

# Chapter 4

- 4.1. SOP form:  $f = \bar{x}_1x_2 + \bar{x}_2x_3$   
 POS form:  $f = (\bar{x}_1 + \bar{x}_2)(x_2 + x_3)$
- 4.2. SOP form:  $f = x_1\bar{x}_2 + x_1x_3 + \bar{x}_2x_3$   
 POS form:  $f = (x_1 + x_3)(x_1 + \bar{x}_2)(\bar{x}_2 + x_3)$
- 4.3. SOP form:  $f = \bar{x}_1x_2x_3\bar{x}_4 + x_1x_2\bar{x}_3x_4 + \bar{x}_2x_3x_4$   
 POS form:  $f = (\bar{x}_1 + x_4)(x_2 + x_3)(\bar{x}_2 + \bar{x}_3 + \bar{x}_4)(x_2 + x_4)(x_1 + x_3)$
- 4.4. SOP form:  $f = \bar{x}_2\bar{x}_3 + \bar{x}_2\bar{x}_4 + x_2x_3x_4$   
 POS form:  $f = (\bar{x}_2 + x_3)(x_2 + \bar{x}_3 + \bar{x}_4)(\bar{x}_2 + x_4)$
- 4.5. SOP form:  $f = \bar{x}_3\bar{x}_5 + \bar{x}_3x_4 + x_2x_4\bar{x}_5 + \bar{x}_1x_3\bar{x}_4x_5 + x_1x_2\bar{x}_4x_5$   
 POS form:  $f = (\bar{x}_3 + x_4 + x_5)(\bar{x}_3 + \bar{x}_4 + \bar{x}_5)(x_2 + \bar{x}_3 + \bar{x}_4)(x_1 + x_3 + x_4 + \bar{x}_5)(\bar{x}_1 + x_2 + x_4 + \bar{x}_5)$
- 4.6. SOP form:  $f = \bar{x}_2x_3 + \bar{x}_1x_5 + \bar{x}_1x_3 + \bar{x}_3\bar{x}_4 + \bar{x}_2x_5$   
 POS form:  $f = (\bar{x}_1 + \bar{x}_2 + \bar{x}_3)(\bar{x}_1 + \bar{x}_2 + \bar{x}_4)(x_3 + \bar{x}_4 + x_5)$
- 4.7. SOP form:  $f = x_3\bar{x}_4\bar{x}_5 + \bar{x}_3\bar{x}_4x_5 + x_1x_4x_5 + x_1x_2x_4 + x_3x_4x_5 + \bar{x}_2x_3x_4 + x_2\bar{x}_3x_4\bar{x}_5$   
 POS form:  $f = (x_3 + x_4 + x_5)(\bar{x}_3 + x_4 + \bar{x}_5)(x_1 + \bar{x}_2 + \bar{x}_3 + \bar{x}_4 + x_5)$
- 4.8.  $f = \sum m(0, 7)$   
 $f = \sum m(1, 6)$   
 $f = \sum m(2, 5)$   
 $f = \sum m(0, 1, 6)$   
 $f = \sum m(0, 2, 5)$   
 etc.
- 4.9.  $f = x_1x_2x_3 + x_1x_2x_4 + x_1x_3x_4 + x_2x_3x_4$
- 4.10. SOP form:  $f = x_1x_2\bar{x}_3 + x_1\bar{x}_2x_4 + x_1x_3\bar{x}_4 + \bar{x}_1x_2x_3 + \bar{x}_1x_3x_4 + x_2\bar{x}_3x_4$   
 POS form:  $f = (x_1 + x_2 + x_3)(x_1 + x_2 + x_4)(x_1 + x_3 + x_4)(x_2 + x_3 + x_4)(\bar{x}_1 + \bar{x}_2 + \bar{x}_3 + \bar{x}_4)$   
 The POS form has lower cost.
- 4.11. The statement is false. As a counter example consider  $f(x_1, x_2, x_3) = \sum m(0, 5, 7)$ .  
 Then, the minimum-cost SOP form  $f = x_1x_3 + \bar{x}_1\bar{x}_2\bar{x}_3$  is unique.  
 But, there are two minimum-cost POS forms:  
 $f = (x_1 + \bar{x}_3)(\bar{x}_1 + x_3)(x_1 + \bar{x}_2)$  and  
 $f = (x_1 + \bar{x}_3)(\bar{x}_1 + x_3)(\bar{x}_2 + x_3)$

4.12. If each circuit is implemented separately:

$$f = \overline{x_1}\overline{x_4} + \overline{x_1}x_2x_3 + x_1\overline{x_2}x_4 \quad \text{Cost} = 15$$

$$g = \overline{x_1}\overline{x_3}\overline{x_4} + \overline{x_2}x_3\overline{x_4} + x_1\overline{x_3}x_4 + x_1x_2x_4 \quad \text{Cost} = 21$$

In a combined circuit:

$$f = \overline{x_2}x_3\overline{x_4} + \overline{x_1}\overline{x_3}\overline{x_4} + x_1\overline{x_2}\overline{x_3}x_4 + \overline{x_1}x_2x_3$$

$$g = \overline{x_2}x_3\overline{x_4} + \overline{x_1}\overline{x_3}\overline{x_4} + x_1\overline{x_2}\overline{x_3}x_4 + x_1x_2x_4$$

The first 3 product terms are shared, hence the total cost is 31.

4.13. If each circuit is implemented separately:

$$f = \overline{x_1}x_2x_4 + x_2x_4x_5 + x_3\overline{x_4}\overline{x_5} + \overline{x_1}\overline{x_2}\overline{x_4}x_5 \quad \text{Cost} = 22$$

$$g = \overline{x_3}\overline{x_5} + \overline{x_4}\overline{x_5} + \overline{x_1}\overline{x_2}\overline{x_4} + \overline{x_1}x_2x_4 + x_2x_4x_5 \quad \text{Cost} = 24$$

In a combined circuit:

$$f = \overline{x_1}x_2x_4 + x_2x_4x_5 + x_3\overline{x_4}\overline{x_5} + \overline{x_1}\overline{x_2}\overline{x_4}x_5$$

$$g = \overline{x_1}x_2x_4 + x_2x_4x_5 + x_3\overline{x_4}\overline{x_5} + \overline{x_1}\overline{x_2}\overline{x_4}x_5 + \overline{x_3}\overline{x_5}$$

The first 4 product terms are shared, hence the total cost is 31. Note that in this implementation  $f \subseteq g$ , thus  $g$  can be realized as  $g = f + \overline{x_3}\overline{x_5}$ , in which case the total cost is lowered to 28.

4.14.  $f = (x_3 \uparrow g) \uparrow ((g \uparrow g) \uparrow x_4)$  where  $g = (x_1 \uparrow (x_2 \uparrow x_2)) \uparrow ((x_1 \uparrow x_1) \uparrow x_2)$

4.15.  $\overline{f} = (((x_3 \downarrow x_3) \downarrow g) \downarrow ((g \downarrow g) \downarrow (x_4 \downarrow x_4)))$ , where  
 $g = ((x_1 \downarrow x_1) \downarrow x_2) \downarrow (x_1 \downarrow (x_2 \downarrow x_2))$ . Then,  $f = \overline{\overline{f}} \downarrow \overline{f}$ .

4.16.  $f = (g \uparrow k) \uparrow ((g \uparrow g) \uparrow (k \uparrow k))$ , where  $g = (x_1 \uparrow x_1) \uparrow (x_2 \uparrow x_2) \uparrow (x_5 \uparrow x_5)$   
and  $k = (x_3 \uparrow (x_4 \uparrow x_4)) \uparrow ((x_3 \uparrow x_3) \uparrow x_4)$

4.17.  $\overline{f} = (g \downarrow k) \downarrow ((g \downarrow g) \downarrow (k \downarrow k))$ , where  $g = x_1 \downarrow x_2 \downarrow x_5$   
and  $k = ((x_3 \downarrow x_3) \downarrow x_4) \downarrow (x_3 \downarrow (x_4 \downarrow x_4))$ . Then,  $f = \overline{\overline{f}} \downarrow \overline{f}$ .

4.18.  $f = \overline{x_1}(x_2 + x_3)(x_4 + x_5) + x_1(\overline{x_2} + x_3)(\overline{x_4} + x_5)$

4.19.  $f = x_1\overline{x_3}\overline{x_4} + x_2\overline{x_3}\overline{x_4} + x_1x_3x_4 + x_2x_3x_4 = (x_1 + x_2)\overline{x_3}\overline{x_4} + (x_1 + x_2)x_3x_4$   
This requires 2 OR and 2 AND gates.

4.20.  $f = x_1 \cdot g + \overline{x_1} \cdot \overline{g}$ , where  $g = \overline{x_3}x_4 + x_3\overline{x_4}$

4.21.  $f = g \cdot h + \overline{g} \cdot \overline{h}$ , where  $g = x_1x_2$  and  $h = x_3 + x_4$

4.22. Let  $D(0, 20)$  be 0 and  $D(15, 26)$  be 1. Then decomposition yields:

$$g = x_5(\overline{x_1} + x_2)$$

$$f = (\overline{x_3}\overline{x_4} + x_3x_4)g + \overline{x_3}x_4\overline{g} = x_3x_4g + \overline{x_3}\overline{x_4}g + \overline{x_3}x_4\overline{g}$$

$$\text{Cost} = 9 + 18 = 27$$

The optimal SOP form is:

$$f = \overline{x}_3 x_4 \overline{x}_5 + \overline{x}_1 x_3 x_4 x_5 + x_1 \overline{x}_2 \overline{x}_3 x_4 + \overline{x}_1 \overline{x}_3 \overline{x}_4 x_5 + x_2 \overline{x}_3 \overline{x}_4 x_5 + x_2 x_3 x_4 x_5$$

$$\text{Cost} = 7 + 29 = 36$$

4.23. The prime implicants are generated as follows:

List 1			List 2		
0	0 0 0 0	✓	0,2	0 0 x 0	
2	0 0 1 0	✓	0,4	0 x 0 0	
4	0 1 0 0	✓	0,8	x 0 0 0	
8	1 0 0 0	✓	4,5	0 1 0 x	
5	0 1 0 1	✓	8,9	1 0 0 x	
9	1 0 0 1	✓	5,7	0 1 x 1	
7	0 1 1 1	✓	7,15	x 1 1 1	
15	1 1 1 1	✓			

The initial prime implicant table is

Prime implicant	Minterm							
	0	2	4	5	7	8	9	15
$p_1 = 0 0 x 0$	✓	✓						
$p_2 = 0 x 0 0$	✓		✓					
$p_3 = x 0 0 0$	✓						✓	
$p_4 = 0 1 0 x$			✓	✓				
$p_5 = 1 0 0 x$							✓	✓
$p_6 = 0 1 x 1$				✓	✓			
$p_7 = x 1 1 1$					✓			✓

The prime implicants  $p_1$ ,  $p_5$  and  $p_7$  are essential. Removing these prime implicants gives

Prime implicant	Minterm	
	4	5
$p_2$	✓	
$p_3$		
$p_4$	✓	✓
$p_6$		✓

Since  $p_4$  covers both minterms, the final cover is

$$\begin{aligned} C &= \{p_1, p_4, p_5, p_7\} \\ &= \{00x0, 010x, 100x, x111\} \end{aligned}$$

and the function is implemented as

$$f = \overline{x}_1\overline{x}_2\overline{x}_4 + \overline{x}_1x_2\overline{x}_3 + x_1\overline{x}_2\overline{x}_3 + x_2x_3x_4$$

4.24. The prime implicants are generated as follows:

List 1			List 2			List 3		
0	0 0 0 0	✓	0,4	0 x 0 0		3,7,11,15	x x 1 1	
4	0 1 0 0	✓	0,8	x 0 0 0		9,11,13,15	1 x x 1	
8	1 0 0 0	✓	4,6	0 1 x 0				
			8,9	1 0 0 x				
3	0 0 1 1	✓	3,7	0 x 1 1	✓			
6	0 1 1 0	✓	3,11	x 0 1 1	✓			
9	1 0 0 1	✓	6,7	0 1 1 x				
7	0 1 1 1	✓	9,11	1 0 x 1	✓			
11	1 0 1 1	✓	9,13	1 x 0 1	✓			
13	1 1 0 1	✓						
15	1 1 1 1	✓	7,15	x 1 1 1	✓			
			11,15	1 x 1 1	✓			
			13,15	1 1 x 1	✓			

The initial prime implicant table is

Prime implicant	Minterm					
	0	4	6	8	9	15
$p_1 = 0 x 0 0$	✓	✓				
$p_2 = x 0 0 0$	✓			✓		
$p_3 = 0 1 x 0$		✓	✓			
$p_4 = 1 0 0 x$				✓	✓	
$p_5 = 0 1 1 x$			✓			
$p_6 = x x 1 1$						✓
$p_7 = 1 x x 1$					✓	✓

There are no essential prime implicants. Prime implicant  $p_3$  dominates  $p_5$  and their costs are the same, so remove  $p_5$ . Similarly,  $p_7$  dominates  $p_6$ , so remove  $p_6$ . This gives

Prime implicant	Minterm					
	0	4	6	8	9	15
$p_1$	✓	✓				
$p_2$	✓			✓		
$p_3$		✓	✓			
$p_4$				✓	✓	
$p_7$					✓	✓

Now,  $p_3$  and  $p_7$  are essential, which leaves

Prime implicant	Minterm 0 8	
$p_1$	✓	
$p_2$	✓	✓
$p_4$		✓

Choosing  $p_2$  results in the minimum cost cover

$$\begin{aligned}
 C &= \{p_2, p_3, p_7\} \\
 &= \{x000, 01x0, 1xx1\}
 \end{aligned}$$

and the function is implemented as

$$f = \overline{x}_2\overline{x}_3\overline{x}_4 + \overline{x}_1x_2\overline{x}_4 + x_1x_4$$

4.25. The prime implicants are generated as follows:

List 1			List 2			List 3		
0	0 0 0 0	✓	0,4 0,8	0 x 0 0 x 0 0 0	✓	0,4,8,12	x x 0 0	
4	0 1 0 0	✓	4,5	0 1 0 x	✓	4,5,12,13	x 1 0 x	
8	1 0 0 0	✓	4,12	x 1 0 0	✓	8,9,12,13	1 x 0 x	
3	0 0 1 1	✓	8,9	1 0 0 x	✓			
5	0 1 0 1	✓	8,12	1 x 0 0	✓			
9	1 0 0 1	✓						
12	1 1 0 0	✓						
7	0 1 1 1	✓	3,7	0 x 1 1				
11	1 0 1 1	✓	3,11	x 0 1 1				
13	1 1 0 1	✓	5,7	0 1 x 1				
14	1 1 1 0	✓	5,13	x 1 0 1	✓			
			9,11	1 0 x 1				
			9,13	1 x 0 1	✓			
			12,13	1 1 0 x	✓			
			12,14	1 1 x 0				

The initial prime implicant table is

Prime implicant	Minterm					
	0	3	4	5	7	9 11
$p_1 = 0 \ x \ 1 \ 1$		✓			✓	
$p_2 = x \ 0 \ 1 \ 1$		✓				✓
$p_3 = 0 \ 1 \ x \ 1$				✓	✓	
$p_4 = 1 \ 0 \ x \ 1$						✓ ✓
$p_5 = x \ x \ 0 \ 0$	✓		✓			
$p_6 = x \ 1 \ 0 \ x$			✓	✓		
$p_7 = 1 \ x \ 0 \ x$						✓
$p_8 = 1 \ 1 \ x \ 0$						

Prime implicant  $p_5$  is essential, so remove columns 0 and 4 to get

Prime implicant	Minterm				
	3	5	7	9	11
$p_1$	✓		✓		
$p_2$	✓				✓
$p_3$		✓	✓		
$p_4$				✓	✓
$p_6$		✓			
$p_7$				✓	

Here,  $p_3$  dominates  $p_6$ , and  $p_4$  dominates  $p_7$ ; but costs of  $p_3$  and  $p_4$  are greater than the costs of  $p_6$  and  $p_7$ , respectively. So, use branching. First choose  $p_3$  to be in the final cover, which leads to

Prime implicant	Minterm		
	3	9	11
$p_1$	✓		
$p_2$	✓		✓
$p_4$		✓	✓
$p_6$			
$p_7$		✓	

Now, choose  $p_2$  and  $p_7$  (lower cost than  $p_4$ ) to cover the remaining minterms. The resulting cover is

$$\begin{aligned}
 C &= \{p_2, p_3, p_5, p_7\} \\
 &= \{x011, 01x1, xx00, 1x0x\}
 \end{aligned}$$

Next, assume that  $p_3$  is not included in the final cover, which leads to

Prime implicant	Minterm				
	3	5	7	9	11
$p_1$	✓		✓		
$p_2$	✓				✓
$p_4$				✓	✓
$p_6$		✓			
$p_7$				✓	

Then  $p_6$  is essential. Also, column 3 dominates 7, hence remove 3 giving

Prime implicant	Minterm		
	7	9	11
$p_1$	✓		
$p_2$			✓
$p_4$		✓	✓
$p_7$		✓	

Choose  $p_1$  and  $p_4$ , which results in the cover

$$\begin{aligned} C &= \{p_1, p_4, p_5, p_6\} \\ &= \{0x11, 10x1, xx00, x10x\} \end{aligned}$$

Both covers have the same cost, so choosing the first cover the function can be implemented as

$$f = \overline{x}_2 x_3 x_4 + \overline{x}_1 x_2 x_4 + \overline{x}_3 \overline{x}_4 + x_1 \overline{x}_3$$

Observe that if we had not taken the cost of prime implicants (rows) into account and pursued the dominance of  $p_3$  over  $p_6$  and  $p_4$  over  $p_7$ , then we would have removed  $p_6$  and  $p_7$  giving the following table

Prime implicant	Minterm				
	3	5	7	9	11
$p_1$	✓		✓		
$p_2$	✓				✓
$p_3$		✓	✓		
$p_4$				✓	✓

Now  $p_3$  and  $p_4$  are essential. Also choose  $p_1$ , so that

$$\begin{aligned} C &= \{p_1, p_3, p_4, p_5\} \\ &= \{0x11, 01x1, 10x1, xx00\} \end{aligned}$$

The cost of this cover is greater by one literal compared to both covers derived above.

4.26. Note that  $X \# Y = X \cdot \overline{Y}$ . Therefore,

$$\begin{aligned}(A \cdot B) \# C &= A \cdot B \cdot \overline{C} \\ (A \# C) \cdot (B \# C) &= A \cdot \overline{C} \cdot B \cdot \overline{C} \\ &= A \cdot B \cdot \overline{C}\end{aligned}$$

Similarly,

$$\begin{aligned}(A + B) \# C &= (A + B) \cdot \overline{C} \\ &= A \cdot \overline{C} + B \cdot \overline{C} \\ (A \# C) + (B \# C) &= A \cdot \overline{C} + B \cdot \overline{C}\end{aligned}$$

4.27. The initial cover is  $C^0 = \{0000, 0011, 0100, 0101, 0111, 1000, 1001, 1111\}$ .

Using the \*-product get the prime implicants

$$P = \{00x0, 0x00, x000, 010x, 01x1, 100x, x111\}.$$

The minimum cover is  $C_{\text{minimum}} = \{00x0, 010x, 100x, x111\}$ , which corresponds to  $f = \overline{x}_1\overline{x}_2\overline{x}_4 + \overline{x}_1x_2\overline{x}_3 + x_1\overline{x}_2\overline{x}_3 + x_2x_3x_4$ .

4.28. The initial cover is  $C^0 = \{0x0x0, 110xx, x1101, 1001x, 11110, 01x10, 0x011\}$ .

Using the \*-product get the prime implicants

$$P = \{0x0x0, xx01x, x1x10, 110xx, x10x0, 11x01, x1101\}.$$

The minimum cover is  $C_{\text{minimum}} = \{0x0x0, xx01x, x1x10, 110xx, x1101\}$ , which corresponds to  $f = \overline{x}_1\overline{x}_3\overline{x}_5 + \overline{x}_3x_4 + x_2x_4\overline{x}_5 + x_1x_2\overline{x}_3 + x_2x_3\overline{x}_4x_5$ .

4.29. The initial cover is  $C^0 = \{00x0, 100x, x010, 1111, 00x1, 011x\}$ .

Using the \*-product get the prime implicants  $P = \{00xx, 0x1x, x00x, x0x0, x111\}$ .

The minimum-cost cover is  $C_{\text{minimum}} = \{x00x, x0x0, x111\}$ , which corresponds to  $f = \overline{x}_2\overline{x}_3 + \overline{x}_2\overline{x}_4 + x_2x_3x_4$ .

4.30. Expansion of  $\overline{x}_1\overline{x}_2\overline{x}_3$  gives  $\overline{x}_1$ .

Expansion of  $\overline{x}_1\overline{x}_2x_3$  gives  $\overline{x}_1$ .

Expansion of  $\overline{x}_1x_2\overline{x}_3$  gives  $\overline{x}_1$ .

Expansion of  $x_1x_2x_3$  gives  $x_2x_3$ .

The set of prime implicants comprises  $\overline{x}_1$  and  $x_2x_3$ .

4.31. Expansion of  $\overline{x}_1x_2\overline{x}_3x_4$  gives  $x_2\overline{x}_3x_4$  and  $\overline{x}_1x_2x_4$ .

Expansion of  $x_1x_2\overline{x}_3x_4$  gives  $x_2\overline{x}_3x_4$ .

Expansion of  $x_1x_2x_3\overline{x}_4$  gives  $x_3\overline{x}_4$ .

Expansion of  $\overline{x}_1x_2x_3$  gives  $\overline{x}_1x_3$ .

Expansion of  $\overline{x}_2x_3$  gives  $\overline{x}_2x_3$ .

The set of prime implicants comprises  $x_2\overline{x}_3x_4$ ,  $\overline{x}_1x_2x_4$ ,  $x_3\overline{x}_4$ ,  $\overline{x}_1x_3$ , and  $\overline{x}_2x_3$ .

4.32. Representing both functions in the form of Karnaugh map, it is easy to show that  $f = g$ . The minimum cost SOP expression is

$$f = g = \overline{x}_2\overline{x}_3\overline{x}_5 + \overline{x}_2x_3\overline{x}_4 + x_1x_3x_4 + x_1x_2x_4x_5.$$



- 4.33. Representing both functions in the form of Karnaugh map, it is easy to show that  $f = g$ . The minimum cost SOP expression is

$$f = g = x_2x_4 + x_1\bar{x}_2\bar{x}_4 + \bar{x}_1x_2x_3 + \bar{x}_2\bar{x}_3\bar{x}_4.$$

- 4.34. Representing both functions in the form of Karnaugh map, it is easy to show that  $f$  and  $g$  do not represent the same function. In particular:  $f(1, 1, 0, 1, 0) = 1$  while  $g(1, 1, 0, 1, 0) = 0$  and  $f(1, 1, 1, 1, 1) = 0$  while  $g(1, 1, 1, 1, 1) = 1$ .

- 4.35. Implementing the circuit as

$$f = \bar{x}_2x_3\bar{x}_4 + x_1x_2x_3x_4 + \bar{x}_2\bar{x}_3$$

$$g = \bar{x}_2x_3\bar{x}_4 + x_1x_2x_3x_4 + x_1\bar{x}_3x_4 + \bar{x}_1x_3x_4$$

there are 7 gates and 22 inputs for a cost of 29.

- 4.36. Implementing the circuit as

$$f = (x_1 + \bar{x}_2 + x_3 + x_4)(x_1 + \bar{x}_2 + \bar{x}_3 + \bar{x}_4)(x_1 + x_3 + \bar{x}_4)(\bar{x}_1 + x_2 + \bar{x}_3)(\bar{x}_1 + x_2 + x_4)$$

$$g = (x_1 + \bar{x}_2 + x_3 + x_4)(x_1 + \bar{x}_2 + \bar{x}_3 + \bar{x}_4)(\bar{x}_1 + \bar{x}_2 + \bar{x}_4)(\bar{x}_1 + x_2 + \bar{x}_3)$$

there are 9 gates and 32 inputs for a cost of 41.

- 4.37. Assuming that the condition where all sensors produce the output of 0 is a don't care, the complement of the desired function is

$$\bar{f} = \bar{x}_1\bar{x}_2\bar{x}_3 + \bar{x}_1\bar{x}_2\bar{x}_4 + \bar{x}_1\bar{x}_3\bar{x}_4 + \bar{x}_2\bar{x}_3\bar{x}_4$$

Then,  $f = \bar{\bar{f}}$ .

- 4.38. Assuming that the condition where all sensors produce the output of 0 is a don't care, the complement of the desired function is

$$\begin{aligned} \bar{f} = & \bar{x}_1\bar{x}_2\bar{x}_3\bar{x}_4\bar{x}_5\bar{x}_6 + \bar{x}_1\bar{x}_2\bar{x}_3\bar{x}_4\bar{x}_5\bar{x}_7 + \bar{x}_1\bar{x}_2\bar{x}_3\bar{x}_4\bar{x}_6\bar{x}_7 + \bar{x}_1\bar{x}_2\bar{x}_3\bar{x}_5\bar{x}_6\bar{x}_7 + \\ & \bar{x}_1\bar{x}_2\bar{x}_4\bar{x}_5\bar{x}_6\bar{x}_7 + \bar{x}_1\bar{x}_3\bar{x}_4\bar{x}_5\bar{x}_6\bar{x}_7 + \bar{x}_2\bar{x}_3\bar{x}_4\bar{x}_5\bar{x}_6\bar{x}_7 \end{aligned}$$

Then,  $f = \bar{\bar{f}}$ .

- 4.39. Implement first the complement of  $f$  as

$$\begin{aligned} \bar{f} &= x_1x_3 + x_2x_4 \\ &= (x_1 \uparrow x_3) \uparrow (x_2 \uparrow x_4) \end{aligned}$$

Then  $f = \bar{f} \uparrow \bar{f}$ .

4.40. Implement first the complement of  $f$  as

$$\begin{aligned}\bar{f} &= \bar{x}_1\bar{x}_3 + x_2x_4 + x_1x_3 \\ &= (\bar{x}_1\bar{x}_3 + x_2x_4) + (x_1x_3 + x_1x_3) \\ &= ((\bar{x}_1 \uparrow \bar{x}_3) \uparrow (x_2 \uparrow x_4)) \uparrow ((x_1 \uparrow x_3) \uparrow (x_1 \uparrow x_3))\end{aligned}$$

Then  $f = \bar{f} \uparrow \bar{f}$ .

4.41. Implement first the complement of  $f$  as

$$\begin{aligned}\bar{f} &= (\bar{x}_1 + x_4)(\bar{x}_2 + \bar{x}_3) \\ &= (\bar{x}_1 \downarrow x_4) \downarrow (\bar{x}_2 \downarrow \bar{x}_3)\end{aligned}$$

Then  $f = \bar{f} \downarrow \bar{f}$ .

4.42. Implement first the complement of  $f$  as

$$\begin{aligned}\bar{f} &= (\bar{x}_1 + \bar{x}_4)(\bar{x}_2 + x_3)(x_2 + \bar{x}_3) \\ &= ((\bar{x}_1 + \bar{x}_4)(\bar{x}_2 + x_3))((x_2 + \bar{x}_3)(x_2 + \bar{x}_3)) \\ &= ((\bar{x}_1 \downarrow \bar{x}_4) \downarrow (\bar{x}_2 \downarrow x_3)) \downarrow ((x_2 \downarrow \bar{x}_3) \downarrow (x_2 \downarrow \bar{x}_3))\end{aligned}$$

Then  $f = \bar{f} \downarrow \bar{f}$ .

4.43. The cost of the circuit in Figure P4.2 is 11 gates and 30 inputs, for a total of 41. The functions implemented by the circuit can also be realized as

$$\begin{aligned}f &= \bar{x}_1\bar{x}_2\bar{x}_4 + x_2\bar{x}_3\bar{x}_4 + \bar{x}_1x_3x_4 + x_1x_4 \\ g &= \bar{x}_1\bar{x}_2\bar{x}_4 + x_2\bar{x}_3\bar{x}_4 + \bar{x}_1x_3x_4 + \bar{x}_2x_4 + x_3\bar{x}_4\end{aligned}$$

The first three product terms in  $f$  and  $g$  are the same; therefore, they can be shared. Then, the cost of implementing  $f$  and  $g$  is 8 gates and 24 inputs, for a total of 32.

4.44. The cost of the circuit in Figure P4.3 is 11 gates and 26 inputs, for a total of 37. The functions implemented by the circuit can also be realized as

$$\begin{aligned}f &= (\bar{x}_2 \uparrow x_4) \uparrow (\bar{x}_1 \uparrow x_2 \uparrow x_3) \uparrow (x_1 \uparrow \bar{x}_2 \uparrow x_3) \uparrow (\bar{x}_2 \uparrow \bar{x}_3) \\ g &= (\bar{x}_2 \uparrow x_4) \uparrow (\bar{x}_1 \uparrow x_2 \uparrow x_3) \uparrow (x_1 \uparrow \bar{x}_2 \uparrow x_3) \uparrow (\bar{x}_1 \uparrow \bar{x}_1)\end{aligned}$$

The first three NAND terms in  $f$  and  $g$  are the same; therefore, they can be shared. Then, the cost of implementing  $f$  and  $g$  is 7 gates and 20 inputs, for a total of 27.

4.45. Using the logic expressions derived in Example 4.7, the circuit in Figure 4.25*b* can be specified as

```
LIBRARY ieee ;
USE ieee.std_logic_1164.all ;

ENTITY prob35 IS
    PORT ( x1, x2, x3, x4 : IN    STD_LOGIC ;
           f              : OUT  STD_LOGIC ) ;
END prob35 ;

ARCHITECTURE LogicFunc OF prob35 IS
    SIGNAL g, k : STD_LOGIC ;
BEGIN
    g <= x1 OR x2 OR x3 ;
    k <= (NOT x3 AND x4) OR (x3 AND NOT x4) ;
    f <= (k AND g) OR (NOT k AND NOT g) ;
END LogicFunc ;
```

4.46. The circuit in Figure 4.27*c* can be specified as follows:

```
LIBRARY ieee ;
USE ieee.std_logic_1164.all ;

ENTITY prob36 IS
    PORT ( x1, x2, x3, x4, x5, x6, x7 : IN    STD_LOGIC ;
           f                          : OUT  STD_LOGIC ) ;
END prob36 ;

ARCHITECTURE LogicFunc OF prob36 IS
    SIGNAL g, h, k : STD_LOGIC ;
BEGIN
    g <= (x1 NAND x1) NAND (x2 NAND x3) ;
    h <= (x4 NAND ((x5 NAND x5) NAND (x6 NAND x6))) NAND (x7 NAND x7) ;
    k <= g NAND h ;
    f <= k NAND k ;
END LogicFunc ;
```

4.47. The circuit in Figure 4.28b can be specified as follows:

```

LIBRARY ieee ;
USE ieee.std_logic_1164.all ;

ENTITY prob37 IS
    PORT ( x1, x2, x3, x4, x5, x6, x7 : IN    STD_LOGIC ;
          f                               : OUT STD_LOGIC ) ;
END prob37 ;

ARCHITECTURE LogicFunc OF prob37 IS
    SIGNAL g, h : STD_LOGIC ;
BEGIN
    g <= x1 NOR ((x2 NOR x2) NOR (x3 NOR x3)) ;
    h <= ((x4 NOR x4) NOR (x5 NOR x6)) NOR x7 ;
    f <= g NOR h ;
END LogicFunc ;

```

4.48. Using the POS expression

$$f = (x_1 + x_2 + \bar{x}_3 + \bar{x}_4)(x_1 + \bar{x}_2 + \bar{x}_3 + x_4)(\bar{x}_1 + x_2 + \bar{x}_3 + x_4)(\bar{x}_1 + \bar{x}_2 + x_3 + \bar{x}_4)$$

the function can be implemented using the code

```

LIBRARY ieee ;
USE ieee.std_logic_1164.all ;

ENTITY prob38 IS
    PORT ( x1, x2, x3, x4 : IN    STD_LOGIC ;
          f               : OUT STD_LOGIC ) ;
END prob38 ;

ARCHITECTURE LogicFunc OF prob38 IS
BEGIN
    f <= (x1 OR x2 OR NOT x3 OR NOT x4) AND (x1 OR NOT x2 OR NOT x3 OR x4) AND
        (NOT x1 OR x2 OR NOT x3 OR x4) AND (NOT x1 OR NOT x2 OR x3 OR NOT x4) ;
END LogicFunc ;

```

4.49. The simplest expression is

$$f = \overline{x_1}\overline{x_3} + x_2x_3(x_1 + x_4)$$

which can be implemented using the code

```

LIBRARY ieee ;
USE ieee.std_logic_1164.all ;

ENTITY prob39 IS
    PORT ( x1, x2, x3, x4 : IN    STD_LOGIC ;
          f               : OUT  STD_LOGIC ) ;
END prob39 ;

ARCHITECTURE LogicFunc OF prob39 IS
BEGIN
    f <= (NOT x1 AND NOT x3) OR ((x2 AND x3) AND (x1 OR x4)) ;
END LogicFunc ;

```

4.50. The simplest expression is

$$f = (\overline{x_1} + x_3)(x_1 + \overline{x_2} + \overline{x_3} + x_4)$$

which can be implemented using the code

```

LIBRARY ieee ;
USE ieee.std_logic_1164.all ;

ENTITY prob40 IS
    PORT ( x1, x2, x3, x4 : IN    STD_LOGIC ;
          f               : OUT  STD_LOGIC ) ;
END prob40 ;

ARCHITECTURE LogicFunc OF prob40 IS
BEGIN
    f <= (NOT x1 OR x3) AND (x1 OR NOT x2 OR NOT x3 OR x4) ;
END LogicFunc ;

```

4.51. The simplest expression is

$$f = (x_2 + \overline{x_3})(\overline{x_1} + \overline{x_3} + x_4)$$

which can be implemented using the code

```

LIBRARY ieee ;
USE ieee.std_logic_1164.all ;

ENTITY prob41 IS
    PORT ( x1, x2, x3, x4 : IN    STD_LOGIC ;
          f               : OUT  STD_LOGIC ) ;
END prob41 ;

ARCHITECTURE LogicFunc OF prob41 IS
BEGIN
    f <= (x2 OR NOT x3) AND (NOT x1 OR NOT x3 OR x4) ;
END LogicFunc ;

```