Algorytmy tekstowe

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Struktury Sufiksowe i Dopasowywanie Wzorców

Struktura projektu

Pliki załączone do tego sprawozdania składają się z:

- Testów (tests)
- Plików tekstowych używanych do analizy (text_samples)
- Plików ze zmodyfikowanymi algorytmami z poprzednich laboratoriów (utils)
- Plików z algorytmami Ukkonena, tablicami sufiksów oraz problemami sufiksowymi
- Notebooków Jupyter zawierających analizę poszczególnych zagadnień tego laboratorium z omówieniem i komentarzami (w języku angielskim dla wygody)
- Plików PDF, które są exportem Notebooków Jupytera (pdfs)

Algorytm Ukkonena jest opisany w tym pliku w sekcji o tej samej nazwie.

Analiza porównawcza algorytmów znajdowania najdłuższego wspólnego podciągu wraz z analizą algorytmów dla problemu najdłuższego wspólnego podciągu dla wielu ciągów znaków oraz najdłuższego palindromicznego podciągu znajdują się w Notebooku longest_common_substring_compare.ipynb (lub PDF).

Analiza porównawcza struktur sufiksowych znajduje się w Notebooku suffix structure compare.ipynb (lub PDF).

Analiza porównawcza algorytmów wyszukiwania wzorca w tekście znajduje się w Notebooku pattern_matching_analysis.ipynb (lub PDF).

Algorytm Ukkonena

```
class End:
    def __init__(self, value):
        self.value = value
class Node:
    def __init__(self, start, end):
        self.children = {}
        self.suffix_link = None
        self.start = start
        self.end = end
        self.id = -1
    def __repr__(self):
        return f"{self.id}"
class SuffixTree:
   def __init__(self, text: str):
        Construct a suffix tree for the given text using Ukkonen's algorithm.
        Args:
            text: The input text for which to build the suffix tree
        self.text = text + "$"
        self.root = Node(None, None)
        self.active_node = self.root
        self.active_edge = 0
        self.active_length = 0
        self.remainder = ∅
        self.build_tree()
        self.count_compares = False
    def build_tree(self):
        Build the suffix tree using Ukkonen's algorithm.
        ID = 0
        END = End(0)
        def split_node(node : Node, last_split : Node = None) -> Node:
           nonlocal ID, END
           parent_node = Node(node.start, node.start + self.active_length)
            self.active_node.children[self.text[parent_node.start]] = parent_node
            # Rule 2
            if last_split is not None:
                last_split.suffix_link = parent_node
            node.start = parent_node.end
            new_leaf = Node(END.value - 1, END)
            new_leaf.id = ID
```

```
ID += 1
            parent node.children[self.text[node.start]] = node
            parent_node.children[self.text[new_leaf.start]] = new_leaf
            last_split = parent_node
            return last_split
        def add_child(node : Node, char, start) -> None:
            nonlocal ID, END
            child = Node(start, END)
            child.id = ID
            ID += 1
            node.children[char] = child
        def get_edge_text_len(node : Node):
            return node.end.value - node.start if isinstance(node.end, End) else node.end -
node.start
        n = len(self.text)
        for i in range(n):
            char = self.text[i]
            self.remainder += 1
            END.value = i + 1
            last_split = None
            while self.remainder > 0:
                if self.active_length == 0:
                    self.active_edge = i
                edge_char = self.text[self.active_edge]
                if edge_char not in self.active_node.children.keys():
                    add_child(self.active_node, char, i)
                    if last_split is not None:
                        last_split.suffix_link = self.active_node
                        last_split = None
                    if self.active_node is not self.root:
                        self.active node = self.active node.suffix link or self.root
                else:
                    node = self.active_node.children[edge_char]
                    edge_len = get_edge_text_len(node)
                    if self.active_length >= edge_len:
                        self.active_node = node
                        self.active_length -= edge_len
                        self.active_edge += edge_len
                        continue
                    if self.text[node.start + self.active_length] == char:
                        self.active_length += 1
                        if last_split:
                            last_split.suffix_link = self.active_node
```

```
last split = split node(node, last split)
                if self.active_node is self.root:
                    self.active_length -= 1
                    self.active_edge += 1
                else:
                    self.active_node = self.active_node.suffix_link or self.root
            self.remainder -= 1
def find_pattern(self, pattern: str):
    Find all occurrences of the pattern in the text.
        pattern: The pattern to search for
    Returns:
       A list of positions where the pattern occurs in the text
    node = self.root
    i = 0
    n = len(pattern)
    compares = 0
    while i < n:
        if pattern[i] not in node.children.keys():
            return ([], compares) if self.count_compares else []
        compares += 1
        child = node.children[pattern[i]]
        edge_end = child.end.value if isinstance(child.end, End) else child.end
        edge_text = self.text[child.start : edge_end]
        j = 0
        m = len(edge_text)
        while j < m and i < n:
            if pattern[i] != edge_text[j]:
                return ([], compares) if self.count_compares else []
            compares += 1
            j += 1
            i += 1
        node = child
    results = []
    def collect DFS(node : Node):
        if not node.children:
            results.append(node.id)
            return
        for child in node.children.values():
            collect_DFS(child)
```

```
collect_DFS(node)
if self.count_compares: return results, compares
return results
```

Implementacja algorytmu Ukkonena służącego do budowy drzew sufiksowych zawierająca łącza sufiksowe i osobną strukturę do wskazywania końca buduje drzewo sufiksów w czasie O(n) dla tekstu o długości n, ponieważ wykonuje operację zmiany zmiennych i dodania zaległych sufiksów w czasie O(1), a dodaje sufiksy do drzewa O(n) razy.

Algorytmy wspólnych podciągów

Analiza porównawcza z wykresami znajduje się w Notebooku longest common substring compare.ipynb (lub PDF).

```
from ukkonen import SuffixTree, Node
from suffix_array import SuffixArray
import random, string
def longest_common_substring_dp(str1: str, str2: str) -> str:
    Find the longest common substring of two strings using dynamic programming.
    Args:
        str1: First string
        str2: Second string
    Returns:
        The longest common substring
    n = len(str1)
    m = len(str2)
    max_len, end_pos = 0, 0
    DP = [[0 \text{ for } \_ \text{ in } range(n + 1)] \text{ for } \_ \text{ in } range(m + 1)]
    for i in range(1, m + 1):
        for j in range(1, n + 1):
            if str1[j - 1] == str2[i - 1]:
                DP[i][j] = DP[i - 1][j - 1] + 1
                if DP[i][j] > max_len:
                     max_len = DP[i][j]
                     end_pos = j
            else:
                DP[i][j] = 0
    return str1[end_pos - max_len: end_pos]
def longest_common_substring_sa(str1: str, str2: str) -> str:
    Find the longest common substring of two strings using a suffix array.
    Args:
        str1: First string
        str2: Second string
    Returns:
        The longest common substring
    # Concatenate the strings with a unique separator
    combined = str1 + "#" + str2 + "$"
    seperator_index = len(str1)
```

```
sa = SuffixArray(combined)
    n = len(sa.suffixes)
    lcp = [0] * (n-1)
    rank = [0] * n
    for i in range(n):
        rank[sa.suffixes[i]] = i
    for i in range(n):
        if rank[i] == n - 1:
            k = 0
            continue
        j = sa.suffixes[rank[i] + 1]
        while i + k < n and j + k < n and sa.text[i + k] == sa.text[j + k]:
            k += 1
        lcp[rank[i]] = k
        if k > 0:
            k -= 1
    max_len = 0
    position = 0
    for i in range(1, n):
       s1 = sa.suffixes[i]
        s2 = sa.suffixes[i - 1]
        if (s1 < seperator_index) != (s2 < seperator_index):</pre>
            if lcp[i - 1] > max_len:
                max_len = lcp[i - 1]
                position = s1
    return combined[position : position + max_len]
def longest_common_substring_st(str1: str, str2: str) -> str:
    Find the longest common substring of two strings using a suffix tree.
       str1: First string
       str2: Second string
    Returns:
       The longest common substring
    combined = str1 + "#" + str2 + "$"
    seperator_index = len(str1)
    longest_substring = ""
    st = SuffixTree(combined)
    def DFS(node : Node, path : list):
```

```
nonlocal longest_substring
        bits = set()
        if not node.children:
            if node.id < seperator_index:</pre>
                bits.add(0)
            elif node.id > seperator_index:
                bits.add(1)
            return bits
        for child in node.children.values():
            edge end = child.end.value if hasattr(child.end, 'value') else child.end
            edge_text = st.text[child.start : edge_end]
            bits.update(DFS(child, path + [edge_text]))
        if 0 in bits and 1 in bits:
            substring = "".join(path)
            if len(substring) > len(longest_substring):
                longest_substring = substring
        return bits
    DFS(st.root, [])
    return longest_substring
def longest_common_substring_multiple(strings: list[str]) -> str:
    Find the longest common substring among multiple strings using suffix structures.
        strings: List of strings to compare
    Returns:
        The longest common substring that appears in all strings
    n = len(strings)
    if n == 1: return strings[0]
    combined = ""
    seperator_indexes = []
    available separators = list(set(string.punctuation) - set(["".join(s) for s in strings]))
    seperators = random.sample(available_separators, n)
    j = 0
    for i, str in enumerate(strings):
        combined += str + seperators[i]
        seperator_indexes.append((j, j + len(str)))
        j += len(str) + 1
    longest_substring = ""
```

```
st = SuffixTree(combined)
    def DFS(node : Node, path : list):
        nonlocal longest_substring
        bits = set()
        if not node.children:
            for k, (start, end) in enumerate(seperator_indexes):
                if start <= node.id < end:</pre>
                    bits.add(k)
                    break
            return bits
        for child in node.children.values():
            edge_end = child.end.value if hasattr(child.end, 'value') else child.end
            edge_text = st.text[child.start : edge_end]
            bits.update(DFS(child, path + [edge_text]))
        if len(bits) == n:
            substring = "".join(path)
            if len(substring) > len(longest_substring):
                longest_substring = substring
        return bits
    DFS(st.root, [])
    return longest_substring
def longest_palindromic_substring(text: str) -> str:
    Find the longest palindromic substring in a given text using suffix structures.
    Args:
        text: Input text
    Returns:
        The longest palindromic substring
    revtext = text[::-1]
    combined = text + "#" + revtext + "$"
    st = SuffixTree(combined)
    n = len(text)
    lps = ""
    def DFS(node : Node, path : list):
        nonlocal lps
        bits = set()
        positions = set()
```

```
if not node.children:
        if node.id < n:</pre>
            bits.add(0)
            positions.add(node.id)
        elif node.id > n:
            bits.add(1)
            positions.add(node.id - (n+1))
        return bits, positions
    for child in node.children.values():
        edge_end = child.end.value if hasattr(child.end, 'value') else child.end
        edge_text = st.text[child.start : edge_end]
        child_bits, child_positions = DFS(child, path + [edge_text])
        bits.update(child_bits)
        positions.update(child_positions)
    if 0 in bits and 1 in bits:
        substring = "".join(path)
        1 = len(substring)
        for pos in positions:
            revindex = n - (pos + 1)
            if revindex in positions and 1 > len(lps):
                lps = substring
    return bits, positions
DFS(st.root, [])
return lps
```

Porównanie struktur sufiksowych

Analiza porównawcza z wykresami znajduje się w Notebooku suffix_structure_compare.ipynb (lub PDF).

```
from ukkonen import SuffixTree, Node
from suffix_array import SuffixArray, Suffix
import time, psutil, os, sys
def get_memory_usage():
    process = psutil.Process(os.getpid())
    return process.memory_info().rss / 1024 # in KB
def get_suffix_tree_size(st: SuffixTree):
    size_of_node = sys.getsizeof(Node(None, None))
    def dfs(node: Node):
        total = size_of_node
        for child in node.children.values():
            total += dfs(child)
        return total
    return dfs(st.root)
def get_suffix_array_size(sa: SuffixArray):
    size_of_suffix = sys.getsizeof(Suffix("", None))
    return len(sa.suffixes) * size_of_suffix
def compare_suffix_structures(text: str) -> dict:
    Compare suffix array and suffix tree data structures.
    Args:
        text: The input text for which to build the structures
    Returns:
       A dictionary containing:
        - Construction time for both structures
        - Memory usage for both structures
        - Size (number of nodes/elements) of both structures
    prior_mem = get_memory_usage()
    start_time = time.time()
    sa = SuffixArray(text)
    end_time = time.time()
    sarray_construction = (end_time - start_time) * 1000
    mem_after_sa = get_memory_usage()
    sarray_mem_usage = mem_after_sa - prior_mem
    start_time = time.time()
    st = SuffixTree(text)
    end time = time.time()
```

```
stree_construction = (end_time - start_time) * 1000
mem_after_st = get_memory_usage()
stree_mem_usage = mem_after_st - mem_after_sa

# Measure size of structures in bytes
sarray_size = get_suffix_array_size(sa)
stree_size = get_suffix_tree_size(st)

return {
    "suffix_array": {
        "construction_time_ms": sarray_construction,
        "memory_usage_kb": sarray_mem_usage,
        "size": sarray_size
    },
    "suffix_tree": {
        "construction_time_ms": stree_construction,
        "memory_usage_kb": stree_mem_usage,
        "size": stree_size
    }
}
```

Porównanie algorytmów wyszukiwania wzorca

A porównawcza z wykresami znajduje się w Notebooku **pattern_matching_analysis.ipynb** (lub **PDF**).

```
def compare_pattern_matching_algorithms(text: str, pattern: str) -> dict:
    Compare the performance of different pattern matching algorithms.
    Args:
       text: The text to search in
       pattern: The pattern to search for
    Returns:
       A dictionary containing the results of each algorithm:
       - Execution time in milliseconds
       - Memory usage in kilobytes
       - Number of character comparisons made
        - Positions where the pattern was found
   prior_mem = get_memory_usage()
    start_time = time.time()
    naive_result, naive_compares = naive_pattern_match(text, pattern)
    end_time = time.time()
   mem_after_naive = get_memory_usage()
    naive time exec = (end time - start time) * 1000
    naive_mem_usage = mem_after_naive - prior_mem
   sa = SuffixArray(text)
   start_time = time.time()
    sa.count_compares = True
    sa_result, sa_compares = sa.find_pattern(pattern)
    end time = time.time()
   mem_after_sa = get_memory_usage()
   sarray_time_exec = (end_time - start_time) * 1000
   sarray_mem_usage = mem_after_sa - naive_mem_usage
    st = SuffixTree(text)
    start_time = time.time()
```

```
st.count_compares = True
st_result, st_compares = st.find_pattern(pattern)
end time = time.time()
mem_after_st = get_memory_usage()
stree_time_exec = (end_time - start_time) * 1000
stree_mem_usage = mem_after_st - mem_after_sa
start time = time.time()
kmp_result, kmp_compares = kmp_pattern_match(text, pattern)
end time = time.time()
mem after_kmp = get_memory_usage()
kmp_time_exec = (end_time - start_time) * 1000
kmp_mem_usage = mem_after_kmp - mem_after_st
start_time = time.time()
bm_result, bm_compares = boyer_moore_pattern_match(text, pattern)
end_time = time.time()
mem_after_bm = get_memory_usage()
bm_time_exec = (end_time - start_time) * 1000
bm_mem_usage = mem_after_bm - mem_after_kmp
start time = time.time()
rk_result, rk_compares = rabin_karp_pattern_match(text, pattern)
end_time = time.time()
mem_after_rk = get_memory_usage()
rk_time_exec = (end_time - start_time) * 1000
rk_mem_usage = mem_after_rk - mem_after_bm
ac = AhoCorasick([pattern])
start time = time.time()
ac_result, ac_compares = ac.search(text)
end_time = time.time()
mem_after_ac = get_memory_usage()
ac_time_exec = (end_time - start_time) * 1000
ac_mem_usage = mem_after_ac - mem_after_rk
sa result.sort()
st_result.sort()
kmp_result.sort()
bm_result.sort()
rk_result.sort()
ac_result = [index for index, _ in ac_result]
ac_result.sort()
return {
    "Naive": {
        "execution_time_ms": naive_time_exec,
        "memory_usage_kb": naive_mem_usage,
        "compares": naive_compares,
        "results": naive_result
    },
    "Suffix array": {
```

```
"execution_time_ms": sarray_time_exec,
    "memory_usage_kb": sarray_mem_usage,
    "compares": sa_compares,
    "results": sa_result
},
"Suffix tree": {
   "execution_time_ms": stree_time_exec,
    "memory_usage_kb": stree_mem_usage,
    "compares": st_compares,
    "results": st_result
},
"Knuth-Morris-Pratt": {
    "execution_time_ms": kmp_time_exec,
    "memory_usage_kb": kmp_mem_usage,
    "compares": kmp_compares,
    "results": kmp_result
"Boyer-Moore": {
    "execution_time_ms": bm_time_exec,
    "memory_usage_kb": bm_mem_usage,
    "compares": bm compares,
    "results": bm_result
"Rabin-Karp": {
    "execution_time_ms": rk_time_exec,
    "memory_usage_kb": rk_mem_usage,
    "compares": rk_compares,
    "results": rk_result
"Aho-Corasick": {
    "execution_time_ms": ac_time_exec,
    "memory_usage_kb": ac_mem_usage,
    "compares": ac_compares,
    "results": ac_result
```

Algorytmy z poprzednich laboratoriów

Algorytm naiwny

```
def naive_pattern_match(text: str, pattern: str) -> list[int]:
   Implementation of the naive pattern matching algorithm.
   Args:
       text: The text to search in
       pattern: The pattern to search for
   Returns:
      A list of starting positions (0-indexed) where the pattern was found in the text
   if pattern == "": return []
   def compare(str1, str2):
       if len(str1) != len(str2): return False, 0
       compares = 0
       for i in range(len(str1)):
           compares += 1
           if str1[i] != str2[i]:
               return False, compares
       return True, compares
   n = len(text)
   m = len(pattern)
   result = []
   compares = 0
   for i in range(n):
       comp_result, add_to_compares = compare(text[i:i+m], pattern)
       if comp_result:
           result.append(i)
       compares += add_to_compares
   return result, compares
```

Aho-Corasick

```
from collections import deque
from typing import List, Tuple, Optional
class AhoCorasickNode:
   def __init__(self, char):
       self.char = char
       self.children = dict()
        self.fail_link = None
       self.outputs = []
class AhoCorasick:
   def __init__(self, patterns: List[str]):
       self.root = AhoCorasickNode(None)
       self.patterns = list(filter(lambda x: len(x) != 0, patterns))
   def _build_trie(self):
        """Builds the trie structure for the given patterns."""
        for pattern in self.patterns:
            node = self.root
            for c in pattern:
                if c not in node.children:
                    node.children[c] = AhoCorasickNode(c)
                node = node.children[c]
            node.outputs.append(pattern)
   def _build_failure_links(self):
        """Builds failure links and propagates outputs through them."""
       Q = deque([])
        for node in self.root.children.values():
            node.fail_link = self.root
            Q.append(node)
       while Q:
            node = Q.popleft()
            for c, child in node.children.items():
                fnode = node.fail_link
                while fnode is not None and c not in fnode.children:
                    fnode = fnode.fail_link
                child.fail link = fnode.children[c] if fnode else self.root
                child.outputs += child.fail_link.outputs if child.fail_link else []
                Q.append(child)
```

```
def search(self, text: str) -> Tuple[List[Tuple[int, str]], int]:
       Searches for all occurrences of patterns in the given text.
       Returns:
          List of tuples (start_index, pattern).
       if len(text) == 0: return [], 0
       results = []
       self._build_trie()
       self._build_failure_links()
       node = self.root
       compares = 0
       for i, c in enumerate(text):
           compares += 1
           while node is not None and c not in node.children:
               compares += 1
               node = node.fail_link
           if node is None:
               node = self.root
               continue
           node = node.children[c]
           for pattern in node.outputs:
               results.append((i - len(pattern) + 1, pattern))
       return results, compares
```

Boyer-Moore

```
def compute_bad_character_table(pattern: str) -> dict:
    """
    Compute the bad character table for the Boyer-Moore algorithm.

Args:
    pattern: The pattern string

Returns:
    A dictionary with keys as characters and values as the rightmost position of the character in the pattern (0-indexed)
    """
    table = {}
    for i, c in enumerate(pattern):
        table[c] = i
    return table

def compute_good_suffix_table(pattern: str) -> list[int]:
    """
    Compute the good suffix table for the Boyer-Moore algorithm.

Args:
        pattern: The pattern string
```

```
Returns:
       A list where shift[i] stores the shift required when a mismatch
       happens at position i of the pattern
   m = len(pattern)
   shift = [0] * (m + 1)
   border_pos = [0] * (m + 1)
   compares = 0
   i = m
   j = m + 1
   border_pos[i] = j
   while i > 0:
       compares += 1
       while j <= m and pattern[i - 1] != pattern[j - 1]:</pre>
            compares += 1
            if shift[j] == 0:
                shift[j] = j - i
           j = border_pos[j]
       j -= 1
       border_pos[i] = j
   j = border_pos[0]
   for i in range(m + 1):
       if shift[i] == 0:
            shift[i] = j
       if i == j:
           j = border_pos[j]
   return shift, compares
def boyer_moore_pattern_match(text: str, pattern: str) -> tuple[list[int], int]:
   Implementation of the Boyer-Moore pattern matching algorithm.
   Args:
       text: The text to search in
       pattern: The pattern to search for
   Returns:
      A list of starting positions (0-indexed) where the pattern was found in the text
   if not text or not pattern: return [], 0
   BCT, (GST, compares) = compute_bad_character_table(pattern),
compute_good_suffix_table(pattern)
   result = []
   n, m = len(text), len(pattern)
   while i + m <= n:
       while j >= 0 and pattern[j] == text[i+j]:
```

```
compares += 1
        j -= 1
compares += 1

if j < 0:
        result.append(i)
        i += GST[0]

else:
        BC_shift = j - BCT.get(text[i+j], -1)
        GS_shift = GST[j + 1]
        i += max(BC_shift, GS_shift, 1)

return result, compares</pre>
```

Knuth-Morris-Pratt

```
def compute_lps_array(pattern: str) -> list[int]:
   Compute the Longest Proper Prefix which is also Suffix array for KMP algorithm.
       pattern: The pattern string
   Returns:
       The LPS array
   n = len(pattern)
   result = [0] * n
   1 = 0
   for i in range(1,n):
       while 1 > 0 and pattern[i] != pattern[1]:
           l = result[l - 1]
       if pattern[i] == pattern[l]:
           1 += 1
       result[i] = 1
   return result
def kmp_pattern_match(text: str, pattern: str) -> tuple[list[int], int]:
   Implementation of the Knuth-Morris-Pratt pattern matching algorithm.
   Args:
       text: The text to search in
       pattern: The pattern to search for
   Returns:
      A list of starting positions (0-indexed) where the pattern was found in the text
   if pattern == "" or text == "": return [], 0
   LPS = compute_lps_array(pattern)
   n = len(text)
   m = len(pattern)
   result = []
   ti = 0
   pi = 0
   compares = ∅
   while ti < n:
       if pattern[pi] == text[ti]:
           ti += 1
           pi += 1
           compares += 1
       if pi == m:
           result.append(ti - pi)
```

```
pi = LPS[pi - 1]

elif ti < n and pattern[pi] != text[ti]:
    compares += 1
    if pi != 0:
        pi = LPS[pi - 1]
    else:
        ti += 1

return result, compa</pre>
```

Rabin-Karp

```
def rabin_karp_pattern_match(text: str, pattern: str, prime: int = 101) ->
tuple[list[int], int]:
   Implementation of the Rabin-Karp pattern matching algorithm.
   Args:
       text: The text to search in
       pattern: The pattern to search for
       prime: A prime number used for the hash function
   Returns:
    A list of starting positions (0-indexed) where the pattern was found in the text
   def compare(str1, str2):
        if len(str1) != len(str2): return False, 0
       compares = ∅
       for i in range(len(str1)):
            compares += 1
            if str1[i] != str2[i]:
                return False, compares
        return True, compares
   def hash(old_hash, oldc):
        return (old_hash + oldc) % prime
   def unhash(old_hash, oldc):
        return (old_hash - oldc + prime) % prime
   def hash_string(string):
       hash res = 0
       for char in string:
            hash_res = hash(hash_res, ord(char))
        return hash res
   if text == "" or pattern == "": return [], 0
   result = []
   n = len(text)
   m = len(pattern)
   hp = hash_string(pattern)
   htw = hash_string(text[:m])
   compares = 0
   i = 0
   while True:
        if htw == hp and text[i:i+m] == pattern:
            comp_result, add_to_compares = compare(text[i:i+m], pattern)
            compares += add_to_compares
            if comp result:
                result.append(i)
        if n <= i + m:
            break
```

```
htw = hash(htw, ord(text[i+m]))
htw = unhash(htw, ord(text[i]))
i += 1
return result, compares
```

Testy

Test algorytmu Ukkonena tworzenia drzew sufiksowych (tests/test_ukkonen.py)

```
import pytest
from ukkonen import SuffixTree
import sys
sys.setrecursionlimit(10000)
class TestUkkonen:
   def test_ukkonen_by_number_of_leaves(self):
       short texts = [
           "abc",
           "aaa",
           "abab",
           "racecar",
           "banana",
           "abcabc",
           "aaaaa",
           "xyzzy",
           "aabbaa",
           "skibidi",
           "addadda"
       medium_texts = [
           "abracadabrabracabracad",
           "mississippiisreallycool",
           "abcabcabcabcabcabcabcabc",
           "zzzyyyxxxyyyzzzzyx",
           "abababababababababab",
           "abcdabcdabcdabcd",
           "xyzxyzxyzxyzxyzxyzxyz",
           "thequickbrownfoxjumpsoverthelazydog",
           "loremipsumdolorsitametconsectetur",
           "aaaaabbbbbcccccdddddeeeeefffff"
       long_texts = [
           "ab" * 150,
           "abcde" * 60,
           "banana" * 50,
           "loremipsum" * 25,
           "zzzyyyxxxwwwvvvuuutttsssrrrqqqpppooonnnmmmlllkkkjjjiii",
           "skibiditoilet" * 20,
           "helloworld" * 30,
           "abcdefghijklmnopqrstuvwxyz" * 4,
```

```
very_long_text = [
           "abababab" * 70,
           "xyz" * 170 + "end",
           "".join(["abcde"[i % 5] for i in range(600)]), # cyclic
           "mississippimississippimis",
           "mississippimississippimississippimississippimississippimississippi",
           "Nory was a Catholic because her mother was a Catholic, and Nory's mother was a Catholic
because her father was a Catholic, and her father was a Catholic because his mother was a Catholic, or
had been."
       def check leaves(tree : SuffixTree) -> bool:
           def count_leaves(node):
               if not node.children:
               return sum(count_leaves(child) for child in node.children.values())
           return len(tree.text) == count_leaves(tree.root)
       list = short_texts + medium_texts + long_texts + very_long_text
       for text in list:
           suffix_tree = SuffixTree(text)
           assert check_leaves(suffix_tree)
   def test_ukkonen_finding_pattern(self):
       def test_batch(text : str, patterns : dict) -> bool:
           suffix tree = SuffixTree(text)
           for pattern, expected in patterns.items():
               result = suffix_tree.find_pattern(pattern)
               result.sort()
               assert result == expected, f"Failed for pattern '{pattern}': got {result}, expected
{expected}"
           print(f"PASSED pattern matching test")
           return True
       assert test_batch("abracadabra", {
           "abra": [0, 7],
           "cad": [4],
           "a": [0, 3, 5, 7, 10],
           "ra": [2, 9],
           "xyz": [],
       })
       assert test_batch("banana", {
           "na": [2, 4],
           "ana": [1, 3],
           "nana": [2],
           "banana": [0],
           "bananas": [],
```

```
assert test_batch("LLLorem im daolor sit amet, consectetur adipLorema elit, sed do eiusmod
Loremc incididunt ut labore ", {"Lorem" : [2, 44, 72]})
assert test_batch("abracad" * 50, {"abra" : list(range(0,350,7))})
assert test_batch("Nory was a Catholic because her mother was a Catholic, and Nory's" \
" mother was a Catholic because her father was a Catholic, and her father was a Catholic" \
" because his mother was a Catholic, or had been.", {"Catholic" : [11, 45, 79, 113, 144, 178]})
```

Test znajdowania wzorca za pomocą tablicy sufiksów (tests/test_suffix_array.py)

```
import pytest
from suffix_array import SuffixArray
class TestSuffixArray:
   def test_array_finding_pattern(self):
        def test_batch(text : str, patterns : dict) -> bool:
            suffix_array = SuffixArray(text)
            for pattern, expected in patterns.items():
                result = suffix_array.find_pattern(pattern)
                result.sort()
                assert result == expected, f"Failed for pattern '{pattern}': got {result}, expected
{expected}"
            print(f"PASSED pattern matching test")
            return True
        assert test_batch("abracadabra", {
            "abra": [0, 7],
            "cad": [4],
            "a": [0, 3, 5, 7, 10],
            "ra": [2, 9],
            "xyz": [],
        })
        assert test_batch("banana", {
            "na": [2, 4],
            "ana": [1, 3],
            "nana": [2],
            "banana": [0],
            "bananas": [],
        })
        assert test_batch("abracad" * 50, {"abra" : list(range(0,350,7))})
        <code>assert test_batch("Nory was a Catholic because her mother was a Catholic, and Nory's" \</code>
        " mother was a Catholic because her father was a Catholic, and her father was a Catholic" \setminus
        " because his mother was a Catholic, or had been.", {"Catholic" : [11, 45, 79, 113, 144,
178]})
```

Test dla problemów sufiksowych (tests/test_substring_problems.py)

```
import pytest
from substring_problems import longest_common_substring_sa, longest_common_substring_st,
longest common substring dp, longest palindromic substring, longest common substring multiple
class TestSubstringProblems:
   def test lcs dp(self):
        test_cases = [
            ("abcdef", "abc", "abc"),
            ("xyz", "abc", ""),
            ("banana", "banana", "banana"),
            ("hello", "lo", "lo"),
            ("startmatch", "start", "start"), ("abracadabra", "racad", "racad"),
            ("bajojajo bajojajo", "ja ci dam pajacu bajojajo", " bajojajo"),
            ("abababababa", "bababababa", "bababababa"),
            ("Zażółć gęślą jaźń", "gęślą", "gęślą"),
            ("abcdef", "", ""),
            ("", "", ""),
            ("x", "x", "x"),
            ("abcxyz123", "xyx789abc", "abc"),
            ("aaaaabbbbcccc", "xxxbbbbyyyy", "bbbb"),
            ("testlongest", "longest", "longest"),
            ("prefixmatch", "matchpostfix", "match"),
            ("hello, world!", "world!", "world!"),
            ("hello≌world", "≅wo", "≅wo"),
        ]
        for s1, s2, expected in test_cases:
            result = longest_common_substring_dp(s1,s2)
            assert result == expected
    def test_lcs_sa(self):
        test_cases = [
            ("abcdef", "abc", "abc"),
            ("xyz", "abc", ""),
            ("banana", "banana"),
            ("hello", "lo", "lo"),
            ("startmatch", "start", "start"), ("abracadabra", "racad", "racad"),
            ("bajojajo bajojajo", "ja ci dam pajacu bajojajo", " bajojajo"),
            ("abababababa", "bababababa", "bababababa"),
            ("Zażółć gęślą jaźń", "gęślą", "gęślą"),
            ("abcdef", "", ""),
            ("", "", ""),
            ("x", "x", "x"),
            ("abcxyz123", "xyx789abc", "abc"),
            ("aaaaabbbbcccc", "xxxbbbbyyyy", "bbbb"),
            ("testlongest", "longest", "longest"),
            ("prefixmatch", "matchpostfix", "match"),
            ("hello, world!", "world!", "world!"),
            ("hello⇔world", "⇔wo", "⇔wo"),
```

```
for s1, s2, expected in test_cases:
         result = longest_common_substring_sa(s1,s2)
         assert result == expected
def test_lcs_st(self):
    test_cases = [
         ("abcdef", "abc", "abc"),
         ("xyz", "abc", ""),
         ("banana", "banana", "banana"),
         ("hello", "lo", "lo"),
         ("startmatch", "start", "start"),
         ("abracadabra", "racad", "racad"),
         ("bajojajo bajojajo", "ja ci dam pajacu bajojajo", " bajojajo"),
         ("abababababab", "bababababa", "bababababa"),
         ("Zażółć gęślą jaźń", "gęślą", "gęślą"),
         ("abcdef", "", ""),
         ("", "", ""),
         ("x", "x", "x"),
         ("abcxyz123", "xyx789abc", "abc"),
         ("aaaaabbbbcccc", "xxxbbbbyyyy", "bbbb"),
         ("testlongest", "longest", "longest"),
("prefixmatch", "matchpostfix", "match"),
         ("hello, world!", "world!", "world!"),
         ("hello<mark>⇔</mark>world", "<mark>⇔</mark>wo", "<mark>⇔</mark>wo"),
    for s1, s2, expected in test_cases:
         result = longest_common_substring_st(s1,s2)
         assert result == expected
def test_multiple_lcs(self):
    test_cases = {
        ("abc", "abc"): "abc",
("abc", "def"): "",
("abcde", "cdefg", "defgh"): "de",
         ("ababab", "babab", "abab"): "abab",
         ("aaaa", "aa", "aaa"): "aa",
         ("racecar", "myrace", "racing"): "rac",
         ("abcXXXXcba", "abc123cba", "zzzcbayy"): "cba",
         ("a#b@c", "#b@c$d", "@c$def"): "@c",
         ("123!abc", "!abc456", "xyz!abc"): "!abc",
         ("",): "",
         ("abc",): "abc",
         (): "",
         ("banana", "ananas", "bandana"): "ana",
("thequickbrownfox", "quickfox", "lazyquickfox"): "quick",
         ("123456", "456789", "0456123"): "456",
         ("ഐhello,", ", ", ", "hello", "sayhello,"): "hello",
    for strings, expected in test_cases.items():
         result = longest_common_substring_multiple(strings)
         assert result == expected
```

```
def test_palindrome(self):
    test_cases = {
        "cbbd": "bb",
        "a": "a",
        "racecar": "racecar",
        "abacdfgdcaba": "aba",
        "": "",
        "abcddcbazzz": "abcddcba",
        "banana": "anana",
        "level": "level",
        "aaabaaaa": "aaabaaaa",
        "xyzzyx": "xyzzyx",
        "bajojab cd ad d": "bajojab",
    }
    for text, expected in test_cases.items():
        result = longest_palindromic_substring(text)
        assert result == expected
```

Wyniki testów