Homework 1: Strings and languages

Theory of Computing (CSE 30151), Spring 2024

Due: 2024-01-26 5pm

Instructions

- Create a PDF file (or files) containing your solutions. You can write your solutions by hand, but please scan them into a PDF.
- Please name your PDF file(s) as follows to ensure that the graders give you credit for all of your work:
 - If you're making a complete submission, name it netid-hw1.pdf, where netid is replaced with your NetID.
 - If you're submitting some problems now and want to submit other problems later, name it netid-hw1-part123.pdf, where 123 is replaced with the problem number(s) you are submitting at this time.
- Submit your PDF file(s) in Canvas under HW1.

Problems (10 points each)

1. Proof practice.

(a) Convert this paragraph proof to a statement–reason proof. Please be sure to write which statement(s) each statement depends on.

To show: If s is a string, every substring of a substring of s is a substring of s.

Proof: Let y be a substring of s, that is, s = xyz for some x, z; and let v be a substring of y, that is, y = uvw for some u, w. Then s = xuvwz, so v is a substring of s.

(b) Convert this statement–reason proof to a paragraph proof.

To show: If w is a string, every prefix of a suffix of w is a suffix of a prefix of w.

1. v is a suffix of w Given 2. y is a prefix of v Given 3. $\exists x$ s.t. xv = w (1), def. suffix 4. $\exists z$ s.t. yz = v (2), def. prefix 5. xyz = w (3), (4), substitution 6. xy is a prefix of w (5), def. prefix 7. y is a suffix of xy (6), def. suffix

2. **String homomorphisms.** If Σ and Γ are finite alphabets, define a *string homomorphism* to be a function $f \colon \Sigma^* \to \Gamma^*$ that has the property that for any $u, v \in \Sigma^*$,

$$f(uv) = f(u) f(v).$$

An example of a string homomorphism is the function that converts hexadecimal numbers into binary numbers, which operates digit-by-digit:

$$f_{
m hb}\colon \{{
m 0},\dots,{
m 9},{
m A},\dots,{
m F}\}^* o \{{
m 0},{
m 1}\}^*$$
 $f_{
m hb}(arepsilon)=arepsilon$ $f_{
m hb}({
m 1})={
m 0001}$ $f_{
m hb}({
m A})={
m 1010}$ $f_{
m hb}({
m 1A1A})={
m 0001101000011010}.$

Prove that, in general, every string homomorphism operates by replacing each symbol with a (possibly empty) string. That is, prove that if f is a string homomorphism, then for any $w = w_1 \cdots w_n$ (where $n \geq 0$ and, for $j = 1, \ldots, n$, $w_j \in \Sigma$), we have

$$f(w) = f(w_1) \cdots f(w_n). \tag{*}$$

Use induction on n.

- (a) State and prove the base case (n = 0).
- (b) Assume that (??) is true for n = i and prove (??) for n = i + 1.

You may assume the following facts about strings:

- Identity: For all $x \in \Sigma^*$, $x\varepsilon = x$ and $\varepsilon x = x$.
- Right cancellation: For all $x, y, z \in \Sigma^*$, if xz = yz then x = y.
- Left cancellation: For all $x, y, z \in \Sigma^*$, if xy = xz then y = z.
- 3. Finite and cofinite. Let $\Sigma = \{a, b\}$. Define FINITE to be the set of all finite languages over Σ , and let coFINITE be the set of all languages over Σ whose complement is finite:

$$\mathsf{coFINITE} = \{L \subseteq \Sigma^* \mid \overline{L} \in \mathsf{FINITE}\}$$

(where $\overline{L} = \Sigma^* \setminus L$). For example, Σ^* is in cofinite because its complement is \emptyset , which is finite. (Please think carefully about this definition, and note that cofinite isn't the same thing as $\overline{\mathsf{FINITE}}$.)

- (a) If $L \in \mathsf{FINITE}$, what data structure could you use to represent L, and given a string w, how would you decide whether $w \in L$?
- (b) If $L \in \mathsf{coFINITE}$, what data structure could you use to represent L, and given a string w, how would you decide whether $w \in L$?
- (c) Are there any languages in $\mathsf{FINITE} \cap \mathsf{coFINITE}?$ Prove your answer.
- (d) Are there any languages not in FINITE \cup coFINITE? Prove your answer.