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L Basic

1.1 vimrc

```
=== .vimrc ===

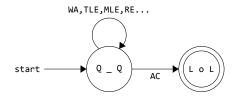
set et nu cin ls=2 ts=4 sw=4 sts=4 ttm=100
syntax on

nn <F4> :w ! cat -n \| lpr <CR>
nn <F7> :w <bar> :!vim %<_.in<left><left><left><left>
nn <F8> :w <bar> :!g++ % -o %< -std=c++11
\ -fsanitize=undefined -Wall -Wextra -Wshadow -DFOX &&
\ for i in %<_*.in; do echo == && ./%< < $i; done <CR>
nn <F9> :w <bar> :!g++ % -o %< -std=c++11
\ -fsanitize=undefined -Wall -Wextra -Wshadow -DFOX &&
\ echo == && ./%<
```

1.2 default code

```
#pragma GCC optimize("Ofast")
#include <bits/stdc++.h>
#include <sys/time.h>
#include <sys/resource.h>
using namespace std;
void setstack(){
    // Set soft limit and hard limit to max
    const rlimit tmp {RLIM_INFINITY,RLIM_INFINITY};
    setrlimit(RLIMIT_STACK,&tmp);
}
int main(){
    #define name ""
    #ifndef FOX
    freopen(name".in","r",stdin);
freopen(name".out","w",stdout);
    #endif
    static_assert(strlen(name));
    ios::sync_with_stdio(0);
    cin.tie(0), cout.tie(0);
```

1.3 state machine





2 Flow

2.1 Dinic

```
(a) Bounded Maxflow Construction:
1. add two node ss, tt
2. add_edge(ss, tt, INF)
3. for each edge u -> v with capacity [1, r]:
        add_edge(u, tt, 1)
        add_edge(ss, v, 1)
        add_edge(u, v, r-1)
4. see (b), check if it is possible.
5. answer is maxflow(ss, tt) + maxflow(s, t)
(b) Bounded Possible Flow:
1. same construction method as (a)
run maxflow(ss, tt)
3. for every edge connected with ss or tt:
        rule: check if their rest flow is exactly 0
4. answer is possible if every edge do satisfy the rule
5. otherwise, it is NOT possible.
(c) Bounded Minimum Flow:
1. same construction method as (a)
2. answer is maxflow(ss, tt)
(d) Bounded Minimum Cost Flow:
* the concept is somewhat like bounded possible flow.
1. same construction method as (a)
2. answer is maxflow(ss, tt) + (\sum 1 * cost for every)
   edge)
______
(e) Minimum Cut:

    run maxflow(s, t)

2. run cut(s)
3. ss[i] = 1: node i is at the same side with s.
const long long INF = 1LL<<60;</pre>
struct Dinic { //O(VVE), with minimum cut
    static const int MAXN = 5003;
    struct Edge{
        int u, v;
        long long cap, rest;
    };
    int n, m, s, t, d[MAXN], cur[MAXN];
    vector<Edge> edges;
    vector<int> G[MAXN];
    void init(){
        edges.clear();
        for ( int i = 0 ; i < MAXN ; i++ ) G[i].clear()</pre>
    // min cut start
    bool side[MAXN];
    void cut(int u) {
        side[u] = 1;
        for ( int i : G[u] ) {
            if ( !side[ edges[i].v ] && edges[i].rest )
                 cut(edges[i].v);
        }
    // min cut end
    void add_edge(int u, int v, long long cap){
  edges.push_back( {u, v, cap, cap} );
        edges.push_back( {v, u, 0, 0LL} );
        m = edges.size();
        G[u].push_back(m-2);
        G[v].push_back(m-1);
    }
```

```
bool bfs(){
         memset(d, -1, sizeof(d));
         queue<int> que;
         que.push(s); d[s]=0;
         while (!que.empty()){
             int u = que.front(); que.pop();
             for (int ei : G[u]){
                 Edge &e = edges[ei];
                 if (d[e.v] < 0 && e.rest > 0){
                     d[e.v] = d[u] + 1;
                     que.push(e.v);
                 }
             }
         return d[t] >= 0;
    long long dfs(int u, long long a){
         if ( u == t || a == 0 ) return a;
         long long flow = 0, f;
         for ( int &i=cur[u]; i < (int)G[u].size(); i++</pre>
             Edge &e = edges[ G[u][i] ];
             if ( d[u] + 1 != d[e.v] ) continue;
             f = dfs(e.v, min(a, e.rest));
             if (f > 0) {
                 e.rest -= f;
                 edges[ G[u][i]^1 ].rest += f;
                 flow += f;
                 a -= f;
                 if ( a == 0 )break;
         return flow;
    long long maxflow(int _s, int _t){
         s = _s, t = _t;
         long long flow = 0, mf;
         while ( bfs() ){
             memset(cur, 0, sizeof(cur));
             while ( (mf = dfs(s, INF)) ) flow += mf;
         return flow;
    }
} dinic;
```

2.2 GomoryHu tree

```
Construct of Gomory Hu Tree
1. make sure the whole graph is clear
2. set node 0 as root, also be the parent of other
   nodes.
3. for every node i > 0, we run maxflow from i to
    parent[i]
4. hense we know the weight between i and parent[i]
5. for each node j > i, if j is at the same side with i
  make the parent of j as i
int e[MAXN][MAXN];
int p[MAXN];
Dinic D; // original graph
void gomory_hu() {
   fill(p, p+n, 0);
    fill(e[0], e[n], INF);
    for ( int s = 1 ; s < n ; s++ ) {</pre>
        int t = p[s];
        Dinic F = D;
        int tmp = F.max_flow(s, t);
        for ( int i = 1 ; i < s ; i++ )
```

```
e[s][i] = e[i][s] = min(tmp, e[t][i]);

for ( int i = s+1 ; i <= n ; i++ )
        if ( p[i] == t && F.side[i] ) p[i] = s;
}
}</pre>
```

2.3 min cost flow

```
// long long version
typedef pair<long long, long long> pll;
struct CostFlow {
    static const int MAXN = 350;
    static const long long INF = 1LL<<60;</pre>
    struct Edge {
        int to, r;
        long long rest, c;
    };
    int n, pre[MAXN], preL[MAXN]; bool inq[MAXN];
    long long dis[MAXN], fl, cost;
    vector<Edge> G[MAXN];
    void init() {
        for ( int i = 0 ; i < MAXN ; i++) G[i].clear();</pre>
    void add_edge(int u, int v, long long rest, long
        long c) {
        G[u].push_back({v, (int)G[v].size(), rest, c});
        G[v].push back({u, (int)G[u].size()-1, 0, -c});
    pll flow(int s, int t) {
        fl = cost = 0;
        while (true) {
             fill(dis, dis+MAXN, INF);
             fill(inq, inq+MAXN, 0);
            dis[s] = 0;
             queue<int> que;
             que.push(s);
             while ( !que.empty() ) {
                int u = que.front(); que.pop();
                inq[u] = 0;
                for ( int i = 0 ; i < (int)G[u].size()</pre>
                     ; i++) {
                     int v = G[u][i].to;
                     long long w = G[u][i].c;
                     if ( G[u][i].rest > 0 && dis[v] >
                         dis[u] + w) {
                         pre[v] = u; preL[v] = i;
                         dis[v] = dis[u] + w;
                         if (!inq[v]) {
                             inq[v] = 1;
                             que.push(v);
                         }
                     }
                }
            }
            if (dis[t] == INF) break;
            long long tf = INF;
            for (int v = t, u, 1; v != s; v = u) {
                u = pre[v]; l = preL[v];
                tf = min(tf, G[u][1].rest);
            for (int v = t, u, 1; v != s; v = u) {
                u = pre[v]; l = preL[v];
                G[u][1].rest -= tf;
                G[v][G[u][1].r].rest += tf;
            cost += tf * dis[t];
            fl += tf;
        return {fl, cost};
} flow;
```

2.4 SW mincut 全點對最小割

```
// all pair min cut
// global min cut
struct SW{ // O(V^3)
  static const int MXN = 514;
  int n,vst[MXN],del[MXN];
  int edge[MXN][MXN],wei[MXN];
  void init(int _n){
    n = _n; FZ(edge); FZ(del);
  void addEdge(int u, int v, int w){
    edge[u][v] += w; edge[v][u] += w;
  void search(int &s, int &t){
    FZ(vst); FZ(wei);
    s = t = -1;
    while (true){
      int mx=-1, cur=0;
       for (int i=0; i<n; i++)</pre>
         if (!del[i] && !vst[i] && mx<wei[i])</pre>
          cur = i, mx = wei[i];
      if (mx == -1) break;
      vst[cur] = 1;
       s = t; t = cur;
      for (int i=0; i<n; i++)</pre>
         if (!vst[i] && !del[i]) wei[i] += edge[cur][i];
    }
  int solve(){
    int res = 2147483647;
    for (int i=0,x,y; i<n-1; i++){</pre>
      search(x,y);
      res = min(res,wei[y]);
       del[y] = 1;
       for (int j=0; j<n; j++)</pre>
         edge[x][j] = (edge[j][x] += edge[y][j]);
    return res;
  }
}graph;
```

3 Matching

3.1 Hungarian

```
// Maximum Cardinality Bipartite Matching
// Worst case O(nm)
struct Graph{
    static const int MAXN = 5003;
    vector<int> G[MAXN];
    int n, match[MAXN], vis[MAXN];
    void init(int _n){
        n = n;
        for (int i=0; i<n; i++) G[i].clear();</pre>
    bool dfs(int u){
        for (int v:G[u]){
             if (vis[v]) continue;
            vis[v]=true;
            if (match[v]==-1 || dfs(match[v])){
                 match[v] = u;
                 match[u] = v;
                 return true;
            }
        }
        return false;
    }
    int solve(){
        int res = 0:
```

for (int i = 0 ; i < (int)G[u].size(); i++) {</pre>

nr = max(nl , nr);

```
memset(match,-1,sizeof(match));
                                                                  while(m--) {
                                                                       int x , y;
        for (int i=0; i<n; i++){</pre>
             if (match[i]==-1){
                                                                       11 w;
                 memset(vis,0,sizeof(vis));
                                                                       RI(x, y, w);
                 if ( dfs(i) ) res++;
                                                                       W[x][y] = w;
                                                                       lx[x] = max(lx[x], w);
        }
        return res;
                                                                  REP1(i , 1 , nl) {
                                                                       REP1(x , 1 , nl) vx[x] = 0;
                                                                       REP1(y , 1 , nr) \ vy[y] = 0 , slack[y] = INF64;
} graph;
                                                                       match(i);
                                                                  11 \text{ ans} = 0LL;
3.2 KM
                                                                  REP1(x , 1 , nl) ans += W[x][mx[x]];
                                                                  PL(ans);
const int MAX_N = 400 + 10;
                                                                  \label{eq:rep1} \texttt{REP1}(\texttt{x , 1 , nl}) \ \texttt{printf}(\text{"$d\%c$",$W[x][mx[x]] ? mx[x]}
const 11 INF64 = 0x3f3f3f3f3f3f3f3f1L;
                                                                       : 0," \n"[x == nl]);
int nl , nr;
                                                                  return 0;
int pre[MAX_N];
                                                              }
11 slack[MAX_N];
11 W[MAX_N][MAX_N];
11 1x[MAX_N], 1y[MAX_N];
int mx[MAX_N], my[MAX_N];
bool vx[MAX_N], vy[MAX_N];
                                                              3.3 Matching.txt
void augment(int u) {
                                                             |最大匹配 + 最小邊覆蓋 = V
    if(!u) return;
                                                              最大獨立集 + 最小點覆蓋 = V
    augment(mx[pre[u]]);
                                                              最大匹配 = 最小點覆蓋
    mx[pre[u]] = u;
                                                              最小路徑覆蓋數 = V - 最大匹配數
    my[u] = pre[u];
inline void match(int x) {
    queue<int> que;
                                                                     Maximum General Matching
    que.push(x);
    while(1) {
        while(!que.empty()) {
                                                             // Maximum Cardinality Matching
             x = que.front();
                                                              struct Graph {
             que.pop();
             vx[x] = 1;
                                                                vector<int> G[MAXN];
                                                                int pa[MAXN], match[MAXN], st[MAXN], S[MAXN], vis[
             REP1(y , 1 , nr) {
                 if(vy[y]) continue;
                 11 t = 1x[x] + 1y[y] - W[x][y];
                                                                int t, n;
                 if(t > 0) {
                     if(slack[y] >= t) slack[y] = t ,
                                                                void init(int _n) {
                         pre[y] = x;
                                                                  n = n;
                     continue;
                                                                  for ( int i = 1 ; i <= n ; i++ ) G[i].clear();</pre>
                 }
                 pre[y] = x;
                                                                void add edge(int u, int v) {
                 if(!my[y])
                                                                  G[u].push_back(v);
                     augment(y);
                                                                  G[v].push_back(u);
                     return;
                                                                int lca(int u, int v){
                 vy[y] = 1;
                                                                  for ( ++t ; ; swap(u, v) ) {
                                                                    if ( u == 0 ) continue;
                 que.push(my[y]);
            }
                                                                    if ( vis[u] == t ) return u;
                                                                    vis[u] = t;
        11 t = INF64;
                                                                    u = st[ pa[ match[u] ] ];
        REP1(y, 1, nr) if(!vy[y]) t = min(t, slack[y])
        REP1(x , 1 , nl) if(vx[x]) lx[x] -= t;
                                                                void flower(int u, int v, int l, queue<int> &q) {
        REP1(y , 1 , nr) {
                                                                  while ( st[u] != 1 ) {
             if(vy[y]) ly[y] += t;
                                                                    pa[u] = v;
                                                                     if ( S[ v = match[u] ] == 1 ) {
             else slack[y] -= t;
                                                                      q.push(v);
        REP1(y , 1 , nr) {
                                                                      S[v] = 0;
             if(vy[y] || slack[y]) continue;
             if(!my[y]) {
                                                                    st[u] = st[v] = 1;
                 augment(y);
                                                                    u = pa[v];
                 return;
                                                                  }
                                                                bool bfs(int u){
             vy[y] = 1;
                                                                  for ( int i = 1 ; i <= n ; i++ ) st[i] = i;
memset(S, -1, sizeof(S));</pre>
             que.push(my[y]);
        }
                                                                  queue<int>q;
    }
                                                                  q.push(u);
int main() {
                                                                  S[u] = 0;
    int m;
                                                                  while ( !q.empty() ) {
    RI(nl , nr , m);
                                                                    u = q.front(); q.pop();
```

```
int v = G[u][i];
        if ( S[v] == -1 ) {
          pa[v] = u;
          S[v] = 1;
          if ( !match[v] ) {
            for ( int lst ; u ; v = lst, u = pa[v] ) {
              lst = match[u];
              match[u] = v;
              match[v] = u;
            }
            return 1;
          q.push(match[v]);
          S[ match[v] ] = 0;
        } else if ( !S[v] && st[v] != st[u] ) {
          int l = lca(st[v], st[u]);
           flower(v, u, 1, q);
          flower(u, v, l, q);
        }
      }
    }
    return 0;
  int solve(){
    memset(pa, 0, sizeof(pa));
    memset(match, 0, sizeof(match));
    int ans = 0;
    for ( int i = 1 ; i <= n ; i++ )</pre>
      if ( !match[i] && bfs(i) ) ans++;
    return ans;
  }
} graph;
```

3.5 Minimum General Weighted Matching

```
// Minimum Weight Perfect Matching (Perfect Match)
struct Graph {
    static const int MAXN = 105;
    int n, e[MAXN][MAXN];
    int match[MAXN], d[MAXN], onstk[MAXN];
    vector<int> stk;
    void init(int _n) {
        n = _n;
        for( int i = 0 ; i < n ; i ++ )</pre>
             for( int j = 0 ; j < n ; j ++ )</pre>
                 e[i][j] = 0;
    void add_edge(int u, int v, int w) {
        e[u][v] = e[v][u] = w;
    bool SPFA(int u){
        if (onstk[u]) return true;
        stk.push_back(u);
        onstk[u] = 1;
        for ( int v = 0 ; v < n ; v++ ) {
             if (u != v && match[u] != v && !onstk[v] )
                 int m = match[v];
                 if ( d[m] > d[u] - e[v][m] + e[u][v] )
                     d[m] = d[u] - e[v][m] + e[u][v];
                     onstk[v] = 1;
                     stk.push_back(v);
                     if (SPFA(m)) return true;
                     stk.pop_back();
                     onstk[v] = 0;
                 }
            }
        onstk[u] = 0;
        stk.pop_back();
        return false;
    int solve() {
        for ( int i = 0 ; i < n ; i += 2 ) {</pre>
```

```
match[i] = i+1;
             match[i+1] = i;
        while (true){
             int found = 0;
             for ( int i = 0 ; i < n ; i++ )</pre>
                 onstk[ i ] = d[ i ] = 0;
             for ( int i = 0 ; i < n ; i++ ) {
                 stk.clear();
                 if ( !onstk[i] && SPFA(i) ) {
                     found = 1;
                     while ( stk.size() >= 2 ) {
                          int u = stk.back(); stk.
                              pop_back();
                          int v = stk.back(); stk.
                              pop_back();
                         match[u] = v;
                         match[v] = u;
                     }
                 }
             if (!found) break;
         int ret = 0;
         for ( int i = 0 ; i < n ; i++ )</pre>
            ret += e[i][match[i]];
         ret /= 2;
        return ret;
} graph;
```

4 Graph

- Maximum Independent Set
 - General: [NPC] maximum clique of complement of G
 - Bipartite Graph: [P] Maximum Cardinality Bipartite Matching
 - Tree: [P] dp
- Minimum Dominating Set
 - General: [NPC]
 - Bipartite Graph: [NPC]
 - Tree: [P] DP
- Minimum Vertex Cover
 - General: [NPC] (?)maximum clique of complement of G
 - Bipartite Graph: [P] Maximum Cardinality Bipartite Matching
 - Tree: [P] Greedy, from leaf to root
- Minimum Edge Cover
 - General: [P] V Maximum Matching
 - Bipartite Graph: [P] Greedy, strategy: cover small degree node first.
 - (Min/Max)Weighted: [P]: Minimum/Minimum Weight Matching

4.1 BCC edge

邊雙連通

任意 兩點間至少有兩條不重疊的路徑連接,找法:

- 1. 標記出所有的橋
- 2. 對全圖進行 DFS,不走橋,每一次 DFS 就是一個新的邊雙 連通

```
// from BCW
```

```
struct BccEdge {
   static const int MXN = 100005;
   struct Edge { int v,eid; };
   int n,m,step,par[MXN],dfn[MXN],low[MXN];
   vector<Edge> E[MXN];
```

```
DisjointSet djs;
  void init(int _n) {
    n = n; m = 0;
    for (int i=0; i<n; i++) E[i].clear();</pre>
    djs.init(n);
  void add_edge(int u, int v) {
    E[u].PB({v, m});
    E[v].PB(\{u, m\});
    m++;
  void DFS(int u, int f, int f_eid) {
    par[u] = f;
    dfn[u] = low[u] = step++;
    for (auto it:E[u]) {
      if (it.eid == f_eid) continue;
      int v = it.v;
      if (dfn[v] == -1) {
        DFS(v, u, it.eid);
        low[u] = min(low[u], low[v]);
      } else {
        low[u] = min(low[u], dfn[v]);
    }
  void solve() {
    step = 0;
    memset(dfn, -1, sizeof(int)*n);
    for (int i=0; i<n; i++) {</pre>
      if (dfn[i] == -1) DFS(i, i, -1);
    djs.init(n);
    for (int i=0; i<n; i++) {</pre>
      if (low[i] < dfn[i]) djs.uni(i, par[i]);</pre>
  }
}graph;
```

4.2 Dijkstra

```
typedef struct Edge{
    int v; long long len;
    bool operator > (const Edge &b)const { return len>b
} State;
const long long INF = 1LL<<60;</pre>
void Dijkstra(int n, vector<Edge> G[], long long d[],
    int s, int t=-1){
    static priority_queue<State, vector<State>, greater
         <State> > pq;
    while ( pq.size() )pq.pop();
    for (int i=1; i<=n; i++)d[i]=INF;</pre>
    d[s]=0; pq.push( (State){s,d[s]} );
    while ( pq.size() ){
        auto x = pq.top(); pq.pop();
        int u = x.v;
        if (d[u]<x.len)continue;</pre>
        if (u==t)return;
        for (auto &e:G[u]){
             if (d[e.v] > d[u]+e.len){
                 d[e.v] = d[u]+e.len;
                 pq.push( (State) {e.v,d[e.v]} );
        }
    }
}
```

4.3 Domination.txt

```
| Maximum Independent Set
| General: [NPC] maximum clique of complement of G
```

```
Tree: [P] Greedy
Bipartite Graph: [P] Maximum Cardinality Bipartite
Minimum Dominating Set
General: [NPC]
Tree: [P] DP
Bipartite Graph: [NPC]
Minimum Vertex Cover
General: [NPC] (?)maximum clique of complement of G
Tree: [P] Greedy, from leaf to root
Bipartite Graph: [P] Maximum Cardinality Bipartite
   Matching
Minimum Edge Cover
General: [P] V - Maximum Matching
Bipartite Graph: [P] Greedy, strategy: cover small
    degree node first.
(Min/Max)Weighted: [P]: Minimum/Minimum Weight Matching
```

4.4 max clique

```
const int MAXN = 105;
int best;
int n;
int num[MAXN];
int path[MAXN];
int G[MAXN][MAXN];
bool dfs( int *adj, int total, int cnt ){
    int t[MAXN];
     if (total == 0){
         if( best < cnt ){</pre>
             best = cnt;
             return true;
         return false;
     for(int i = 0; i < total; i++){</pre>
         if( cnt+(total-i) <= best ) return false;</pre>
         if( cnt+num[adj[i]] <= best ) return false;</pre>
         int k=0;
         for(int j=i+1; j<total; j++)</pre>
             if(G[ adj[i] ][ adj[j] ])
                 t[k++] = adj[j];
         if (dfs(t, k, cnt+1)) return true;
    }
    return false:
int MaximumClique(){
     int adj[MAXN];
     if (n <= 0) return 0;
     best = 0;
     for(int i = n-1; i >= 0; i--){
         int k=0;
         for(int j = i+1; j < n; j++)</pre>
             if (g[i][j]) adj[k++] = j;
         dfs( adj, k, 1 );
         num[i] = best;
    return best;
}
```

4.5 min mean cycle

```
// from BCW

/* minimum mean cycle */
const int MAXE = 1805;
const int MAXN = 35;
const double inf = 1029384756;
const double eps = 1e-6;
```

```
struct Edge {
  int v,u;
  double c;
int n,m,prv[MAXN][MAXN], prve[MAXN][MAXN], vst[MAXN];
vector<int> edgeID, cycle, rho;
double d[MAXN][MAXN];
inline void bellman_ford() {
  for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {</pre>
    fill(d[i+1], d[i+1]+n, inf);
    for(int j=0; j<m; j++) {</pre>
       int v = e[j].v, u = e[j].u;
       if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
         d[i+1][u] = d[i][v]+e[j].c;
         prv[i+1][u] = v;
         prve[i+1][u] = j;
       }
    }
  }
double karp_mmc() {
  // returns inf if no cycle, mmc otherwise
  double mmc=inf;
  int st = -1:
  bellman_ford();
  for(int i=0; i<n; i++) {</pre>
    double avg=-inf;
    for(int k=0; k<n; k++) {</pre>
       if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i])</pre>
           /(n-k));
       else avg=max(avg,inf);
    if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
  for(int i=0; i<n; i++) vst[i] = 0;</pre>
  edgeID.clear(); cycle.clear(); rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
    vst[st]++;
    edgeID.PB(prve[i][st]);
    rho.PB(st);
  while (vst[st] != 2) {
    int v = rho.back(); rho.pop_back();
    cycle.PB(v);
    vst[v]++;
  }
  reverse(ALL(edgeID));
  edgeID.resize(SZ(cycle));
  return mmc;
}
```

4.6 SSSP related concepts

```
最短路問題分類:
```

三個工具 Bellman-Ford, Floyd, Dijkstra,

- 1. 可以把 Dijkstra Priority Queue 裡面存的東西想成「狀態」,他可以拿來統計甚至轉移。
- 2. 當遇到邊權會扣掉走的人的血量(或油量之類的),當不能 有負值的時候,就要使用 Bellman-Ford 來做,
- 一開始可以把起點設為最初的血量(油量),拿去做 Bellman-Ford,當做了 n-1 次之後,還能轉移,那就是有負環或正 環(端看如何轉移 Bellman-Ford,這部分的轉移式很自由 可以依照題目敘述亂改。)
- 3. 特別注意如果要判到某一個點的長度是不是無限小,可在做了 n-1 次之後,發現 u->v 可以更新,那我可以去看 v 是否可以到另一點 k,如果是聯通的,代表 k 這個點的長度是無限小。

4.7 Tarjan.cpp

```
割點
點 u 為割點 if and only if 滿足 1. or 2.
1. u 爲樹根,且 u 有多於一個子樹。
2. u 不爲樹根,且滿足存在 (u,v) 爲樹枝邊 (或稱父子邊,
    即 u 爲 v 在搜索樹中的父親),使得 DFN(u) <= Low(v)
一條無向邊 (u,v) 是橋 if and only if (u,v) 爲樹枝邊,且
    滿足 DFN(u) < Low(v)。
// 0 base
struct TarjanSCC{
  static const int MAXN = 1000006;
  int n, dfn[MAXN], low[MAXN], scc[MAXN], scn, count;
  vector<int> G[MAXN];
  stack<int> stk;
  bool ins[MAXN];
  void tarjan(int u){
    dfn[u] = low[u] = ++count;
    stk.push(u);
    ins[u] = true;
    for(auto v:G[u]){
      if(!dfn[v]){
        tarjan(v);
        low[u] = min(low[u], low[v]);
      }else if(ins[v]){
        low[u] = min(low[u], dfn[v]);
    if(dfn[u] == low[u]){
     int v;
      do {
      v = stk.top();
      stk.pop();
      scc[v] = scn;
      ins[v] = false;
      } while(v != u);
      scn++;
   }
  }
  void getSCC(){
    memset(dfn,0,sizeof(dfn));
    memset(low,0,sizeof(low));
    memset(ins,0,sizeof(ins));
    memset(scc,0,sizeof(scc));
    count = scn = 0;
    for(int i = 0 ; i < n ; i++ ){</pre>
      if(!dfn[i]) tarjan(i);
  }
}SCC;
```

4.8 2-SAT

```
const int MAXN = 2020;
struct TwoSAT{
    static const int MAXv = 2*MAXN;
    vector<int> GO[MAXv], BK[MAXv], stk;
    bool vis[MAXv];
    int SC[MAXv];

    void imply(int u,int v){ // u imply v
        GO[u].push_back(v);
        BK[v].push_back(u);
```

```
int dfs(int u,vector<int>*G,int sc){
         vis[u]=1, SC[u]=sc;
         for (int v:G[u])if (!vis[v])
             dfs(v,G,sc);
         if (G==GO)stk.push_back(u);
    int scc(int n=MAXv){
         memset(vis,0,sizeof(vis));
         for (int i=0; i<n; i++)if (!vis[i])</pre>
             dfs(i,G0,-1);
         memset(vis,0,sizeof(vis));
         int sc=0;
         while (!stk.empty()){
             if (!vis[stk.back()])
                 dfs(stk.back(),BK,sc++);
             stk.pop_back();
         }
}SAT;
int main(){
    SAT.scc(2*n);
    bool ok=1;
    for (int i=0; i<n; i++){</pre>
         if (SAT.SC[2*i]==SAT.SC[2*i+1])ok=0;
    if (ok){
         for (int i=0; i<n; i++){</pre>
             if (SAT.SC[2*i]>SAT.SC[2*i+1]){
                 cout << i << endl;</pre>
         }
    else puts("NO");
}
void warshall(){
    bitset<2003> d[2003];
    for (int k=0; k<n; k++){</pre>
         for (int i=0; i<n; i++) if (d[i][k]) {</pre>
             d[i] \mid = d[k];
         }
    }
}
```

4.9 平面圖判定

//skydog

```
#include <iostream>
#include <cstdio>
#include <cstdlib>
#include <iomanip>
#include <vector>
#include <cstring>
#include <string>
#include <queue>
#include <deque>
#include <stack>
#include <man>
#include <set>
#include <utility>
#include <list>
#include <cmath>
#include <algorithm>
#include <cassert>
#include <bitset>
#include <complex>
#include <climits>
#include <functional>
using namespace std;
typedef long long 11;
```

```
typedef pair<int, int> ii;
typedef pair<ll, 11> 14;
#define mp make_pair
#define pb push back
#define debug(x) cerr << #x << " = " << x << " "
const int N=400+1;
struct Planar
    int n,m,hash[N],fa[N],deep[N],low[N],ecp[N];
    vector<int> g[N],son[N];
    set< pair<int,int> > SDlist[N],proots[N];
    int nxt[N][2],back[N],rev[N];
    deque<int> q;
    void dfs(int u)
    {
        hash[u]=1; q.pb(u);
        ecp[u]=low[u]=deep[u];
        int v;
        for (int i = 0; i < g[u].size(); ++i)</pre>
            if(!hash[v=g[u][i]])
            {
                fa[v]=u;
                deep[v]=deep[u]+1;
                dfs(v);
                low[u]=min(low[u],low[v]);
                SDlist[u].insert(mp(low[v],v));
            else ecp[u]=min(ecp[u],deep[v]);
        low[u]=min(low[u],ecp[u]);
    int visited[N];
    void addtree(int u,int t1,int v,int t2)
    {
        nxt[u][t1]=v; nxt[v][t2]=u;
    void findnxt(int u,int v,int& u1,int& v1)
        u1=nxt[u][v^1];
        if(nxt[u1][0]==u) v1=0;
        else v1=1;
    }
    void walkup(int u,int v)
        back[v]=u;
        int v1=v, v2=v, u1=1, u2=0, z;
        for (;;)
            if(hash[v1]==u || hash[v2]==u) break;
            hash[v1]=u;hash[v2]=u; z=max(v1,v2);
            if(z>n)
                int p=fa[z-n];
                if(p!=u)
                     proots[p].insert(mp(-low[z-n], z));
                     v1=p, v2=p, u1=0, u2=1;
                else break;
            }
            else
                findnxt(v1,u1,v1,u1);
                findnxt(v2,u2,v2,u2);
            }
        }
    }
    int topstack;
    pair<int,int> stack[N];
```

```
int outer(int u,int v)
{
    return ecp[v]<deep[u] || (SDlist[v].size() &&</pre>
        SDlist[v].begin()->first<deep[u]);</pre>
}
int inside(int u,int v)
    return proots[v].size()>0 || back[v]==u;
int active(int u,int v)
{
    return inside(u,v) || outer(u,v);
}
void push(int a,int b)
{
    stack[++topstack]=mp(a,b);
void mergestack()
{
    int v1,t1,v2,t2,s,s1;
    v1=stack[topstack].first;t1=stack[topstack].
        second;
    topstack--;
    v2=stack[topstack].first;t2=stack[topstack].
        second;
    topstack--;
    s=nxt[v1][t1^1];
    s1=(nxt[s][1]==v1);
    nxt[s][s1]=v2;
    nxt[v2][t2]=s;
    SDlist[v2].erase( make_pair(low[v1-n],v1-n) );
    proots[v2].erase( make_pair(-low[v1-n],v1) );
void findnxtActive(int u,int t,int& v,int& w1,int S
    )
{
    findnxt(u,t,v,w1);
    while(u!=v && !active(S,v))
        findnxt(v,w1,v,w1);
}
void walkdown(int S,int u)
{
    topstack=0;
    int t1,v=S,w1,x2,y2,x1,y1,p;
    for(t1=0;t1<2;++t1)</pre>
        findnxt(S,t1^1,v,w1);
        while(v!=S)
        {
            if(back[v]==u)
                 while(topstack>0) mergestack();
                addtree(S,t1,v,w1); back[v]=0;
            if(proots[v].size())
                push(v,w1);
                p=proots[v].begin()->second;
                 findnxtActive(p,1,x1,y1,u);
                findnxtActive(p,0,x2,y2,u);
                if(active(u,x1) && !outer(u,x1))
                    v=x1,w1=y1;
                else if(active(u,x2) && !outer(u,x2
                     ))
                    v=x2, w1=y2;
                else if(inside(u,x1) || back[x1]==u
                     v=x1, w1=y1;
```

```
else v=x2,w1=y2;
                 push(p,v==x2);
             else if(v>n || ( ecp[v]>=deep[u] && !
                 outer(u,v)))
                 findnxt(v,w1,v,w1);
             else if(v<=n && outer(u,v) && !topstack</pre>
                 addtree(S,t1,v,w1); break;
             else break;
        }
    }
}
int work(int u)
    int v;
    for (int i = 0; i < g[u].size(); ++i)</pre>
        if(fa[v=g[u][i]]==u)
            son[u].push_back(n+v);
            proots[n+v].clear();
             addtree(n+v,1,v,0);
            addtree(n+v,0,v,1);
    for (int i = 0; i < g[u].size(); ++i)</pre>
        if(deep[v=g[u][i]]>deep[u]+1)
             walkup(u,v);
    topstack=0;
    for (int i = 0; i < son[u].size(); ++i)</pre>
        walkdown(son[u][i], u);
    for (int i = 0; i < g[u].size(); ++i)</pre>
        if(deep[v=g[u][i]]>deep[u]+1 && back[v])
             return 0:
    return 1;
}
void init(int _n)
{
    n = _n;
    m = 0;
    for(int i=1;i<=2*n;++i)</pre>
        g[i].clear();
        SDlist[i].clear();
        son[i].clear();
        proots[i].clear();
        nxt[i][0]=nxt[i][1]=0;
        fa[i]=0;
        hash[i]=0;low[i]=ecp[i]=deep[i]=back[i]=0;
        q.clear();
    }
}
void add(int u, int v)
    g[u].pb(v); g[v].pb(u);
bool check_planar()
    if(m>3*n-5)
        return false;
    // memset(hash,0,sizeof hash);
    for(int i=1;i<=n;++i)</pre>
        if(!hash[i])
        {
             deep[i]=1;
             dfs(i);
    memset(hash,0,sizeof hash);
    //memset(hash, 0, (2*n+1)*sizeof(hash[0]));
    // originally only looks at last n element
    assert(q.size() == n);
    while (!q.empty())
    {
```

```
if (!work(q.back()))
                                                            | }
                return false;
            q.pop_back();
        return true;
    }
} base, new;
vector<ii> edges;
int n, m;
inline void build(int n, Planar &_new)
     new.init(n);
    for (auto e : edges)
        _new.add(e.first, e.second);
void end()
{
    puts("-1");
    exit(0);
bool vis[N];
const int maxp = 5;
int path[maxp], tp=0;
void dfs(int cur)
{
    vis[cur] = true;
    path[tp++] = cur;
    if (tp == maxp)
    auto it = lower_bound(base.g[cur].begin(), base.g[
        cur].end(), path[0]);
        if ( it != base.g[cur].end() && *it == path[0])
            //a cycle
            int x = n+1;
            for (int i = 0; i < 5; ++i) edges.pb(mp(x,
                 path[i]));
            build(x, _new);
            if (_new.check_planar())
                 for (int i = 0; i < maxp; ++i) printf("</pre>
                     %d%c", path[i], i==maxp-1?'\n':' ')
                 exit(0);
            for (int i = 0; i < 5; ++i) edges.pop_back</pre>
                 ();
        }
    }
    else.
    {
        for (auto e : base.g[cur]) if (!vis[e]) dfs(e);
    vis[cur] = false;
    --tp;
int main()
    scanf("%d %d", &n, &m);
    if (n <= 4)
      assert(false);
  puts("0"); return 0;
    for (int i = 0; i < m; ++i)
        int u, v; scanf("%d %d", &u, &v);
        edges.pb(mp(u, v));
    build(n, base);
    if (!base.check_planar()) end();
    for (int i = 1; i <= n; ++i)</pre>
        sort(base.g[i].begin(), base.g[i].end());
    for (int i = 1; i <= n; ++i)</pre>
        dfs(i);
    end();
```

5 Math

```
• Stirling number of second kind S(n,m): n 個相異球,放到 m 個相同的相子,每個箱子至少 1 = m \times S(n-1,m) + S(n-1,m-1) = \frac{1}{m!} \sum_{j=0}^m {m \choose j} (m-j)^n (-1)^j
```

- Stirling number of first kind s(n,m): n 個相異球,分配到 m 個有向環,每個環至少 1 $s(n+1,m)=n\times s(n,m)+s(n,m-1)$ $s(n,m)\equiv {\lfloor n/2\rfloor\choose m-\lfloor n/2\rfloor}$ mod 2
- Pick's Theorem (Bangkok regional 2016 pD) 多邊形頂點都在整數點上 多邊形面積 = 內部整數點個數 + 邊上格子點個數/2 1 A=i+b/2-1

5.1 ax+by=gcd(a,b)

```
pair<int,int> extgcd(int a, int b){
   if (b==0) return {1,0};
   int k = a/b;
   pair<int,int> p = extgcd(b,a-k*b);
   return { p.second, p.first - k*p.second };
}
```

5.2 FFT

```
// use llround() to avoid EPS
typedef double Double;
const Double PI = acos(-1);
// STL complex may TLE
typedef complex<Double> Complex;
#define x real()
#define y imag()
template<typename Iter> // Complex*
void BitReverse(Iter a, int n){
    for (int i=1, j=0; i<n; i++){</pre>
        for (int k = n>>1; k>(j^=k); k>>=1);
        if (i<j) swap(a[i],a[j]);</pre>
    }
}
template<typename Iter> // Complex*
void FFT(Iter a, int n, int rev=1){ // rev = 1 or -1
    assert( (n&(-n)) == n ); // n is power of 2
    BitReverse(a,n);
    Iter A = a;
    for (int s=1; (1<<s)<=n; s++){</pre>
        int m = (1 << s);
        Complex wm( cos(2*PI*rev/m), sin(2*PI*rev/m));
        for (int k=0; k<n; k+=m){</pre>
             Complex w(1,0);
             for (int j=0; j<(m>>1); j++){
                 Complex t = w * A[k+j+(m>>1)];
                 Complex u = A[k+j];
                 A[k+j] = u+t;
                 A[k+j+(m>>1)] = u-t;
                 w = w*wm;
            }
        }
    }
    if (rev==-1){
        for (int i=0; i<n; i++){</pre>
            A[i] /= n;
```

```
| }
|}
```

5.3 GaussElimination

```
// by bcw_codebook
const int MAXN = 300;
const double EPS = 1e-8;
int n;
double A[MAXN][MAXN];
void Gauss() {
  for(int i = 0; i < n; i++) {</pre>
    bool ok = 0;
    for(int j = i; j < n; j++) {</pre>
       if(fabs(A[j][i]) > EPS) {
         swap(A[j], A[i]);
         ok = 1;
         break;
      }
    if(!ok) continue;
    double fs = A[i][i];
    for(int j = i+1; j < n; j++) {</pre>
       double r = A[j][i] / fs;
       for(int k = i; k < n; k++) {</pre>
         A[j][k] -= A[i][k] * r;
    }
  }
}
```

5.4 inverse

```
const int MAXN = 1000006;
int inv[MAXN];
void invTable(int bound, int p){
   inv[1] = 1;
   for (int i=2; i<bound; i++){
        inv[i] = (long long)inv[p%i] * (p-p/i) %p;
   }
}
int inv(int b, int p){
   if (b==1) return 1;
   return (long long)inv(p%b,p) * (p-p/b) %p;
}</pre>
```

5.5 Miller-Rabin

```
typedef long long LL;
inline LL bin_mul(LL a, LL n,const LL& MOD){
    LL re=0;
    while (n>0){
        if (n&1) re += a;
        a += a; if (a>=MOD) a-=MOD;
        n>>=1;
    }
    return re%MOD;
}
inline LL bin_pow(LL a, LL n,const LL& MOD){
    LL re=1;
    while (n>0){
        if (n&1) re = bin_mul(re,a,MOD);
        a = bin_mul(a,a,MOD);
        n>>=1;
```

```
return re:
}
bool is prime(LL n){
 //static LL sprp[3] = { 2LL, 7LL, 61LL};
  static LL sprp[7] = { 2LL, 325LL, 9375LL,
    28178LL, 450775LL, 9780504LL,
    1795265022LL };
  if (n==1 || (n&1)==0 ) return n==2;
  int u=n-1, t=0;
  while ( (u&1)==0 ) u>>=1, t++;
  for (int i=0; i<3; i++){</pre>
    LL x = bin_pow( sprp[i]%n, u, n);
    if (x==0 || x==1 || x==n-1)continue;
    for (int j=1; j<t; j++){</pre>
      x=x*x%n;
      if (x==1 || x==n-1)break;
   if (x==n-1)continue;
   return 0;
 }
  return 1;
```

5.6 Mobius

5.7 pollardRho

```
// from PEC
// does not work when n is prime
Int f(Int x, Int mod){
  return add(mul(x, x, mod), 1, mod);
Int pollard_rho(Int n) {
  if ( !(n & 1) ) return 2;
  while (true) {
    Int y = 2, x = rand()%(n-1) + 1, res = 1;
    for ( int sz = 2 ; res == 1 ; sz *= 2 ) {
      for ( int i = 0 ; i < sz && res <= 1 ; i++) {</pre>
        x = f(x, n);
        res = \_gcd(abs(x-y), n);
      }
      y = x;
    if ( res != 0 && res != n ) return res;
  }
}
```

5.8 SG

```
|Anti Nim (取走最後一個石子者敗)
|
|先手必勝 if and only if
```

```
1. 「所有」堆的石子數都為 1 且遊戲的 SG 值為 0。
2. 「有些」堆的石子數大於 1 且遊戲的 SG 值不為 0。
Anti-SG (決策集合為空的遊戲者贏)
定義 SG 值為 0 時,遊戲結束,
則先手必勝 if and only if
1. 遊戲中沒有單一遊戲的 SG 函數大於 1 且遊戲的 SG 函數
   為 0。
2. 遊戲中某個單一遊戲的 SG 函數大於 1 且遊戲的 SG 函數
   不為 0。
Sprague-Grundy
1. 雙人、回合制
2. 資訊完全公開
3. 無隨機因素
4. 可在有限步內結束
5. 沒有和局
6. 雙方可採取的行動相同
SG(S) 的值為 0:後手(P)必勝
不為 0: 先手(N) 必勝
int mex(set S) {
 // find the min number >= 0 that not in the S
 // e.g. S = \{0, 1, 3, 4\} mex(S) = 2
state = []
int SG(A) {
 if (A not in state) {
   S = sub_states(A)
   if( len(S) > 1 ) state[A] = reduce(operator.xor, [
      SG(B) for B in S1)
   else state[A] = mex(set(SG(B) for B in next_states(
      A)))
 return state[A]
}
```

5.9 theorem

```
Lucas's Theorem
 For non-negative integer n,m and prime P,
 C(m,n) \mod P = C(m/M,n/M) * C(m%M,n%M) \mod P
 = mult_i ( C(m_i,n_i) )
 where m_i is the i-th digit of m in base P.
______
Kirchhoff's theorem
 A_{\{ii\}} = deg(i), A_{\{ij\}} = (i,j) \setminus in E ? -1 : 0
 Deleting any one row, one column, and cal the det(A)
Nth Catalan recursive function:
C_0 = 1, C_{n+1} = C_n * 2(2n + 1)/(n+2)
Mohius Formula
u(n) = 1 , if n = 1
(-1)^m , 若 n 無平方數因數,且 n = p1*p2*p3
           *...*pk
              , 若 n 有大於 1 的平方數因數
- Property
1. (積性函數) u(a)u(b) = u(ab)
2. \sum \{d|n\} \ u(d) = [n == 1]
Mobius Inversion Formula
if
      f(n) = \sum \{d \mid n\} \ g(d)
       g(n) = \sum \{d/n\} \ u(n/d)f(d)
            = \sum \{d|n\} \ u(d)f(n/d)
- Application
```

```
the number/power of gcd(i, j) = k
- Trick
分塊, O(sqrt(n))
Chinese Remainder Theorem (m_i 兩兩互質)
 x = a_1 \pmod{m_1}
 x = a_2 \pmod{m_2}
 x = a_i \pmod{m_i}
construct a solution:
 Let M = m_1 * m_2 * m_3 * ... * m_n
 Let M_i = M / m_i
 ti = 1 / Mi
 t_i * M_i = 1 \pmod{m_i}
 solution x = a_1 * t_1 * M_1 + a_2 * t_2 * M_2 + ...
     + a_n * t_n * M_n + k * M
 = k*M + \sum a_i * t_i * M_i, k is positive integer.
 under mod M, there is one solution x = \sum a_i * t_i *
     M_i
Burnside's Lemma
|G| * |X/G| = sum(|X^g|) where g in G
總方法數: 每一種旋轉下不動點的個數總和 除以 旋轉的方法
```

6 Geometry

6.1 2D point template

```
typedef double Double;
struct Point {
 Double x,y;
  bool operator < (const Point &b)const{</pre>
    //return\ tie(x,y) < tie(b.x,b.y);
    //return atan2(y,x) < atan2(b.y,b.x);
    assert(0 && "choose compare");
  Point operator + (const Point &b)const{
   return {x+b.x,y+b.y};
  Point operator - (const Point &b)const{
   return {x-b.x,y-b.y};
  Point operator * (const Double &d)const{
    return {d*x,d*y};
  Point operator / (const Double &d)const{
    return {x/d,y/d};
  Double operator * (const Point &b)const{
   return x*b.x + y*b.y;
  Double operator % (const Point &b)const{
   return x*b.y - y*b.x;
  friend Double abs2(const Point &p){
    return p.x*p.x + p.y*p.y;
 friend Double abs(const Point &p){
   return sqrt( abs2(p) );
};
typedef Point Vector;
struct Line{
 Point P; Vector v;
```

```
bool operator < (const Line &b)const{
   return atan2(v.y,v.x) < atan2(b.v.y,b.v.x);
  }
};</pre>
```

6.2 circumcentre

```
#include "2Dpoint.cpp"

Point circumcentre(Point &p0, Point &p1, Point &p2){
    Point a = p1-p0;
    Point b = p2-p0;
    Double c1 = abs2(a)*0.5;
    Double c2 = abs2(b)*0.5;
    Double d = a % b;
    Double x = p0.x + ( c1*b.y - c2*a.y ) / d;
    Double y = p0.y + ( c2*a.x - c1*b.x ) / d;
    return {x,y};
}
```

6.3 ConvexHull

```
#include "2Dpoint.cpp"
// retunr H, 第一個點會在 H 出現兩次
void ConvexHull(vector<Point> &P, vector<Point> &H){
    int n = P.size(), m=0;
    sort(P.begin(),P.end());
    H.clear();
    for (int i=0; i<n; i++){</pre>
        while (m>=2 && (P[i]-H[m-2]) % (H[m-1]-H[m-2])
            <0)H.pop_back(), m--;
        H.push_back(P[i]), m++;
    for (int i=n-2; i>=0; i--){
        while (m>=2 && (P[i]-H[m-2]) % (H[m-1]-H[m-2])
            <0)H.pop_back(), m--;
        H.push_back(P[i]), m++;
    }
}
```

6.4 half plane intersection

```
bool OnLeft(const Line& L,const Point& p){
  return Cross(L.v,p-L.P)>0;
Point GetIntersection(Line a, Line b){
  Vector u = a.P-b.P;
  Double t = Cross(b.v,u)/Cross(a.v,b.v);
  return a.P + a.v*t;
int HalfplaneIntersection(Line* L,int n,Point* poly){
  sort(L,L+n);
  int first,last;
  Point *p = new Point[n];
  Line *q = new Line[n];
  q[first=last=0] = L[0];
  for(int i=1;i<n;i++){</pre>
    while(first < last && !OnLeft(L[i],p[last-1])) last</pre>
    while(first < last && !OnLeft(L[i],p[first])) first</pre>
        ++;
    q[++last]=L[i];
    if(fabs(Cross(q[last].v,q[last-1].v))<EPS){</pre>
      last--:
      if(OnLeft(q[last],L[i].P)) q[last]=L[i];
```

6.5 Intersection of two circle

```
vector<Point> interCircle(Point o1, Double r1, Point o2
    , Double r2) {
    Double d2 = abs2(o1 - o2);
    Double d = sqrt(d2);
    Point u = (o1+o2)*0.5 + (o1-o2)*(r2*r2-r1*r1)/(2.0*d2);
        if (abs((r1+r2)*(r1+r2) - d2) < 1e-6) return {u};
        if (d < fabs(r1-r2) || r1+r2 < d) return {};
        Double A = sqrt((r1+r2+d) * (r1-r2+d) * (r1+r2-d) * (-r1+r2+d));
        Point v = Point{o1.y-o2.y, -o1.x+o2.x} * A / (2.0*d2);
        return {u+v, u-v};
}</pre>
```

6.6 Intersection of two lines

6.7 Smallest Circle

```
#include "circumcentre.cpp"
pair<Point,Double> SmallestCircle(int n, Point _p[]){
  Point *p = new Point[n];
  memcpy(p,_p,sizeof(Point)*n);
  random_shuffle(p,p+n);
  Double r2=0;
  Point cen;
  for (int i=0; i<n; i++){</pre>
    if ( abs2(cen-p[i]) <= r2)continue;</pre>
    cen = p[i], r2=0;
    for (int j=0; j<i; j++){</pre>
      if ( abs2(cen-p[j]) <= r2)continue;</pre>
      cen = (p[i]+p[j])*0.5;
      r2 = abs2(cen-p[i]);
      for (int k=0; k<j; k++){</pre>
        if ( abs2(cen-p[k]) <= r2)continue;</pre>
         cen = circumcentre(p[i],p[j],p[k]);
        r2 = abs2(cen-p[k]);
      }
    }
  }
```

```
delete[] p;
  return {cen,r2};
}
// auto res = SmallestCircle(,);
```

7 String

7.1 AC automaton

```
// remember make_fail() !!!
// notice MLE
const int sigma = 62;
const int MAXC = 200005;
inline int idx(char c){
    if ('A'<= c && c <= 'Z')return c-'A';</pre>
    if ('a' \le c \&\& c \le 'z') return c - 'a' + 26;
    if ('0'<= c && c <= '9')return c-'0' + 52;
}
struct ACautomaton{
    struct Node{
        Node *next[sigma], *fail;
        int cnt; // dp
        Node(){
             memset(next,0,sizeof(next));
             fail=0;
             cnt=0:
    } buf[MAXC], *bufp, *ori, *root;
    void init(){
        bufp = buf;
        ori = new (bufp++) Node();
        root = new (bufp++) Node();
    void insert(int n, char *s){
        Node *ptr = root;
        for (int i=0; s[i]; i++){
             int c = idx(s[i]);
             if (ptr->next[c]==NULL)
                 ptr->next[c] = new (bufp++) Node();
             ptr = ptr->next[c];
        ptr->cnt=1;
    }
    Node* trans(Node *o, int c){
        while (o->next[c]==NULL) o = o->fail;
        return o->next[c];
    void make_fail(){
        static queue<Node*> que;
        for (int i=0; i<sigma; i++)</pre>
             ori->next[i] = root;
        root->fail = ori;
        que.push(root);
        while ( que.size() ){
             Node *u = que.front(); que.pop();
             for (int i=0; i<sigma; i++){</pre>
                 if (u->next[i]==NULL)continue;
                 u->next[i]->fail = trans(u->fail,i);
                 que.push(u->next[i]);
             u->cnt += u->fail->cnt;
        }
    }
} ac;
```

7.2 KMP

```
template<typename T>
void build_KMP(int n, T *s, int *f){ // 1 base
  f[0]=-1, f[1]=0;
  for (int i=2; i<=n; i++){</pre>
    int w = f[i-1];
    while (w>=0 \&\& s[w+1]!=s[i])w = f[w];
    f[i]=w+1;
  }
}
template<typename T>
int KMP(int n, T *a, int m, T *b){
  build_KMP(m,b,f);
  int ans=0;
  for (int i=1, w=0; i<=n; i++){</pre>
    while ( w \ge 0 \&\& b[w+1]! = a[i] )w = f[w];
    w++;
    if (w==m){
      ans++;
      w=f[w];
    }
  }
  return ans;
}
```

7.3 palindromic tree

```
// remember init()
// remember make_fail() !!!
// insert s need 1 base !!!
// notice MLE
const int sigma = 62;
const int MAXC = 1000006;
inline int idx(char c){
    if ('a'<= c && c <= 'z')return c-'a';</pre>
    if ('A'<= c && c <= 'Z')return c-'A'+26;
if ('0'<= c && c <= '9')return c-'0'+52;
}
struct PalindromicTree{
    struct Node{
         Node *next[sigma], *fail;
         int len, cnt; // for dp
         Node(){
             memset(next,0,sizeof(next));
             fail=0;
             len = cnt = 0;
    } buf[MAXC], *bufp, *even, *odd;
    void init(){
         bufp = buf;
         even = new (bufp++) Node();
         odd = new (bufp++) Node();
         even->fail = odd;
         odd \rightarrow len = -1;
    void insert(char *s){
         Node* ptr = even;
         for (int i=1; s[i]; i++){
             ptr = extend(ptr,s+i);
         }
    Node* extend(Node *o, char *ptr){
         int c = idx(*ptr);
         while ( *ptr != *(ptr-1-o->len) )o=o->fail;
         Node *&np = o->next[c];
         if (!np){
              np = new (bufp++) Node();
             np \rightarrow len = o \rightarrow len + 2;
             Node *f = o->fail;
```

7.4 SAM

```
// par : fail link
// val : a topological order ( useful for DP )
// go[x] : automata edge ( x is integer in [0,26) )
struct SAM{
  struct State{
    int par, go[26], val;
    State () : par(0), val(0){ FZ(go); }
    State (int _val) : par(0), val(_val){ FZ(go); }
  vector<State> vec;
  int root, tail;
  void init(int arr[], int len){
    vec.resize(2);
    vec[0] = vec[1] = State(0);
    root = tail = 1;
    for (int i=0; i<len; i++)</pre>
      extend(arr[i]);
  void extend(int w){
    int p = tail, np = vec.size();
    vec.PB(State(vec[p].val+1));
    for ( ; p && vec[p].go[w]==0; p=vec[p].par)
      vec[p].go[w] = np;
    if (p == 0){
      vec[np].par = root;
    } else {
      if (vec[vec[p].go[w]].val == vec[p].val+1){
        vec[np].par = vec[p].go[w];
      } else {
        int q = vec[p].go[w], r = vec.size();
        vec.PB(vec[q]);
        vec[r].val = vec[p].val+1;
        vec[q].par = vec[np].par = r;
        for ( ; p && vec[p].go[w] == q; p=vec[p].par)
          vec[p].go[w] = r;
      }
    tail = np;
};
```

7.5 smallest rotation

```
string mcp(string s){
  int n = s.length();
  s += s;
  int i=0, j=1;
  while (i<n && j<n){
    int k = 0;
    while (k < n && s[i+k] == s[j+k]) k++;
    if (s[i+k] <= s[j+k]) j += k+1;
    else i += k+1;
    if (i == j) j++;
}</pre>
```

```
int ans = i < n ? i : j;
return s.substr(ans, n);
}
Contact GitHub API Training Shop Blog About</pre>
```

7.6 suffix array

```
/*he[i]保存了在後綴數組中相鄰兩個後綴的最長公共前綴長度
 *sa[i]表示的是字典序排名為i的後綴是誰(字典序越小的排
     名越靠前)
 *rk[i]表示的是後綴我所對應的排名是多少
const int MAX = 1020304;
int ct[MAX], he[MAX], rk[MAX];
int sa[MAX], tsa[MAX], tp[MAX][2];
void suffix_array(char *ip){
  int len = strlen(ip);
  int alp = 256;
  memset(ct, 0, sizeof(ct));
  for(int i=0;i<len;i++) ct[ip[i]+1]++;</pre>
  for(int i=1;i<alp;i++) ct[i]+=ct[i-1];</pre>
  for(int i=0;i<len;i++) rk[i]=ct[ip[i]];</pre>
  for(int i=1;i<len;i*=2){</pre>
    for(int j=0;j<len;j++){</pre>
      if(j+i>=len) tp[j][1]=0;
      else tp[j][1]=rk[j+i]+1;
      tp[j][0]=rk[j];
    memset(ct, 0, sizeof(ct));
    for(int j=0;j<len;j++) ct[tp[j][1]+1]++;</pre>
    for(int j=1;j<len+2;j++) ct[j]+=ct[j-1];</pre>
    for(int j=0;j<len;j++) tsa[ct[tp[j][1]]++]=j;</pre>
    memset(ct, 0, sizeof(ct));
    for(int j=0;j<len;j++) ct[tp[j][0]+1]++;</pre>
    for(int j=1;j<len+1;j++) ct[j]+=ct[j-1];</pre>
    for(int j=0;j<len;j++)</pre>
      sa[ct[tp[tsa[j]][0]]++]=tsa[j];
    rk[sa[0]]=0;
    for(int j=1;j<len;j++){</pre>
      if( tp[sa[j]][0] == tp[sa[j-1]][0] &&
        tp[sa[j]][1] == tp[sa[j-1]][1])
        rk[sa[j]] = rk[sa[j-1]];
        rk[sa[j]] = j;
  for(int i=0,h=0;i<len;i++){</pre>
    if(rk[i]==0) h=0;
    else{
      int j=sa[rk[i]-1];
      h=max(0,h-1);
      for(;ip[i+h]==ip[j+h];h++);
    he[rk[i]]=h;
  }
}
```

7.7 Z value

```
| z[0] = 0;

for ( int bst = 0, i = 1; i < len ; i++ ) {

   if ( z[bst] + bst <= i ) z[i] = 0;

   else z[i] = min(z[i - bst], z[bst] + bst - i);

   while ( str[i + z[i]] == str[z[i]] ) z[i]++;

   if ( i + z[i] > bst + z[bst] ) bst = i;

}

// 回文版

void Zpal(const char *s, int len, int *z) {

   // Only odd palindrome len is considered
```

7.8 BWT (Burrows-Wheeler Transform)

```
string BWT(string); // by suffix array
string iBWT(string &s, int start=0){
   int n = (int) s.size();
   string ret(n,' ');
   vector<int> next(n,0), box[256];

   for (int i=0; i<n; i++) // bucket sort
      box[ (int)s[i]].push_back(i);

   for (int i=0, j=0; i<256; i++)
      for (int x:box[i])
        next[j++] = x;

   for (int i=0, p=start; i<n; i++)
      ret[i] = s[ p=next[p]];

   return ret;
}</pre>
```

8 Data structure

8.1 2D range tree

```
// remember sort x !!!!!
typedef int T;
const int LGN = 20;
const int MAXN = 100005;
struct Point{
    T x, y;
    friend bool operator < (Point a, Point b){</pre>
        return tie(a.x,a.y) < tie(b.x,b.y);</pre>
};
struct TREE{
    Point pt;
    int toleft;
}tree[LGN][MAXN];
struct SEG{
    T mx, Mx;
    int sz;
    TREE *st;
}seg[MAXN*4];
vector<Point> P;
void build(int 1, int r, int o, int deep){
    seg[o].mx = P[1].x;
    seg[o].Mx = P[r].x;
    seg[o].sz = r-l+1;;
    if(1 == r){
        tree[deep][r].pt = P[r];
        tree[deep][r].toleft = 0;
        seg[o].st = &tree[deep][r];
```

```
return;
    int mid = (l+r)>>1;
    build(1,mid,o+o,deep+1);
    build(mid+1,r,o+o+1,deep+1);
    TREE *ptr = &tree[deep][1];
    TREE *pl = &tree[deep+1][l], *nl = &tree[deep+1][
        mid+1];
    TREE *pr = &tree[deep+1][mid+1], *nr = &tree[deep
        +1][r+1];
    int cnt = 0;
    while(pl != nl && pr != nr) {
        *(ptr) = pl->pt.y <= pr->pt.y ? cnt++, *(pl++):
             *(pr++);
        ptr -> toleft = cnt; ptr++;
    while(pl != nl) *(ptr) = *(pl++), ptr -> toleft =
        ++cnt, ptr++;
    while(pr != nr) *(ptr) = *(pr++), ptr -> toleft =
        cnt, ptr++;
int main(){
   int n; cin >> n;
    for(int i = 0 ;i < n; i++){</pre>
        T x,y; cin >> x >> y;
        P.push_back((Point){x,y});
    sort(P.begin(),P.end());
    build(0,n-1,1,0);
```

8.2 ext heap

```
#include <bits/extc++.h>
typedef __gnu_pbds::priority_queue<int> heap_t;
heap_t a,b;
int main() {
  a.clear();
  b.clear();
  a.push(1);
  a.push(3);
  b.push(2);
  b.push(4);
  assert(a.top() == 3);
  assert(b.top() == 4);
  // merge two heap
  a.join(b);
  assert(a.top() == 4);
  assert(b.empty());
  return 0;
| }
```

8.3 KD tree

```
// from BCW
const int MXN = 100005;
struct KDTree {
    struct Node {
        int x,y,x1,y1,x2,y2;
        int id,f;
        Node *L, *R;
    }tree[MXN];
    int n;
    Node *root;
long long dis2(int x1, int y1, int x2, int y2) {
```

```
long long dx = x1-x2;
    long long dy = y1-y2;
    return dx*dx+dy*dy;
  static bool cmpx(Node& a, Node& b){ return a.x<b.x; }</pre>
  static bool cmpy(Node& a, Node& b){ return a.y<b.y; }</pre>
  void init(vector<pair<int,int>> ip) {
    n = ip.size();
    for (int i=0; i<n; i++) {</pre>
      tree[i].id = i;
      tree[i].x = ip[i].first;
      tree[i].y = ip[i].second;
    root = build_tree(0, n-1, 0);
  Node* build_tree(int L, int R, int dep) {
    if (L>R) return nullptr;
    int M = (L+R)/2;
    tree[M].f = dep%2;
    nth_element(tree+L, tree+M, tree+R+1, tree[M].f ?
        cmpy : cmpx);
    tree[M].x1 = tree[M].x2 = tree[M].x;
    tree[M].y1 = tree[M].y2 = tree[M].y;
    tree[M].L = build_tree(L, M-1, dep+1);
    if (tree[M].L) {
      tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
      tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
      tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
    tree[M].R = build_tree(M+1, R, dep+1);
    if (tree[M].R) {
      tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
      tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
      tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
    return tree+M;
  int touch(Node* r, int x, int y, long long d2){
    long long dis = sqrt(d2)+1;
    if (x<r->x1-dis || x>r->x2+dis || y<r->y1-dis || y>
        r->y2+dis)
      return 0;
    return 1;
  void nearest(Node* r, int x, int y, int &mID, long
      long &md2) {
    if (!r || !touch(r, x, y, md2)) return;
    long long d2 = dis2(r->x, r->y, x, y);
    if (d2 < md2 || (d2 == md2 && mID < r->id)) {
      mID = r \rightarrow id;
      md2 = d2;
    // search order depends on split dim
    if ((r->f == 0 && x < r->x) ||
        (r->f == 1 \&\& y < r->y)) {
      nearest(r->L, x, y, mID, md2);
      nearest(r->R, x, y, mID, md2);
    } else {
      nearest(r->R, x, y, mID, md2);
      nearest(r->L, x, y, mID, md2);
  int query(int x, int y) {
    int id = 1029384756;
    long long d2 = 102938475612345678LL;
    nearest(root, x, y, id, d2);
    return id:
  }
}tree;
```

8.4 Link-Cut tree

```
// from bcw codebook
const int MXN = 100005;
const int MEM = 100005;
struct Splay {
  static Splay nil, mem[MEM], *pmem;
  Splay *ch[2], *f;
  int val, rev, size;
  Splay () : val(-1), rev(0), size(0) {
   f = ch[0] = ch[1] = &nil;
  Splay (int _val) : val(_val), rev(0), size(1) {
    f = ch[0] = ch[1] = &nil;
  bool isr() {
    return f->ch[0] != this && f->ch[1] != this;
  int dir() {
   return f->ch[0] == this ? 0 : 1;
  void setCh(Splay *c, int d) {
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
  void push() {
    if (rev) {
      swap(ch[0], ch[1]);
      if (ch[0] != &nil) ch[0]->rev ^= 1;
      if (ch[1] != &nil) ch[1]->rev ^= 1;
      rev=0;
    }
  void pull() {
    size = ch[0] -> size + ch[1] -> size + 1;
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
} Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::
   mem:
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
  Splay *p = x - > f
  int d = x - > dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
  p->setCh(x->ch[!d], d);
 x->setCh(p, !d);
 p->pull(); x->pull();
vector<Splay*> splayVec;
void splay(Splay *x) {
  splayVec.clear();
  for (Splay *q=x;; q=q->f) {
    splayVec.push_back(q);
    if (q->isr()) break;
 reverse(begin(splayVec), end(splayVec));
 for (auto it : splayVec) it->push();
 while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir()==x->f->dir()) rotate(x->f),rotate
        (x);
    else rotate(x),rotate(x);
 }
Splay* access(Splay *x) {
 Splay *q = nil;
  for (;x!=nil;x=x->f) {
    splay(x);
    x->setCh(q, 1);
```

```
return q;
void evert(Splay *x) {
  access(x);
  splay(x);
  x->rev ^= 1;
  x->push(); x->pull();
void link(Splay *x, Splay *y) {
// evert(x);
  access(x);
  splay(x);
  evert(y);
  x->setCh(y, 1);
void cut(Splay *x, Splay *y) {
// evert(x);
  access(y);
  splay(y);
  y->push();
  y->ch[0] = y->ch[0]->f = nil;
int N, Q;
Splay *vt[MXN];
int ask(Splay *x, Splay *y) {
  access(x);
  access(y);
  splay(x);
  int res = x->f->val;
  if (res == -1) res=x->val;
  return res;
int main(int argc, char** argv) {
  scanf("%d%d", &N, &Q);
for (int i=1; i<=N; i++)</pre>
    vt[i] = new (Splay::pmem++) Splay(i);
  while (Q--) {
    char cmd[105];
    int u, v;
scanf("%s", cmd);
if (cmd[1] == 'i') {
       scanf("%d%d", &u, &v);
       link(vt[v], vt[u]);
    } else if (cmd[0] == 'c') {
       scanf("%d", &v);
       cut(vt[1], vt[v]);
    } else {
       scanf("%d%d", &u, &v);
       int res=ask(vt[u], vt[v]);
       printf("%d \setminus n", res);
    }
  }
  return 0;
}
```

8.5 Treap Lin

```
#include <cstdio>
#include <cstdlib>
#include <algorithm>
#include <string.h>
using namespace std;
const int INF = 999999999;
int ran(){
    static unsigned x = 20170928;
    return x = 0xdefaced*x+1;
struct Treap{
    Treap *1,*r;
    int num,m,sz,tag,ra,ad;
    Treap(int a){
```

```
(February 18, 2019) 18
        1=r=NULL;
        num=m=a:
        sz=1;
        tag=ad=0;
        ra = ran();
    }
}*head,*tp;
int size(Treap *a){
    return a ? a->sz : 0;
int min(Treap *a){
    return a ? a->m+a->ad : INF;
int add(Treap *a){
    return a ? a->ad : 0;
void push(Treap *a){
    if(!a) return;
    if(a->tag){
        swap(a->1,a->r);
        if(a->1)a->1->tag ^= 1;
        if(a->r)a->r->tag ^= 1;
        a->tag=0;
    if(a->1)a->1->ad += a->ad;
    if(a->r)a->r->ad += a->ad;
    a->num += a->ad;
    a->m += a->ad;
    a \rightarrow ad = 0;
void pull(Treap *a){
    if(!a) return;
    a->sz=1+size(a->l)+size(a->r);
    a->m = min(a->num, min(min(a->1), min(a->r));
Treap* merge(Treap *a, Treap *b){
    if(!a || !b) return a ? a : b;
    if(a->ra > b->ra){
        push(a);
        a->r = merge(a->r,b);
        pull(a);
        return a;
    }else{
        push(b);
        b \rightarrow 1 = merge(a, b \rightarrow 1);
        pull(b);
        return b;
    }
void split (Treap *o, Treap *&a, Treap *&b,int k){
    if(!k) a=NULL, b=o;
    else if(size(o)==k) a=o, b=NULL;
    else{
        push(o);
        if(k <= size(o->1)){
            b = o;
            split(o->1, a, b->1,k);
            pull(b);
        }else{
            a = o;
            split(o->r, a->r, b, k-size(o->l)-1);
            pull(a);
        }
    }
int main(){
    int n, tmp;
    scanf("%d",&n);
    for(int i = 0 ;i < n ;i++){</pre>
        scanf("%d",&tmp);
        tp = new Treap(tmp);
        head = merge(head, tp);
    int Q;
    scanf("%d n", &Q);
```

```
char ss[50];
    int a, b, c;
    Treap *ta, *tb, *tc, *td;
    while(Q--){
        scanf("%s",ss);
        if(strcmp(ss,"ADD")==0){
             scanf("%d %d %d",&a,&b,&c);
             split(head,tb,tc,b);
             split(tb,ta,tb,a-1);
             tb -> ad += c;
             head = merge(ta, merge(tb, tc));
        }else if(strcmp(ss,"REVERSE")==0){
             scanf("%d %d",&a,&b);
             split(head,tb,tc,b);
             split(tb,ta,tb,a-1);
             tb -> tag ^= 1;
             head = merge(ta, merge(tb, tc));
        }else if(strcmp(ss,"REVOLVE")==0){
             scanf("%d %d %d",&a,&b,&c);
             split(head,tb,tc,b);
             split(tb,ta,tb,a-1);
             int szz = size(tb);
             c %= szz;
             split(tb,tb,td,szz-c);
             tb=merge(td,tb);
             head = merge(ta, merge(tb, tc));
        }else if(strcmp(ss,"INSERT")==0){
             scanf("%d %d",&a,&b);
             split(head,ta,tc,a);
             tb = new Treap(b);
             head = merge(ta, merge(tb, tc));
        }else if(strcmp(ss,"DELETE")==0){
             scanf("%d",&a);
             split(head,ta,tc,a-1);
             split(tc,tb,tc,1);
             delete tb;
             head = merge(ta,tc);
        }else if(strcmp(ss,"MIN")==0){
    scanf("%d %d",&a,&b);
             split(head,tb,tc,b);
             split(tb,ta,tb,a-1);
             printf("%d \ n", min(tb));
             head = merge(ta, merge(tb, tc));
    }
}
```

9 Other

生成树计数

9.1 count spanning tree

```
新的方法介绍
下面我们介绍一种新的方法——Matrix-Tree定理(Kirchhoff矩
  阵-树定理)。
Matrix-Tree定理是解决生成树计数问题最有力的武器之一。它
  首先于1847年被Kirchhoff证明。在介绍定理之前,我们首
  先明确几个概念:
1、G的度数矩阵D[G]是一个n*n的矩阵,并且满足: 当i≠j时,
  dij=0;当i=j时,dij等于vi的度数。
2、G的邻接矩阵A[G]也是一个n*n的矩阵, 并且满足:如果vi
   、vj之间有边直接相连,则aij=1,否则为0。
我们定义G的Kirchhoff矩阵(也称为拉普拉斯算子)C[G]为C[G]=
则Matrix-Tree定理可以描述为:G的所有不同的生成树的个数
  等于其Kirchhoff矩阵C[G]任何一个n-1阶主子式的行列式
  的绝对值。
所谓n-1阶主子式,就是对于r(1≤r≤n),将C[G]的第r行、第r列
  同时去掉后得到的新矩阵,用Cr[G]表示。
```

```
算法步骤:
1、 构建拉普拉斯矩阵
     Matrix[i][j] =
degree(i) , i==j
          -1, i-j有边
           0,其他情况
2、 去掉第r行,第r列(r任意)
3、 计算矩阵的行列式
/* ********************************
        : Chen Fan
MYTD
LANG
        : G++
PROG
        : Count_Spaning_Tree_From_Kuangbin
#include <stdio.h>
#include <string.h>
#include <algorithm>
#include <iostream>
#include <math.h>
using namespace std;
const double eps = 1e-8;
const int MAXN = 110;
int sgn(double x)
    if(fabs(x) < eps)return 0;</pre>
    if(x < 0) return -1;
    else return 1;
double b[MAXN][MAXN];
double det(double a[][MAXN],int n)
    int i, j, k, sign = 0;
    double ret = 1;
    for(i = 0;i < n;i++)</pre>
    for(j = 0;j < n;j++) b[i][j] = a[i][j];</pre>
    for(i = 0; i < n; i++)
        if(sgn(b[i][i]) == 0)
            for(j = i + 1; j < n; j++)
            if(sgn(b[j][i]) != 0) break;
            if(j == n)return 0;
            for(k = i; k < n; k++) swap(b[i][k], b[j][k]);
            sign++;
        ret *= b[i][i];
        for(k = i + 1;k < n;k++) b[i][k]/=b[i][i];</pre>
        for(j = i+1;j < n;j++)</pre>
        for(k = i+1;k < n;k++) b[j][k] -= b[j][i]*b[i][</pre>
            k];
    if(sign & 1)ret = -ret;
    return ret;
double a[MAXN][MAXN];
int g[MAXN][MAXN];
int main()
    int T;
    int n,m;
    int u,v;
    scanf("%d",&T);
    while(T--)
        scanf("%d%d",&n,&m);
        memset(g,0,sizeof(g));
        while(m--)
            scanf("%d%d",&u,&v);
            u--;v--;
            g[u][v] = g[v][u] = 1;
        memset(a,0,sizeof(a));
        for(int i = 0;i < n;i++)</pre>
        for(int j = 0; j < n; j++)</pre>
        if(i != j && g[i][j])
```

```
a[i][i]++;
a[i][j] = -1;
}
double ans = det(a,n-1);
printf("%.0lf\n",ans);
}
return 0;
}
```

9.2 C++11 random

```
void init(){
   std::random_device rd;
   std::default_random_engine gen( rd() );
   std::uniform_int_distribution <unsigned long long>
        dis(0,ULLONG_MAX);

for (int i=0; i<MAXN; i++){
        h[i] = dis(gen);
   }
}</pre>
```

9.3 Digit Counting

```
int dfs(int pos, int state1, int state2 ...., bool
    limit, bool zero) {
    if (pos == -1) return 是否符合條件;
    int &ret = dp[pos][state1][state2][....];
    if ( ret != -1 && !limit ) return ret;
    int ans = 0;
    int upper = limit ? digit[pos] : 9;
    for ( int i = 0 ; i <= upper ; i++ ) {</pre>
        ans += dfs(pos - 1, new_state1, new_state2,
            limit & ( i == upper), ( i == 0) && zero);
    if ( !limit ) ret = ans;
   return ans;
int solve(int n) {
   int it = 0;
    for ( ; n ; n /= 10 ) digit[it++] = n % 10;
    return dfs(it - 1, 0, 0, 1, 1);
```

9.4 DP optimization

```
Monotonicity & 1D/1D DP & 2D/1D DP
Definition xD/yD
1D/1D DP[j] = min(0 \le i < j) \{ DP[i] + w(i, j) \}; DP[0] = k
2D/1D DP[i][j] = min(i < k \le j) \{ DP[i][k - 1] + DP[k][j] \}
    + w(i, j); DP[i][i] = 0
Monotonicity
a \mid w(a, c) w(a, d)
b \mid w(b, c) w(b, d)
Monge Condition
Concave(凹四邊形不等式): w(a, c) + w(b, d) >= w(a, d) +
     w(b, c)
Convex (凸四邊形不等式): w(a, c) + w(b, d) <= w(a, d) +
     w(b, c)
Totally Monotone
Concave(凹單調): w(a, c) <= w(b, d) ----> w(a, d) <= w
Convex (凸單調): w(a, c) >= w(b, d) ----> w(a, d) >= w
    (b, c)
```

9.5 DP 1D/1D

```
#include<bits/stdc++.h>
int t, n, L;
int p;
char s[MAXN][35];
11 \text{ sum}[MAXN] = \{0\};
long double dp[MAXN] = {0};
int prevd[MAXN] = {0};
long double pw(long double a, int n) {
    if ( n == 1 ) return a;
    long double b = pw(a, n/2);
    if ( n & 1 ) return b*b*a;
    else return b*b;
long double f(int i, int j) {
     cout << (sum[i] - sum[j]+i-j-1-L) << endl;</pre>
    return pw(abs(sum[i] - sum[j]+i-j-1-L), p) + dp[j];
struct INV {
    int L, R, pos;
INV stk[MAXN*10];
int top = 1, bot = 1;
void update(int i) {
    while ( top > bot && i < stk[top].L && f(stk[top].L</pre>
         , i) < f(stk[top].L, stk[top].pos) ) {</pre>
        stk[top - 1].R = stk[top].R;
        top--:
    int lo = stk[top].L, hi = stk[top].R, mid, pos =
        stk[top].pos;
    //if ( i >= lo ) lo = i + 1;
    while ( lo != hi ) {
        mid = lo + (hi - lo) / 2;
        if ( f(mid, i) < f(mid, pos) ) hi = mid;</pre>
        else lo = mid + 1;
    if ( hi < stk[top].R ) {</pre>
        stk[top + 1] = (INV) { hi, stk[top].R, i };
        stk[top++].R = hi;
    }
}
int main() {
    cin >> t;
    while ( t-- ) {
        cin >> n >> L >> p;
        dp[0] = sum[0] = 0;
        for ( int i = 1 ; i <= n ; i++ ) {</pre>
            cin >> s[i];
             sum[i] = sum[i-1] + strlen(s[i]);
             dp[i] = numeric_limits<long double>::max();
        stk[top] = (INV) {1, n + 1, 0};
        for ( int i = 1 ; i <= n ; i++ ) {
             if ( i >= stk[bot].R ) bot++;
            dp[i] = f(i, stk[bot].pos);
            update(i);
//
              cout << (ll) f(i, stk[bot].pos) << endl;</pre>
        if ( dp[n] > 1e18 ) {
             cout << "Too hard to arrange" << endl;</pre>
        } else {
            vector<PI> as;
```

```
cout << (l1)dp[n] << endl;
}
}
return 0;
}</pre>
```

9.6 stable marriage

```
// normal stable marriage problem
// input:
//3
//Albert Laura Nancy Marcy
//Brad Marcy Nancy Laura
//Chuck Laura Marcy Nancy
//Laura Chuck Albert Brad
//Marcy Albert Chuck Brad
//Nancy Brad Albert Chuck
#include < bits / stdc++.h>
using namespace std;
const int MAXN = 505;
int n:
int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id
int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank
int current[MAXN]; // current[boy_id] = rank; boy_id
    will pursue current[boy_id] girl.
int girl_current[MAXN]; // girl[girl_id] = boy_id;
void initialize() {
 for ( int i = 0 ; i < n ; i++ ) {</pre>
    current[i] = 0;
    girl_current[i] = n;
    order[i][n] = n;
 }
map<string, int> male, female;
string bname[MAXN], gname[MAXN];
int fit = 0;
void stable_marriage() {
  queue<int> que;
  for ( int i = 0 ; i < n ; i++ ) que.push(i);</pre>
  while ( !que.empty() ) {
    int boy_id = que.front();
    que.pop();
    int girl_id = favor[boy_id][current[boy_id]];
    current[boy_id] ++;
    if ( order[girl_id][boy_id] < order[girl_id][</pre>
        girl_current[girl_id]] ) {
      if ( girl_current[girl_id] < n ) que.push(</pre>
          girl_current[girl_id]); // if not the first
          time
      girl_current[girl_id] = boy_id;
    } else {
      que.push(boy_id);
    }
  }
int main() {
 cin >> n:
  for ( int i = 0 ; i < n; i++ ) {</pre>
    string p, t;
    cin >> p;
    male[p] = i;
    bname[i] = p;
    for ( int j = 0 ; j < n ; j++ ) {</pre>
```

```
cin >> t;
    if ( !female.count(t) ) {
       gname[fit] = t;
      female[t] = fit++;
    favor[i][j] = female[t];
}
for ( int i = 0 ; i < n ; i++ ) {</pre>
  string p, t;
  cin >> p;
  for ( int j = 0 ; j < n ; j++ ) {</pre>
    cin >> t;
    order[female[p]][male[t]] = j;
}
initialize();
stable_marriage();
for ( int i = 0 ; i < n ; i++ ) {</pre>
  cout << bname[i] << " " << gname[favor[i][current[i</pre>
      ] - 1]] << endl;
```

9.7 Mo's algorithm

```
int 1 = 0, r = 0, nowAns = 0, BLOCK_SIZE, n, m;
int ans[];
struct QUE{
    int 1, r, id;
    friend bool operator < (QUE a, QUE b){</pre>
        if(a.l / BLOCK_SIZE != b.l / BLOCK_SIZE)
            return a.1 / BLOCK SIZE < b.1 / BLOCK SIZE;</pre>
        return a.r < b.r;</pre>
    }
}querys[];
inline void move(int pos, int sign) {
    // update nowAns
}
void solve() {
    BLOCK_SIZE = int(ceil(pow(n, 0.5)));
    sort(querys, querys + m);
    for (int i = 0; i < m; ++i) {</pre>
        const QUE &q = querys[i];
        while (1 > q.1) move(--1, 1);
        while (r < q.r) move(r++, 1);
        while (1 < q.1) move(1++, -1);
        while (r > q.r) move(--r, -1);
        ans[q.id] = nowAns;
    }
```

9.8 Parser

```
~Node(){ delete L; delete R; }
    enum { op, op1, num } type;
    LL val;
    Node *L, *R;
} *root;
char getOp1(int LEVEL, istream& is){
    is >>ws;
    for (auto& x : Op1s){
        auto& op = x.first;
        auto& lev = x.second;
        if (LEVEL == lev && is.peek() == op)
            return is.get();
    return 0;
template <int LEVEL> void parse(Node*& x, istream& is){
    char op1 = getOp1(LEVEL, is);
    parse<LEVEL+1>(x, is);
    if (op1) x = new Node{Node::op1, op1, x, nullptr};
    auto& ops = Ops[LEVEL];
    while (is>>ws && count(ops.begin(), ops.end(), is.
        peek())){
        x = new Node{Node::op, is.get(), x, nullptr};
        parse<LEVEL+1>(x->R, is);
template <> void parse<MAXLEVEL>(Node*& x, istream& is)
    char op1 = getOp1(MAXLEVEL, is);
    is>>ws;
    if (is.peek()>='0' && is.peek()<='9'){</pre>
        LL t; is >>t;
        x = new Node{Node::num, t, nullptr, nullptr};
    } else if (is.peek() == '('){
        is.get();
        parse<0>(x, is);
        is>>ws;
        if (is.get()!=')') throw 0;
    } else throw 0;
    if (op1) x = new Node{Node::op1, op1, x, nullptr};
// throw when error occur !!!!!
void build(istream& is){
    parse<0>(root, is);
    if ((is>>ws).peek() != EOF) throw 0;
```

9.9 java cheat sheet

```
import java.util.*;
import java.math.*;
import java.io.*;
public class java{
    static class Comp implements Comparator<Integer>{
        public int compare(Integer lhs, Integer rhs){
            return lhs - rhs;
        }
    static class Yee implements Comparable<Yee>{
        public int compareTo(Yee y){
            return 0;
    static class Reader{
        private BufferedReader br;
        private StringTokenizer st;
        public Reader(){
            br = new BufferedReader(new
                InputStreamReader(System.in));
        boolean hasNext() throws IOException{
                                                          }
            String s;
            while (st == null || !st.hasMoreElements())
                {
```

```
if ((s = br.readLine())==null) return
                false:
            st = new StringTokenizer(s);
        return true;
    String next() throws IOException{
        while (st == null || !st.hasMoreElements())
            st = new StringTokenizer(br.readLine())
        return st.nextToken();
    int nextInt() throws IOException{
        return Integer.parseInt(next());
    }// Long.parseLong, Double.parseDouble, br.
        readLine
public static void main(String args[])throws
    IOException{
    Reader cin = new Reader();
    //Scanner cin = new Scanner(System.in);
    PrintWriter cout = new PrintWriter(System.out);
    //Scanner cin = new Scanner(new File("t.in"));
    //PrintWriter cout = new PrintWriter(new File("
        t.out"));
       **** cout.close() or cout.flush() is needed
    // 2D array: int[][] a = new int[10][10];
    // input, EOF, Graph
    int n = cin.nextInt();
    // nextFloat, nextLine, next
    ArrayList<ArrayList<Integer>> G = new ArrayList
        <>();
    for (int i=0; i<n; i++) G.add(new ArrayList<>()
    while (cin.hasNext()){ // EOF
        int u = cin.nextInt(), v = cin.nextInt();
        G.get(u).add(v);
    // Math: E, PI, min, max, random(double 0~1),
    // Collections(List a): swap(a,i,j), sort(a[,
        comp]), min(a), binarySearch(a,val[,comp])
    // set
    Set<Integer> set = new TreeSet<>();
    set.add(87); set.remove(87);
    if (!set.contains(87)) cout.println("no 87");
    Map<String, Integer> map = new HashMap<>();
    map.put("0", 1); map.put("2", 3);
    for ( Map.Entry<String,Integer> i : map.
        entrySet() )
        cout.println(i.getKey() + " " + i.getValue
            () + " wry");
    cout.println( map.get("1") );
    // Big Number: TEN ONE ZERO, modInverse
        isProbablePrime modInverse modPow
    // add subtract multiply divide remainder, and
        or xor not shiftLeft shiftRight
    // queue: add, peek(==null), poll
    PriorityQueue<Integer> pq = new PriorityQueue<
        Integer>(Collections.reverseOrder());
    Queue<Integer> q = new ArrayDeque<Integer>();
    // stack: push, empty, pop
    Stack<Integer> s = new Stack<Integer>();
    cout.close();
}
```

9.10 python cheat sheet

```
#!/usr/bin/env python3
# import
import math
from math import *
import math as M
from math import sqrt
# input
n = int( input() )
a = [ int(x) for x in input().split() ]
# EOF
while True:
    try:
        solve()
    except:
        break;
# output
print( x, sep=' ')
print( ''.join( str(x)+' ' for x in a ) )
print( '{:5d}'.format(x) )
# sort
a.sort()
sorted(a)
# list
a = [ x for x in range(n) ]
a.append(x)
# Basic operator
a, b = 10, 20
a/b # 0.5
a//b # 0
a%b # 10
a**b # 10^20
# if, else if, else
if a==0:
    print('zero')
elif a>0:
    print('postive')
else:
    print('negative')
# loop
while a==b and b==c:
for i in LIST:
               # C++
# stack
stack = [3,4,5]
stack.append(6) # push()
stack.pop()
                # pop()
                # top()
stack[-1]
                # size() 0(1)
len(stack)
                # C++
# queue
from collections import deque
queue = deque([3,4,5])
queue.append(6) # push()
queue.popleft() # pop()
queue[0]
                # front()
               # size() 0(1)
len(queue)
# random
from random import *
randrange(L,R,step) # [L,R) L+k*step
randint(L,R) # int from [L,R]
choice(list) # pick 1 item from list
choices(list,k) # pick k item
shuffle(list)
Uniform(L,R) # float from [L,R]
```

```
# Decimal
from fractions import Fraction
from decimal import Decimal, getcontext
getcontext().prec = 250 # set precision
itwo = Decimal(0.5)
two = Decimal(2)
N = 200
def angle(cosT):
  """given cos(theta) in decimal return theta"""
 for i in range(N):
    cosT = ((cosT + 1) / two) ** itwo
 sinT = (1 - cosT * cosT) ** itwo
return sinT * (2 ** N)
pi = angle(Decimal(-1))
# file IO
r = open("filename.in")
a = r.read() # read whole content into one string
w = open("filename.out", "w")
w.write('123\n')
# IO redirection
import sys
sys.stdin = open('filename.in')
sys.stdout = open('filename.out', 'w')
```

