# The Number System of the Permutations Generated by Cyclic Shift

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#### Abstract

A number system coding for the permutations generated by cyclic shift is described. The system allows to find the rank of a permutation given how it has been generated, and to determine a permutation given its rank. It defines a code describing the symmetry properties of the set of permutations generated by cyclic shift. This code is conjectured to be a combinatorial Gray code listing the set of permutations: this corresponds to an Hamiltonian path of minimal weight in an appropriate regular digraph.

#### 1 Introduction

Since the work of Laisant in 1888 [4] – and even since Fischer and Krause in 1812 [1] according to Hall and Knuth [2] –, it is known that the factorial number system codes for the permutations generated in lexicographic order. More precisely, when the set of all permutations on n symbols is ordered by lexicographic order, the rank of a permutation written in the factorial system provides a code determining the permutation. The code specifies which interchanges of the symbols according to lexicographic order have to be performed to generate the permutation. Conversely, the rank of a permutation can be computed from its code. This coding has been rediscovered several times since (e.g., Lehmer [5]).

In this study, we describe a number system on the finite ring  $\mathbb{Z}_{n!}$  coding for the permutations generated by cyclic shift. When the set  $\mathcal{S}_n$  of permutations is ordered according to generation by cyclic shift, the rank of a permutation written in this number system entirely specifies how the permutation has been generated. Conversely, the rank can be computed from the code. This number system is a special case of a large class of methods presented by Knuth [3] for generating  $\mathcal{S}_n$ .

We shall describe properties of  $S_n$  generated by cyclic shift:

- 1. A decomposition into k-orbits;
- 2. The symmetries:

- 3. An infinite family of regular digraphs associated with  $\{S_n; n \geq 1\}$ ;
- 4. A conjectured combinatorial Gray code generating the permutations on n symbols. The adjacency rule associated with this code is that the last symbols of each permutation match the first symbols of the next optimally.

#### 2 Number system

For any positive integer a, the ring  $(\mathbb{Z}/a\mathbb{Z}, +, \times)$  of integers modulo a is denoted  $\mathbb{Z}_a$ . The set  $\mathbb{Z}_a$  is identified with a subset of the set  $\mathbb{N}$  of natural integers.

**Proposition 1.** For  $n \geq 2$ , any element  $\alpha \in \mathbb{Z}_{n!}$  can be uniquely represented as

$$\alpha = \sum_{i=0}^{n-2} \alpha_i \varpi_{n,i}, \quad \alpha_i \in \mathbb{Z}_{n-i},$$

with the base elements

$$\varpi_{n,0} = 1, \qquad \varpi_{n,i} = n(n-1)\cdots(n-i+1), \qquad i = 1,\ldots,n-2.$$

The  $\alpha_i$ 's are the digits of  $\alpha$  in this number system, which we call the  $\varpi$ -system. Any element of  $\mathbb{Z}_{n!}$  can be written uniquely

$$\alpha = \alpha_{n-2} \cdots \alpha_1 \alpha_{0_{\overline{m}}}.$$

Unless  $\alpha_{n-2} = 1$ , the rightmost digits are set to 0, so that the sum always involves n-1 elements, indexed  $0, \ldots, n-2$ .

For example, in  $\mathbb{Z}_{5!}$  the base is  $\{\varpi_{5,0}=1, \varpi_{5,1}=5, \varpi_{5,2}=20, \varpi_{5,3}=60\}$ . The element 84 writes

$$84 = 1 \times 60 + 1 \times 20 + 0 \times 5 + 4 \times 1 = 1104_{\pi}$$

and the element 35 writes

$$35 = 0 \times 60 + 1 \times 20 + 3 \times 5 + 0 \times 1 = 0130_{\pi}$$

*Proof.* For simplicity, we denote  $\varpi_i = \varpi_{n,i}$ . For n = 2, there is a single base element,  $\varpi_0 = 1$ , and the result clearly holds. For  $n \geq 3$ , and  $\alpha \in \mathbb{Z}_{n!}$ , we set

$$\alpha^{(0)} = \alpha$$
.

$$\alpha_i = \alpha^{(i)} \mod (n-i), \quad \alpha^{(i+1)} = \alpha^{(i)} \dim (n-i), \quad i = 0, \dots, n-2,$$

where div denotes the integer division. These relations imply

$$\alpha^{(i)} = (n-i)\alpha^{(i+1)} + \alpha_i, \quad i = 0, \dots, n-2.$$

We multiply by  $\varpi_i$  on both sides. For  $i = 0, \ldots, n-3$ , we use the identity  $\varpi_{i+1} = (n-i)\varpi_i$ , and for i = n-2, we use the identity  $2\varpi_{n-2} = 0$  in  $\mathbb{Z}_2$ , to get

$$\varpi_i \alpha^{(i)} - \varpi_{i+1} \alpha^{(i+1)} = \alpha_i \varpi_i, \quad i = 0, \dots, n-3, 
\varpi_{n-2} \alpha^{(n-2)} = \alpha_{n-2} \varpi_{n-2}.$$

Adding these relations together, and accounting for telescoping cancellation on the left side,

$$\varpi_0 \alpha^{(0)} = \alpha_{n-2} \varpi_{n-2} + \dots + \alpha_0 \varpi_0.$$

We obtain the representation

$$\alpha = \alpha_{n-2}\varpi_{n-2} + \dots + \alpha_0\varpi_0.$$

By construction,  $\alpha_i \in \mathbb{Z}_{n-i}$  for  $i = 0, \dots, n-2$ . The digits  $\alpha_i$  are uniquely determined, so that the representation is unique.

Arithmetics can be performed in the ring  $(\mathbb{Z}_{n!}, +, \times)$  endowed with the  $\varpi$ -system. The computation of the sum and product works in the usual way of positional number systems, using the ring structure of  $\mathbb{Z}_{n-i}$  for the operations on the digits of the operands. There is no carry to propagate after the rightmost digit.

**Lemma 1.** The base elements verify

$$\overline{\omega}_{n,i+k} = \overline{\omega}_{n-k,i}\overline{\omega}_{n,k},\tag{1}$$

$$\sum_{i=0}^{k-1} (n-i-1)\varpi_{n,i} = \varpi_{n,k} - 1, \quad k \in \{1, \dots, n-2\},$$
(2)

$$\sum_{i=0}^{n-2} (n-i-1)\overline{\omega}_{n,i} = -1.$$
(3)

*Proof.* The verification of the first relation is straightforward. For the two other relations, let

$$\xi = \sum_{i=0}^{k-1} \alpha_i \overline{\omega}_{n,i}, \quad \alpha_i = n - i - 1 \in \mathbb{Z}_{n-i}.$$

In  $\mathbb{Z}_{n-i}$ ,  $\alpha_i + 1 = 0$ . Therefore, when computing  $\xi + 1$ , the carry propagates from  $\alpha_0$  up to  $\alpha_{k-1}$ , and the  $\alpha_i$ 's are set to 0. If  $k \leq n-2$ , the digit  $\alpha_k = 0$  gets the carry and is replaced by 1. In this case,  $\xi + 1 = \varpi_{n,k}$ . If k = n-1, there is no carry to propagate after the rightmost digit, and all digits of  $\xi + 1$  are set to 0. In this case,  $\xi + 1 = 0$ .

Corollary 1. For  $\alpha, \alpha' \in \mathbb{Z}_{n!}$ , with digits  $\alpha_i, \alpha_i' \in \mathbb{Z}_{n-i}$ ,

$$\alpha + \alpha' = -1 \iff \alpha_i + \alpha'_i = -1, \quad i = 0, \dots, n-2.$$

*Proof.* We write

$$\alpha + \alpha' = \sum_{i=0}^{n-2} (\alpha_i + \alpha_i') \varpi_{n,i}.$$

In  $\mathbb{Z}_{n-i}$ ,  $\alpha_i + \alpha'_i = -1$  if and only if  $\alpha_i + \alpha'_i = n - i - 1$ . By uniqueness of the decomposition in the  $\varpi$ -system, the result follows from (3).

It can be noted that (3) leads in  $\mathbb{N}$  to the identity

$$\sum_{i=0}^{n-2} (n-i-1) \frac{n!}{(n-i)!} = n! - 1,$$

which is related to the identity

$$\sum_{i=1}^{n-1} i \cdot i! = n! - 1,\tag{4}$$

associated with the factorial number system. Identities (2) and (3) are instances of general identities of mixed radix number systems.

#### 3 Code

The set of permutations on n symbols  $x_1, \ldots, x_n$  is denoted  $S_n$ . From a permutation q on the n-1 symbols  $x_1, \ldots, x_{n-1}$ , n permutations on n symbols are generated by inserting  $x_n$  to the right and cyclically permuting the symbols. The insertion of  $x_n$  to the right defines an injection

$$S_{n-1} \xrightarrow{\iota} S_n$$
 $q = (a_1 \cdots a_{n-1}) \longmapsto (a_1 \cdots a_{n-1} x_n) = \tilde{q}.$ 

We define the cyclic shift  $S: \mathcal{S}_{n-1} \to \mathcal{S}_n$  by  $S = C \circ \iota$ , where  $C: \mathcal{S}_n \to \mathcal{S}_n$  is the circular permutation, so that

$$S^{0}q = (a_{1}a_{2} \cdots a_{n-1}x_{n}) = C^{0}\tilde{q} = \tilde{q},$$

$$S^{1}q = (a_{2} \cdots a_{n-1}x_{n}a_{1}) = C^{1}\tilde{q},$$

$$\vdots$$

$$S^{n-1}q = (x_{n}a_{1}a_{2} \cdots a_{n-1}) = C^{n-1}\tilde{q}.$$

The set  $\mathcal{O}(q) = \{S^0 q, \dots, S^{n-1} q\}$  is the *orbit* of q. As  $S^i = S^j$  is equivalent to  $i = j \mod n$ , the exponents of the cyclic shift are elements of  $\mathbb{Z}_n$ .

**Lemma 2.** The set of permutations  $S_n$  is the disjoint union of the orbits O(q) for  $q \in S_{n-1}$ .

Proof. If  $q, r \in \mathcal{S}_{n-1}$ , their orbits are disjoint subsets of  $\mathcal{S}_n$ . Indeed, if  $S^i q = S^j r$  there exists  $k \in \mathbb{Z}_n$  such that  $S^k q = S^0 r = \tilde{r}$ . The only possibility is k = 0, implying  $S^0 q = \tilde{q} = \tilde{r}$ , and q = r. There are (n-1)! disjoint orbits, each of size n, so that they span  $\mathcal{S}_n$ .

According to Lemma 2, the set  $S_n$  can be generated by cyclic shift. The generation by cyclic shift defines an order on the set of permutations,  $S_n = \{p_0, \ldots, p_{n!-1}\}$ , indexed from 0 (a cyclic order in fact). For this order, the rank  $\alpha$  of a permutation  $p_{\alpha} \in S_n$  is an element of  $\mathbb{Z}_{n!}$ .

The generation by cyclic shift of  $p \in \mathcal{S}_n$  from  $(1) \in \mathcal{S}_1$  can be schematized:

$$\left\{ \begin{array}{c}
p^{(1)} = (1) \xrightarrow{\alpha_{n-2}} p^{(2)} \longrightarrow \cdots \xrightarrow{\alpha_2} p^{(n-2)} \xrightarrow{\alpha_1} p^{(n-1)} \xrightarrow{\alpha_0} p^{(n)} = p, \\
p^{(n-i)} = S_{n-i}^{\alpha_i} p^{(n-i-1)},
\end{array} \right\}$$
(5)

where  $p^{(n-i)} \in \mathcal{S}_{n-i}$  is generated from  $p^{(n-i-1)} \in \mathcal{S}_{n-i-1}$  by the cyclic shift

$$S_{n-i}: \mathcal{S}_{n-i-1} \longrightarrow \mathcal{S}_{n-i}$$

with the exponent  $\alpha_i \in \mathbb{Z}_{n-i}$ .

**Definition 1.** The sequence of exponents associated with successive cyclic shifts leading from  $(1) \in S_1$  to  $p \in S_n$  is the code of p in the  $\varpi$ -system:

$$\alpha = \alpha_{n-2} \cdots \alpha_{0_{\varpi}} \in \mathbb{Z}_{n!}.$$

**Theorem 1.** The rank of a permutation on n symbols generated by cyclic shift is given by its code. A permutation on n symbols generated by cyclic shift is determined by writing its rank in the  $\varpi$ -system.

For example, the permutation  $p_{84} = (51324) \in \mathcal{S}_5$  is generated:

$$(1) \xrightarrow{\alpha_3=1} (21) \xrightarrow{\alpha_2=1} (132) \xrightarrow{\alpha_1=0} (1324) \xrightarrow{\alpha_0=4} (51324).$$

Its code is  $1104_{\varpi} = 84$ .

Proof. We use induction on n. For n=2, in  $\mathcal{S}_2=\{(12),(21)\}$ , the rank of the permutation (12) is  $0=0_{\varpi}$ , and the rank of the permutation (21) is  $1=1_{\varpi}$ . For n>2, let  $p=p_{\alpha}\in\mathcal{S}_n$  of rank  $\alpha$ , generated by cyclic shift from  $q=q_{\beta}\in\mathcal{S}_{n-1}$  of rank  $\beta$ . Then  $p=S_n^{\alpha_0}q$  for some  $\alpha_0\in\mathbb{Z}_n$ ,  $\alpha_0$  being the rank of p within the orbit of q. As the orbits contain n elements and as  $\beta$  is the rank of q in  $\mathcal{S}_{n-1}$ , the rank of p in  $\mathcal{S}_n$  is

$$\alpha = \beta n + \alpha_0 = \beta \varpi_{n,1} + \alpha_0 \varpi_{n,0}.$$

By induction hypothesis, the rank  $\beta$  of q is given by the code

$$\beta = \sum_{i=0}^{n-3} \beta_i \varpi_{n-1,i}, \quad \beta_i \in \mathbb{Z}_{n-1-i}.$$

For k = 1, Eq. (1) gives

$$\varpi_{n,i+1} = \varpi_{n-1,i}\varpi_{n,1},$$

so that

$$\beta \varpi_{n,1} = \sum_{i=0}^{n-3} \beta_i \varpi_{n-1,i} \varpi_{n,1} = \sum_{i=0}^{n-3} \beta_i \varpi_{n,i+1} = \sum_{i=1}^{n-2} \beta_{i-1} \varpi_{n,i}.$$

Let  $\alpha_i = \beta_{i-1}$  for  $i = 1, \dots, n-2$ . As  $\beta_i \in \mathbb{Z}_{n-1-i}$ ,  $\alpha_i \in \mathbb{Z}_{n-i}$ . We obtain that

$$\alpha = \beta \varpi_{n,1} + \alpha_0 \varpi_{n,0} = \sum_{i=0}^{n-2} \alpha_i \varpi_{n,i}, \quad \alpha_i \in \mathbb{Z}_{n-i},$$

is the code of  $p_{\alpha}$ . Conversely, let  $p_{\alpha} \in \mathcal{S}_n$ . We write the rank  $\alpha$  in the  $\varpi$ -system,  $\alpha = \alpha_{n-2} \cdots \alpha_{0_{\varpi}}$ , and use scheme (5) – from right to left – with the exponents  $\alpha_0, \ldots, \alpha_{n-2}$  to determine  $p_{\alpha}$ .

We end the section by a package of algorithms performing the correspondance rank  $\leftrightarrow$  permutation of Theorem 1. Permutations are represented by strings indexed from 1. Algorithm C in Knuth [3] generates  $S_n$  by cyclic shift in a simple version of the scheme described in this section.

```
Int2Num(n, \alpha) { conversion from integer to \varpi-system }
for i \leftarrow 0 to n-2 do
   A[i] \leftarrow \alpha \mod (n-i)
   \alpha \leftarrow \alpha \operatorname{div} (n-i)
end for
return A
Num2Int(n, A) { conversion from \varpi-system to integer }
\alpha \leftarrow 0
base \leftarrow 1
for i \leftarrow 0 to n-2 do
   \alpha \leftarrow \alpha + A[i] * base
   base \leftarrow base * (n - i)
end for
return \alpha
CIRC(m, k, p) { Circular permutation of exponent k on m symbols }
for i \leftarrow 1 to k do
   c \leftarrow p[1]
   for j \leftarrow 2 to m do
      p[j-1] \leftarrow p[j]
   end for
   p[m] \leftarrow c
end for
return p
Pos(m, p) { Position of x_m in a permutation p on m symbols }
for j \leftarrow 1 to m do
   if p[j] = x_m then
       return m-j
   end if
end for
PERM2RANK(n, p) { Find the rank of a given permutation p }
for i \leftarrow 0 to n-2 do
   m \leftarrow n - i
   A[i] \leftarrow \operatorname{Pos}(m, p)
   p \leftarrow \text{CIRC}(m, m - A[i], p)
end for
\alpha \leftarrow \text{Num2Int}(n, A)
return A
Rank2Perm(n, \alpha) { Determine a permutation given its rank \alpha }
A \leftarrow \text{Int2Num}(n, \alpha)
```

```
\begin{array}{l} p \leftarrow x_1 \\ \textbf{for } i \leftarrow n-2 \ \text{downto} \ 0 \ \textbf{do} \\ m \leftarrow n-i \\ p \leftarrow \text{CIRC}(m,A[i],p+x_m) \\ \textbf{end for} \\ \textbf{return} \quad p \\ \\ \\ \text{SETPERM}(n) \ \{ \ \text{Generation of the permutations on} \ n \ \text{symbols} \ \} \\ \textbf{for} \ \alpha \leftarrow 0 \ \text{to} \ n!-1 \ \textbf{do} \\ p \leftarrow \text{RANK2PERM}(n,\alpha) \\ \textbf{end for} \end{array}
```

In the sequel, we assume that the set of permutations  $S_n$  is ordered according to generation by cyclic shift.

$\alpha$	$p_{\alpha}$	$\alpha_2$	$\alpha_1$	$\alpha_0$
0	1234	0	0	0
1	2341	0	0	1
2	3412	0	0	2
3	4123	0	0	3
4	2314	0	1	0
5	3142	0	1	1
6	1423	0	1	2
7	4231	0	1	3
8	3124	0	2	0
9	1243	0	2	1
10	2431	0	2	2
11	4312	0	2	3
12	2134	1	0	0
13	1342	1	0	1
14	3421	1	0	2
15	4213	1	0	3
16	1324	1	1	0
17	3241	1	1	1
18	2413	1	1	2
19	4132	1	1	3
20	3214	1	2	0
21	2143	1	2	1
22	1432	1	$^2$	2
23	4321	1	2	3

Table 1: The codes of the permutations of  $\{1, 2, 3, 4\}$  generated by cyclic shift.

#### 4 k-orbits

In this section, structural properties of  $\mathcal{S}_n$  are described using the  $\varpi$ -system.

**Proposition 2.** Let  $k \in \{0, ..., n-2\}$  and  $p_{\alpha} \in \mathcal{S}_n$  with code  $\alpha \in \mathbb{Z}_{n!}$ . There exists a permutation  $q_{\beta} \in \mathcal{S}_{n-k}$  with code  $\beta \in \mathbb{Z}_{(n-k)!}$  such that

$$\alpha = \beta \varpi_{n,k} + \gamma, \quad \gamma \in \{0, \dots, \varpi_{n,k} - 1\}. \tag{6}$$

The code  $\beta$  is made of the n-k-1 leftmost digits of  $\alpha$ , and  $\gamma$  is made of the k rightmost digits of  $\alpha$ .

*Proof.* We have the decomposition

$$\alpha = \alpha_{n-2} \cdots \alpha_{0\varpi} = \alpha_{n-2} \cdots \alpha_k 0 \cdots 0_{\varpi} + 0 \cdots 0 \alpha_{k-1} \cdots \alpha_{0\varpi} = \tilde{\alpha} + \gamma.$$

Let  $\beta_i = \alpha_{i+k}$  for  $i = 0, \dots, n-k-2$ , so that the  $\beta$ 's are the n-k-1 leftmost digits of  $\alpha$ . As  $\alpha_i \in \mathbb{Z}_{n-i}$ ,  $\beta_i = \alpha_{i+k} \in \mathbb{Z}_{n-k-i}$ . Hence

$$\beta = \sum_{i=0}^{n-k-2} \beta_i \varpi_{n-k,i}, \quad \beta_i \in \mathbb{Z}_{n-k-i},$$

is an element of  $\mathbb{Z}_{(n-k)!}$  which is the code of a permutation  $q_{\beta} \in \mathcal{S}_{n-k}$ . Using relation (1), we obtain

$$\tilde{\alpha} = \sum_{i=k}^{n-2} \alpha_i \varpi_{n,i} = \sum_{i=0}^{n-k-2} \alpha_{i+k} \varpi_{n,i+k} = \sum_{i=0}^{n-k-2} \alpha_{i+k} \varpi_{n-k,i} \varpi_{n,k} = \left(\sum_{i=0}^{n-k-2} \beta_i \varpi_{n-k,i}\right) \varpi_{n,k}.$$

The term

$$\gamma = \sum_{i=0}^{k-1} \alpha_i \varpi_{n,i}$$

is made of the k rightmost digits of  $\alpha$ . It is an element of  $\mathbb{Z}_{n-k+1} \times \cdots \times \mathbb{Z}_n$  ranging from 0 to  $\sum_{i=0}^{k-1} (n-i-1)\varpi_{n,i}$ , which equals  $\varpi_{n,k}-1$  by (2). We obtain

$$\alpha = \tilde{\alpha} + \gamma = \beta \varpi_{n,k} + \gamma.$$

**Definition 2.** For  $k \in \{0, ..., n-2\}$ , and  $q_{\beta} \in \mathcal{S}_{n-k}$ , the k-orbit of  $q_{\beta}$  in  $\mathcal{S}_n$  is the subset

$$\mathcal{O}_{n,k}(q_{\beta}) = \{ p_{\alpha} \in S_n; \quad \alpha = \beta \varpi_{n,k} + \gamma, \quad \gamma = 0, \dots, \varpi_{n,k} - 1 \}.$$

For k = 0, the 0-orbit of  $q \in \mathcal{S}_n$  is  $\{q\}$ . Indeed, for k = 0,  $\varpi_{n,0} = 1$ ,  $\gamma = 0$ , and  $q = p_{\alpha}$ . For  $k \geq 1$ , a k-orbit  $\mathcal{O}_{n,k}(q)$  can be described as the subset of  $\mathcal{S}_n$  generated from  $q \in \mathcal{S}_{n-k}$  by k successive cyclic shifts. Indeed, by Proposition 2, the code of  $p_{\alpha} \in \mathcal{O}_{n,k}(q_{\beta})$  is obtained by appending  $\alpha_{k-1} \cdots \alpha_0$  to the code  $\beta_{n-k-2} \cdots \beta_{0_{\varpi}}$  of  $q_{\beta}$ . By scheme (5), the digits  $\alpha_{k-1}, \ldots, \alpha_0$  describe the generation of  $p_{\alpha}$  from  $q_{\beta}$ . In particular, for k = 1, the 1-orbit  $\mathcal{O}_{n,1}(q)$  of  $q \in \mathcal{S}_{n-1}$  is the orbit  $\mathcal{O}(q)$ . We may further define the (n-1)-orbit  $\mathcal{O}_{n,n-1}(q)$  as the whole set  $\mathcal{S}_n$ , with  $q = (1) \in \mathcal{S}_1$ .

We have the following generalization of Lemma 2:

**Proposition 3.** For  $k \in \{0, ..., n-2\}$ , the set of permutations  $S_n$  is the disjoint union of the k-orbits  $O_{n,k}(q)$  for  $q \in S_{n-k}$ .

*Proof.* The k-orbits are disjoint by uniqueness of the decomposition (6). They are in number (n-k)! and contain  $\varpi_{n,k}$  elements each. As  $(n-k)!\varpi_{n,k}=n!$  in  $\mathbb{N}$ , the k-orbits span  $\mathcal{S}_n$ .  $\square$ 

In decomposition (6),  $\beta$  specifies to which k-orbit  $p_{\alpha}$  belongs and  $\gamma$  specifies the rank of  $p_{\alpha}$  within the k-orbit. The first element of the k-orbit has rank  $\alpha^{first} = \beta \varpi_{n,k}$  (i.e.,  $\gamma = 0$ ). The last element has rank  $\alpha^{last} = \beta \varpi_{n,k} + \varpi_{n,k} - 1$  (i.e.,  $\gamma = \varpi_{n,k} - 1$ ). The element next to the last has rank  $\alpha^{last} + 1 = \beta \varpi_{n,k} + \varpi_{n,k} = (\beta + 1)\varpi_{n,k}$ . It is the first element of the next k-orbit  $\mathcal{O}_{n,k}(q_{\beta+1})$ , where  $q_{\beta+1}$  is the element next to  $q_{\beta}$  in  $\mathcal{S}_{n-k}$ .

**Proposition 4.** For  $k \in \{0, ..., n-2\}$ , the digit  $\alpha_k$  of the code of  $p_\alpha$  is the rank of the k-orbit within the (k+1)-orbit containing  $p_\alpha$ .

For example, Table 1 shows that  $S_4$  contains two 2-orbits within the 3-orbit  $S_4$ . The ranks 0, 1 of these 2-orbits in  $S_4$  are specified by the digit  $\alpha_2$ .

*Proof.* The number of k-orbits within a (k+1)-orbit is n-k (indeed, (n-k)!/(n-(k+1))! = n-k). When performing  $\beta \to \beta + 1$ , the digit  $\beta_0 = \alpha_k$  ranges from 0 to n-k-1 in  $\mathbb{Z}_{n-k}$ . It specifies the rank of the k-orbit within the (k+1)-orbit.

**Lemma 3.** Let  $p_{\alpha} \in \mathcal{S}_n$ . There exists a largest  $k \in \{0, ..., n-2\}$  and  $q_{\beta} \in \mathcal{S}_{n-k}$  such that  $p_{\alpha}$  is the last element of the k-orbit  $\mathcal{O}_{n,k}(q_{\beta})$ , and not the last element of the (k+1)-orbit containing this k-orbit.

Proof. If  $p_{\alpha}$  is not the last element of the 1-orbit it belongs to, it is the last element of the 0-orbit  $\{p_{\alpha}\}$ . In this trivial case, k=0 and  $p_{\alpha}=q_{\beta}$ . Otherwise the last digit of  $p_{\alpha}$  is  $\alpha_0=n-1$ . There exists a largest  $k\geq 1$  such that  $\alpha_i=n-i-1$  for  $i=0,\ldots,k-1$ , and  $\alpha_k\neq n-i-1$ . This means that  $p_{\alpha}$  is the last element of nested j-orbits,  $j=1,\ldots,k$ .  $\square$ 

### 5 Symmetries

For compatibility with the cyclic shift, we adopt the convention that the positions of the symbols in a permutation are computed from the right and are considered as elements of  $\mathbb{Z}_n$  (the position of the last symbol is 0 and the position of the first symbol is n-1).

According to scheme (5), symbol  $x_{n-i}$  ( $i \ge 2$ ) appears at step n-i with the digit  $\alpha_i$  as exponent of the cyclic shift. Its position in the generated permutation  $p^{(n-i)}$  is therefore

$$pos_{n-i}(x_{n-i}, p^{(n-i)}) = \alpha_i.$$

In particular,

$$pos_n(x_n, p^{(n)}) = \alpha_0.$$

For a permutation  $p = (a_1 a_2 \cdots a_{n-1} a_n) \in \mathcal{S}_n$ , we introduce the *mirror image* of p,  $\overline{p} = (a_n a_{n-1} \cdots a_2 a_1)$ .

**Proposition 5.** The permutations  $p_{\alpha}$  and  $p_{\alpha'}$  are the mirror image of one another if and only if

$$\alpha + \alpha' = -1$$
.

For example, in  $\mathbb{Z}_{5!}$  we have 84 + 35 = -1, and in  $S_5$ ,  $p_{84} = (51324)$  is the mirror image of  $p_{35} = (42315)$ .

*Proof.* The proof is by induction on n. For n = 2,  $p_0 = (12)$ ,  $p_1 = (21)$ , and 0 + 1 = 1 = -1 in  $\mathbb{Z}_2$ . Let n > 2. By Proposition 6,

$$\alpha = \beta \varpi_{n,1} + \alpha_0 \varpi_{n,0}, \quad \alpha' = \beta' \varpi_{n,1} + \alpha'_0 \varpi_{n,0}, \quad q_\beta, q_{\beta'} \in \mathcal{S}_{n-1}, \quad \alpha_0, \alpha'_0 \in \mathbb{Z}_n.$$

By Corollary 1, the condition  $\alpha + \alpha' = -1$  is equivalent to  $\beta + \beta' = -1$  and  $\alpha_0 + \alpha'_0 = -1$ . By induction hypothesis,  $q_{\beta}$  is the mirror image of  $q_{\beta'}$  in  $\mathcal{S}_{n-1}$  if and only if  $\beta + \beta' = -1$ . The condition  $\alpha + \alpha'_0 = -1$  is equivalent to  $\alpha'_0 = n - 1 - \alpha_0$ , i.e., the ranks of  $\alpha_0$  and  $\alpha'_0$  are symmetrical in  $\mathbb{Z}_n$ . As these ranks are the positions of symbol  $x_n$  when  $p_{\alpha}$  and  $p_{\alpha'}$  are generated by cyclic shift from  $q_{\beta}$  and  $q_{\beta'}$  respectively, we obtain the result.

**Corollary 2.** The word constructed by concatenating the symbols of the permutations generated by cyclic shift is a palindrome.

*Proof.* Let  $p_{\alpha} \in \mathcal{S}_n$ . The rank symmetrical to  $\alpha$  in  $\mathbb{Z}_{n!}$  is  $(n!-1)-\alpha=-(\alpha+1)$ . By Proposition 5,  $p_{-(\alpha+1)}$  is the mirror image of  $p_{\alpha}$ .

The set  $S_n$  has in fact deeper symmetries, coming from the recursive structure of the k-orbits.

According to Theorem 1, the generation of  $S_n$  by cyclic shift is obtained by performing  $\alpha \to \alpha + 1$  for  $\alpha \in \mathbb{Z}_{n!}$ , and writing  $\alpha$  in the  $\varpi$ -system. This determines each permutation  $p_{\alpha}$ . As  $\alpha$  runs through  $\mathbb{Z}_{n!}$ ,  $p_{\alpha}$  runs through the k-orbits of  $S_n$ . For a fixed k, and by Proposition 4,  $p_{\alpha}$  leaves a k-orbit to enter the next when, in the computation of  $\alpha + 1$ , the carry propagates up to the digit  $\alpha_k$ , incrementing the rank  $\beta$  of the k-orbit. This occurs when  $\alpha = \beta \varpi_{n,k} + \varpi_{n,k} - 1$ .

**Proposition 6.** Any two successive permutations of  $S_n$  are written as

$$p_{\alpha} = \overline{A}B, \qquad p_{\alpha+1} = BA,$$

with an integer  $k \in \{0, ..., n-2\}$  such that

$$|A| = k + 1.$$

For example, in  $S_5$ ,  $p_{39} = (542\underline{31})$  and  $p_{40} = (\underline{31}245)$ , with  $39 = 0134_{\varpi}$  and  $40 = 0200_{\varpi}$ .

*Proof.* If  $p_{\alpha}$  and  $p_{\alpha+1}$  are in the same 1-orbit then

$$p_{\alpha} = (a_1 a_2 \cdots a_n), \qquad p_{\alpha+1} = (a_2 \cdots a_n a_1).$$

The result holds with  $A = (a_1)$ ,  $B = (a_2 \cdots a_n)$ , and this corresponds to k = 0. Otherwise, by Lemma 3, there exists a largest  $k \ge 1$  such that  $p_{\alpha}$  is the last element of a k-orbit, and not

the last element of a (k+1)-orbit. The elements of a k-orbit are generated by successively inserting the symbols  $x_{n-k+1}, \ldots, x_n$  from a permutation  $q_{\beta} \in \mathcal{S}_{n-k}$ . The last element is

$$(x_n \cdots x_{n-k+1} b_1 \cdots b_{n-k}),$$

where  $q_{\beta} = (b_1 \cdots b_{n-k})$  is a permutation of the symbols  $x_1, \ldots, x_{n-k}$ . The first element of the next k-orbit is

$$(c_1 \cdots c_{n-k} x_{n-k+1} \cdots x_n),$$

where  $q_{\beta+1} = (c_1 \cdots c_{n-k})$ . As  $S_{n-k}$  is generated by cyclic shift,  $q_{\beta+1} = C_{n-k}q_{\beta}$ , with  $C_{n-k}$  the circular permutation in  $S_{n-k}$ . We can now write

$$p_{\alpha} = (x_n \cdots x_{n-k+1} b_1 b_2 \cdots b_{n-k}) = \overline{A}B$$
  

$$p_{\alpha+1} = (b_2 \cdots b_{n-k} b_1 x_{n-k+1} \cdots x_n) = BA,$$

where  $A = (b_1 x_{n-k+1} \cdots x_n)$  contains k+1 symbols.

According to the Proposition, k+1 symbols have to be erased to the left of  $p_{\alpha}$  so that the last symbols of  $p_{\alpha}$  match the first symbols of  $p_{\alpha+1}$ . We define the weight  $e_n(\alpha) \in \{1, \ldots, n-1\}$  of the transition  $\alpha \to \alpha + 1$  as the number of symbols of A in the above decomposition of  $p_{\alpha}$  and  $p_{\alpha+1}$ .

We define the  $\varpi$ -ruler sequence as

$$E_n = \{e_n(\alpha); \quad \alpha = 0, \dots, n! - 2\}.$$

**Proposition 7.** The  $\varpi$ -ruler sequence is a palindrome.

Proof. If the ranks of  $\alpha$  and  $\alpha'$  are symmetrical in  $\mathbb{Z}_{n!}$ ,  $\alpha + \alpha' = -1$ , and  $\alpha_i + \alpha'_i = -1$  for  $i = 0, \ldots, n-2$  by Corollary 1. By the definition of  $e_n(\alpha)$ , we want to show that  $e_n(\alpha) = e_n(\alpha' - 1)$ . If  $p_{\alpha}$  is the last element of a k-orbit, then  $\alpha_i = -1$  for  $i = 0, \ldots, k-1$ , so that  $\alpha'_i = 0$  for  $i = 0, \ldots, k-1$ :  $p_{\alpha'}$  is the first element of a k-orbit and  $p_{\alpha'-1}$  is the last element of the previous k-orbit. Hence  $e_n(\alpha) = e_n(\alpha' - 1) = k + 1$ . If  $p_{\alpha}$  is not the last element of a k-orbit, then  $\alpha_0 \neq -1$ ,  $\alpha'_0 \neq 0$ ,  $p_{\alpha'}$  is not the first element of a 1-orbit. In this case  $e_n(\alpha) = e_n(\alpha' - 1) = 1$ .

**Proposition 8.** The number of terms of the  $\varpi$ -ruler sequence such that  $e_n(\alpha) = k$  is

$$(n-k)(n-k)!$$
.

The sum of its n! - 1 terms is

$$W_n = 1! + 2! + \ldots + n! - n.$$

*Proof.* We have  $e_n(\alpha) = k \ge 1$  if and only if  $p_\alpha$  is the last element of a (k-1)-orbit, and not the last element of a k-orbit. The number of (k-1)-orbits within a k-orbit is n-k+1 (see Proposition 4). We exclude the last (k-1)-orbit, giving n-k possibilities. The number of k-orbits is (n-k)! so that there are (n-k)(n-k)! possibilities for  $e_n(\alpha) = k$ .

The formula for the sum is shown by induction. We have  $W_2 = 1 = 1! + 2! - 2$ , and for n > 2,

$$W_n = \sum_{k=1}^{n-1} k(n-k)(n-k)! = \sum_{k=0}^{n-2} (k+1)(n-1-k)(n-1-k)!$$

$$= \sum_{k=1}^{n-2} k(n-1-k)(n-1-k)! + \sum_{k=0}^{n-2} (n-1-k)(n-1-k)!$$

$$= W_{n-1} + \sum_{i=1}^{n-1} i \cdot i! = 1! + \dots + (n-1)! - (n-1) + n! - 1 = 1! + \dots + n! - n.$$

In the last line, we have used identity (4) and the induction hypothesis.

The  $\varpi$ -ruler sequence is analogous to the ruler sequence (sequence A001511 in Sloane [6]). The difference is that the number of intermediate ticks increases with n (Table 2).

n	$E_n$
2	1
3	$1^{2}21^{2}$
4	$1^3 21^3 21^3 31^3 21^3$
5	$1^{4} 2$

Table 2: The  $\varpi$ -ruler sequence for n = 2, 3, 4, 5 (1<sup>j</sup> denotes 1 repeated j times).

#### 6 Combinatorial Gray code

A combinatorial Gray code is a method for generating combinatorial objects so that successive objects differ by some pre-specified adjacency rule involving a minimality criterion (Savage [7]). Such a code can be formulated as an Hamiltonian path or cycle in a graph whose vertices are the combinatorial objects to be generated. Two vertices are joined by an edge if they differ from each other in the pre-specified way.

The code associated with the  $\varpi$ -system corresponds to an Hamiltonian path in a weighted directed graph  $G_n$ .

**Definition 3.** The vertices of the digraph  $G_n$  are the elements of  $S_n$ . For two permutations (vertices)  $p_{\alpha}$  and  $p_{\alpha'}$ , there is an arc from  $p_{\alpha}$  to  $p_{\alpha'}$  if and only if the last symbols of  $p_{\alpha}$  match the first symbols of  $p_{\alpha'}$  (there is no arc when there is no match). Let  $p_{\alpha}, p_{\alpha'} \in S_n$  be two connected vertices in  $G_n$ . The weight  $f_n(\alpha, \alpha') \in \{1, \ldots, n-1\}$  associated with the arc  $(p_{\alpha}, p_{\alpha'})$  is the number of symbols that have to be erased to the left of  $p_{\alpha}$  so that the last symbols of  $p_{\alpha}$  match the first symbols of  $p_{\alpha'}$ .

By Proposition 6, for each  $\alpha$ , there is an arc of weight  $e_n(\alpha) = f_n(\alpha, \alpha + 1)$  joining  $p_\alpha$  to  $p_{\alpha+1}$ . This allows to define the Hamiltonian path

$$\mathbf{w}_n = \{(p_{\alpha}, p_{\alpha+1}); \ \alpha = 0, \dots, n! - 2\}$$

joining successive permutations. This path has total weight  $W_n = 1! + ... + n! - n$  by Proposition 8. The path  $\mathbf{w}_n$  can be closed into an Hamiltonian cycle by joining the last permutation  $p_{n!-1}$  to the first  $p_0$  by an arc of weight n-1:

$$(x_n \cdots x_2 x_1) \xrightarrow{n-1} (x_1 x_2 \cdots x_n).$$

Hence, an oriented path exists from any vertex to any other, so that  $G_n$  is strongly connected. Table 3 displays the weighted adjacency matrix of the digraph  $G_4$  and the Hamiltonian path  $\mathbf{w}_4$ .

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23
0	0	1	2	3	0	0	0	3	0	0	0	3	0	0	2	3	0	0	0	3	0	0	0	3
1	3	0	1	2	0	0	3	0	0	3	0	0	0	3	0	0	3	0	0	2	0	0	3	0
2	2	3	0	1	3	0	0	0	0	2	3	0	3	0	0	0	0	0	3	0	0	3	0	0
3	1	2	3	0	2	3	0	0	3	0	0	0	0	0	3	0	0	3	0	0	3	0	0	0
4	0	0	0	3	0	1	2	3	0	0	0	3	0	0	0	3	0	0	0	3	0	0	2	3
5	0	3	0	O	3	0	1	2	0	0	3	0	3	0	0	2	0	0	3	0	0	3	0	0
6	0	2	3	0	2	3	0	1	3	0	0	0	0	0	3	0	0	3	0	0	3	0	0	0
7	3	0	0	0	1	2	3	0	2	3	0	0	0	3	0	0	3	0	0	0	0	0	3	0
8	0	0	0	3	0	0	0	3	0	1	2	3	0	0	0	3	0	0	2	3	0	0	0	3
9	0	0	3	0	0	3	0	0	3	0	1	2	0	0	3	0	0	3	0	0	3	0	0	2
10	3	0	0	O	0	2	3	0	2	3	0	1	0	3	0	0	3	0	0	0	0	0	3	0
11	2	3	0	0	3	0	0	0	1	2	3	0	3	0	0	0	0	0	3	0	0	3	0	0
12	0	0	2	3	0	0	0	3	0	0	0	3	0	1	2	3	0	0	0	3	0	0	0	3
13	0	3	0	0	3	0	0	2	0	0	3	0	3	0	1	2	0	0	3	0	0	3	0	0
14	3	0	0	0	0	0	3	0	0	3	0	0	2	3	0	1	3	0	0	0	0	2	3	0
15	0	0	3	0	0	3	0	0	3	0	0	0	1	2	3	0	2	3	0	0	3	0	0	0
16	0	0	0	3	0	0	0	3	0	0	2	3	0	0	0	3	0	1	2	3	0	0	0	3
17	3	0	0	2	0	0	3	0	0	3	0	0	0	3	0	0	3	0	1	2	0	0	3	0
18	0	0	3	0	0	3	0	0	3	0	0	0	0	2	3	0	2	3	0	1	3	0	0	0
19	0	3	0	0	3	0	0	0	0	0	3	0	3	0	0	0	1	2	3	0	2	3	0	0
20	0	0	0	3	0	0	2	3	0	0	0	3	0	0	0	3	0	0	0	3	0	1	2	3
21	0	0	3	O	0	3	0	0	3	0	0	2	0	0	3	0	0	3	0	0	3	0	1	2
22	0	3	0	0	3	0	0	0	0	0	3	0	3	0	0	0	0	2	3	0	2	3	0	1
23	3	0	0	0	0	0	3	0	0	3	0	0	2	3	0	0	3	0	0	0	1	2	3	0

Table 3: The adjacency matrix of the digraph  $G_4$ . Lines delineate the 1-orbits. Double lines delineate the 2-orbits. Bold entries on the upper diagonal indicate the Hamiltonian path corresponding to the  $\varpi$ -system code, and forming the  $\varpi$ -ruler sequence.

**Proposition 9.** Each vertex of  $G_n$  has exactly j! in-arcs of weight j and j! out-arcs of weight j, j = 1, ..., n - 1. Hence the vertices of  $G_n$  have  $L = 1! + \cdots + (n - 1)!$  as in- and out-degree, and  $G_n$  is L-regular.

*Proof.* Let us consider the arcs of weight  $j \in \{1, ..., n-1\}$  joining a vertex to another in  $G_n$ :

$$(a_1 \cdots a_j b_1 \cdots b_{n-j}) \xrightarrow{j} (b_1 \cdots b_{n-j} c_1 \cdots c_j),$$

where the c's are a permutation of the a's. There are j! possibilities for the c's, the a's and the b's being fixed. Hence j! arcs of weight j leave each vertex. Similarly, there are j! possibilities for the a's, the b's and the c's being fixed, so that j! arcs of weight j enter each vertex.

We conjecture that the Hamiltonian path  $\mathbf{w}_n$  joining successive permutations in the digraph  $G_n$  is of minimal total weight. Assuming this conjecture we may state:

The  $\varpi$ -system is a combinatorial Gray code listing the permutations generated by cyclic shift. The adjacency rule is that the minimal number of symbols is erased to the left of a permutation so that the last symbols of the permutation match the first symbols of the next permutation.

#### 7 Acknowledgements

Marco Castera initiated the problem which motivated this study. Philippe Paclet discovered the weighted directed graph described in section 6.

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2000 Mathematics Subject Classification: Primary 05A05.

Keywords: permutations, cyclic shift, number system, palindrome, combinatorial Gray code.