NCTU_Oimo

Contents

1 Basic

1.1 Vimrc

```
set mouse=a
set t_Co=256
set cursorline
syntax on
syntax enable
set number
map <F9> : !g++ % -o %<.out <CR>
map <F5> : !./%<.out < %<.in <CR>
set bg=light
set shiftwidth=2
set tabstop=2
set ai
set nu
set ruler
set incsearch
filetype indent on
hi Comment ctermfg=darkcyan
```

2 DataStructure

2.1 KDTree

```
// from BCW

const int MXN = 100005;

struct KDTree {
  struct Node {
    int x,y,x1,y1,x2,y2;
    int id,f;
    Node *L, *R;
  }tree[MXN];
  int n;
  Node *root;
```

```
  long long dis2(int x1, int y1, int x2, int y2) {
    long long dx = x1-x2;
    long long dy = y1-y2;
    return dx*dx+dy*dy;
  }
  static bool cmpx(Node& a, Node& b){ return a.x<b.x; }
  static bool cmpy(Node& a, Node& b){ return a.y<b.y; }
  void init(vector<pair<int,int>> ip) {
    n = ip.size();
    for (int i=0; i<n; i++) {
      tree[i].id = i;
      tree[i].x = ip[i].first;
      tree[i].y = ip[i].second;
    }
    root = build_tree(0, n-1, 0);
  }
  Node* build_tree(int L, int R, int dep) {
    if (L>R) return nullptr;
    int M = (L+R)/2;
    tree[M].f = dep%2;
    nth_element(tree+L, tree+M, tree+R+1, tree[M].f ?
        cmpy : cmpx);
    tree[M].x1 = tree[M].x2 = tree[M].x;
    tree[M].y1 = tree[M].y2 = tree[M].y;

    tree[M].L = build_tree(L, M-1, dep+1);
    if (tree[M].L) {
      tree[M].x1 = min(tree[M].x1, tree[M].L->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].L->x2);
      tree[M].y1 = min(tree[M].y1, tree[M].L->y1);
      tree[M].y2 = max(tree[M].y2, tree[M].L->y2);
    }

    tree[M].R = build_tree(M+1, R, dep+1);
    if (tree[M].R) {
      tree[M].x1 = min(tree[M].x1, tree[M].R->x1);
      tree[M].x2 = max(tree[M].x2, tree[M].R->x2);
      tree[M].y1 = min(tree[M].y1, tree[M].R->y1);
      tree[M].y2 = max(tree[M].y2, tree[M].R->y2);
    }

    return tree+M;
  }
  int touch(Node* r, int x, int y, long long d2){
    long long dis = sqrt(d2)+1;
    if (x<r->x1-dis || x>r->x2+dis || y<r->y1-dis || y>
        r->y2+dis)
      return 0;
    return 1;
  }
  void nearest(Node* r, int x, int y, int &mID, long
        long &md2) {
    if (!r || !touch(r, x, y, md2)) return;
    long long d2 = dis2(r->x, r->y, x, y);
    if (d2 < md2 || (d2 == md2 && mID < r->id)) {
      mID = r->id;
      md2 = d2;
    }
    // search order depends on split dim
    if ((r->f == 0 && x < r->x) ||
        (r->f == 1 && y < r->y)) {
      nearest(r->L, x, y, mID, md2);
      nearest(r->R, x, y, mID, md2);
    } else {
      nearest(r->R, x, y, mID, md2);
      nearest(r->L, x, y, mID, md2);
    }
  }
  int query(int x, int y) {
    int id = 1029384756;
    long long d2 = 102938475612345678LL;
    nearest(root, x, y, id, d2);
    return id;
  }
}tree;
```

2.2 SparseTable

```
const int MAXN = 200005;
const int lgN = 20;

struct SP{ //sparse table
  int Sp[MAXN][lgN];
  function<int(int,int)> opt;
  void build(int n, int *a){ // 0 base
    for (int i=0 ;i<n; i++) Sp[i][0]=a[i];

    for (int h=1; h<lgN; h++){
      int len = 1<<(h-1), i=0;
      for (; i+len<n; i++)
        Sp[i][h] = opt( Sp[i][h-1] , Sp[i+len][h-1] );
      for (; i<n; i++)
        Sp[i][h] = Sp[i][h-1];
    }
  }
  int query(int l, int r){
    int h = __lg(r-l+1);
    int len = 1<<h;
    return opt( Sp[l][h] , Sp[r-len+1][h] );
  }
};
```

2.3 Treap

```
#include<bits/stdc++.h>
using namespace std;
template<class T,unsigned seed>class treap{
  public:
    struct node{
      T data;
      int size;
      node *l,*r;
      node(T d){
        size=1;
        data=d;
        l=r=NULL;
      }
      inline void up(){
        size=1;
        if(l)size+=l->size;
        if(r)size+=r->size;
      }
      inline void down(){
      }
    }*root;
    inline int size(node *p){return p?p->size:0;}
    inline bool ran(node *a,node *b){
      static unsigned x=seed;
      x=0xdefaced*x+1;
      unsigned all=size(a)+size(b);
      return (x%all+all)%all<size(a);
    }
    void clear(node *&p){
      if(p)clear(p->l),clear(p->r),delete p,p=NULL;
    }
    ~treap(){clear(root);}
    void split(node *o,node *&a,node *&b,int k){
      if(!k)a=NULL,b=o;
      else if(size(o)==k)a=o,b=NULL;
      else{
        o->down();
        if(k<=size(o->l)){
          b=o;
          split(o->l,a,b->l,k);
          b->up();
        }else{
          a=o;
          split(o->r,a->r,b,k-size(o->l)-1);
          a->up();
        }
      }
```

```
    }
    void merge(node *&o,node *a,node *b){
      if(!a||!b)o=a?a:b;
      else{
        if(ran(a,b)){
          a->down();
          o=a;
          merge(o->r,a->r,b);
        }else{
          b->down();
          o=b;
          merge(o->l,a,b->l);
        }
        o->up();
      }
    }
    void build(node *&p,int l,int r,T *s){
      if(l>r)return;
      int mid=(l+r)>>1;
      p=new node(s[mid]);
      build(p->l,l,mid-1,s);
      build(p->r,mid+1,r,s);
      p->up();
    }
    inline int rank(T data){
      node *p=root;
      int cnt=0;
      while(p){
        if(data<=p->data)p=p->l;
        else cnt+=size(p->l)+1,p=p->r;
      }
      return cnt;
    }
    inline void insert(node *&p,T data,int k){
      node *a,*b,*now;
      split(p,a,b,k);
      now=new node(data);
      merge(a,a,now);
      merge(p,a,b);
    }
    inline void remove(node *&p, int k) {
      node *a, *b, *res, *die;
      split(p, a, res, k);
      if (res == NULL) return;
      split(res, die, b, 1);
      merge(a, a, b);
      if (size(a) > size(b)) p = a;
      else p = b;
      clear(die);
    }
};
treap<T ,20141223>bst;
int main(){
  bst.remove(bst.root, bst.rank(E));
  bst.insert(bst.root, E, bst.rank(E));
}
```

2.4 Link Cut Tree

```
// from bcw codebook

const int MXN = 100005;
const int MEM = 100005;

struct Splay {
  static Splay nil, mem[MEM], *pmem;
  Splay *ch[2], *f;
  int val, rev, size;
  Splay () : val(-1), rev(0), size(0) {
    f = ch[0] = ch[1] = &nil;
  }
  Splay (int _val) : val(_val), rev(0), size(1) {
    f = ch[0] = ch[1] = &nil;
  }
  bool isr() {
    return f->ch[0] != this && f->ch[1] != this;
  }
  int dir() {
    return f->ch[0] == this ? 0 : 1;
  }
  void setCh(Splay *c, int d) {
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
  }
  void push() {
    if (rev) {
      swap(ch[0], ch[1]);
      if (ch[0] != &nil) ch[0]->rev ^= 1;
      if (ch[1] != &nil) ch[1]->rev ^= 1;
      rev=0;
    }
  }
  void pull() {
    size = ch[0]->size + ch[1]->size + 1;
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
  }
} Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::
    mem;
Splay *nil = &Splay::nil;

void rotate(Splay *x) {
  Splay *p = x->f;
  int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
  p->setCh(x->ch[!d], d);
  x->setCh(p, !d);
  p->pull(); x->pull();
}

vector<Splay*> splayVec;
void splay(Splay *x) {
  splayVec.clear();
  for (Splay *q=x;; q=q->f) {
    splayVec.push_back(q);
    if (q->isr()) break;
  }
  reverse(begin(splayVec), end(splayVec));
  for (auto it : splayVec) it->push();
  while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir()==x->f->dir()) rotate(x->f),rotate
        (x);
    else rotate(x),rotate(x);
  }
}

Splay* access(Splay *x) {
  Splay *q = nil;
  for (;x!=nil;x=x->f) {
    splay(x);
    x->setCh(q, 1);
    q = x;
  }
  return q;
}
void evert(Splay *x) {
  access(x);
  splay(x);
  x->rev ^= 1;
  x->push(); x->pull();
}
void link(Splay *x, Splay *y) {
//  evert(x);
  access(x);
  splay(x);
  evert(y);
  x->setCh(y, 1);
}
void cut(Splay *x, Splay *y) {
//  evert(x);
```

```
  access(y);
  splay(y);
  y->push();
  y->ch[0] = y->ch[0]->f = nil;
}

int N, Q;
Splay *vt[MXN];

int ask(Splay *x, Splay *y) {
  access(x);
  access(y);
  splay(x);
  int res = x->f->val;
  if (res == -1) res=x->val;
  return res;
}
int main(int argc, char** argv) {
  scanf("%d%d", &N, &Q);
  for (int i=1; i<=N; i++)
    vt[i] = new (Splay::pmem++) Splay(i);
  while (Q--) {
    char cmd[105];
    int u, v;
    scanf("%s", cmd);
    if (cmd[1] == 'i') {
      scanf("%d%d", &u, &v);
      link(vt[v], vt[u]);
    } else if (cmd[0] == 'c') {
      scanf("%d", &v);
      cut(vt[1], vt[v]);
    } else {
      scanf("%d%d", &u, &v);
      int res=ask(vt[u], vt[v]);
      printf("%d\n", res);
    }
  }

  return 0;
}
```

2.5 Pb Ds Heap

```
#include <bits/extc++.h>
typedef __gnu_pbds::priority_queue<int> heap_t;
heap_t a,b;
int main() {
  a.clear();b.clear();
  a.push(1);a.push(3);b.push(2);b.push(4);
  // merge two heap
  a.join(b);
  assert(a.top() == 4);
  assert(b.empty());
  return 0;
}
```

2.6 Pb Ds Rbtree

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
#define ordered_set tree<int, null_type,less<int>,
    rb_tree_tag,tree_order_statistics_node_update>
int main() {
    ordered_set o_set;
    o_set.insert(5);
    o_set.insert(1);
    o_set.insert(2);
    cout << *(o_set.find_by_order(1)) << endl; // 2
    cout << o_set.order_of_key(4) << endl;     // 2
    cout << o_set.order_of_key(5) << endl;     // 2
    if (o_set.find(2) != o_set.end())
        o_set.erase(o_set.find(2));
    cout << *(o_set.find_by_order(1)) << endl; // 5
    cout << o_set.order_of_key(4) << endl;     // 1
}
```

3 Flow

3.1 Dinic

```
(a) Bounded Maxflow Construction:
1. add two node ss, tt
2. add_edge(ss, tt, INF)
3. for each edge u -> v with capacity [l, r]:
        add_edge(u, tt, l)
        add_edge(ss, v, l)
        add_edge(u, v, r-l)
4. see (b), check if it is possible.
5. answer is maxflow(ss, tt) + maxflow(s, t)
-------------------------------------------------------
(b) Bounded Possible Flow:
1. same construction method as (a)
2. run maxflow(ss, tt)
3. for every edge connected with ss or tt:
        rule: check if their rest flow is exactly 0
4. answer is possible if every edge do satisfy the rule
    ;
5. otherwise, it is NOT possible.
-------------------------------------------------------
(c) Bounded Minimum Flow:
1. same construction method as (a)
2. answer is maxflow(ss, tt)
-------------------------------------------------------
(d) Bounded Minimum Cost Flow:
* the concept is somewhat like bounded possible flow.
1. same construction method as (a)
2. answer is maxflow(ss, tt) + (∑ l * cost for every
    edge)
-------------------------------------------------------
(e) Minimum Cut:
1. run maxflow(s, t)
2. run cut(s)
3. ss[i] = 1: node i is at the same side with s.
-------------------------------------------------------

const long long INF = 1LL<<60;
struct Dinic {  //O(VVE), with minimum cut
    static const int MAXN = 5003;
    struct Edge{
        int u, v;
        long long cap, rest;
    };

    int n, m, s, t, d[MAXN], cur[MAXN];
    vector<Edge> edges;
    vector<int> G[MAXN];

    void init(){
        edges.clear();
        for ( int i = 0 ; i < n ; i++ ) G[i].clear();
        n = 0;
    }

    // min cut start
    bool side[MAXN];
    void cut(int u) {
        side[u] = 1;
        for ( int i : G[u] ) {
            if ( !side[ edges[i].v ] && edges[i].rest )
                 cut(edges[i].v);
        }
    }
    // min cut end

    int add_node(){
        return n++;
```

```
    }

    void add_edge(int u, int v, long long cap){
        edges.push_back( {u, v, cap, cap} );
        edges.push_back( {v, u, 0, 0LL} );
        m = edges.size();
        G[u].push_back(m-2);
        G[v].push_back(m-1);
    }

    bool bfs(){
        fill(d,d+n,-1);
        queue<int> que;
        que.push(s); d[s]=0;
        while (!que.empty()){
            int u = que.front(); que.pop();
            for (int ei : G[u]){
                Edge &e = edges[ei];
                if (d[e.v] < 0 && e.rest > 0){
                    d[e.v] = d[u] + 1;
                    que.push(e.v);
                }
            }
        }
        return d[t] >= 0;
    }

    long long dfs(int u, long long a){
        if ( u == t || a == 0 ) return a;
        long long flow = 0, f;
        for ( int &i=cur[u]; i < (int)G[u].size() ; i++
             ) {
            Edge &e = edges[ G[u][i] ];
            if ( d[u] + 1 != d[e.v] ) continue;
            f = dfs(e.v, min(a, e.rest) );
            if ( f > 0 ) {
                e.rest -= f;
                edges[ G[u][i]^1 ].rest += f;
                flow += f;
                a -= f;
                if ( a == 0 )break;
            }
        }
        return flow;
    }

    long long maxflow(int _s, int _t){
        s = _s, t = _t;
        long long flow = 0, mf;
        while ( bfs() ){
            fill(cur,cur+n,0);
            while ( (mf = dfs(s, INF)) ) flow += mf;
        }
        return flow;
    }
} dinic;
```

3.2 Gomory Hu

```
Construct of Gomory Hu Tree

1. make sure the whole graph is clear
2. set node 0 as root, also be the parent of other
     nodes.
3. for every node i > 0, we run maxflow from i to
     parent[i]
4. hense we know the weight between i and parent[i]
5. for each node j > i, if j is at the same side with i
   ,
   make the parent of j as i
-------------------------------------------------------

int e[MAXN][MAXN];
int p[MAXN];

Dinic D; // original graph
void gomory_hu() {
    fill(p, p+n, 0);
    fill(e[0], e[n], INF);
    for ( int s = 1 ; s < n ; s++ ) {
        int t = p[s];
        Dinic F = D;
        int tmp = F.max_flow(s, t);

        for ( int i = 1 ; i < s ; i++ )
            e[s][i] = e[i][s] = min(tmp, e[t][i]);

        for ( int i = s+1 ; i <= n ; i++ )
            if ( p[i] == t && F.side[i] ) p[i] = s;

    }
}
```

3.3 Min Cost Flow

```
#include<bits/stdc++.h>
using namespace std;
#define int long long
typedef pair<int,int> P;
struct edge{
    edge(){}
    edge(int a,int b,int c,int d):to(a),cap(b),cost(c),
        rev(d){}
    int to,cap,cost,rev;
};
#define V 210
#define inf 1000000000000000
vector<edge> g[V];
int h[V],dist[V],prev_v[V],prev_e[V];
void add_edge(int from,int to,int cap,int cost){
    g[from].push_back(edge(to,cap,cost,g[to].size()));
    g[to].push_back(edge(from,0,-cost,g[from].size()-1)
        );
}
int min_costflow(int s,int t,int f){
    int res=0;
    memset(h,0,sizeof(h));
    while(f>0){
        priority_queue<P,vector<P>,greater<P> >que;
        fill(dist,dist+V,inf);
        dist[s]=0;
        que.push(P(dist[s],s));
        while(!que.empty()){
            P p=que.top();
            que.pop();
            int v=p.second;
            if(dist[v]<p.first)continue;
            for(int i=0;i<g[v].size();++i){
                edge &e=g[v][i];
                if(e.cap>0&&dist[e.to]>dist[v]+e.cost+h
                    [v]-h[e.to]){
                    dist[e.to]=dist[v]+e.cost+h[v]-h[e.
                        to];
                    prev_v[e.to]=v;
                    prev_e[e.to]=i;
                    que.push(P(dist[e.to],e.to));
                }
            }
        }
        if(dist[t]==inf) return -1;
        for(int v=0;v<V;++v)h[v]+=dist[v];
        int d=f;
        for(int v=t;v!=s;v=prev_v[v]) d=min(d,g[prev_v[
            v]][prev_e[v]].cap);
        f-=d;
        res+=d*h[t];
        for(int v=t;v!=s;v=prev_v[v]){
            edge &e=g[prev_v[v]][prev_e[v]];
            e.cap-=d;
            g[v][e.rev].cap+=d;
        }
```

```
    }
    return res;
}
#undef int
int main()
{
#define int long long
    int T,n,m,cost,l,s,t,ans;
    cin>>T;
    while(T--){
      cin>>n>>m;
        for(int q=0;q<V;++q)g[q].clear();
        s=m+n;
        t=m+n+1;
        for(int i=0;i<n;++i)
          for(int j=0;j<m;++j){
            cin>>cost;
            if(cost>0)
              add_edge(n+j,i,1,cost);
          }
        for(int i=0;i<m;++i){
          cin>>l;
          add_edge(s,n+i,l,0);
        }
        for(int i=0;i<n;++i)
          add_edge(i,t,1,0);
        ans=min_costflow(s,t,n);
        cout<<ans<<endl;
    }
    return 0;
}
```

3.4 SW-mincut

```
// all pair min cut
// global min cut
struct SW{ // O(V^3)
  static const int MXN = 514;
  int n,vst[MXN],del[MXN];
  int edge[MXN][MXN],wei[MXN];
  void init(int _n){
    n = _n; FZ(edge); FZ(del);
  }
  void addEdge(int u, int v, int w){
    edge[u][v] += w; edge[v][u] += w;
  }
  void search(int &s, int &t){
    FZ(vst); FZ(wei);
    s = t = -1;
    while (true){
      int mx=-1, cur=0;
      for (int i=0; i<n; i++)
        if (!del[i] && !vst[i] && mx<wei[i])
          cur = i, mx = wei[i];
      if (mx == -1) break;
      vst[cur] = 1;
      s = t; t = cur;
      for (int i=0; i<n; i++)
        if (!vst[i] && !del[i]) wei[i] += edge[cur][i];
    }
  }
  int solve(){
    int res = 2147483647;
    for (int i=0,x,y; i<n-1; i++){
      search(x,y);
      res = min(res,wei[y]);
      del[y] = 1;
      for (int j=0; j<n; j++)
        edge[x][j] = (edge[j][x] += edge[y][j]);
    }
    return res;
  }
}graph;
```

4 Geometry

4.1 2Dpoint

```
typedef double Double;
struct Point {
  Double x,y;

  bool operator < (const Point &b)const{
    //return tie(x,y) < tie(b.x,b.y);
    //return atan2(y,x) < atan2(b.y,b.x);
    assert(0 && "choose compare");
  }
  Point operator + (const Point &b)const{
    return (Point){x+b.x,y+b.y};
  }
  Point operator - (const Point &b)const{
    return (Point){x-b.x,y-b.y};
  }
  Point operator * (const Double &d)const{
    return Point(d*x,d*y);
  }
  Double operator * (const Point &b)const{
    return x*b.x + y*b.y;
  }
  Double operator % (const Point &b)const{
    return x*b.y - y*b.x;
  }
  friend Double abs2(const Point &p){
    return p.x*p.x + p.y*p.y;
  }
  friend Double abs(const Point &p){
    return sqrt( abs2(p) );
  }
};
typedef Point Vector;

struct Line{
  Point P; Vector v;
  bool operator < (const Line &b)const{
    return atan2(v.y,v.x) < atan2(b.v.y,b.v.x);
  }
};
```

4.2 ConvexHull

```
#include "2Dpoint.cpp"

// retunr H, 第一個點會在 H 出現兩次
void ConvexHull(vector<Point> &P, vector<Point> &H){
    int n = P.size(), m=0;
    sort(P.begin(),P.end());
    H.clear();

    for (int i=0; i<n; i++){
        while (m>=2 && (P[i]-H[m-2]) % (H[m-1]-H[m-2])
            <0)H.pop_back(), m--;
        H.push_back(P[i]), m++;
    }

    for (int i=n-2; i>=0; i--){
        while (m>=2 && (P[i]-H[m-2]) % (H[m-1]-H[m-2])
            <0)H.pop_back(), m--;
        H.push_back(P[i]), m++;
    }
}
```

4.3 Intersection Of Two Circle

```
vector<Double> interCircle(Double o1, Double r1, Double
     o2, Double r2) {
  Double d2 = abs2(o1 - o2);
```

```
  Double d = sqrt(d2);
  if (d < fabs(r1-r2) || r1+r2 < d) return {};
  Double u = 0.5*(o1+o2) + ((r2*r2-r1*r1)/(2.0*d2))*(o1
      -o2);
  Double A = sqrt((r1+r2+d) * (r1-r2+d) * (r1+r2-d) *
      (-r1+r2+d));
  Double v = A / (2.0*d2) * Double(o1.S-o2.S, -o1.F+o2.
      F);
  return {u+v, u-v};
}
```

4.4 Intersection Of Two Lines

```
Point interPnt(Point p1, Point p2, Point q1, Point q2,
    bool &res){
  Double f1 = cross(p2, q1, p1);
  Double f2 = -cross(p2, q2, p1);
  Double f = (f1 + f2);

  if(fabs(f) < EPS) {
    res = false;
    return {};
    }

  res = true;
  return (f2 / f) * q1 + (f1 / f) * q2;
}
```

4.5 Smallest Circle

```
#include "circumcentre.cpp"
pair<Point,Double> SmallestCircle(int n, Point _p[]){
  Point *p = new Point[n];
  memcpy(p,_p,sizeof(Point)*n);
  random_shuffle(p,p+n);

  Double r2=0;
  Point cen;
  for (int i=0; i<n; i++){
    if ( abs2(cen-p[i]) <= r2)continue;
    cen = p[i], r2=0;
    for (int j=0; j<i; j++){
      if ( abs2(cen-p[j]) <= r2)continue;
      cen = (p[i]+p[j])*0.5;
      r2 = abs2(cen-p[i]);
      for (int k=0; k<j; k++){
        if ( abs2(cen-p[k]) <= r2)continue;
        cen = circumcentre(p[i],p[j],p[k]);
        r2 = abs2(cen-p[k]);
      }
    }
  }

  delete[] p;
  return {cen,r2};
}
// auto res = SmallestCircle(,);
```

4.6 Circumcentre

```
#include "2Dpoint.cpp"

Point circumcentre(Point &p0, Point &p1, Point &p2){
  Point a = p1-p0;
  Point b = p2-p0;
  Double c1 = abs2(a)*0.5;
  Double c2 = abs2(b)*0.5;
  Double d = a % b;
  Double x = p0.x + ( c1*b.y - c2*a.y ) / d;
  Double y = p0.y + ( c2*a.x - c1*b.x ) / d;
  return {x,y};
}
```

4.7 Half Plane Intersection

```
bool OnLeft(const Line& L,const Point& p){
  return Cross(L.v,p-L.P)>0;
}
Point GetIntersection(Line a,Line b){
  Vector u = a.P-b.P;
  Double t = Cross(b.v,u)/Cross(a.v,b.v);
  return a.P + a.v*t;
}
int HalfplaneIntersection(Line* L,int n,Point* poly){
  sort(L,L+n);

  int first,last;
  Point *p = new Point[n];
  Line *q = new Line[n];
  q[first=last=0] = L[0];
  for(int i=1;i<n;i++){
    while(first < last && !OnLeft(L[i],p[last-1])) last
        --;
    while(first < last && !OnLeft(L[i],p[first])) first
        ++;
    q[++last]=L[i];
    if(fabs(Cross(q[last].v,q[last-1].v))<EPS){
      last--;
      if(OnLeft(q[last],L[i].P)) q[last]=L[i];
    }
    if(first < last) p[last-1]=GetIntersection(q[last
        -1],q[last]);
  }
  while(first<last && !OnLeft(q[first],p[last-1])) last
      --;
  if(last-first<=1) return 0;
  p[last]=GetIntersection(q[last],q[first]);

  int m=0;
  for(int i=first;i<=last;i++) poly[m++]=p[i];
  return m;
}
```

5 Graph

5.1 BCC Edge

```
邊雙連通

任意兩點間至少有兩條不重疊的路徑連接，找法：
1. 標記出所有的橋
2. 對全圖進行 DFS，不走橋，每一次 DFS 就是一個新的邊雙
    連通

// from BCW

struct BccEdge {
  static const int MXN = 100005;
  struct Edge { int v,eid; };
  int n,m,step,par[MXN],dfn[MXN],low[MXN];
  vector<Edge> E[MXN];
  DisjointSet djs;
  void init(int _n) {
    n = _n; m = 0;
    for (int i=0; i<n; i++) E[i].clear();
    djs.init(n);
  }
  void add_edge(int u, int v) {
    E[u].PB({v, m});
    E[v].PB({u, m});
    m++;
  }
  void DFS(int u, int f, int f_eid) {
    par[u] = f;
    dfn[u] = low[u] = step++;
    for (auto it:E[u]) {
```

```
      if (it.eid == f_eid) continue;
      int v = it.v;
      if (dfn[v] == -1) {
        DFS(v, u, it.eid);
        low[u] = min(low[u], low[v]);
      } else {
        low[u] = min(low[u], dfn[v]);
      }
    }
  }
  void solve() {
    step = 0;
    memset(dfn, -1, sizeof(int)*n);
    for (int i=0; i<n; i++) {
      if (dfn[i] == -1) DFS(i, i, -1);
    }
    djs.init(n);
    for (int i=0; i<n; i++) {
      if (low[i] < dfn[i]) djs.uni(i, par[i]);
    }
  }
}graph;
```

5.2 Dijkstra

```
from heapq import *
INF = 2*10**10000
t = input()
for pp in range(t):
  n, m = map(int, raw_input().split())
  g, d, q = [[] for _ in range(n+1)], [0] + [INF] * n,
      [(0, 0)]
  #for i in range(1, m):
  #  a[i], b[i], c[i], l[i], o[i] = map(int, input().
      split())
  for _ in range(m):
    u, v, c, l, o = map(int, raw_input().split())
    g[u] += [(o, v, c, l)]
  while q:
    u = heappop(q)[1]
    for e in g[u]:
      k = d[u] / e[2]
      if k < 0:
        k = 0
      else:
        k = k * e[3]
      t, v = d[u] + e[0] + k, e[1]
      if t < d[v]:
        d[v] = t
        heappush(q, (d[v], v))
  print(d[n])
```

5.3 Directed MST

```
template<typename T>
struct zhu_liu{
  static const int MAXN=110,MAXM=10005;
  struct node{
    int u,v;
    T w,tag;
    node *l,*r;
    node(int u=0,int v=0,T w=0):u(u),v(v),w(w),tag(0),l
        (0),r(0){}
    void down(){
      w+=tag;
      if(l)l->tag+=tag;
      if(r)r->tag+=tag;
      tag=0;
    }
  }mem[MAXM];//靜態記憶體
  node *pq[MAXN*2],*E[MAXN*2];
  int st[MAXN*2],id[MAXN*2],m;
  void init(int n){
    for(int i=1;i<=n;++i){
      pq[i]=E[i]=0;
      st[i]=id[i]=i;
    }m=0;
  }
  node *merge(node *a,node *b){//skew heap
    if(!a||!b)return a?a:b;
    a->down(),b->down();
    if(b->w<a->w)return merge(b,a);
    swap(a->l,a->r);
    a->l=merge(b,a->l);
    return a;
  }
  void add_edge(int u,int v,T w){
    if(u!=v)pq[v]=merge(pq[v],&(mem[m++]=node(u,v,w)));
  }
  int find(int x,int *st){
    return st[x]==x?x:st[x]=find(st[x],st);
  }
  T build(int root,int n){
    T ans=0;int N=n,all=n;
    for(int i=1;i<=N;++i){
      if(i==root||!pq[i])continue;
      while(pq[i]){
        pq[i]->down(),E[i]=pq[i];
        pq[i]=merge(pq[i]->l,pq[i]->r);
        if(find(E[i]->u,id)!=find(i,id))break;
      }
      if(find(E[i]->u,id)==find(i,id))continue;
      ans+=E[i]->w;
      if(find(E[i]->u,st)==find(i,st)){
        if(pq[i])pq[i]->tag-=E[i]->w;
        pq[++N]=pq[i],id[N]=N;
        for(int u=find(E[i]->u,id);u!=i;u=find(E[u]->u,
            id)){
          if(pq[u])pq[u]->tag-=E[u]->w;
          id[find(u,id)]=N;
          pq[N]=merge(pq[N],pq[u]);
        }
        st[N]=find(i,st);
        id[find(i,id)]=N;
      }else st[find(i,st)]=find(E[i]->u,st),--all;
    }
    return all==1?ans:-INT_MAX;//圖不連通就無解
  }
};
```

5.4 LCA

```
//lv紀錄深度
//father[多少冪次][誰]
//已經建好每個人的父親是誰 (father[0][i]已經建好)
//已經建好深度 (lv[i]已經建好)
void makePP(){
  for(int i = 1; i < 20; i++){
    for(int j = 2; j <= n; j++){
      father[i][j]=father[i-1][ father[i-1][j] ];
    }
  }
}
int find(int a, int b){
  if(lv[a] < lv[b]) swap(a,b);
  int need = lv[a] - lv[b];
  for(int i = 0; need!=0; i++){
    if(need&1) a=father[i][a];
    need >>= 1;
  }
  for(int i = 19 ;i >= 0 ;i--){
    if(father[i][a] != father[i][b]){
      a=father[i][a];
      b=father[i][b];
    }
  }
  return a!=b?father[0][a] : a;
}
```

5.5 MaximumClique

```
const int MAXN = 105;
int best;
int m ,n;
int num[MAXN];
// int x[MAXN];
int path[MAXN];
int g[MAXN][MAXN];

bool dfs( int *adj, int total, int cnt ){
    int i, j, k;
    int t[MAXN];
    if( total == 0 ){
        if( best < cnt ){
            // for( i = 0; i < cnt; i++) path[i] = x[i
                ];
            best = cnt; return true;
        }
        return false;
    }
    for( i = 0; i < total; i++){
        if( cnt+(total-i) <= best ) return false;
        if( cnt+num[adj[i]] <= best ) return false;
        // x[cnt] = adj[i];
        for( k = 0, j = i+1; j < total; j++ )
            if( g[ adj[i] ][ adj[j] ] )
                t[ k++ ] = adj[j];
        if( dfs( t, k, cnt+1 ) ) return true;
    } return false;
}
int MaximumClique(){
    int i, j, k;
    int adj[MAXN];
    if( n <= 0 ) return 0;
    best = 0;
    for( i = n-1; i >= 0; i-- ){
        // x[0] = i;
        for( k = 0, j = i+1; j < n; j++ )
            if( g[i][j] ) adj[k++] = j;
        dfs( adj, k, 1 );
        num[i] = best;
    }
    return best;
}
```

5.6 Min Mean Cycle

```
// from BCW

/* minimum mean cycle */
const int MAXE = 1805;
const int MAXN = 35;
const double inf = 1029384756;
const double eps = 1e-6;
struct Edge {
  int v,u;
  double c;
};
int n,m,prv[MAXN][MAXN], prve[MAXN][MAXN], vst[MAXN];
Edge e[MAXE];
vector<int> edgeID, cycle, rho;
double d[MAXN][MAXN];
inline void bellman_ford() {
  for(int i=0; i<n; i++) d[0][i]=0;
  for(int i=0; i<n; i++) {
    fill(d[i+1], d[i+1]+n, inf);
    for(int j=0; j<m; j++) {
      int v = e[j].v, u = e[j].u;
      if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
        d[i+1][u] = d[i][v]+e[j].c;
        prv[i+1][u] = v;
        prve[i+1][u] = j;
      }
    }
  }
}
double karp_mmc() {
  // returns inf if no cycle, mmc otherwise
  double mmc=inf;
  int st = -1;
  bellman_ford();
  for(int i=0; i<n; i++) {
    double avg=-inf;
    for(int k=0; k<n; k++) {
      if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i])
          /(n-k));
      else avg=max(avg,inf);
    }
    if (avg < mmc) tie(mmc, st) = tie(avg, i);
  }
  for(int i=0; i<n; i++) vst[i] = 0;
  edgeID.clear(); cycle.clear(); rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
    vst[st]++;
    edgeID.PB(prve[i][st]);
    rho.PB(st);
  }
  while (vst[st] != 2) {
    int v = rho.back(); rho.pop_back();
    cycle.PB(v);
    vst[v]++;
  }
  reverse(ALL(edgeID));
  edgeID.resize(SZ(cycle));
  return mmc;
}
```

5.7 MinimumSteinerTree

```
// Minimum Steiner Tree
// O(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
  int n , dst[V][V] , dp[1 << T][V] , tdst[V];
  void init( int _n ){
    n = _n;
    for( int i = 0 ; i < n ; i ++ ){
      for( int j = 0 ; j < n ; j ++ )
        dst[ i ][ j ] = INF;
      dst[ i ][ i ] = 0;
    }
  }
  void add_edge( int ui , int vi , int wi ){
    dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
    dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
  }
  void shortest_path(){
    for( int k = 0 ; k < n ; k ++ )
      for( int i = 0 ; i < n ; i ++ )
        for( int j = 0 ; j < n ; j ++ )
          dst[ i ][ j ] = min( dst[ i ][ j ],
                dst[ i ][ k ] + dst[ k ][ j ] );
  }
  int solve( const vector<int>& ter ){
    int t = (int)ter.size();
    for( int i = 0 ; i < ( 1 << t ) ; i ++ )
      for( int j = 0 ; j < n ; j ++ )
        dp[ i ][ j ] = INF;
    for( int i = 0 ; i < n ; i ++ )
      dp[ 0 ][ i ] = 0;
    for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){
      if( msk == ( msk & (-msk) ) ){
        int who = __lg( msk );
        for( int i = 0 ; i < n ; i ++ )
          dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];
        continue;
      }
      for( int i = 0 ; i < n ; i ++ )
```

```
        for( int submsk = ( msk - 1 ) & msk ; submsk ;
                 submsk = ( submsk - 1 ) & msk )
          dp[ msk ][ i ] = min( dp[ msk ][ i ],
                                dp[ submsk ][ i ] +
                                dp[ msk ^ submsk ][ i ] );
      for( int i = 0 ; i < n ; i ++ ){
        tdst[ i ] = INF;
        for( int j = 0 ; j < n ; j ++ )
          tdst[ i ] = min( tdst[ i ],
                           dp[ msk ][ j ] + dst[ j ][ i ] );
      }
      for( int i = 0 ; i < n ; i ++ )
        dp[ msk ][ i ] = tdst[ i ];
    }
    int ans = INF;
    for( int i = 0 ; i < n ; i ++ )
      ans = min( ans , dp[ ( 1 << t ) - 1 ][ i ] );
    return ans;
  }
} solver;
```

5.8 Tarjan

```
割點
點 u 為割點 if and only if 滿足 1. or 2.
1. u 爲樹根，且 u 有多於一個子樹。
2. u 不爲樹根，且滿足存在 (u,v) 爲樹枝邊 (或稱父子邊，
    即 u 爲 v 在搜索樹中的父親)，使得 DFN(u) <= Low(v)
    。

-------------------------------------------------------
橋
一條無向邊 (u,v) 是橋 if and only if (u,v) 爲樹枝邊，且
    滿足 DFN(u) < Low(v)。

// 0 base
struct TarjanSCC{
  static const int MAXN = 1000006;
  int n, dfn[MAXN], low[MAXN], scc[MAXN], scn, count;
  vector<int> G[MAXN];
  stack<int> stk;
  bool ins[MAXN];

  void tarjan(int u){
    dfn[u] = low[u] = ++count;
    stk.push(u);
    ins[u] = true;

    for(auto v:G[u]){
      if(!dfn[v]){
        tarjan(v);
        low[u] = min(low[u], low[v]);
      }else if(ins[v]){
        low[u] = min(low[u], dfn[v]);
      }
    }

    if(dfn[u] == low[u]){
      int v;
      do {
      v = stk.top();
      stk.pop();
      scc[v] = scn;
      ins[v] = false;
      } while(v != u);
      scn++;
    }
  }

  void getSCC(){
    memset(dfn,0,sizeof(dfn));
    memset(low,0,sizeof(low));
    memset(ins,0,sizeof(ins));
    memset(scc,0,sizeof(scc));
    count = scn = 0;
    for(int i = 0 ; i < n ; i++ ){
      if(!dfn[i]) tarjan(i);
    }
  }

}SCC;
```

5.9 TwoSAT

```
const int MAXN = 2020;

struct TwoSAT{
    static const int MAXv = 2*MAXN;
    vector<int> GO[MAXv],BK[MAXv],stk;
    bool vis[MAXv];
    int SC[MAXv];

    void imply(int u,int v){ // u imply v
        GO[u].push_back(v);
        BK[v].push_back(u);
    }
    int dfs(int u,vector<int>*G,int sc){
        vis[u]=1, SC[u]=sc;
        for (int v:G[u])if (!vis[v])
            dfs(v,G,sc);
        if (G==GO)stk.push_back(u);
    }
    int scc(int n=MAXv){
        memset(vis,0,sizeof(vis));
        for (int i=0; i<n; i++)if (!vis[i])
            dfs(i,GO,-1);
        memset(vis,0,sizeof(vis));
        int sc=0;
        while (!stk.empty()){
            if (!vis[stk.back()])
                dfs(stk.back(),BK,sc++);
            stk.pop_back();
        }
    }
}SAT;

int main(){
    SAT.scc(2*n);
    bool ok=1;
    for (int i=0; i<n; i++){
        if (SAT.SC[2*i]==SAT.SC[2*i+1])ok=0;
    }
    if (ok){
        for (int i=0; i<n; i++){
            if (SAT.SC[2*i]>SAT.SC[2*i+1]){
                cout << i << endl;
            }
        }
    }
    else puts("NO");
}
```

6 Matching

6.1 KM

```
#define MAXN 100
#define INF INT_MAX
int g[MAXN][MAXN],lx[MAXN],ly[MAXN],slack_y[MAXN];
int px[MAXN],py[MAXN],match_y[MAXN],par[MAXN];
int n;
void adjust(int y){//把增廣路上所有邊反轉
  match_y[y]=py[y];
  if(px[match_y[y]]!=-2)
    adjust(px[match_y[y]]);
}
bool dfs(int x){//DFS找增廣路
```

```
  for(int y=0;y<n;++y){
    if(py[y]!=-1)continue;
    int t=lx[x]+ly[y]-g[x][y];
    if(t==0){
      py[y]=x;
      if(match_y[y]==-1){
        adjust(y);
        return 1;
      }
      if(px[match_y[y]]!=-1)continue;
      px[match_y[y]]=y;
      if(dfs(match_y[y]))return 1;
    }else if(slack_y[y]>t){
      slack_y[y]=t;
      par[y]=x;
    }
  }
  return 0;
}
inline int km(){
  memset(ly,0,sizeof(int)*n);
  memset(match_y,-1,sizeof(int)*n);
  for(int x=0;x<n;++x){
    lx[x]=-INF;
    for(int y=0;y<n;++y){
      lx[x]=max(lx[x],g[x][y]);
    }
  }
  for(int x=0;x<n;++x){
    for(int y=0;y<n;++y)slack_y[y]=INF;
    memset(px,-1,sizeof(int)*n);
    memset(py,-1,sizeof(int)*n);
    px[x]=-2;
    if(dfs(x))continue;
    bool flag=1;
    while(flag){
      int cut=INF;
      for(int y=0;y<n;++y)
        if(py[y]==-1&&cut>slack_y[y])cut=slack_y[y];
      for(int j=0;j<n;++j){
        if(px[j]!=-1)lx[j]-=cut;
        if(py[j]!=-1)ly[j]+=cut;
        else slack_y[j]-=cut;
      }
      for(int y=0;y<n;++y){
        if(py[y]==-1&&slack_y[y]==0){
          py[y]=par[y];
          if(match_y[y]==-1){
            adjust(y);
            flag=0;
            break;
          }
          px[match_y[y]]=y;
          if(dfs(match_y[y])){
            flag=0;
            break;
          }
        }
      }
    }
  }
  int ans=0;
  for(int y=0;y<n;++y)if(g[match_y[y]][y]!=-INF)ans+=g[
      match_y[y]][y];
  return ans;
}
```

6.2 Maximum General Matching

```
// Maximum Cardinality Matching

struct Graph {
  vector<int> G[MAXN];
  int pa[MAXN], match[MAXN], st[MAXN], S[MAXN], vis[
      MAXN];
  int t, n;

  void init(int _n) {
    n = _n;
    for ( int i = 1 ; i <= n ; i++ ) G[i].clear();
  }
  void add_edge(int u, int v) {
    G[u].push_back(v);
    G[v].push_back(u);
  }
  int lca(int u, int v){
    for ( ++t ; ; swap(u, v) ) {
      if ( u == 0 ) continue;
      if ( vis[u] == t ) return u;
      vis[u] = t;
      u = st[ pa[ match[u] ] ];
    }
  }
  void flower(int u, int v, int l, queue<int> &q) {
    while ( st[u] != l ) {
      pa[u] = v;
      if ( S[ v = match[u] ] == 1 ) {
        q.push(v);
        S[v] = 0;
      }
      st[u] = st[v] = l;
      u = pa[v];
    }
  }
  bool bfs(int u){
    for ( int i = 1 ; i <= n ; i++ ) st[i] = i;
    memset(S, -1, sizeof(S));
    queue<int>q;
    q.push(u);
    S[u] = 0;
    while ( !q.empty() ) {
      u = q.front(); q.pop();
      for ( int i = 0 ; i < (int)G[u].size(); i++) {
        int v = G[u][i];
        if ( S[v] == -1 ) {
          pa[v] = u;
          S[v] = 1;
          if ( !match[v] ) {
            for ( int lst ; u ; v = lst, u = pa[v] ) {
              lst = match[u];
              match[u] = v;
              match[v] = u;
            }
            return 1;
          }
          q.push(match[v]);
          S[ match[v] ] = 0;
        } else if ( !S[v] && st[v] != st[u] ) {
          int l = lca(st[v], st[u]);
          flower(v, u, l, q);
          flower(u, v, l, q);
        }
      }
    }
    return 0;
  }
  int solve(){
    memset(pa, 0, sizeof(pa));
    memset(match, 0, sizeof(match));
    int ans = 0;
    for ( int i = 1 ; i <= n ; i++ )
      if ( !match[i] && bfs(i) ) ans++;
    return ans;
  }
} graph;
```

6.3 Minimum General Weighted Matching

```
// Minimum Weight Perfect Matching (Perfect Match)

struct Graph {
    static const int MAXN = 105;
```

```
    int n, e[MAXN][MAXN];
    int match[MAXN], d[MAXN], onstk[MAXN];
    vector<int> stk;
    void init(int _n) {
        n = _n;
        for( int i = 0 ; i < n ; i ++ )
            for( int j = 0 ; j < n ; j ++ )
                e[i][j] = 0;
    }
    void add_edge(int u, int v, int w) {
        e[u][v] = e[v][u] = w;
    }
    bool SPFA(int u){
        if (onstk[u]) return true;
        stk.push_back(u);
        onstk[u] = 1;
        for ( int v = 0 ; v < n ; v++ ) {
            if (u != v && match[u] != v && !onstk[v] )
                {
                int m = match[v];
                if ( d[m] > d[u] - e[v][m] + e[u][v] )
                    {
                    d[m] = d[u] - e[v][m] + e[u][v];
                    onstk[v] = 1;
                    stk.push_back(v);
                    if (SPFA(m)) return true;
                    stk.pop_back();
                    onstk[v] = 0;
                }
            }
        }
        onstk[u] = 0;
        stk.pop_back();
        return false;
    }
    int solve() {
        for ( int i = 0 ; i < n ; i += 2 ) {
            match[i] = i+1;
            match[i+1] = i;
        }
        while (true){
            int found = 0;
            for ( int i = 0 ; i < n ; i++ )
                onstk[ i ] = d[ i ] = 0;
            for ( int i = 0 ; i < n ; i++ ) {
                stk.clear();
                if ( !onstk[i] && SPFA(i) ) {
                    found = 1;
                    while ( stk.size() >= 2 ) {
                        int u = stk.back(); stk.
                            pop_back();
                        int v = stk.back(); stk.
                            pop_back();
                        match[u] = v;
                        match[v] = u;
                    }
                }
            }
            if (!found) break;
        }
        int ret = 0;
        for ( int i = 0 ; i < n ; i++ )
            ret += e[i][match[i]];
        ret /= 2;
        return ret;
    }
} graph;
```

6.4 Stable Marriage

```
#define F(n) Fi(i, n)
#define Fi(i, n) Fl(i, 0, n)
#define Fl(i, l, n) for(int i = l ; i < n ; ++i)
#include <bits/stdc++.h>
using namespace std;
int D, quota[205], weight[205][5];
int S, scoretodep[12005][205], score[5];
int P, prefer[12005][85], iter[12005];
int ans[12005];
typedef pair<int, int> PII;
map<int, int> samescore[205];
typedef priority_queue<PII, vector<PII>, greater<PII>>
    QQQ;
QQQ pri[205];
void check(int d) {
  PII t = pri[d].top();
  int v;
  if (pri[d].size() - samescore[d][t.first] + 1 <=
      quota[d]) return;
  while (pri[d].top().first == t.first) {
    v = pri[d].top().second;
    ans[v] = -1;
    --samescore[d][t.first];
    pri[d].pop();
  }
}
void push(int s, int d) {
  if (pri[d].size() < quota[d]) {
    pri[d].push(PII(scoretodep[s][d], s));
    ans[s] = d;
    ++samescore[s][scoretodep[s][d]];
  } else if (scoretodep[s][d] >= pri[d].top().first) {
    pri[d].push(PII(scoretodep[s][d], s));
    ans[s] = d;
    ++samescore[s][scoretodep[s][d]];
    check(d);
  }
}
void f() {
  int over;
  while (true) {
    over = 1;
    Fi (q, S) {
      if (ans[q] != -1 || iter[q] >= P) continue;
      push(q, prefer[q][iter[q]++]);
      over = 0;
    }
    if (over) break;
  }
}
main() {
  ios::sync_with_stdio(false);
  cin.tie(NULL);
  int sadmit, stof, dexceed, dfew;
  while (cin >> D, D) { // Beware of the input format
      or judge may troll us.
    sadmit = stof = dexceed = dfew = 0;
    memset(iter, 0, sizeof(iter));
    memset(ans, 0, sizeof(ans));
    Fi (q, 205) {
      pri[q] = QQQ();
      samescore[q].clear();
    }
    cin >> S >> P;
    Fi (q, D) {
      cin >> quota[q];
      Fi (w, 5) cin >> weight[q][w];
    }
    Fi (q, S) {
      Fi (w, 5) cin >> score[w];
      Fi (w, D) {
        scoretodep[q][w] = 0;
        F (5) scoretodep[q][w] += weight[w][i] * score[
            i];
      }
    }
    Fi (q, S) Fi (w, P) {
      cin >> prefer[q][w];
      --prefer[q][w];
    }
    f();
    Fi (q, D) sadmit += pri[q].size();
    Fi (q, S) if (ans[q] == prefer[q][0]) ++stof;
```

```
    Fi (q, D) if (pri[q].size() > quota[q]) ++dexceed;
    Fi (q, D) if (pri[q].size() < quota[q]) ++dfew;
    cout << sadmit << ' ' << stof << ' ' << dexceed <<
        ' ' << dfew << '\n';
  }
}
```

7 Math

7.1 FFT

```
// use llround() to avoid EPS
typedef double Double;
const Double PI = acos(-1);

// STL complex may TLE
typedef complex<Double> Complex;
#define x real()
#define y imag()

template<typename Iter> // Complex*
void BitReverse(Iter a, int n){
    for (int i=1, j=0; i<n; i++){
        for (int k = n>>1; k>(j^=k); k>>=1);
        if (i<j) swap(a[i],a[j]);
    }
}

template<typename Iter> // Complex*
void FFT(Iter a, int n, int rev=1){ // rev = 1 or -1
    assert( (n&(-n)) == n ); // n is power of 2
    BitReverse(a,n);
    Iter A = a;

    for (int s=1; (1<<s)<=n; s++){
        int m = (1<<s);

        Complex wm( cos(2*PI*rev/m), sin(2*PI*rev/m) );
        for (int k=0; k<n; k+=m){
            Complex w(1,0);
            for (int j=0; j<(m>>1); j++){
                Complex t = w * A[k+j+(m>>1)];
                Complex u = A[k+j];
                A[k+j] = u+t;
                A[k+j+(m>>1)] = u-t;
                w = w*wm;
            }
        }
    }

    if (rev==-1){
        for (int i=0; i<n; i++){
            A[i] /= n;
        }
    }
}
```

7.2 GaussElimination

```
// by bcw_codebook

const int MAXN = 300;
const double EPS = 1e-8;

int n;
double A[MAXN][MAXN];

void Gauss() {
  for(int i = 0; i < n; i++) {
    bool ok = 0;
    for(int j = i; j < n; j++) {
      if(fabs(A[j][i]) > EPS) {
        swap(A[j], A[i]);
        ok = 1;
        break;
      }
    }
    if(!ok) continue;

    double fs = A[i][i];
    for(int j = i+1; j < n; j++) {
      double r = A[j][i] / fs;
      for(int k = i; k < n; k++) {
        A[j][k] -= A[i][k] * r;
      }
    }
  }
}
```

7.3 Karatsuba

```
// N is power of 2
template<typename Iter>
void DC(int N, Iter tmp, Iter A, Iter B, Iter res){
    fill(res,res+2*N,0);
    if (N<=32){
        for (int i=0; i<N; i++){
            for (int j=0; j<N; j++){
                res[i+j] += A[i]*B[j];
            }
        }
        return;
    }
    int n = N/2;
    auto a = A+n, b = A;
    auto c = B+n, d = B;
    DC(n,tmp+N,a,c,res+2*N);
    for (int i=0; i<N; i++){
        res[i+N] += res[2*N+i];
        res[i+n] -= res[2*N+i];
    }
    DC(n,tmp+N,b,d,res+2*N);
    for (int i=0; i<N; i++){
        res[i] += res[2*N+i];
        res[i+n] -= res[2*N+i];
    }

    auto x = tmp;
    auto y = tmp+n;
    for (int i=0; i<n; i++) x[i] = a[i]+b[i];
    for (int i=0; i<n; i++) y[i] = c[i]+d[i];
    DC(n,tmp+N,x,y,res+2*N);
    for (int i=0; i<N; i++){
        res[i+n] += res[2*N+i];
    }
}
// DC(1<<16,tmp.begin(),A.begin(),B.begin(),res.begin()
    );
```

7.4 LinearPrime

```
const int MAXP = 100; //max prime
vector<int> P;  // primes
void build_prime(){
  static bitset<MAXP> ok;
  int np=0;
  for (int i=2; i<MAXP; i++){
    if (ok[i]==0)P.push_back(i), np++;
    for (int j=0; j<np && i*P[j]<MAXP; j++){
      ok[ i*P[j] ] = 1;
      if ( i%P[j]==0 )break;
    }
  }
}
```

7.5 Miller-Rabin

```
typedef long long LL;

inline LL bin_mul(LL a, LL n,const LL& MOD){
  LL re=0;
  while (n>0){
    if (n&1) re += a;
    a += a; if (a>=MOD) a-=MOD;
    n>>=1;
  }
  return re%MOD;
}

inline LL bin_pow(LL a, LL n,const LL& MOD){
  LL re=1;
  while (n>0){
    if (n&1) re = bin_mul(re,a,MOD);
    a = bin_mul(a,a,MOD);
    n>>=1;
  }
  return re;
}

bool is_prime(LL n){
  //static LL sprp[3] = { 2LL, 7LL, 61LL};
  static LL sprp[7] = { 2LL, 325LL, 9375LL,
    28178LL, 450775LL, 9780504LL,
    1795265022LL };
  if (n==1 || (n&1)==0 ) return n==2;
  int u=n-1, t=0;
  while ( (u&1)==0 ) u>>=1, t++;
  for (int i=0; i<3; i++){
    LL x = bin_pow( sprp[i]%n, u, n);
    if (x==0 || x==1 || x==n-1)continue;

    for (int j=1; j<t; j++){
      x=x*x%n;
      if (x==1 || x==n-1)break;
    }
    if (x==n-1)continue;
    return 0;
  }
  return 1;
}
```

7.6 Mobius

```
void mobius() {
    fill(isPrime, isPrime + MAXN, 1);
    mu[1] = 1, num = 0;
    for (int i = 2; i < MAXN; ++i) {
        if (isPrime[i]) primes[num++] = i, mu[i] = -1;
        static int d;
        for (int j = 0; j < num && (d = i * primes[j])
            < MAXN; ++j) {
            isPrime[d] = false;
            if (i % primes[j] == 0) {
                mu[d] = 0; break;
            } else mu[d] = -mu[i];
        }
    }
}
```

7.7 Simplex

```
// Two-phase simplex algorithm for solving linear
    programs of the form
//
//     maximize     c^T x
//     subject to   Ax <= b
//                  x >= 0
//
// INPUT: A -- an m x n matrix
//        b -- an m-dimensional vector
//        c -- an n-dimensional vector
//        x -- a vector where the optimal solution will
    be stored
//
// OUTPUT: value of the optimal solution (infinity if
    unbounded
//         above, nan if infeasible)
//
// To use this code, create an LPSolver object with A,
    b, and c as
// arguments.  Then, call Solve(x).

#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>

using namespace std;

typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;

const DOUBLE EPS = 1e-9;

struct LPSolver {
  int m, n;
  VI B, N;
  VVD D;

  LPSolver(const VVD &A, const VD &b, const VD &c) :
    m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2,
        VD(n + 2)) {
    for (int i = 0; i < m; i++) for (int j = 0; j < n;
        j++) D[i][j] = A[i][j];
    for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n]
         = -1; D[i][n + 1] = b[i]; }
    for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -
        c[j]; }
    N[n] = -1; D[m + 1][n] = 1;
  }

  void Pivot(int r, int s) {
    double inv = 1.0 / D[r][s];
    for (int i = 0; i < m + 2; i++) if (i != r)
      for (int j = 0; j < n + 2; j++) if (j != s)
        D[i][j] -= D[r][j] * D[i][s] * inv;
    for (int j = 0; j < n + 2; j++) if (j != s) D[r][j]
         *= inv;
    for (int i = 0; i < m + 2; i++) if (i != r) D[i][s]
         *= -inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
  }

  bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
    while (true) {
      int s = -1;
      for (int j = 0; j <= n; j++) {
        if (phase == 2 && N[j] == -1) continue;
        if (s == -1 || D[x][j] < D[x][s] || D[x][j] ==
            D[x][s] && N[j] < N[s]) s = j;
      }
      if (D[x][s] > -EPS) return true;
      int r = -1;
      for (int i = 0; i < m; i++) {
        if (D[i][s] < EPS) continue;
        if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n +
            1] / D[r][s] ||
          (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r
              ][s]) && B[i] < B[r]) r = i;
      }
```

```
      if (r == -1) return false;
      Pivot(r, s);
    }
  }

  DOUBLE Solve(VD &x) {
    int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][
        n + 1]) r = i;
    if (D[r][n + 1] < -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return
           -numeric_limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1;
        for (int j = 0; j <= n; j++)
          if (s == -1 || D[i][j] < D[i][s] || D[i][j]
              == D[i][s] && N[j] < N[s]) s = j;
        Pivot(i, s);
      }
    }
    if (!Simplex(2)) return numeric_limits<DOUBLE>::
        infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] =
         D[i][n + 1];
    return D[m][n + 1];
  }
};

int main() {

  const int m = 4;
  const int n = 3;
  DOUBLE _A[m][n] = {
    { 6, -1, 0 },
    { -1, -5, 0 },
    { 1, 5, 1 },
    { -1, -5, -1 }
  };
  DOUBLE _b[m] = { 10, -4, 5, -5 };
  DOUBLE _c[n] = { 1, -1, 0 };

  VVD A(m);
  VD b(_b, _b + m);
  VD c(_c, _c + n);
  for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] +
      n);

  LPSolver solver(A, b, c);
  VD x;
  DOUBLE value = solver.Solve(x);

  cerr << "VALUE: " << value << endl; // VALUE: 1.29032
  cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
  for (size_t i = 0; i < x.size(); i++) cerr << " " <<
      x[i];
  cerr << endl;
  return 0;
}
```

7.8 Sprague-Grundy

```
Anti Nim (取走最後一個石子者敗)

先手必勝 if and only if
1. 「所有」堆的石子數都為 1 且遊戲的 SG 值為 0。
2. 「有些」堆的石子數大於 1 且遊戲的 SG 值不為 0。

-------------------------------------------------------
Anti-SG (決策集合為空的遊戲者贏)

定義 SG 值為 0 時，遊戲結束，
則先手必勝 if and only if
1. 遊戲中沒有單一遊戲的 SG 函數大於 1 且遊戲的 SG 函數
    為 0。
2. 遊戲中某個單一遊戲的 SG 函數大於 1 且遊戲的 SG 函數
    不為 0。

-------------------------------------------------------
Sprague-Grundy

1. 雙人、回合制
2. 資訊完全公開
3. 無隨機因素
4. 可在有限步內結束
5. 沒有和局
6. 雙方可採取的行動相同

SG(S) 的值為 0：後手(P)必勝
不為 0：先手(N)必勝

int mex(set S) {
  // find the min number >= 0 that not in the S
  // e.g. S = {0, 1, 3, 4} mex(S) = 2
}

state = []
int SG(A) {
  if (A not in state) {
    S = sub_states(A)
    if( len(S) > 1 ) state[A] = reduce(operator.xor, [
        SG(B) for B in S])
    else state[A] = mex(set(SG(B) for B in next_states(
        A)))
  }
  return state[A]
}
```

7.9 Ax+by=gcd

```
pair<int,int> extgcd(int a, int b){
    if (b==0) return {1,0};
    int k = a/b;
    pair<int,int> p = extgcd(b,a-k*b);
    return { p.second, p.first - k*p.second };
}
int inv[maxN];
LL invtable(int n,LL P){
  inv[1]=1;
  for(int i=2;i<n;++i)
    inv[i]=(P-(P/i))*inv[P%i]%P;
}
```

7.10 PollardRho

```
// does not work when n is prime
inline LL f(LL x , LL mod) {
return (x * x % mod + 1) % mod;
}
inline LL pollard_rho(LL n) {
  if(!(n&1)) return 2;
  while(true) {
    LL y = 2 , x = rand() % (n - 1) + 1 , res = 1;
    for(int sz = 2; res == 1; sz *= 2) {
      for(int i = 0; i < sz && res <= 1; i++) {
        x = f(x , n);
        res = __gcd(abs(x - y) , n);
      }
      y = x;
    }
    if (res != 0 && res != n) return res;
  }
}
```

7.11 Theorem

```
/*
Lucas's Theorem
  For non-negative integer n,m and prime P,
  C(m,n) mod P = C(m/M,n/M) * C(m%M,n%M) mod P
  = mult_i ( C(m_i,n_i) )
  where m_i is the i-th digit of m in base P.
-------------------------------------------------------
Kirchhoff's theorem
  A_{ii} = deg(i), A_{ij} = (i,j) \in E ? -1 : 0
  Deleting any one row, one column, and cal the det(A)
-------------------------------------------------------
Nth Catalan recursive function:
C_0 = 1, C_{n+1} = C_n * 2(2n + 1)/(n+2)
-------------------------------------------------------
Mobius Formula
u(n) = 1          , if n = 1
       (-1)^m     , 若 n 無平方數因數，且 n = p1*p2*p3
           *...*pk
       0          , 若 n 有大於 1 的平方數因數
- Property
1. (積性函數) u(a)u(b) = u(ab)
2. ∑_{d|n} u(d) = [n == 1]
-------------------------------------------------------
Mobius Inversion Formula
if      f(n) = ∑_{d|n} g(d)
then    g(n) = ∑_{d|n} u(n/d)f(d)
             = ∑_{d|n} u(d)f(n/d)
- Application
the number/power of gcd(i, j) = k
- Trick
分塊, O(sqrt(n))
-------------------------------------------------------
Chinese Remainder Theorem (m_i 兩兩互質)

  x = a_1 (mod m_1)
  x = a_2 (mod m_2)
  ....
  x = a_i (mod m_i)

construct a solution:

  Let M = m_1 * m_2 * m_3 * ... * m_n
  Let M_i = M / m_i

  t_i = 1 / M_i
  t_i * M_i = 1 (mod m_i)

  solution x = a_1 * t_1 * M_1 + a_2 * t_2 * M_2 + ...
      + a_n * t_n * M_n + k * M
  = k*M + ∑ a_i * t_i * M_i, k is positive integer.

  under mod M, there is one solution x = ∑ a_i * t_i *
      M_i
-------------------------------------------------------
Burnside's lemma
|G| * |X/G|  = sum( |X^g| ) where g in G
總方法數: 每一種旋轉下不動點的個數總和 除以 旋轉的方法
    數
-------------------------------------------------------
Lagrange multiplier
f(x,y) 求極值。必須滿足 g(x,y) = 0。

湊得 f(x,y) = f(x,y) + λ g(x,y)
定義 s(x,y,λ) = f(x,y) + λ g(x,y)

f(x,y) 的極值，等同 s(x,y,λ) = f(x,y) + λ g(x,y) 的極
    值。
欲求極值：
對 x 偏微分，讓斜率是 0。
對 y 偏微分，讓斜率是 0。

不管 λ 如何變化，λ g(x,y) 都是零，s(x,y,λ) 永遠不變。
欲求永遠不變的地方：
對 λ 偏微分，讓斜率是 0。

三道偏微分方程式聯立之後，其解涵蓋了（不全是）所有符合
    約束條件的極值。
{ ∂/∂x s(x,y,λ) = 0
{ ∂/∂y s(x,y,λ) = 0
{ ∂/∂λ s(x,y,λ) = 0
*/
```

8 Other

8.1 CYK

```
// 2016 NCPC from sunmoon

// 轉換

#define MAXN 55
struct CNF{
  int s,x,y;//s->xy | s->x, if y==-1
  int cost;
  CNF(){}
  CNF(int s,int x,int y,int c):s(s),x(x),y(y),cost(c){}
};
int state;//規則數量
map<char,int> rule;//每個字元對應到的規則，小寫字母為終
    端字符
vector<CNF> cnf;
inline void init(){
  state=0;
  rule.clear();
  cnf.clear();
}
inline void add_to_cnf(char s,const string &p,int cost)
    {
  if(rule.find(s)==rule.end())rule[s]=state++;
  for(auto c:p)if(rule.find(c)==rule.end())rule[c]=
      state++;
  if(p.size()==1){
    cnf.push_back(CNF(rule[s],rule[p[0]],-1,cost));
  }else{
    int left=rule[s];
    int sz=p.size();
    for(int i=0;i<sz-2;++i){
      cnf.push_back(CNF(left,rule[p[i]],state,0));
      left=state++;
    }
    cnf.push_back(CNF(left,rule[p[sz-2]],rule[p[sz-1]],
        cost));
  }
}

// 計算

vector<long long> dp[MAXN][MAXN];
vector<bool> neg_INF[MAXN][MAXN];//如果花費是負的可能會
    有無限小的情形
inline void relax(int l,int r,const CNF &c,long long
    cost,bool neg_c=0){
  if(!neg_INF[l][r][c.s]&&(neg_INF[l][r][c.x]||cost<dp[
      l][r][c.s])){
    if(neg_c||neg_INF[l][r][c.x]){
      dp[l][r][c.s]=0;
      neg_INF[l][r][c.s]=true;
    }else dp[l][r][c.s]=cost;
  }
}
inline void bellman(int l,int r,int n){
  for(int k=1;k<=state;++k)
    for(auto c:cnf)
      if(c.y==-1)relax(l,r,c,dp[l][r][c.x]+c.cost,k==n)
          ;
}
```

```cpp
inline void cyk(const vector<int> &tok){
  for(int i=0;i<(int)tok.size();++i){
    for(int j=0;j<(int)tok.size();++j){
      dp[i][j]=vector<long long>(state+1,INT_MAX);
      neg_INF[i][j]=vector<bool>(state+1,false);
    }
    dp[i][i][tok[i]]=0;
    bellman(i,i,tok.size());
  }
  for(int r=1;r<(int)tok.size();++r){
    for(int l=r-1;l>=0;--l){
      for(int k=l;k<r;++k)
        for(auto c:cnf)
          if(~c.y)relax(l,r,c,dp[l][k][c.x]+dp[k+1][r][
              c.y]+c.cost);
      bellman(l,r,tok.size());
    }
  }
}
```

8.2 DP-optimization

```
Monotonicity & 1D/1D DP & 2D/1D DP
------------------------------------------------------------
Definition xD/yD
1D/1D DP[j] = min(0≤i<j) { DP[i] + w(i, j) }; DP[0] = k
2D/1D DP[i][j] = min(i<k≤j) { DP[i][k - 1] + DP[k][j] }
     + w(i, j); DP[i][i] = 0
------------------------------------------------------------
Monotonicity
       c       d
  -----------------
a | w(a, c) w(a, d)
b | w(b, c) w(b, d)

Monge Condition
Concave(凹四邊形不等式): w(a, c) + w(b, d) >= w(a, d) +
     w(b, c)
Convex (凸四邊形不等式): w(a, c) + w(b, d) <= w(a, d) +
     w(b, c)

Totally Monotone
Concave(凹單調): w(a, c) <= w(b, d) -----> w(a, d) <= w
    (b, c)
Convex (凸單調): w(a, c) >= w(b, d) -----> w(a, d) >= w
    (b, c)
------------------------------------------------------------
1D/1D DP O(n^2) -> O(nlgn)
**CONSIDER THE TRANSITION POINT**
Solve 1D/1D Concave by Stack
Solve 1D/1D Convex by Deque
------------------------------------------------------------
2D/1D Convex DP (Totally Monotone) O(n^3) -> O(n^2)
h(i, j - 1) ≤ h(i, j) ≤ h(i + 1, j)
```

8.3 DigitCounting

```cpp
int dfs(int pos, int state1, int state2 ....., bool
    limit, bool zero) {
    if ( pos == -1 ) return 是否符合條件;
    int &ret = dp[pos][state1][state2][....];
    if ( ret != -1 && !limit ) return ret;
    int ans = 0;
    int upper = limit ? digit[pos] : 9;
    for ( int i = 0 ; i <= upper ; i++ ) {
        ans += dfs(pos - 1, new_state1, new_state2,
            limit & ( i == upper), ( i == 0) && zero);
    }
    if ( !limit ) ret = ans;
    return ans;
}

int solve(int n) {
    int it = 0;
    for ( ; n ; n /= 10 ) digit[it++] = n % 10;
    return dfs(it - 1, 0, 0, 1, 1);
}
```

8.4 Dp1D1D

```cpp
#include<bits/stdc++.h>

int t, n, L;
int p;
char s[MAXN][35];
ll sum[MAXN] = {0};
long double dp[MAXN] = {0};
int prevd[MAXN] = {0};

long double pw(long double a, int n) {
    if ( n == 1 ) return a;
    long double b = pw(a, n/2);
    if ( n & 1 ) return b*b*a;
    else return b*b;
}
long double f(int i, int j) {
//    cout << (sum[i] - sum[j]+i-j-1-L) << endl;
    return pw(abs(sum[i] - sum[j]+i-j-1-L), p) + dp[j];
}
struct INV {
    int L, R, pos;
};
INV stk[MAXN*10];
int top = 1, bot = 1;
void update(int i) {
    while ( top > bot && i < stk[top].L && f(stk[top].L
        , i) < f(stk[top].L, stk[top].pos) ) {
        stk[top - 1].R = stk[top].R;
        top--;
    }
    int lo = stk[top].L, hi = stk[top].R, mid, pos =
        stk[top].pos;
    //if ( i >= lo ) lo = i + 1;
    while ( lo != hi ) {
        mid = lo + (hi - lo) / 2;
        if ( f(mid, i) < f(mid, pos) ) hi = mid;
        else lo = mid + 1;
    }
    if ( hi < stk[top].R ) {
        stk[top + 1] = (INV) { hi, stk[top].R, i };
        stk[top++].R = hi;
    }
}

int main() {
    cin >> t;
    while ( t-- ) {
        cin >> n >> L >> p;
        dp[0] = sum[0] = 0;
        for ( int i = 1 ; i <= n ; i++ ) {
            cin >> s[i];
            sum[i] = sum[i-1] + strlen(s[i]);
            dp[i] = numeric_limits<long double>::max();
        }
        stk[top] = (INV) {1, n + 1, 0};
        for ( int i = 1 ; i <= n ; i++ ) {
            if ( i >= stk[bot].R ) bot++;
            dp[i] = f(i, stk[bot].pos);
            update(i);
//              cout << (ll) f(i, stk[bot].pos) << endl;
        }
        if ( dp[n] > 1e18 ) {
            cout << "Too hard to arrange" << endl;
        } else {
            vector<PI> as;
            cout << (ll)dp[n] << endl;
        }
    }
    return 0;
```

```
}
```

8.5 ManhattanMST

```
#include <bits/stdc++.h>
using namespace std;

const int MAXN = 100005;
const int OFFSET = 2000; // y-x may < 0, offset it, if
    y-x too large, please write a unique function
const int INF = 0xFFFFFFF;
int n;
int x[MAXN], y[MAXN], p[MAXN];

typedef pair<int, int> pii;
pii bit[MAXN]; // [ val, pos ]

struct P {
    int x, y, id;
    bool operator<(const P&b ) const {
        if ( x == b.x ) return y > b.y;
        else return x > b.x;
    }
};
vector<P> op;

struct E {
    int x, y, cost;
    bool operator<(const E&b ) const {
        return cost < b.cost;
    }
};
vector<E> edges;

int find(int x) {
    return p[x] == x ? x : p[x] = find(p[x]);
}

void update(int i, int v, int p) {
    while ( i ) {
        if ( bit[i].first > v ) bit[i] = {v, p};
        i -= i & (-i);
    }
}

pii query(int i) {
    pii res = {INF, INF};
    while ( i < MAXN ) {
        if ( bit[i].first < res.first ) res = {bit[i].
            first, bit[i].second};
        i += i & (-i);
    }
    return res;
}

void input() {
    cin >> n;
    for ( int i = 0 ; i < n ; i++ ) cin >> x[i] >> y[i
        ], op.push_back((P) {x[i], y[i], i});
}

void mst() {
    for ( int i = 0 ; i < MAXN ; i++ ) p[i] = i;
    int res = 0;
    sort(edges.begin(), edges.end());
    for ( auto e : edges ) {
        int x = find(e.x), y = find(e.y);
        if ( x != y ) {
            p[x] = y;
            res += e.cost;
        }
    }
    cout << res << endl;
}

void construct() {
    sort(op.begin(), op.end());
    for ( int i = 0 ; i < n ; i++ ) {
        pii q = query(op[i].y - op[i].x + OFFSET);
        update(op[i].y - op[i].x + OFFSET, op[i].x + op
            [i].y, op[i].id);
        if ( q.first == INF ) continue;
        edges.push_back((E) {op[i].id, q.second, abs(x[
            op[i].id]-x[q.second]) + abs(y[op[i].id]-y[
            q.second]) });
    }
}

void solve() {

    // [45 ~ 90 deg]
    for ( int i = 0 ; i < MAXN ; i++ ) bit[i] = {INF,
        INF};
    construct();

    // [0 ~ 45 deg]
    for ( int i = 0 ; i < MAXN ; i++ ) bit[i] = {INF,
        INF};
    for ( int i = 0 ; i < n ; i++ ) swap(op[i].x, op[i
        ].y);
    construct();
    for ( int i = 0 ; i < n ; i++ ) swap(op[i].x, op[i
        ].y);

    // [-90 ~ -45 deg]
    for ( int i = 0 ; i < MAXN ; i++ ) bit[i] = {INF,
        INF};
    for ( int i = 0 ; i < n ; i++ ) op[i].y *= -1;
    construct();

    // [-45 ~ 0 deg]
    for ( int i = 0 ; i < MAXN ; i++ ) bit[i] = {INF,
        INF};
    for ( int i = 0 ; i < n ; i++ ) swap(op[i].x, op[i
        ].y);
    construct();

    // mst
    mst();

}

int main () {
    input();
    solve();
    return 0;
}
```

8.6 Count Spanning Tree

```
新的方法介绍
下面我们介绍一种新的方法——Matrix-Tree定理(Kirchhoff矩
    阵-树定理)。

Matrix-Tree定理是解决生成树计数问题最有力的武器之一。它
    首先于1847年被Kirchhoff证明。在介绍定理之前，我们首
    先明确几个概念：
1、G的度数矩阵D[G]是一个n*n的矩阵，并且满足：当i≠j时，
    dij=0；当i=j时，dij等于vi的度数。
2、G的邻接矩阵A[G]也是一个n*n的矩阵， 并且满足：如果vi
    、vj之间有边直接相连，则aij=1，否则为0。
我们定义G的Kirchhoff矩阵(也称为拉普拉斯算子)C[G]为C[G]=
    D[G]-A[G]，
则Matrix-Tree定理可以描述为：G的所有不同的生成树的个数
    等于其Kirchhoff矩阵C[G]任何一个n-1阶主子式的行列式
    的绝对值。
所谓n-1阶主子式，就是对于r(1≤r≤n)，将C[G]的第r行、第r列
    同时去掉后得到的新矩阵，用Cr[G]表示。
```

```
生成树计数
算法步骤：
1、 构建拉普拉斯矩阵
     Matrix[i][j] =
degree(i) , i==j
           -1，i-j有边
            0，其他情况
2、 去掉第r行，第r列 (r任意)
3、 计算矩阵的行列式

/* **************************************************
MYID    : Chen Fan
LANG    : G++
PROG    : Count_Spaning_Tree_From_Kuangbin
************************************************ */
#include <stdio.h>
#include <string.h>
#include <algorithm>
#include <iostream>
#include <math.h>
using namespace std;
const double eps = 1e-8;
const int MAXN = 110;
int sgn(double x)
{
    if(fabs(x) < eps)return 0;
    if(x < 0)return -1;
    else return 1;
}
double b[MAXN][MAXN];
double det(double a[][MAXN],int n)
{
    int i, j, k, sign = 0;
    double ret = 1;
    for(i = 0;i < n;i++)
    for(j = 0;j < n;j++) b[i][j] = a[i][j];
    for(i = 0;i < n;i++)
    {
        if(sgn(b[i][i]) == 0)
        {
            for(j = i + 1; j < n;j++)
            if(sgn(b[j][i]) != 0) break;
            if(j == n)return 0;
            for(k = i;k < n;k++) swap(b[i][k],b[j][k]);
            sign++;
        }
        ret *= b[i][i];
        for(k = i + 1;k < n;k++) b[i][k]/=b[i][i];
        for(j = i+1;j < n;j++)
        for(k = i+1;k < n;k++) b[j][k] -= b[j][i]*b[i][
            k];
    }
    if(sign & 1)ret = -ret;
    return ret;
}
double a[MAXN][MAXN];
int g[MAXN][MAXN];
int main()
{
    int T;
    int n,m;
    int u,v;
    scanf("%d",&T);
    while(T--)
    {
        scanf("%d%d",&n,&m);
        memset(g,0,sizeof(g));
        while(m--)
        {
            scanf("%d%d",&u,&v);
            u--;v--;
            g[u][v] = g[v][u] = 1;
        }
        memset(a,0,sizeof(a));
        for(int i = 0;i < n;i++)
        for(int j = 0;j < n;j++)
        if(i != j && g[i][j])
        {
            a[i][i]++;
            a[i][j] = -1;
        }
        double ans = det(a,n-1);
        printf("%.0lf\n",ans);
    }
    return 0;
}
```

9 String

9.1 AC

```
// remember make_fail() !!!
// notice MLE

const int sigma = 62;
const int MAXC = 200005;

inline int idx(char c){
    if ('A'<= c && c <= 'Z')return c-'A';
    if ('a'<= c && c <= 'z')return c-'a' + 26;
    if ('0'<= c && c <= '9')return c-'0' + 52;
}

struct ACautomaton{
    struct Node{
        Node *next[sigma], *fail;
        int cnt; // dp
        Node(){
            memset(next,0,sizeof(next));
            fail=0;
            cnt=0;
        }
    } buf[MAXC], *bufp, *ori, *root;

    void init(){
        bufp = buf;
        ori = new (bufp++) Node();
        root = new (bufp++) Node();
    }

    void insert(int n, char *s){
        Node *ptr = root;
        for (int i=0; s[i]; i++){
            int c = idx(s[i]);
            if (ptr->next[c]==NULL)
                ptr->next[c] = new (bufp++) Node();
            ptr = ptr->next[c];
        }
        ptr->cnt=1;
    }

    Node* trans(Node *o, int c){
        while (o->next[c]==NULL) o = o->fail;
        return o->next[c];
    }

    void make_fail(){
        static queue<Node*> que;

        for (int i=0; i<sigma; i++)
            ori->next[i] = root;
        root->fail = ori;

        que.push(root);
        while ( que.size() ){
            Node *u = que.front(); que.pop();
            for (int i=0; i<sigma; i++){
                if (u->next[i]==NULL)continue;
                u->next[i]->fail = trans(u->fail,i);
                que.push(u->next[i]);
            }
```

```
            u->cnt += u->fail->cnt;
        }
    }
} ac;
```

9.2 BWT

```
// BWT
const int N = 8;            // 字串長度
int s[N+N+1] = "suffixes";  // 字串，後面預留一倍空間。
int sa[N];                  // 後綴陣列
int pivot;

int cmp(const void* i, const void* j)
{
    return strncmp(s+*(int*)i, s+*(int*)j, N);
}

// 此處便宜行事，採用 O(N²logN) 的後綴陣列演算法。
void BWT()
{
    strncpy(s + N, s, N);
    for (int i=0; i<N; ++i) sa[i] = i;
    qsort(sa, N, sizeof(int), cmp);
    // 當輸入字串的所有字元都相同，必須當作特例處理。
    // 或者改用stable sort。

    for (int i=0; i<N; ++i)
        cout << s[(sa[i] + N-1) % N];

    for (int i=0; i<N; ++i)
        if (sa[i] == 0)
        {
            pivot = i;
            break;
        }
}

// Inverse BWT
const int N = 8;            // 字串長度
char t[N+1] = "xuffessi";   // 字串
int pivot;
int next[N];

void IBWT()
{
    vector<int> index[256];
    for (int i=0; i<N; ++i)
        index[t[i]].push_back(i);

    for (int i=0, n=0; i<256; ++i)
        for (int j=0; j<index[i].size(); ++j)
            next[n++] = index[i][j];

    int p = pivot;
    for (int i=0; i<N; ++i)
        cout << t[p = next[p]];
}
```

9.3 KMP

```
template<typename T>
void build_KMP(int n, T *s, int *f){ // 1 base
  f[0]=-1, f[1]=0;
  for (int i=2; i<=n; i++){
    int w = f[i-1];
    while (w>=0 && s[w+1]!=s[i])w = f[w];
    f[i]=w+1;
  }
}

template<typename T>
int KMP(int n, T *a, int m, T *b){
  build_KMP(m,b,f);
  int ans=0;

  for (int i=1, w=0; i<=n; i++){
    while ( w>=0 && b[w+1]!=a[i] )w = f[w];
    w++;
    if (w==m){
      ans++;
      w=f[w];
    }
  }
  return ans;
}
```

9.4 PalindromicTree

```
// remember init()      !!!
// remember make_fail() !!!
// insert s need 1 base !!!
// notice MLE
const int sigma = 62;
const int MAXC = 1000006;
inline int idx(char c){
    if ('a'<= c && c <= 'z')return c-'a';
    if ('A'<= c && c <= 'Z')return c-'A'+26;
    if ('0'<= c && c <= '9')return c-'0'+52;
}
struct PalindromicTree{
    struct Node{
        Node *next[sigma], *fail;
        int len, cnt; // for dp
        Node(){
            memset(next,0,sizeof(next));
            fail=0;
            len = cnt = 0;
        }
    } buf[MAXC], *bufp, *even, *odd;

    void init(){
        bufp = buf;
        even = new (bufp++) Node();
        odd  = new (bufp++) Node();
        even->fail = odd;
        odd->len = -1;
    }

    void insert(char *s){
        Node* ptr = even;
        for (int i=1; s[i]; i++){
            ptr = extend(ptr,s+i);
        }
    }

    Node* extend(Node *o, char *ptr){
        int c = idx(*ptr);
        while ( *ptr != *(ptr-1-o->len) )o=o->fail;
        Node *&np = o->next[c];
        if (!np){
            np = new (bufp++) Node();
            np->len = o->len+2;
            Node *f = o->fail;
            if (f){
                while ( *ptr != *(ptr-1-f->len) )f=f->
                    fail;
                np->fail = f->next[c];
            }
            else {
                np->fail = even;
            }
            np->cnt = np->fail->cnt;
        }
        np->cnt++;
        return np;
    }
} PAM;
```

9.5 SAM

```
// par : fail link
// val : a topological order ( useful for DP )
// go[x] : automata edge ( x is integer in [0,26) )

struct SAM{
  struct State{
    int par, go[26], val;
    State () : par(0), val(0){ FZ(go); }
    State (int _val) : par(0), val(_val){ FZ(go); }
  };
  vector<State> vec;
  int root, tail;

  void init(int arr[], int len){
    vec.resize(2);
    vec[0] = vec[1] = State(0);
    root = tail = 1;
    for (int i=0; i<len; i++)
      extend(arr[i]);
  }
  void extend(int w){
    int p = tail, np = vec.size();
    vec.PB(State(vec[p].val+1));
    for ( ; p && vec[p].go[w]==0; p=vec[p].par)
      vec[p].go[w] = np;
    if (p == 0){
      vec[np].par = root;
    } else {
      if (vec[vec[p].go[w]].val == vec[p].val+1){
        vec[np].par = vec[p].go[w];
      } else {
        int q = vec[p].go[w], r = vec.size();
        vec.PB(vec[q]);
        vec[r].val = vec[p].val+1;
        vec[q].par = vec[np].par = r;
        for ( ; p && vec[p].go[w] == q; p=vec[p].par)
          vec[p].go[w] = r;
      }
    }
    tail = np;
  }
};
```

9.6 Z-value

```
z[0] = 0;
for ( int bst = 0, i = 1; i < len ; i++ ) {
  if ( z[bst] + bst <= i ) z[i] = 0;
  else z[i] = min(z[i - bst], z[bst] + bst - i);
  while ( str[i + z[i]] == str[z[i]] ) z[i]++;
  if ( i + z[i] > bst + z[bst] ) bst = i;
}
```

```
// 回文版

void Zpal(const char *s, int len, int *z) {
    // Only odd palindrome len is considered
    // z[i] means that the longest odd palindrom
        centered at
    // i is [i-z[i] .. i+z[i]]
    z[0] = 0;
    for (int b=0, i=1; i<len; i++) {
        if (z[b] + b >= i) z[i] = min(z[2*b-i], b+z[b]-
            i);
        else z[i] = 0;
        while (i+z[i]+1 < len and i-z[i]-1 >= 0 and
               s[i+z[i]+1] == s[i-z[i]-1]) z[i] ++;
        if (z[i] + i > z[b] + b) b = i;
    }
}
```

9.7 Smallest Rotation

```
string mcp(string s){
  int n = s.length();
  s += s;
  int i=0, j=1;
  while (i<n && j<n){
    int k = 0;
    while (k < n && s[i+k] == s[j+k]) k++;
    if (s[i+k] <= s[j+k]) j += k+1;
    else i += k+1;
    if (i == j) j++;
  }
  int ans = i < n ? i : j;
  return s.substr(ans, n);
}
```

9.8 Suffix Array

```
/*he[i]保存了在後綴數組中相鄰兩個後綴的最長公共前綴長度
 *sa[i]表示的是字典序排名為i的後綴是誰（字典序越小的排
     名越靠前）
 *rk[i]表示的是後綴我所對應的排名是多少   */

const int MAX = 1020304;
int ct[MAX], he[MAX], rk[MAX];
int sa[MAX], tsa[MAX], tp[MAX][2];
void suffix_array(char *ip){
  int len = strlen(ip);
  int alp = 256;
  memset(ct, 0, sizeof(ct));
  for(int i=0;i<len;i++) ct[ip[i]+1]++;
  for(int i=1;i<alp;i++) ct[i]+=ct[i-1];
  for(int i=0;i<len;i++) rk[i]=ct[ip[i]];
  for(int i=1;i<len;i*=2){
    for(int j=0;j<len;j++){
      if(j+i>=len) tp[j][1]=0;
      else tp[j][1]=rk[j+i]+1;
      tp[j][0]=rk[j];
    }
    memset(ct, 0, sizeof(ct));
    for(int j=0;j<len;j++) ct[tp[j][1]+1]++;
    for(int j=1;j<len+2;j++) ct[j]+=ct[j-1];
    for(int j=0;j<len;j++) tsa[ct[tp[j][1]]++]=j;
    memset(ct, 0, sizeof(ct));
    for(int j=0;j<len;j++) ct[tp[j][0]+1]++;
    for(int j=1;j<len+1;j++) ct[j]+=ct[j-1];
    for(int j=0;j<len;j++)
      sa[ct[tp[tsa[j]][0]]++]=tsa[j];
    rk[sa[0]]=0;
    for(int j=1;j<len;j++){
      if( tp[sa[j]][0] == tp[sa[j-1]][0] &&
        tp[sa[j]][1] == tp[sa[j-1]][1] )
        rk[sa[j]] = rk[sa[j-1]];
      else
        rk[sa[j]] = j;
    }
  }
  for(int i=0,h=0;i<len;i++){
    if(rk[i]==0) h=0;
    else{
      int j=sa[rk[i]-1];
      h=max(0,h-1);
      for(;ip[i+h]==ip[j+h];h++);
    }
    he[rk[i]]=h;
  }
}
```

10 無權邊的生成樹個數 Kirchhoff's Theorem

1. 定義 $n \times m$ 矩陣 $E = (a_{i,j})$，n 為點數，m 為邊數，若 i 點在 j 邊上，i 為小點 $a_{i,j} = 1$，i 為大點 $a_{i,j} = -1$，否則 $a_{i,j} = 0$。
(證明省略)
4. 令 $E(E^T) = Q$，他是一種有負號的 kirchhoff 的矩陣，取 Q 的子矩陣即為 $F(F^T)$
結論：做 Q 取子矩陣算 det 即為所求。(除去第一行第一列 by mz)

11 monge

$i \leq i^{'} < j \leq j^{'}$
$m(i,j) + m(i^{'}, j^{'}) \leq m(i^{'}, j) + m(i, j^{'})$
$k(i,j-1) <= k(i,j) <= k(i+1,j)$

12 四心

$\frac{sa*A+sb*B+sc*C}{sa+sb+sc}$
外心 sin 2A : sin 2B : sin 2C
內心 sin A : sin B : sin C
垂心 tan A : tan B : tan C
重心 1 : 1 : 1

13 Runge-Kutta

$y_{n+1} = y_n + \frac{h}{6}(k_1 + 2k_2 + 2k_3 + k_4)$
$k_1 = f(t_n, y_n)$
$k_2 = f(t_n + \frac{h}{2}, y_n + \frac{h}{2}k_2)$
$k_3 = f(t_n + \frac{h}{2}, y_n + \frac{h}{2}k_3)$
$k_2 = f(t_n + h, y_n + hk_3)$

14 Householder Matrix

$I - 2\frac{vv^T}{v^Tv}$