

NCKU Programming Contest Training Course

Math

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Lin Yun Wen

Department of Computer Science and Information Engineering
National Cheng Kung University
Tainan, Taiwan



Outline

Prime Numbers

Big Number

GCD, Extended Euclid's Algorithm



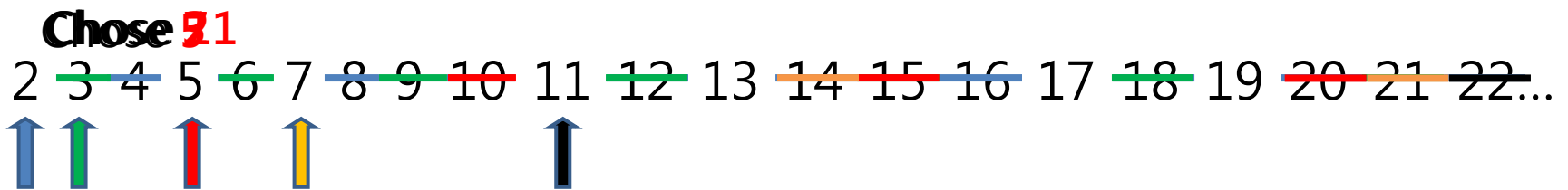
Prime Number

- Sieve of Eratosthenes (埃拉托斯特尼篩法)
 - 由小到大選擇質數，並刪除其倍數
- $6n \pm 1$ Method
 - 拿 2 和 3 這兩個質數先篩過一遍，剩下的數字則用除法驗證是不是質數。



Prime Number

- We use **sieve** to create a prime array
 - Chose the smallest number at each iteration and delete the multiple of this number



Prime Number

- Sieve of Eratosthenes (埃拉托斯特尼篩法)
 - 由小到大選擇質數，並刪除其倍數

```
1  #include <cmath>
2  #include <cstring>
3  #define MAX 10000000
4  bool is_prime[MAX];
5  void eratosthenes()
6  {
7      memset(is_prime, 1, sizeof(is_prime));
8      is_prime[0] = false;
9      is_prime[1] = false;
10
11     for (int i = 2; i <= sqrt(MAX); ++i)
12         if (is_prime[i])
13             for(int j = i+i; j < MAX; j += i)
14                 is_prime[j] = false;
15 }
```



Prime Number

- $6n \pm 1$ Method
 - 2 和 3 的最小公倍數是 6，把所有數字分為 $6n$ 、 $6n+1$ 、 $6n+2$ 、 $6n+3$ 、 $6n+4$ 、 $6n+5$ 六種，可以看出 $6n$ 、 $6n+2$ 、 $6n+3$ 、 $6n+4$ 會是 2 或 3 的倍數，不屬於質數。因此，只要驗證 $6n+1$ 和 $6n+5 (= 6n-1)$ 是不是質數就可以了。



Prime Number

- $6n \pm 1$ Method

```
1  #include <vector>
2  #define MAX 10000000
3  vector<int> prime;
4  bool is_prime(int n) {
5      for (int i = 0; prime[i]*prime[i] <= n; ++i)
6          if (n % prime[i] == 0)
7              return false;
8      return true;
9  }
10 void make_prime() {
11     prime.push_back(2);
12     prime.push_back(3);
13     for (int i = 5, gap = 2; i < MAX; i+=gap, gap = 6 - gap)
14         if (is_prime(i))
15             prime.push_back(i);
16 }
```



Prime Number

- 方法二比方法一慢，但較省空間
- But just remember that the code in previous page is fast enough to solve almost every prime problems
- 其他方法:
 - [演算法 筆記 - Prime](#)



Practice - 1

UVa 10392 - Factoring Large Numbers



Outline

Prime Numbers

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Big Number

- Array
- 習慣上將低位數放在index比較小的位置
 - Ex: 680468975231245

0	1	2	3	4	5	6	7	8	9	...									
5	4	2	1	3	2	5	7	9	8	6	4	0	8	6	_	_	_	---	

- 右方補0



Big Number

- 加法：位數各自相加後，由低至高位依序進位
- 減法：位數各自相減後，由低至高位依序借位
- 乘法：直式乘法
- 除法：長除法



Big Number

- 加法：

```
1 void add(int a[100], int b[100], int c[100]) {  
2     for (int i = 0; i < 100; ++i)  
3         c[i] = a[i] + b[i];  
4  
5     for (int i = 0; i < 100-1; ++i) {  
6         c[i+1] += c[i] / 10;  
7         c[i] %= 10;  
8     }  
9 }
```



Practice - 2

UVa 10106 - Product

Problem Description

The problem is to multiply two integers X, Y .($0 \leq X, Y < 10250$)

Input

The input will consist of a set of pairs of lines. Each line in pair contains one multiplier.

Output

For each input pair of lines the output line should consist one integer the product.



Outline

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Greatest Common Divisor

- 輾轉相除法 (Euclidean Algorithm)

```
11  int gcd(int a, int b) {  
12      ... if (a == 0)  
13      ...     return b;  
14      ... return gcd(b % a, a);  
15  }
```



- $\text{gcd}(462, 1071)$
 - $\text{gcd}(147, 462)$
- $\text{gcd}(21, 147)$
 - $\text{gcd}(0, 7)$

```
11  int gcd(int a, int b) {  
12      ... if (a == 0)  
13      ...  
14      ... r  
15  }
```

- 從1071中不斷減去462直到小於462（可以減2次，即商 $q_0 = 2$ ），餘數是147：
 - $1071 = 2 \times 462 + 147$.
- 然後從462中不斷減去147直到小於147（可以減3次，即 $q_1 = 3$ ），餘數是21：
 - $462 = 3 \times 147 + 21$.
- 再從147中不斷減去21直到小於21（可以減7次，即 $q_2 = 7$ ），沒有餘數：
 - $147 = 7 \times 21 + 0$.
- 此時，餘數是0，所以1071和462的最大公因數是21，



Practice - 3

UVa 408 – Uniform Generator



Extended Euclidean Algorithm

- 找到 $aX + bY = \gcd(a, b)$ 的整數解 X, Y
- Ex (from wiki)
 - $47x + 30y = 1$



Extended Euclidean Algorithm

- $47 = 30 * 1 + 17$
- $30 = 17 * 1 + 13$
- $17 = 13 * 1 + 4$
- $13 = 4 * 3 + 1$
- $4 = 1 * 4 + 0$

gcd(30, 47)
gcd(17, 30)
gcd(13, 17)
gcd(4, 13)
gcd(1, 4)
gcd(0, 1)



Extended Euclidean Algorithm

- $47 = 30 * 1 + 17$
- $30 = 17 * 1 + 13$
- $17 = 13 * 1 + 4$
- $13 = 4 * 3 + 1$
- $4 = 1 * 4 + 0$
- $17 = 47 * 1 + 30 * (-1)$
- $13 = 30 * 1 + 17 * (-1)$
- $4 = 17 * 1 + 13 * (-1)$
- $1 = 13 * 1 + 4 * (-3)$

$$47x + 30y = 1$$



Extended Euclidean Algorithm

- $1 = 13 * 1 + 4 * (-3)$
- $1 = 13 * 1 + [17 * 1 + 13 * (-1)] * (-3)$
- $1 = 17 * (-3) + 13 * 4$
- $1 = 17 * (-3) + [30 * 1 + 17 * (-1)] * 4$
- $1 = 30 * 4 + 17 * (-7)$
- $1 = 30 * 4 + [47 * 1 + 30 * (-1)] * (-7)$
- $1 = 47 * (-7) + 30 * 11$



Extended Euclidean Algorithm

- $\gcd(a, b) = \gcd(b, a \% b)$
- $aX + bY = \gcd(a, b) = \gcd(b, a \% b) = bX' + (a \% b)Y'$
- $aX + bY = bX' + [a - (a/b)b]Y' = aY' + b(X' - (a/b)Y')$
 - $X = Y'$
 - $Y = X' - (a/b)Y'$



Extended Euclidean Algorithm

```
17  int exGCD(int a, int b, int &X, int &Y) {
18      if (b == 0) {
19          X = 1;
20          Y = 0;
21          return a;
22      } else {
23          int gcd = exGCD(b, a % b, X, Y);
24          int tmp = X;
25          X = Y;
26          Y = tmp - (a/b)*Y;
27          return gcd;
28      }
29 }
```



Practice - 4

UVa 10104 - Euclid Problem



Extended Euclidean Algorithm

- $\frac{m!}{n!} \% P$ (P 是一個很大的質數)



Extended Euclidean Algorithm

- $aX + bY = \gcd(a, b)$
 - $a = n!$
 - $b = p$
 - $\gcd(a, b) = 1$

方程式 $ax+by=1$ 有整數解
iff 整數 a 和 b 互質



Extended Euclidean Algorithm

- $n!X + pY = 1$ (use Extended Euclidean Algorithm get (X, Y))
- $n!X + pY = 1 \rightarrow \text{mod } p$
- $(n!X) \% p = 1 \text{ --- (1)}$
- $\frac{m!}{n!} \% p = ans \text{ --- (2)}$
- $(1) * (2)$
 $\rightarrow \left(\frac{m!}{n!} \times n!X \right) \% p = ans \rightarrow (m! \times X) \% p = ans$



Practice - 5

Facebook Hacker Cup 2017 Round I

Beach Umbrellas



Epsilon ϵ

- Float :
 - 數值範圍 : $-3.4\text{e-}38 \sim 3.4\text{e}38$
 - 十位數精確度位數 : 6~7
- Double :
 - 數值範圍 : $-1.7\text{e}308 \sim 1.7\text{e}308$
 - 十位數精確度位數 : 14~15



Epsilon ϵ

- Example

```
1  #include <stdio>
2  #include <cmath>
3
4  int main() {
5      double a = asin(sqrt(2.0) / 2) * 4.0;
6      double b = acos(-1.0);
7
8      printf("a = %.20lf\n", a);
9      printf("b = %.20lf\n", b);
10     printf("a-b = %.20lf\n", a - b);
11     printf("a == b? %s\n", a == b ? "True" : "False");
12 }
```



Epsilon ϵ

- Result

```
linyunwen@Lin-Yun-Wens-MacBook-Air ~/D/L/c/ACM> ./sample_epsilon  
a = 3.14159265358979356009  
b = 3.14159265358979311600  
a-b = 0.00000000000000044409  
a == b? False
```



Epsilon ϵ

- 引入 **eps** 判斷浮點數是否相等
 - $\text{eps} = 1\text{e-}8$

整數	浮點數
$a == b$	$ a - b < \text{eps}$
$a != b$	$ a - b > \text{eps}$
$a < b$	$a - b < -\text{eps}$
$a > b$	$a - b > \text{eps}$



Practice - 6

UVa 906 – Rational Neighbor

