

CS 171 HW2

Q1. a) $\exists N_f \in \mathbb{Z}^+ \forall n > N_f \ f(n) < \frac{1}{2P(n)} \left\{ \begin{array}{l} h(n) = f(n) + g(n) < \frac{1}{2P(n)} \forall n > (\max(N_f, N_g)) \\ \exists N_g \in \mathbb{Z}^+ \forall n > N_g \ g(n) < \frac{1}{2P(n)} \end{array} \right.$

b) $\exists N_f \in \mathbb{Z}^+ \forall n > N_f \ f(n) < \frac{1}{n^c \text{Poly}(n)} \left\{ \begin{array}{l} h(n) = f(n) \cdot P(n) < \frac{1}{\text{Poly}(n)} \\ P(n) < n^c \end{array} \right.$

c) Polynomial

$\therefore n^{-100}$ is polynomial as $n^{-100} < \frac{1}{n^{101}}$ is impossible

d) Negligible

\therefore For large enough n , $n^c < 1.01^n$, so $\frac{1}{1.01^n} < \frac{1}{P(n)}$

e) Polynomial

$\therefore 2^{-(\log_2 n)^2} = \frac{1}{n^2} > \frac{1}{n^3}$

f) Negligible

\therefore Dominant term $e^{-\log n}$ is negligible since $n^c < e^{\log_2 n}$ for large enough n
which means $\frac{1}{e^{\log_2 n}} < \frac{1}{n^c}$

Q2. a) When $a=0$, it is not invertible so Rec leads to error, losing the message.

$\Pr[a=0] = \frac{1}{23}$

b) $\left. \begin{array}{l} c_1 = a \cdot m_1 + b \text{ mod } 23 \\ c_2 = a \cdot m_2 + b \text{ mod } 23 \end{array} \right\} c_1 - c_2 = a \cdot \underbrace{(m_1 - m_2)}_{\text{non-zero}} \text{ mod } 23$

$\therefore (m_1 - m_2)$ has inverse mod 23

\therefore Unique solution for a and thus b

The above is same for m'_1, m'_2

$\therefore \Pr[\text{Enc}(k, m_1) = c_1 \wedge \text{Enc}(k, m_2) = c_2] = \Pr[\text{Enc}(k, m'_1) = c_2 \wedge \text{Enc}(k, m'_2) = c_2]$
 $= \frac{1}{23} \times \frac{1}{22}$