## Homework 2: DFAs and NFAs

Theory of Computing (CSE 30151), Spring 2025

Due 2025-01-31 5pm

## Instructions

- Create a PDF file (or files) containing your solutions. You can write your solutions by hand, but please scan them into a PDF.
- Please name your PDF file(s) as follows to ensure that the graders give you credit for all of your work:
  - If you're making a complete submission, name it netid-hw2.pdf, where netid is replaced with your NetID.
  - If you're submitting some problems now and want to submit other problems later, name it netid-hw2-part123.pdf, where 123 is replaced with the problem numbers you are submitting at this time.
- Submit your PDF file(s) in Canvas.

## Problems (10 points each)

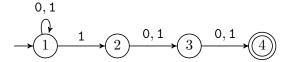
1. **Divisibility tests.** Define, for all k > 0,

 $D_k = \{w \in \{0, \dots, 9\}^* \mid w \text{ is the decimal representation of a multiple of } k\}$ 

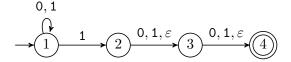
where  $\varepsilon$  is considered to represent the number 0. For example, the strings  $\varepsilon$ , 0, 1234, and 01234 all belong to  $D_2$ , but 99 and 099 do not.

- (a) Prove that  $D_2$  is regular by writing a DFA for  $D_2$ .
- (b) Prove that  $D_3$  is regular by writing a DFA for  $D_3$ .
- (c) Prove that  $D_6$  is regular. An explicit DFA is not necessary.
- (\*) Optional alternative: You can get full credit for all of the above if you can prove that for any k > 0,  $D_k$  is regular, by describing how to write the formal description of a DFA  $M = (Q, \{0, ..., 9\}, \delta, s, F)$  in terms of k.

2. **Nondeterminism.** Consider the following NFA  $N_2$  (same as in Figure 1.31), which accepts a string iff the third-to-last symbol is a 1:



- (a) Use the subset construction (Theorem 1.39) to convert  $N_2$  to a DFA M. You may omit curly braces and commas when naming states; for example, instead of  $\{1, 2, 3, 4\}$  you may write 1234. (Hint: the DFA should be equivalent to the one in Figure 1.32.)
- (b) Why do you think the states in Figure 1.32 are named the way they are?
- (c) In Example 1.30, Sipser asks what happens if you modify  $N_2$  into the following NFA let's call it  $N'_2$ :



Use the subset construction (Theorem 1.39) to convert  $N'_2$  to a DFA M'.

- 3. Procrustean closure properties.
  - (a) For any string  $w = w_1 w_2 \cdots w_{n-1} w_n$ , define

STRETCH
$$(w_1 w_2 \cdots w_{n-1} w_n) = w_1 w_1 w_2 w_2 \cdots w_{n-1} w_{n-1} w_n w_n$$
.

For example,

$$\begin{split} & \text{STRETCH}(\varepsilon) = \varepsilon \\ & \text{STRETCH(abab)} = \text{aabbaabb}. \end{split}$$

This induces an operation on languages,

$$STRETCH(L) = \{STRETCH(w) \mid w \in L\}.$$

For example,

$$STRETCH(\{\varepsilon, abab\}) = \{\varepsilon, aabbaabb\}.$$

Prove that for any regular language L, STRETCH(L) is also regular.

(b) For any  $w = w_1 w_2 \cdots w_{n-1} w_n$  with  $n \geq 2$ , define

$$CHOP(w_1w_2\cdots w_{n-1}w_n)=w_2\cdots w_{n-1}.$$

For example,

$$\mathrm{CHOP}(\mathtt{ab}) = \varepsilon$$
  $\mathrm{CHOP}(\mathtt{abab}) = \mathtt{ba}$   $\mathrm{CHOP}(\mathtt{bbaa}) = \mathtt{ba}.$ 

This induces an operation on languages,

$$CHOP = \{CHOP(w) \mid w \in L \text{ and } |w| \ge 2\}.$$

For example,

$$\begin{split} \mathrm{CHOP}(\{\mathtt{ab},\mathtt{abab},\mathtt{bbaa}\}) &= \{\varepsilon,\mathtt{ba}\} \\ \mathrm{CHOP}(\{\varepsilon,\mathtt{a}\}) &= \emptyset. \end{split}$$

Prove that for any regular language L, CHOP(L) is also regular.