

Assignment - 1

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Computer Fundamentals & Architecture

1) what is computer?

→ Computer is an electronic device

- which is used to perform simple as well as complex operations to solve simple as well as complex problem with the help of input and output devices as well as central processing unit (CPU)

2) what are different components of

computer

→ Hardware is thing which is we can touch and feel its hardware

- following are the different components of computer

1) microprocessor

2) math co-processor

3) CPU

4) storage device

→ cutatidha slotnabait valgmo

5] Input device

Svetgmo ei torlu

6] output device

givab simbola no ai valgmo

7] motherboard

an alqmib mifaq of gzu ei dudu

8] systembusa xalgmo an llaw

ma dudu xalgmo an llaw an alqmib
3] what is storage device and types
of storage devices? llaw an givab
(U9) fin



storage device

to fengmo frangib ero torlu

- Storage responsible for storing the data.

an gzu ei dudu paidt ei gnuhroll

- It can used for storing information for temporary and permanent storage.

- There are 2 types of storage device

1) primary storage device

2) Secondary storage device

givab apoto

1) Primary storage device

- Generally storage devices accessible to microprocessor directly are called primary storage device.
- e.g.: RAM

- It is primary electric device which is used to store data temporarily.

- Every program gets loaded into RAM for execution purpose.

Ram

RAM is electrical primary storage device

- Every Ram becoming a Rom is considered as RAM

2) Secondary storage device

- This storage device are used to store the data permanently.

- Before the data get loaded into secondary it initial store in the primary.

- It is compulsory & mandatory to shift the data from primary to secondary permanently.
- when data is not in RAM we can shift into Secondary.

Hard drive

- It is secondary state of memory.

- used to store the data using magnetism.
- when we create the file & save that file it gets into Secondary storage.

Q] What is meant by Toolchain on x86?

Toolchain

Toolchain is considered as a set of software (tools) which are used to convert human understandable program into machine understandable program (binary).

using toolchain convert human understandable code to machine understandable code

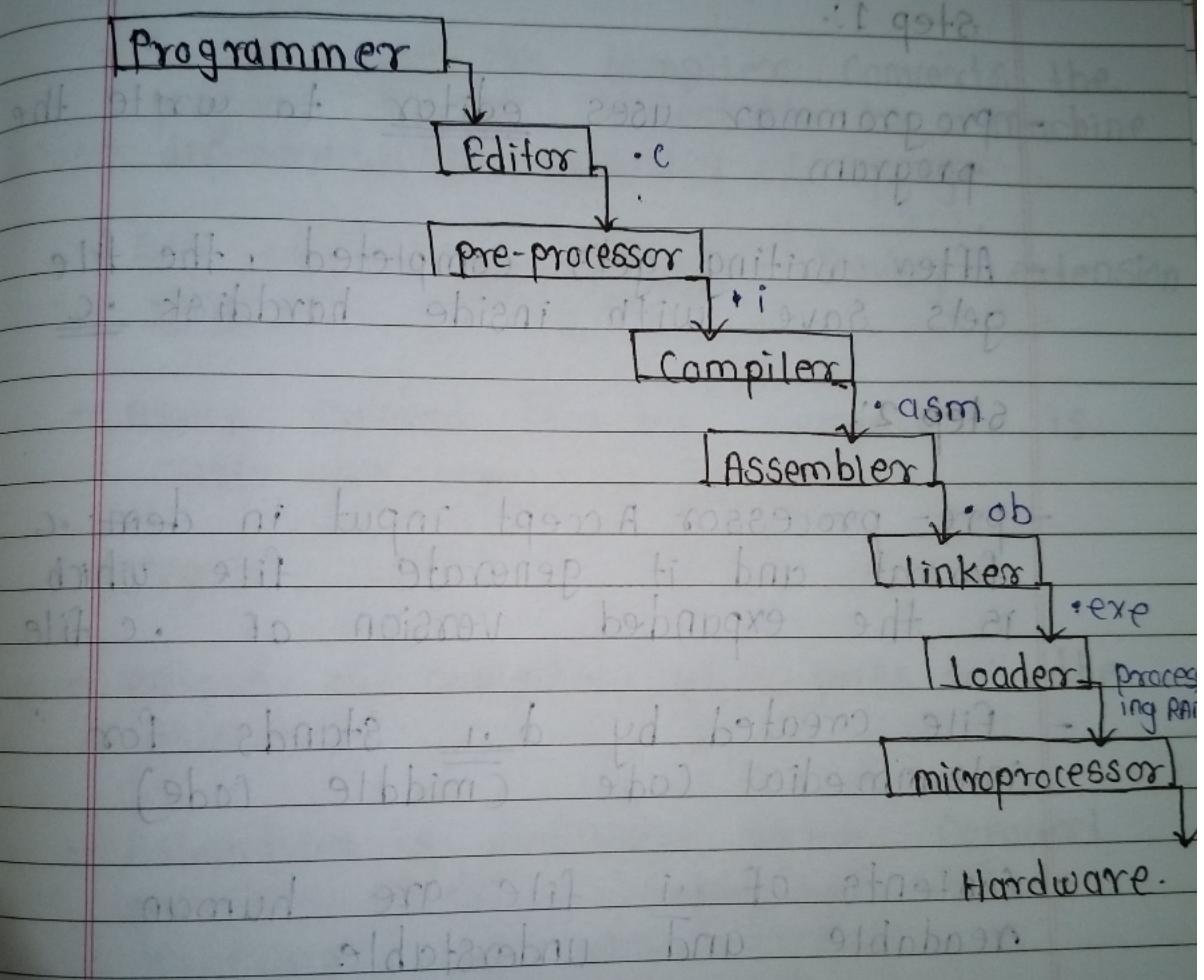


fig: Toolchain X86

Above fig shows the toolchain of x86 architecture of

Step 1:

- programmer uses editor to write the program
- After writing gets completed, the file gets save with inside harddisk .c

Step 2:

pre-processor Accept input in .c format and it generate file which is the expanded version of .c file

- file created by .i stands for intermediate code (middle code)

contents of .i file are human readable and understandable

Step 3:

- output of pre-processor gets provided as input to the compiler

- Compiler is software program from one language to another.
- In our case the compiler converts the program human undestable to machine dependant like Assembly language.
- file created by compiler having extension .asm.
- newly created file by compiler is .asm or .s
- Step 4: ldibbcd abian -b reloc
- output of compiler gets to pass to the Assembler.
- Assembler is Software which converts program machine dependant into machine undestable format.
- output of Assembly having file extension .obj
- .obj contain the code in binary format but it is not directly executable.

Step 5: After assembly of program.

It is now time for linking.

- Linker is responsible to link obj file generate by the Assembler and it generates dependant obj file (e.g.) .
- Linker generates output file with extension .exe

After linking program is formed.

- In our case .exe is output of linker.

Step 6: Now pd if before plug.

Exe is now ready.

- .exe generated by linker currently stored inside harddisk.

To execute any application or program it has to be present or loaded inside RAM.

- After loading the .exe file into RAM it gets considered as process and it gets executed with the help of operating system.

5] what is CPU register? what are types of it? and explain use of each CPU register.

→ CPU register

- CPU registers are small fast storage areas in a computer's central processing (CPU) that temporarily hold data & instructions while the HA (computer) is working.

- following are multiple CPU register

AH, AL = (AL) Arithmetic register

BH, BL = (BL) Base register

CH, CL = (CL) Count register

DH, DL = (DL) Data register

SP = Stack pointer

SI = (source) index register

DI = (destination) index register

BP = base register

ES = extra segment register

CS = code segment register

SS = stack segment register

Arithmetic register (AL) or word [2 bytes to save memory b/w 8 bit to 16 bit]

- AH and AL registers together form 16-bit registers that are paired together to form the 16-bit AL register.

AH is the high byte and AL is the low byte of the 16-bit register.

- AH and AL are 8-bit "char" size registers.

Base register (BL) and pointer.

- BL is known as the base register. It could be used in indexed addressing.
- It can also point to a 16-byte in memory or the beginning of a memory array.

Count register (CH) (CL)

CL is known as the data register. It is also used in input CH, CL store the loop count in iterative operations.

Data register (DL) at anfbaillq

- DL is known as the data register it is also used in input / output operations it is also used with DH and DL registers for multiply and divide operations involving large values

Stack pointer (SP)

- A stack pointer is a register in a processor that stores the memory address of the most recently added item to a stack.
- It's also known as the extended stack pointer.

Source index (SI)

The source index register (SI) is a 16 bit general purpose register in a computer's CPU that's used to identify memory addresses in the data segment that the data register (DS) is addressing.

Destination index (DI) register

The Destination index (DI) register is a 16-bit general purpose register in the microprocessor; it is used as a memory pointer to store the address of data being read or written.

- 6] what is operating system & different tasks of O.S?

→ operating system

operating system is a system software that manages computer hardware & software resources and provides common services for computer program.

- operating system (OS) is the program that after being initially loaded into the computer by a boot program, manages all of the other application programs in a computer.
- following are the tasks of the operating system.

- 1) file management
 - 2) process management
 - 3) memory management
 - 4) CPU scheduling
 - 5) Hardware abstraction
- 7] Explain working of each tool from toolchain i.e. Editor, preprocessor, Compiler, Assembler, linker, loader

→ Editor:

After writing gets completed the file gets save inside harddisk in .c format.

Pre-processor

pre processor Accepts input in .c format it generate file which is the expanded version of .c file

- file created by ~~mpp~~ stands for intermediate code in (middle)
- file generated by pre-processor is human readable & understandable

Compiler

- Compiler Accepts the input from the pre-processor in .i format
- Compiler is software which convert a program from one language to another language
- Compiler convert the program from human understandable format to machine dependant i.e. Assembly language

The output of the compiler is generated in file `.obj` or `.asmlib`

Assembler

- Assembler is software which convert program machine dependant into machine understandable format

Output of Assembly having file extension

• .obj

- .obj file which generated by Assembler is in binary format but not executable.

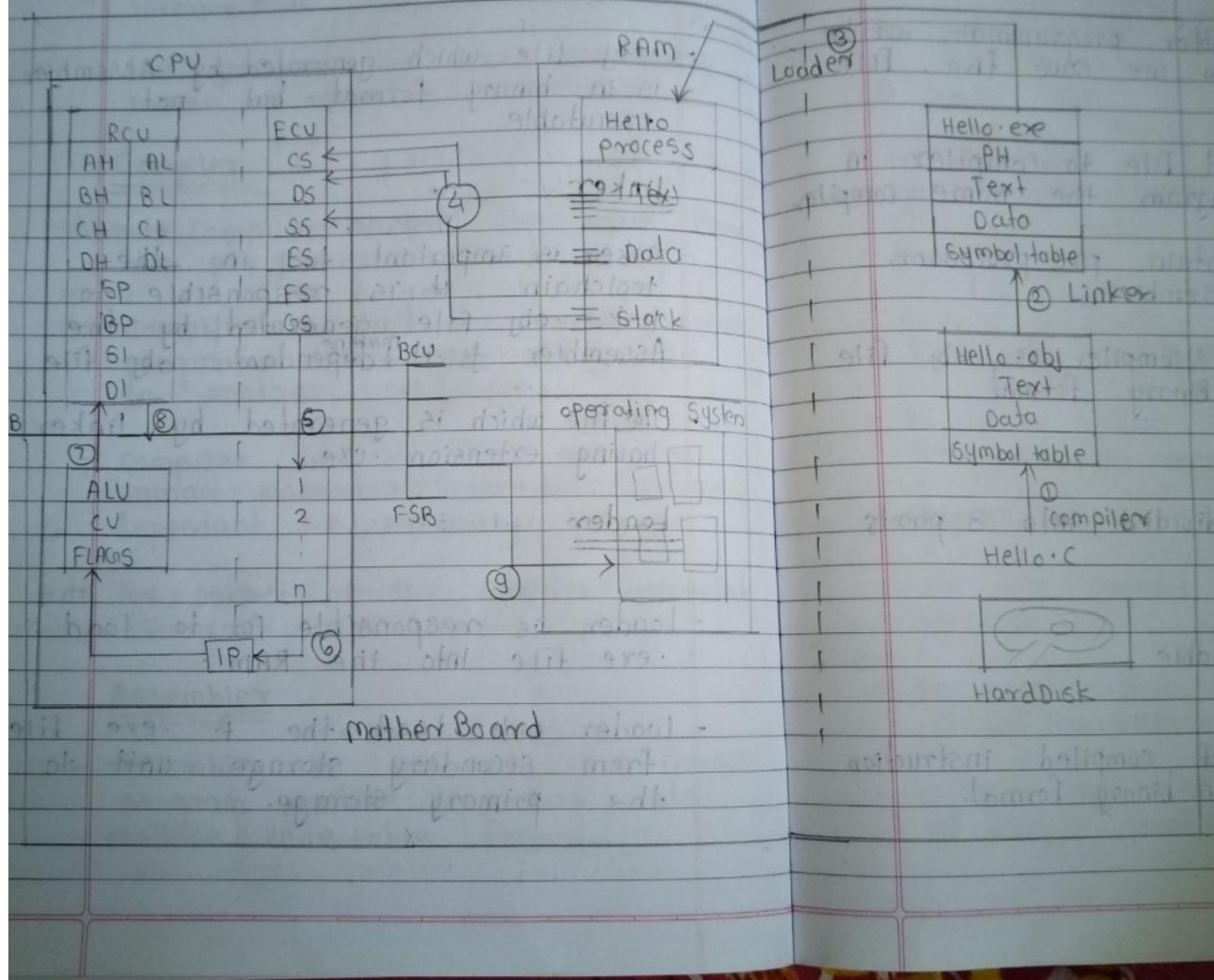
Linker

- Linker is important tool in the toolchain it is responsible to link .obj file generated by the Assembler & it ^{machine} dependant obj file.
- The file which is generated by linker having extension -exe.

Loader

- Loader is responsible for to load .exe file into the RAM
- Loader will load the p. .exe file from secondary storage unit to the primary storage.

8] Explain each step of below diagram



Step 1:-

- By using editor program we write the program we save the file as hello.c
- we pass that file to Compiler in our prog diagram the name compiler internally contain pre-processor, compiler, Assembler
- The output of compiler is .obj file which is in Binary format

Step 2:

- .obj file is divided into 3 parts

- 1) Text
- 2) Data
- 3) Symbol table

1) Text

Content of compiled instruction of a program in binary format.

2) Data:

Data section contains memory for global variable use in program.

3) Symbol table:

It's a table which contains the information about symbol character, special symbols which is used in program.

- Now we pass .obj file to the linker.
- Linker is responsible to link our .obj file with its dependent .obj files.
- The linker creates .exe file which is considered as executable file.

Step 3:-

- .exe file created by linker stored in harddisk.
- To execute that file it should be loaded in RAM.
- To load that file into RAM the loader of OS is responsible.

- Who Inside the .exe file a header is added which contain information about .exe file
- PH contain information about .exe file
- When the loader load that file into RAM it is considered as process
- process into RAM is divided into 3 parts
 - 1) Text
 - 2) Data
 - 3) Stack

The stack section contain information about the function which is written inside our program

- for each function there is a stack frame
- RAM is not responsible for execute the process due to which we pass that process to microprocessor
- inside the microprocessor there is no sufficient memory like RAM due to which direct loading text, data, stack

is not practically possible.

- So that each section is divided into multiple segments
 - Text segment gets copied into CS (Code Segment)
 - Data gets copied into DS (Data Segment)
 - Stack gets copied into SS (Stack Segment)
 - IF the FS is also full then GS are used
 - CFS & CS has no any full form
- Step 5:-
- Now the data from segments gets copied into instruction queue
 - For frame that instruction queue single instruction gets fetch at a time & forward that instruction at the time.

Step 6: Addressing pipeline for ALU

- a) Now the instruction gets forwarded to ALU if instruction is related to the Arithmetic operation.
- b) otherwise bypass step happens first. The instruction gets forwarded to Control unit.

~~Example about an ALU bypass step also.~~

Step 7:-

- In the ALU unit flags indicate the internal instructions after it. instructions gets stored forward to CPU
- Now instruction gets stored in CPU registers.

~~After that there are multiple CPU registers depends on their requirement.~~

Step 8:- Addressing load memory.

- Inside the ALU after CPU register execution of instruction gets performed.

step 8:

- The instructions in the CPU register & the output of instruction will gets forward to IDB about Forwarding

Step 9:

- After getting the output of the instructions are forwarded for the operating system through Bus Control units of the front side bus for sharing data
- After the output getting operating will display the result on console.

q] what are the contents of physical Header?

- physical Header
- By using the physical header operating system will understand to locate the loader or not.

- physical headers store the address of main function.

following include the contents of following physical header field of header

y Text

: P 978

y Data

adt no logno adt pntfop nftn

3) symbol table adt no mftb utani

10] what is meant by Text, Data, and stack section?

→ Text: no thzad adt pntfop nftn

They program into the RAM for EP divided into Text, Data, Stack section

- In the text there is the text format binary format about the executable file is stored

Data: adt no logno adt pntfop nftn

Data section contain memory for global variable in the program

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S stack section mapping are done by compiler to architect are done
The stack section contains information about the function which are written inside our program

- for each function there is stack frame upon so basically it is a function int basic form of function adt

so it can map of form in it's address location so glad edit this into within function of and so many editing or return