

# Homework 1: Strings and Languages

Theory of Computing (CSE 30151), Spring 2026

Due: 2026-01-23 5pm

## Instructions

- Create a PDF file (or files) containing your solutions. You can write your solutions by hand, but please scan them into a PDF.
- Please name your PDF file(s) as follows to ensure that the graders give you credit for all of your work:
  - If you're making a complete submission, name it *netid-hw1.pdf*, where *netid* is replaced with your NetID.
  - If you're submitting some problems now and want to submit other problems later, name it *netid-hw1-part123.pdf*, where 123 is replaced with the problem number(s) you are submitting at this time.
- Submit your PDF file(s) in Canvas under HW1.

## Problems (10 points each)

### 1. Proof practice.

- (a) Convert this paragraph proof to a statement–reason proof. Please be sure to write which statement(s) each statement depends on.

To show: If  $s$  is a string, every substring of a substring of  $s$  is a substring of  $s$ .

Proof: Let  $y$  be a substring of  $s$ , that is,  $s = xyz$  for some  $x, z$ ; and let  $v$  be a substring of  $y$ , that is,  $y = uvw$  for some  $u, w$ . Then  $s = xuvwz$ , so  $v$  is a substring of  $s$ .

- (b) Convert this statement–reason proof to a paragraph proof.

To show: If  $w$  is a string, every prefix of a suffix of  $w$  is a suffix of a prefix of  $w$ .

- |                              |                        |
|------------------------------|------------------------|
| 1. $v$ is a suffix of $w$    | Given                  |
| 2. $y$ is a prefix of $v$    | Given                  |
| 3. $\exists x$ s.t. $xv = w$ | (1), def. suffix       |
| 4. $\exists z$ s.t. $yz = v$ | (2), def. prefix       |
| 5. $xyz = w$                 | (3), (4), substitution |
| 6. $xy$ is a prefix of $w$   | (5), def. prefix       |
| 7. $y$ is a suffix of $xy$   | (6), def. suffix       |

2. **String homomorphisms.** If  $\Sigma$  and  $\Gamma$  are finite alphabets, define a *string homomorphism* to be a function  $f: \Sigma^* \rightarrow \Gamma^*$  that has the property that for any  $u, v \in \Sigma^*$ ,

$$f(uv) = f(u)f(v).$$

An example of a string homomorphism is the function that converts hexadecimal numbers into binary numbers, which operates digit-by-digit:

$$\begin{aligned} f_{\text{hex}}: \{0, \dots, 9, \text{A}, \dots, \text{F}\}^* &\rightarrow \{0, 1\}^* \\ f_{\text{hex}}(\varepsilon) &= \varepsilon \\ f_{\text{hex}}(1) &= 0001 \\ f_{\text{hex}}(\text{A}) &= 1010 \\ f_{\text{hex}}(1\text{A}1\text{A}) &= 0001101000011010. \end{aligned}$$

Prove that all string homomorphisms operate symbol-by-symbol in this way. That is, prove: for any string homomorphism  $f$ , and for any string  $w = w_1 \cdots w_n$  (where  $n \geq 0$  and, for  $j = 1, \dots, n$ ,  $w_j \in \Sigma$ ), we have

$$f(w) = f(w_1) \cdots f(w_n). \quad (*)$$

Use induction on  $n$ .

- (a) State and prove the base case, that is,  $(*)$  for  $n = 0$ .
- (b) Assume that  $(*)$  is true for  $n = i$  and prove  $(*)$  for  $n = i + 1$ .

You may assume the following facts about strings:

- Identity: For all  $x \in \Sigma^*$ ,  $x\varepsilon = x$  and  $\varepsilon x = x$ .
- Right cancellation: For all  $x, y, z \in \Sigma^*$ , if  $xz = yz$  then  $x = y$ .
- Left cancellation: For all  $x, y, z \in \Sigma^*$ , if  $xy = xz$  then  $y = z$ .

3. **Finite and cofinite.** Let  $\Sigma = \{\text{a}, \text{b}\}$ . Define FINITE to be the set of all finite languages over  $\Sigma$ , and let coFINITE be the set of all languages over  $\Sigma$  whose complement is finite:

$$\text{coFINITE} = \{L \subseteq \Sigma^* \mid \bar{L} \in \text{FINITE}\}$$

(where  $\overline{L} = \Sigma^* \setminus L$ ). For example,  $\Sigma^*$  is in **coFINITE** because its complement is  $\emptyset$ , which is finite. (Please think carefully about this definition, and note that **coFINITE** isn't the same thing as  $\overline{\text{FINITE}}$ .)

- (a) If  $L \in \text{FINITE}$ , what data structure could you use to represent  $L$ , and given a string  $w$ , how would you decide whether  $w \in L$ ?
- (b) If  $L \in \text{coFINITE}$ , what data structure could you use to represent  $L$ , and given a string  $w$ , how would you decide whether  $w \in L$ ?
- (c) Are there any languages in  $\text{FINITE} \cap \text{coFINITE}$ ? Prove your answer.
- (d) Are there any languages over  $\Sigma$  that are *not* in  $\text{FINITE} \cup \text{coFINITE}$ ? Prove your answer.

You may assume the following fact about sets:

- The union of two finite sets is finite.