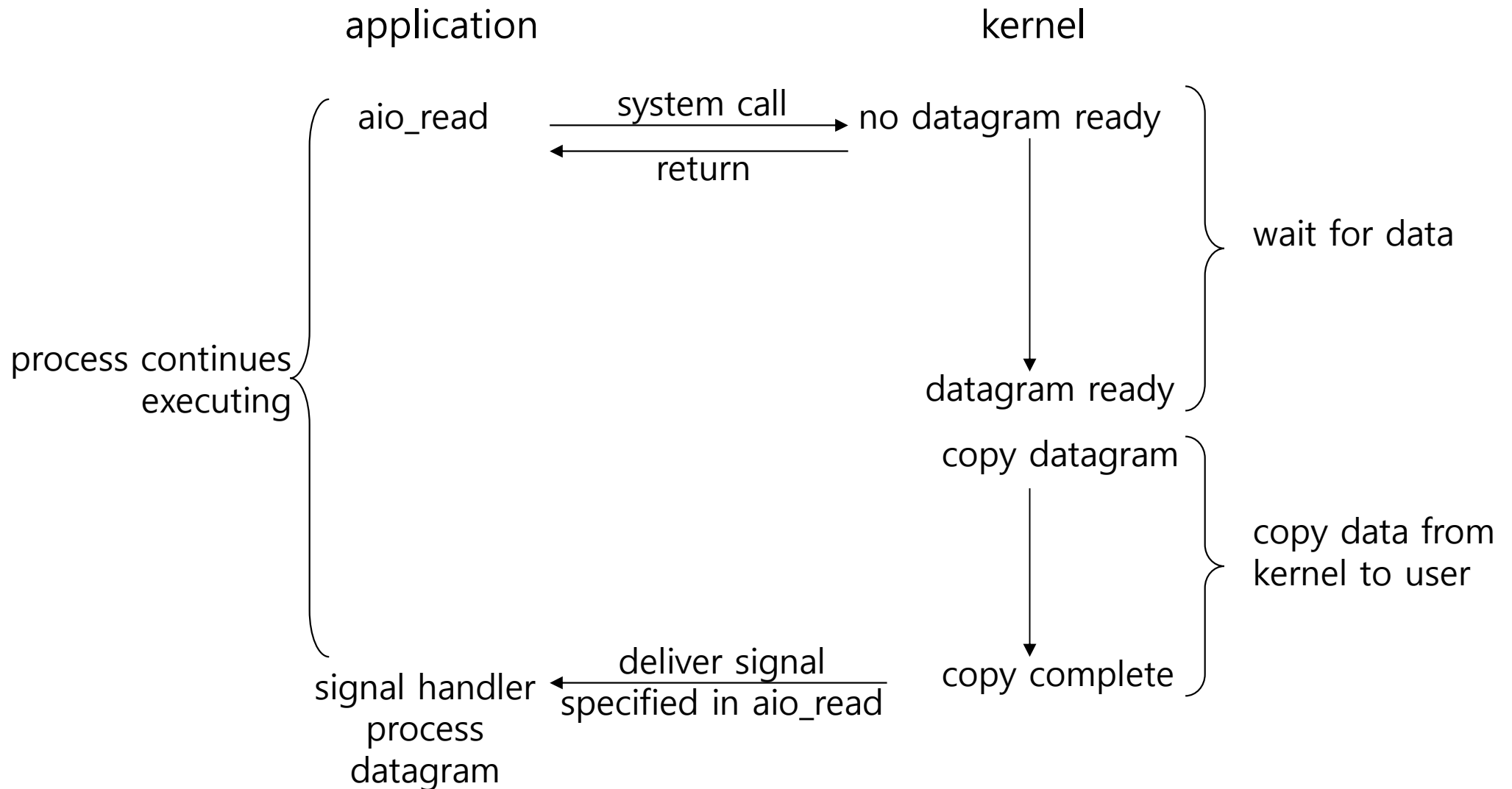
# Signal-Driven I/O Model

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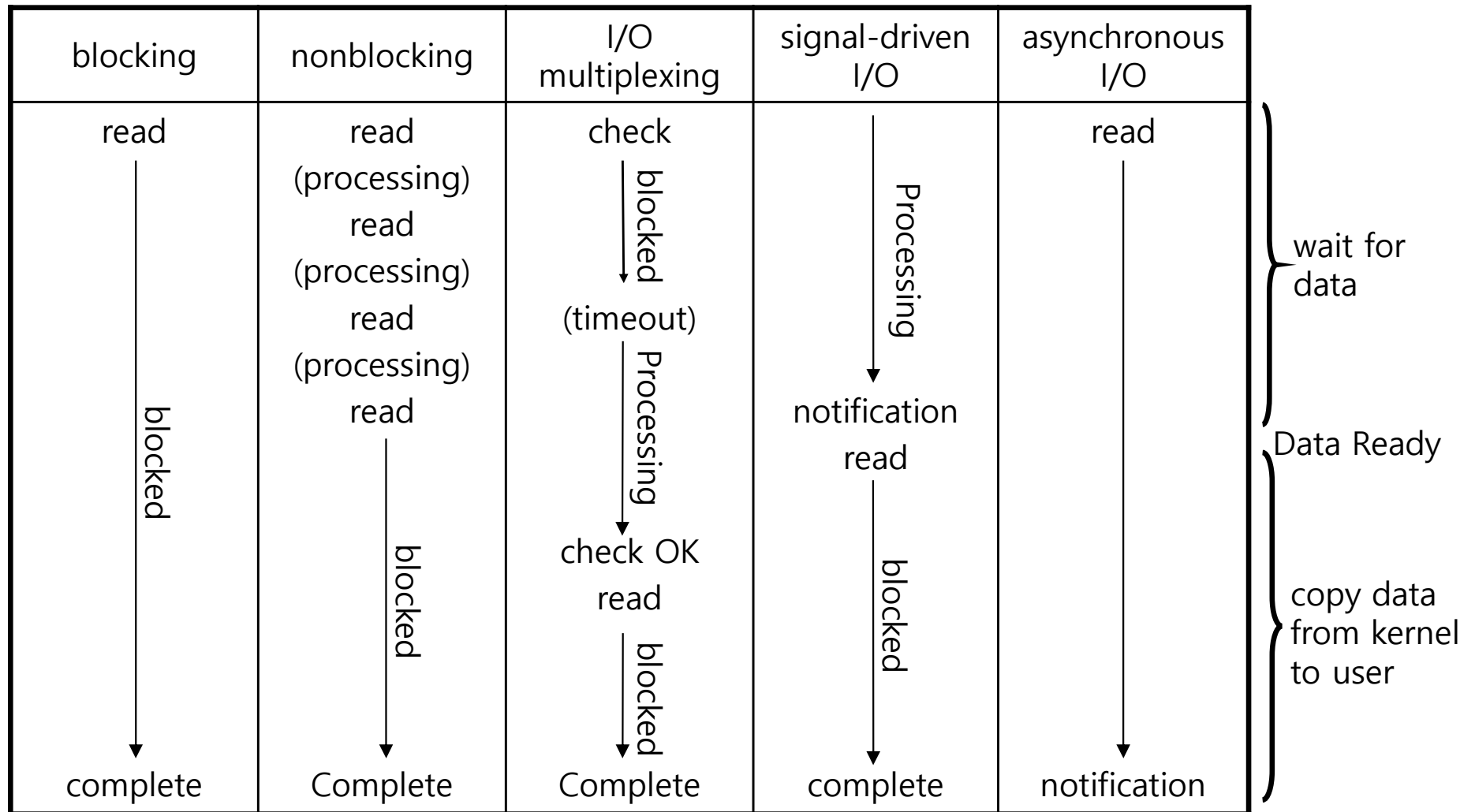


# Asynchronous I/O Model

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# I/O Model의 비교(read)



# I/O Model의 비교(write)

