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# Team Note of Deobureo Minkyu Party

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#### ALL BELOW HERE ARE USELESS IF YOU READ THE STATEMENT WRONG

#### 1 Flows, Matching

## 1.1 Hopcroft-Karp Bipartite Matching

```
const int MAXN = 50005, MAXM = 50005;
vector<int> gph[MAXN];
int dis[MAXN], 1[MAXN], r[MAXM], vis[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int 1, int r){ gph[1].push_back(r); }
bool bfs(int n){
  queue<int> que:
 bool ok = 0;
 memset(dis, 0, sizeof(dis));
 for(int i=0; i<n; i++){</pre>
   if(l[i] == -1 && !dis[i]){
      que.push(i):
     dis[i] = 1;
   }
  while(!que.empty()){
   int x = que.front();
    que.pop();
    for(auto &i : gph[x]){
     if(r[i] == -1) ok = 1;
      else if(!dis[r[i]]){
        dis[r[i]] = dis[x] + 1;
        que.push(r[i]);
   }
 }
 return ok;
bool dfs(int x){
 if(vis[x]) return 0;
 vis[x] = 1:
 for(auto &i : gph[x]){
   if(r[i] == -1 \mid | (!vis[r[i]] \&\& dis[r[i]] == dis[x] + 1 \&\& dfs(r[i]))){
     l[x] = i: r[i] = x:
     return 1;
 }
 return 0;
int match(int n){
 memset(1, -1, sizeof(1));
 memset(r, -1, sizeof(r));
 int ret = 0;
 while(bfs(n)){
   memset(vis, 0, sizeof(vis));
   for(int i=0; i<n; i++) if(l[i] == -1 && dfs(i)) ret++;
 }
 return ret;
bool chk[MAXN + MAXM]:
void rdfs(int x, int n){
 if(chk[x]) return;
```

```
chk[x] = 1:
 for(auto &i : gph[x]){
   chk[i + n] = 1;
   rdfs(r[i], n);
 }
vector<int> getcover(int n, int m){ // solve min. vertex cover
 match(n):
 memset(chk, 0, sizeof(chk));
 for(int i=0; i<n; i++) if(l[i] == -1) rdfs(i, n);
 vector<int> v:
 for(int i=0; i<n; i++) if(!chk[i]) v.push_back(i);</pre>
 for(int i=n; i<n+m; i++) if(chk[i]) v.push_back(i);</pre>
1.2 Dinic's Algorithm
const int MAXN = 505;
struct edg{ int pos, cap, rev; };
vector<edg> gph[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int s, int e, int x){
 gph[s].push_back({e, x, (int)gph[e].size()});
 gph[e].push_back({s, 0, (int)gph[s].size()-1});
int dis[MAXN], pnt[MAXN];
bool bfs(int src, int sink){
 memset(dis, 0, sizeof(dis));
 memset(pnt, 0, sizeof(pnt));
 queue<int> que:
 que.push(src);
 dis[src] = 1;
 while(!que.empty()){
   int x = que.front();
   que.pop();
   for(auto &e : gph[x]){
     if(e.cap > 0 && !dis[e.pos]){
       dis[e.pos] = dis[x] + 1;
        que.push(e.pos);
 }
 return dis[sink] > 0;
int dfs(int x, int sink, int f){
 if(x == sink) return f:
 for(; pnt[x] < gph[x].size(); pnt[x]++){</pre>
   edg e = gph[x][pnt[x]];
   if(e.cap > 0 \&\& dis[e.pos] == dis[x] + 1){
     int w = dfs(e.pos, sink, min(f, e.cap));
     if(w){
        gph[x][pnt[x]].cap -= w;
       gph[e.pos][e.rev].cap += w;
       return w:
```

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```
return 0:
lint match(int src, int sink){
 lint ret = 0:
 while(bfs(src, sink)){
   while((r = dfs(src, sink, 2e9))) ret += r;
 }
 return ret;
1.3 Min Cost Max Flow
const int MAXN = 100:
struct edg{ int pos, cap, rev, cost; };
vector<edg> gph[MAXN];
void clear(){
 for(int i=0; i<MAXN; i++) gph[i].clear();</pre>
void add_edge(int s, int e, int x, int c){
 gph[s].push_back({e, x, (int)gph[e].size(), c});
 gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
int dist[MAXN], pa[MAXN], pe[MAXN];
bool inque[MAXN];
bool spfa(int src. int sink){
 memset(dist, 0x3f, sizeof(dist));
 memset(inque, 0, sizeof(inque));
 queue<int> que;
 dist[src] = 0;
 inque[src] = 1:
 que.push(src);
 bool ok = 0;
  while(!que.empty()){
   int x = que.front();
   que.pop();
   if(x == sink) ok = 1:
   inque[x] = 0;
   for(int i=0; i<gph[x].size(); i++){</pre>
     edg e = gph[x][i];
     if(e.cap > 0 \&\& dist[e.pos] > dist[x] + e.cost){
       dist[e.pos] = dist[x] + e.cost;
       pa[e.pos] = x;
       pe[e.pos] = i;
       if(!inque[e.pos]){
         inque[e.pos] = 1;
          que.push(e.pos);
     }
   }
 }
 return ok;
int match(int src, int sink){
 int ret = 0:
 while(spfa(src, sink)){
   int cap = 1e9;
   for(int pos = sink; pos != src; pos = pa[pos]){
```

```
cap = min(cap, gph[pa[pos]][pe[pos]].cap);
   ret += dist[sink] * cap;
   for(int pos = sink; pos != src; pos = pa[pos]){
     int rev = gph[pa[pos]][pe[pos]].rev;
     gph[pa[pos]][pe[pos]].cap -= cap;
     gph[pos][rev].cap += cap;
   }
 }
 return ret;
1.4 Hell-Joseon style MCMF
const int MAXN = 100;
struct edg{ int pos, cap, rev, cost; };
vector<edg> gph[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int s, int e, int x, int c){
 gph[s].push_back({e, x, (int)gph[e].size(), c});
 gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
int phi[MAXN], inque[MAXN], dist[MAXN];
void prep(int src, int sink){
 memset(phi, 0x3f, sizeof(phi));
 memset(dist, 0x3f, sizeof(dist)):
 queue<int> que;
 que.push(src);
 inque[src] = 1;
 while(!que.empty()){
   int x = que.front():
   que.pop();
   inque[x] = 0;
   for(auto &i : gph[x]){
     if(i.cap > 0 && phi[i.pos] > phi[x] + i.cost){
        phi[i.pos] = phi[x] + i.cost;
        if(!inque[i.pos]){
         inque[i.pos] = 1;
         que.push(i.pos);
     }
   }
 }
 for(int i=0; i<MAXN; i++){</pre>
   for(auto &j : gph[i]){
      if(j.cap > 0) j.cost += phi[i] - phi[j.pos];
 }
 priority_queue<pi, vector<pi>, greater<pi> > pq;
 pq.push(pi(0, src));
 dist[src] = 0;
 while(!pq.empty()){
   auto 1 = pq.top();
   pq.pop();
   if(dist[l.second] != l.first) continue:
   for(auto &i : gph[l.second]){
     if(i.cap > 0 && dist[i.pos] > 1.first + i.cost){
        dist[i.pos] = l.first + i.cost;
```

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```
pq.push(pi(dist[i.pos], i.pos));
     }
   }
 }
bool vis[MAXN]:
int ptr[MAXN];
int dfs(int pos, int sink, int flow){
 vis[pos] = 1;
 if(pos == sink) return flow;
  for(; ptr[pos] < gph[pos].size(); ptr[pos]++){</pre>
   auto &i = gph[pos][ptr[pos]];
   if(!vis[i.pos] && dist[i.pos] == i.cost + dist[pos] && i.cap > 0){
     int ret = dfs(i.pos, sink, min(i.cap, flow));
     if(ret != 0){
       i.cap -= ret;
        gph[i.pos][i.rev].cap += ret;
       return ret;
     }
   }
 }
 return 0;
int match(int src, int sink, int sz){
  prep(src. sink):
 for(int i=0; i<sz; i++) dist[i] += phi[sink] - phi[src];</pre>
 int ret = 0:
  while(true){
   memset(ptr, 0, sizeof(ptr));
   memset(vis, 0, sizeof(vis));
    int tmp = 0;
    while((tmp = dfs(src, sink, 1e9))){
     ret += dist[sink] * tmp;
     memset(vis, 0, sizeof(vis));
    tmp = 1e9:
    for(int i=0; i<sz; i++){</pre>
     if(!vis[i]) continue:
     for(auto &j : gph[i]){
       if(j.cap > 0 && !vis[j.pos]){
          tmp = min(tmp, (dist[i] + j.cost) - dist[j.pos]);
       }
     }
   }
    if(tmp > 1e9 - 200) break;
   for(int i=0; i<sz; i++){</pre>
      if(!vis[i]) dist[i] += tmp;
   }
 }
 return ret;
1.5 Circulation Problem
maxflow mf:
lint lsum:
void clear(){
 lsum = 0:
```

```
mf.clear():
void add_edge(int s, int e, int l, int r){
 lsum += 1:
 mf.add_edge(s + 2, e + 2, r - 1);
 mf.add edge(0, e + 2, 1):
 mf.add_edge(s + 2, 1, 1);
bool solve(int s. int e){
 mf.add_edge(e+2, s+2, 1e9); // to reduce as maxflow with lower bounds, in circulation
 problem skip this line
 return lsum == mf.match(0, 1):
 // to get maximum LR flow, run maxflow from s+2 to e+2 again
1.6 Min Cost Circulation
// Cycle canceling (Dual of successive shortest path)
// Time complexity is ridiculously high (F * maxC * nm^2). But runs reasonably in practice
(V = 70 \text{ in } 1s)
struct edg{ int pos, cap, rev, cost; };
vector<edg> gph[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add edge(int s, int e, int x, int c){
 gph[s].push_back({e, x, (int)gph[e].size(), c});
 gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
int dist[MAXN], par[MAXN], pae[MAXN];
int negative_cycle(int n){
 bool mark[MAXN] = {};
 memset(dist, 0, sizeof(dist)):
 int upd = -1;
 for(int i=0; i<=n; i++){
   for(int j=0; j<n; j++){
     int idx = 0;
     for(auto &k : gph[j]){
        if(k.cap > 0 && dist[k.pos] > dist[j] + k.cost){
         dist[k.pos] = dist[j] + k.cost;
         par[k.pos] = j;
         pae[k.pos] = idx;
         if(i == n){
           upd = j;
           while(!mark[upd]){
             mark[upd] = 1;
             upd = par[upd];
           return upd;
       }
        idx++;
   }
 }
 return -1;
int match(int n){
 int rt = -1;
 int ans = 0:
```

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```
while(~(rt = negative cvcle(n))){
   bool mark[MAXN] = {};
   vector<pi> cyc;
   while(!mark[rt]){
     cyc.push_back(pi(par[rt], pae[rt]));
     mark[rt] = 1:
     rt = par[rt];
   reverse(cyc.begin(), cyc.end());
   int capv = 1e9;
   for(auto &i : cyc){
     auto e = &gph[i.first][i.second];
     capv = min(capv, e->cap);
   for(auto &i : cvc){
     auto e = &gph[i.first][i.second];
     e->cap -= capv;
     gph[e->pos][e->rev].cap += capv;
     ans += e->cost * capv:
 }
 return ans:
1.7 Gomory-Hu Tree
struct edg{ int s, e, x; };
vector<edg> edgs;
maxflow mf:
void clear(){ edgs.clear(); }
void add_edge(int s, int e, int x){ edgs.push_back({s, e, x}); }
bool vis[MAXN]:
void dfs(int x){
 if(vis[x]) return;
 vis[x] = 1:
 for(auto &i : mf.gph[x]) if(i.cap > 0) dfs(i.pos);
vector<pi> solve(int n){ // i - j cut : i - j minimum edge cost. 0 based.
 vector<pi> ret(n); // if i > 0, stores pair(parent,cost)
 for(int i=1: i<n: i++){</pre>
   for(auto &j : edgs){
     mf.add_edge(j.s, j.e, j.x);
     mf.add_edge(j.e, j.s, j.x);
   ret[i].first = mf.match(i, ret[i].second);
   memset(vis, 0, sizeof(vis));
   dfs(i);
   for(int j=i+1; j<n; j++){</pre>
     if(ret[j].second == ret[i].second && vis[j]){
       ret[j].second = i;
   }
   mf.clear():
 return ret;
```

## 1.8 Blossom Algorithm for General Matching

```
const int MAXN = 2020 + 1;
// 1-based Vertex index
int vis[MAXN], par[MAXN], orig[MAXN], match[MAXN], aux[MAXN], t, N;
vector<int> conn[MAXN];
queue<int> Q;
void addEdge(int u, int v) {
 conn[u].push_back(v); conn[v].push_back(u);
void init(int n) {
 N = n: t = 0:
 for(int i=0; i<=n; ++i) {
   conn[i].clear();
   match[i] = aux[i] = par[i] = 0;
 }
void augment(int u, int v) {
 int pv = v, nv;
 do {
   pv = par[v]; nv = match[pv];
   match[v] = pv; match[pv] = v;
   v = nv:
 } while(u != pv);
int lca(int v, int w) {
 ++t;
 while(true) {
     if(aux[v] == t) return v: aux[v] = t:
     v = orig[par[match[v]]];
   }
   swap(v, w);
 }
void blossom(int v. int w. int a) {
 while(orig[v] != a) {
   par[v] = w; w = match[v];
   if(vis[w] == 1) Q.push(w), vis[w] = 0;
   orig[v] = orig[w] = a;
   v = par[w];
 }
bool bfs(int u) {
 fill(vis+1, vis+1+N, -1); iota(orig + 1, orig + N + 1, 1);
 Q = queue < int > (); Q.push(u); vis[u] = 0;
 while(!Q.emptv()) {
   int v = Q.front(); Q.pop();
   for(int x: conn[v]) {
     if(vis[x] == -1) {
       par[x] = v; vis[x] = 1;
       if(!match[x]) return augment(u, x), true:
       Q.push(match[x]); vis[match[x]] = 0;
      else if(vis[x] == 0 && orig[v] != orig[x]) {
       int a = lca(orig[v], orig[x]);
       blossom(x, v, a); blossom(v, x, a);
```

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```
}
   }
 }
  return false;
int Match() {
  int ans = 0;
  // find random matching (not necessary, constant improvement)
  vector<int> V(N-1); iota(V.begin(), V.end(), 1);
  shuffle(V.begin(), V.end(), mt19937(0x94949));
  for(auto x: V) if(!match[x]){
   for(auto v: conn[x]) if(!match[v]) {
      match[x] = y, match[y] = x;
      ++ans: break:
   }
 }
  for(int i=1: i<=N: ++i) if(!match[i] && bfs(i)) ++ans:</pre>
  return ans;
1.9 Blossom Algorithm for Weighted General Matching
// N^3 (but fast in practice)
static const int INF = INT MAX:
static const int N = 514;
struct edge{
  int u,v,w; edge(){}
  edge(int ui,int vi,int wi)
    :u(ui).v(vi).w(wi){}
};
int n.n x:
edge g[N*2][N*2];
int lab[N*2];
int match[N*2],slack[N*2],st[N*2],pa[N*2];
int flo_from[N*2][N+1],S[N*2],vis[N*2];
vector<int> flo[N*2];
aueue<int> a:
int e_delta(const edge &e){
  return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
void update_slack(int u,int x){
  if(!slack[x]||e_delta(g[u][x])<e_delta(g[slack[x]][x]))slack[x]=u;</pre>
void set slack(int x){
  slack[x]=0:
  for(int u=1;u<=n;++u)</pre>
   if(g[u][x].w>0&&st[u]!=x&&S[st[u]]==0)
      update slack(u.x):
}
void q_push(int x){
  if(x \le n)q.push(x);
  else for(size_t i=0;i<flo[x].size();i++)</pre>
    q_push(flo[x][i]);
void set st(int x.int b){
  if(x>n)for(size_t i=0;i<flo[x].size();++i)</pre>
    set_st(flo[x][i],b);
```

```
int get_pr(int b,int xr){
 int pr=find(flo[b].begin(),flo[b].end(),xr)-flo[b].begin();
 if(pr%2==1){
   reverse(flo[b].begin()+1,flo[b].end());
   return (int)flo[b].size()-pr:
 }else return pr;
void set_match(int u,int v){
 match[u]=g[u][v].v;
 if(u<=n) return;</pre>
 edge e=g[u][v];
 int xr=flo_from[u][e.u],pr=get_pr(u,xr);
 for(int i=0;i<pr;++i)set_match(flo[u][i],flo[u][i^1]);</pre>
 set_match(xr,v);
 rotate(flo[u].begin(),flo[u].begin()+pr,flo[u].end());
void augment(int u,int v){
 for(::){
   int xnv=st[match[u]];
   set_match(u,v);
   if(!xnv)return:
   set_match(xnv,st[pa[xnv]]);
   u=st[pa[xnv]],v=xnv;
int get_lca(int u,int v){
 static int t=0:
 for(++t;u||v;swap(u,v)){
   if(u==0)continue:
   if(vis[u]==t)return u;
   vis[u]=t:
   u=st[match[u]];
   if(u)u=st[pa[u]];
 return 0:
void add blossom(int u.int lca.int v){
 int b=n+1:
 while(b \le n_x \& st[b])++b;
 if(b>n x)++n x:
 lab[b]=0,S[b]=0;
 match[b]=match[lca]:
 flo[b].clear():
 flo[b].push_back(lca);
 for(int x=u,y;x!=lca;x=st[pa[y]])
   flo[b].push_back(x),flo[b].push_back(y=st[match[x]]),q_push(y);
 reverse(flo[b].begin()+1,flo[b].end());
 for(int x=v.v:x!=lca:x=st[pa[v]])
   flo[b].push_back(x),flo[b].push_back(y=st[match[x]]),q_push(y);
 set_st(b,b);
 for(int x=1;x<=n_x;++x)g[b][x].w=g[x][b].w=0;
 for(int x=1;x<=n;++x)flo_from[b][x]=0;</pre>
 for(size t i=0:i<flo[b].size():++i){</pre>
   int xs=flo[b][i];
   for(int x=1:x\leq n x:++x)
```

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```
if(g[b][x].w==0||e_delta(g[xs][x]) < e_delta(g[b][x]))
       g[b][x]=g[xs][x],g[x][b]=g[x][xs];
   for(int x=1;x<=n;++x)</pre>
     if(flo_from[xs][x])flo_from[b][x]=xs;
 }
 set slack(b):
void expand_blossom(int b){
 for(size_t i=0;i<flo[b].size();++i)</pre>
   set_st(flo[b][i],flo[b][i]);
 int xr=flo_from[b][g[b][pa[b]].u],pr=get_pr(b,xr);
  for(int i=0:i<pr:i+=2){</pre>
   int xs=flo[b][i],xns=flo[b][i+1];
   pa[xs]=g[xns][xs].u:
   S[xs]=1,S[xns]=0;
   slack[xs]=0,set_slack(xns);
   q_push(xns);
 S[xr]=1,pa[xr]=pa[b];
  for(size_t i=pr+1;i<flo[b].size();++i){</pre>
   int xs=flo[b][i];
   S[xs]=-1.set slack(xs):
 }
 st[b]=0;
bool on_found_edge(const edge &e){
 int u=st[e.u].v=st[e.v]:
 if(S[v]==-1){}
   pa[v]=e.u,S[v]=1;
   int nu=st[match[v]]:
   slack[v]=slack[nu]=0;
   S[nu]=0,q_push(nu);
 else if(S[v]==0){
   int lca=get_lca(u,v);
   if(!lca)return augment(u,v), augment(v,u), true;
   else add blossom(u.lca.v):
 }
 return false:
bool matching(){
 memset(S+1,-1,sizeof(int)*n x):
 memset(slack+1,0,sizeof(int)*n_x);
 q=queue<int>();
 for(int x=1;x<=n_x;++x)</pre>
   if (st[x]==x\&\&!match[x])pa[x]=0,S[x]=0,q_push(x);
 if(q.empty())return false;
 for(;;){
   while(q.size()){
     int u=q.front();q.pop();
     if(S[st[u]]==1)continue;
     for(int v=1; v<=n;++v)</pre>
       if(g[u][v].w>0&&st[u]!=st[v]){
          if(e_delta(g[u][v])==0){
            if(on_found_edge(g[u][v]))return true;
          }else update_slack(u,st[v]);
```

```
int d=INF;
    for(int b=n+1;b<=n_x;++b)</pre>
      if(st[b] == b&&S[b] == 1)d=min(d, lab[b]/2);
    for(int x=1;x<=n_x;++x)</pre>
      if(st[x]==x&&slack[x]){
        if(S[x]==-1)d=min(d,e_delta(g[slack[x]][x]));
        else if(S[x]==0)d=min(d,e_delta(g[slack[x]][x])/2);
    for(int u=1;u<=n;++u){
      if(S[st[u]]==0){
        if(lab[u]<=d)return 0:
        lab[u]-=d;
     }else if(S[st[u]]==1)lab[u]+=d:
   for(int b=n+1;b<=n_x;++b)</pre>
     if(st[b]==b){
        if(S[st[b]]==0)lab[b]+=d*2;
        else if(S[st[b]]==1)lab[b]-=d*2:
    q=queue<int>();
    for(int x=1:x\leq n x:++x)
      if(st[x]==x\&\&slack[x]\&\&st[slack[x]]!=x\&\&e_delta(g[slack[x]][x])==0)
        if(on_found_edge(g[slack[x]][x]))return true;
   for(int b=n+1:b \le n x:++b)
      if(st[b]==b\&\&S[b]==1\&\&lab[b]==0)expand_blossom(b);
 return false:
pair<long long,int> solve(){
 memset(match+1,0,sizeof(int)*n);
 n x=n:
 int n_matches=0;
 long long tot_weight=0;
 for(int u=0;u<=n;++u)st[u]=u,flo[u].clear();</pre>
 int w max=0:
 for(int u=1;u<=n;++u)
   for(int v=1:v<=n:++v){
     flo_from[u][v]=(u==v?u:0);
      w_max=max(w_max,g[u][v].w);
  for(int u=1;u<=n;++u)lab[u]=w_max;</pre>
  while(matching())++n_matches;
 for(int u=1:u<=n:++u)
   if (match[u]&&match[u]<u)
      tot_weight+=g[u][match[u]].w;
 return make_pair(tot_weight,n_matches);
void add edge( int ui , int vi , int wi ){
 g[ui][vi].w = g[vi][ui].w = wi;
void init( int _n ){
 n = _n;
 for(int u=1:u<=n:++u)
   for(int v=1;v<=n;++v)</pre>
      g[u][v]=edge(u,v,0);
```

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```
2 Graph
2.1 2-SAT
strongly_connected scc;
int n; // = number of clauses
void init(int _n){ scc.clear(); n = _n; }
int NOT(int x) { return x \ge n ? (x - n) : (x + n); }
void add_edge(int x, int y){ // input ~x to denote NOT
 if((x >> 31) \& 1) x = (^x) + n;
 if((y >> 31) & 1) y = (~y) + n;
 scc.add_edge(x, y), scc.add_edge(NOT(y), NOT(x));
bool satisfy(vector<bool> &res){
 res.resize(n):
  scc.get_scc(2*n);
  for(int i=0; i<n; i++){</pre>
   if(scc.comp[i] == scc.comp[NOT(i)]) return 0;
   if(scc.comp[i] < scc.comp[NOT(i)]) res[i] = 0;</pre>
   else res[i] = 1:
 }
 return 1;
2.2 BCC
void color(int x, int p){
 if(p){
   bcc[p].push_back(x);
    cmp[x].push_back(p);
  for(auto &i : gph[x]){
    if(cmp[i].size()) continue;
    if(low[i] >= dfn[x]){
     bcc[++c].push_back(x);
      cmp[x].push_back(c);
      color(i, c);
   }
    else color(i, p);
}
2.3 Splay Tree + Link-Cut Tree
// Checklist 1. Is it link cut, or splay?
// Checklist 2. In link cut, is son always root?
void rotate(node *x){
 if(!x->p) return;
  push(x->p); // if there's lazy stuff
 push(x);
  node *p = x->p;
  bool is_left = (p->l == x);
  node *b = (is_left ? x->r : x->l);
  x->p = p->p;
 if(x-p \&\& x-p->1 == p) x-p->1 = x;
  if(x->p && x->p->r == p) x->p->r = x;
  if(is_left){
   if(b) b \rightarrow p = p;
   p->1 = b;
   p->p = x;
```

```
x->r = p;
 }
  else{
   if(b) b \rightarrow p = p;
   p->r = b;
   p->p = x;
   x->1 = p;
 pull(p); // if there's something to pull up
 pull(x);
 if(!x->p) root = x; // IF YOU ARE SPLAY TREE
 if(p->pp){ // IF YOU ARE LINK CUT TREE
   x->pp = p->pp;
   p->pp = NULL;
 }
void splay(node *x){
 while(x->p){
   node *p = x->p;
   node *g = p->p;
   if(g){
      if((p\rightarrow 1 == x) \hat{(g\rightarrow 1 == p)}) rotate(x);
      else rotate(p);
   }
   rotate(x);
 }
void access(node *x){
 splay(x);
 push(x);
 if(x->r){
   x->r->pp = x;
   x->r->p = NULL;
   x->r = NULL;
 pull(x);
 while(x->pp){
   node *nxt = x->pp;
    splay(nxt);
    push(nxt);
   if(nxt->r){
     nxt->r->pp = nxt;
     nxt->r->p = NULL;
     nxt->r = NULL;
   nxt->r = x;
   x->p = nxt;
   x->pp = NULL;
   pull(nxt);
    splay(x);
node *root(node *x){
 access(x):
  while(x->1){
    push(x);
```

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```
x = x -> 1:
 }
  access(x);
 return x;
node *par(node *x){
  access(x);
 if(!x->1) return NULL;
  push(x);
 x = x -> 1;
  while(x->r){
   push(x);
   x = x->r;
  access(x);
 return x;
node *lca(node *s, node *t){
  access(s):
  access(t);
  splay(s);
 if(s->pp == NULL) return s;
 return s->pp;
void link(node *par, node *son){
  access(par);
  access(son):
  son->rev ^= 1: // remove if needed
  push(son);
 son->1 = par;
 par->p = son;
 pull(son);
void cut(node *p){
  access(p);
  push(p):
 if(p->1){
   p->1->p = NULL;
   p->1 = NULL;
  pull(p);
2.4 Offline Dynamic MST
int n, m, q;
int st[MAXN], ed[MAXN], cost[MAXN], chk[MAXN];
pi qr[MAXN];
bool cmp(int &a, int &b){ return pi(cost[a], a) < pi(cost[b], b); }</pre>
void contract(int s, int e, vector<int> v, vector<int> &must_mst, vector<int> &maybe_mst){
  sort(v.begin(), v.end(), cmp);
 vector<pi> snapshot;
  for(int i=s; i<=e; i++) disj.uni(st[qr[i].first], ed[qr[i].first], snapshot);</pre>
  for(auto &i : v) if(disj.uni(st[i], ed[i], snapshot)) must_mst.push_back(i);
  disj.revert(snapshot);
  for(auto &i : must_mst) disj.uni(st[i], ed[i], snapshot);
```

```
for(auto &i : v) if(disj.uni(st[i], ed[i], snapshot)) maybe_mst.push_back(i);
 disj.revert(snapshot);
void solve(int s, int e, vector<int> v, lint cv){
 if(s == e){}
    cost[qr[s].first] = qr[s].second;
   if(st[qr[s].first] == ed[qr[s].first]){
     printf("%lld\n", cv);
     return;
   }
   int minv = qr[s].second;
   for(auto &i : v) minv = min(minv, cost[i]);
   printf("%lld\n".minv + cv):
   return;
 }
 int m = (s+e)/2;
 vector<int> lv = v, rv = v;
 vector<int> must mst. maybe mst:
 for(int i=m+1; i<=e; i++){</pre>
   chk[qr[i].first]--;
   if(chk[qr[i].first] == 0) lv.push_back(qr[i].first);
 vector<pi> snapshot;
 contract(s, m, lv, must_mst, maybe_mst);
 lint lcv = cv;
 for(auto &i : must_mst) lcv += cost[i], disj.uni(st[i], ed[i], snapshot);
 solve(s, m, maybe_mst, lcv);
 disj.revert(snapshot);
 must_mst.clear(); maybe_mst.clear();
 for(int i=m+1; i<=e; i++) chk[qr[i].first]++;</pre>
 for(int i=s; i<=m; i++){</pre>
   chk[qr[i].first]--;
   if(chk[qr[i].first] == 0) rv.push_back(qr[i].first);
 lint rcv = cv:
 contract(m+1, e, rv, must_mst, maybe_mst);
 for(auto &i : must_mst) rcv += cost[i], disj.uni(st[i], ed[i], snapshot);
 solve(m+1, e, maybe_mst, rcv);
 disj.revert(snapshot);
 for(int i=s; i<=m; i++) chk[qr[i].first]++;</pre>
int main(){
 scanf("%d %d",&n,&m);
 vector<int> ve:
 for(int i=0; i<m; i++){</pre>
   scanf("%d %d %d",&st[i],&ed[i],&cost[i]);
 scanf("%d",&g);
 for(int i=0; i<q; i++){</pre>
   scanf("%d %d",&qr[i].first,&qr[i].second);
   gr[i].first--;
   chk[qr[i].first]++;
 }
 disj.init(n);
```

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```
for(int i=0: i<m: i++) if(!chk[i]) ve.push back(i):</pre>
 solve(0, q-1, ve, 0);
2.5 Dominator Tree
vector<int> E[MAXN], RE[MAXN], rdom[MAXN];
int S[MAXN], RS[MAXN], cs;
int par[MAXN], val[MAXN], sdom[MAXN], rp[MAXN], dom[MAXN];
void clear(int n) {
 cs = 0:
 for(int i=0;i<=n;i++) {</pre>
   par[i] = val[i] = sdom[i] = rp[i] = dom[i] = S[i] = RS[i] = 0:
   E[i].clear(); RE[i].clear(); rdom[i].clear();
 }
void add_edge(int x, int y) { E[x].push_back(y); }
void Union(int x, int y) { par[x] = y; }
int Find(int x, int c = 0) {
 if(par[x] == x) return c ? -1 : x;
 int p = Find(par[x], 1);
 if(p == -1) return c ? par[x] : val[x];
 if(sdom[val[x]] > sdom[val[par[x]]]) val[x] = val[par[x]];
 par[x] = p:
 return c ? p : val[x];
void dfs(int x) {
 RS[S[x] = ++cs] = x;
 par[cs] = sdom[cs] = val[cs] = cs:
 for(int e : E[x]) {
   if(S[e] == 0) dfs(e), rp[S[e]] = S[x];
   RE[S[e]].push_back(S[x]);
 }
}
int solve(int s. int *up) { // Calculate idoms
 dfs(s);
 for(int i=cs:i:i--) {
   for(int e : RE[i]) sdom[i] = min(sdom[i], sdom[Find(e)]);
   if(i > 1) rdom[sdom[i]].push_back(i);
   for(int e : rdom[i]) {
     int p = Find(e);
     if(sdom[p] == i) dom[e] = i;
     else dom[e] = p;
   }
   if(i > 1) Union(i, rp[i]);
 for(int i=2;i<=cs;i++) if(sdom[i] != dom[i]) dom[i] = dom[dom[i]];</pre>
 for(int i=2:i<=cs:i++) up[RS[i]] = RS[dom[i]]:</pre>
 return cs;
2.6 Global Min-Cut
int minimum cut phase(int n. int &s. int &t. vector<vector<int>> &adi. vector<int> vis){
 vector<int> dist(n);
 int mincut = 1e9;
 while(true){
```

```
int pos = -1, cur = -1e9:
   for(int i=0; i<n; i++){</pre>
     if(!vis[i] && dist[i] > cur){
        cur = dist[i]:
       pos = i:
   if(pos == -1) break;
   s = t:
   t = pos;
   mincut = cur;
   vis[pos] = 1:
   for(int i=0; i<n; i++){</pre>
     if(!vis[i]) dist[i] += adi[pos][i]:
   }
 }
 return mincut; // optimal s-t cut here is, {t} and V \ {t}
int solve(int n. vector<vector<int>> adi){
 if(n <= 1) return 0:
 vector<int> vis(n);
 int ans = 1e9:
 for(int i=0; i<n-1; i++){</pre>
   int s, t;
   ans = min(ans, minimum_cut_phase(n, s, t, adj, vis));
   vis[t] = 1:
   for(int j=0; j<n; j++){</pre>
     if(!vis[i]){
        adi[s][i] += adi[t][i];
       adj[j][s] += adj[j][t];
   }
   adj[s][s] = 0;
 }
 return ans:
2.7 Edmond's Directed MST
 Should be revised.
// starts from node 0. assumes there exists at least one dmst.
// edge is reversed : if there is edge s -> e, INSERT IN gph[e]
struct edge { int to, cost, id; };
using elist = vector<edge>;
void dmst(vector<elist> &g, vector<int>& res) {
 const int n = g.size();
 vector<edge*> to(n):
 vector<int> u(n, 0);
 for (int i = 1; i < n; ++i) {
   int mn = g[i][0].cost;
   for (int j = 0; j < g[i].size(); ++j)
     mn = min(mn, g[i][j].cost);
   for (int j = 0; j < g[i].size(); ++j) {
     if (g[i][j].cost == mn)
       to[i] = &g[i][j];
     g[i][j].cost -= mn;
```

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```
for (int i = 1; i < n; ++i) {
   if (u[i]) continue;
   int x = i:
   vector<int> order(1, x);
   u[x] = 1:
    while (to[x]->to > 0 && u[to[x]->to] == 0) {
     x = to[x] \rightarrow to;
     u[x] = 1:
     order.push_back(x);
   int y = to[x] \rightarrow to;
   vector<int> cycle(find(order.begin(), order.end(), y), order.end());
   if (cycle.size() == 0) continue;
   vector<int> in_cycle(n);
   for (int j = 0; j < cycle.size(); ++j) in_cycle[cycle[j]] = 1;</pre>
   vector<int> nw_id(n);
   for (int j = 1; j < n; ++j) nw_id[j] = nw_id[j - 1] + !in_cycle[j];
   int nn = n - cvcle.size():
   vector<elist> gn(nn + 1);
   for (int j = 1; j < n; ++j) {
     if (in_cycle[j]) {
       for (int k = 0; k < g[j].size(); ++k)
          if (!in_cvcle[g[j][k].to])
            gn[nn].push_back({nw_id[g[j][k].to], g[j][k].cost, g[j][k].id});
     } else {
       for (int k = 0; k < g[j].size(); ++k)
          if (in_cycle[g[j][k].to])
            gn[nw_id[j]].push_back({nn, g[j][k].cost, g[j][k].id});
            gn[nw_id[j]].push_back({nw_id[g[j][k].to], g[j][k].cost, g[j][k].id});
     }
   }
   dmst(gn, res);
    set<int> used_e(res.begin(), res.end());
    for (int j = 0; j < cycle.size(); ++j) {</pre>
     bool found = false;
     for (int k = 0; k < g[cycle[j]].size(); ++k)</pre>
       found |= used_e.count(g[cycle[j]][k].id);
      if (found) {
       for (int k = 0; k < cycle.size(); ++k) {
         if (k != j)
            res.push_back(to[cycle[k]]->id);
       }
       return;
     }
   }
 for (int i = 1; i < n; ++i) res.push_back(to[i]->id);
2.8 Vizing's Theorem
namespace Vizing{ // returns edge coloring in adjacent matrix G. 1 - based
 int C[MAXN] [MAXN], G[MAXN] [MAXN];
 void clear(int N){
   for(int i=0; i<=N; i++){
```

```
for(int j=0; j<=N; j++) C[i][j] = G[i][j] = 0;</pre>
   }
 }
 void solve(vector<pi> &E, int N, int M){
   int X[MAXN] = {}, a;
   auto update = [\&] (int u) { for(X[u] = 1; C[u][X[u]]; X[u]++); };
   auto color = [&](int u, int v, int c){
     int p = G[u][v];
     G[u][v] = G[v][u] = c;
     C[u][c] = v; C[v][c] = u;
     C[u][p] = C[v][p] = 0;
     if(p) X[u] = X[v] = p;
     else update(u), update(v);
     return p:
   };
    auto flip = [&] (int u, int c1, int c2){
      int p = C[u][c1];
      swap(C[u][c1], C[u][c2]);
     if( p ) G[u][p] = G[p][u] = c2;
     if( !C[u][c1] ) X[u] = c1;
     if( !C[u][c2] ) X[u] = c2;
     return p;
   };
   for(int i = 1; i <= N; i++) X[i] = 1;
   for(int t = 0; t < E.size(); t++){}
      int u = E[t].first, v0 = E[t].second, v = v0, c0 = X[u], c = c0, d;
     vector<pi> L;
     int vst[MAXN] = {}:
      while(!G[u][v0]){
        L.emplace_back(v, d = X[v]);
        if(!C[v][c]) for(a = (int)L.size()-1; a >= 0; a--) c = color(u, L[a].first, c);
        else if(!C[u][d])for(a=(int)L.size()-1;a>=0;a--)color(u,L[a].first,L[a].second);
        else if( vst[d] ) break;
        else vst[d] = 1, v = C[u][d];
      if( !G[u][v0] ){
       for(;v; v = flip(v, c, d), swap(c, d));
        if(C[u][c0]){
         for(a = (int)L.size()-2; a >= 0 && L[a].second != c; a--);
         for(; a >= 0; a--) color(u, L[a].first, L[a].second);
        } else t--:
   }
 }
   Strings
3.1 Aho-Corasick Algorithm
const int MAXN = 100005, MAXC = 26;
int trie[MAXN][MAXC], fail[MAXN], term[MAXN], piv;
void init(vector<string> &v){
 memset(trie, 0, sizeof(trie));
 memset(fail, 0, sizeof(fail));
 memset(term, 0, sizeof(term));
 piv = 0;
 for(auto &i : v){
```

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```
int p = 0:
    for(auto &j : i){
     if(!trie[p][j]) trie[p][j] = ++piv;
     p = trie[p][j];
   term[p] = 1;
 }
  queue<int> que;
  for(int i=0; i<MAXC; i++){</pre>
   if(trie[0][i]) que.push(trie[0][i]);
  while(!que.empty()){
   int x = que.front();
   que.pop():
   for(int i=0; i<MAXC; i++){</pre>
     if(trie[x][i]){
        int p = fail[x];
        while(p && !trie[p][i]) p = fail[p];
        p = trie[p][i];
        fail[trie[x][i]] = p;
        if(term[p]) term[trie[x][i]] = 1;
        que.push(trie[x][i]);
   }
 }
bool query(string &s){
  int p = 0:
  for(auto &i : s){
   while(p && !trie[p][i]) p = fail[p];
   p = trie[p][i];
   if(term[p]) return 1;
 }
 return 0;
}
3.2 Suffix Array
const int MAXN = 500005:
int ord[MAXN], nord[MAXN], cnt[MAXN], aux[MAXN];
void solve(int n, char *str, int *sfx, int *rev, int *lcp){
 int p = 1;
 memset(ord, 0, sizeof(ord));
 for(int i=0; i<n; i++){</pre>
   sfx[i] = i;
   ord[i] = str[i];
 int pnt = 1;
  while(1){
   memset(cnt, 0, sizeof(cnt));
   for(int i=0; i<n; i++) cnt[ord[min(i+p, n)]]++;</pre>
   for(int i=1; i<=n || i<=255; i++) cnt[i] += cnt[i-1];
   for(int i=n-1: i>=0: i--)
     aux[--cnt[ord[min(i+p, n)]]] = i;
    memset(cnt, 0, sizeof(cnt));
   for(int i=0; i<n; i++) cnt[ord[i]]++;</pre>
    for(int i=1; i<=n || i<=255; i++) cnt[i] += cnt[i-1];
   for(int i=n-1; i>=0; i--)
```

```
sfx[--cnt[ord[aux[i]]]] = aux[i]:
   if(pnt == n) break;
   pnt = 1;
   nord[sfx[0]] = 1;
   for(int i=1; i<n; i++){
     if(ord[sfx[i-1]] != ord[sfx[i]] || ord[sfx[i-1] + p] != ord[sfx[i] + p]){
     nord[sfx[i]] = pnt;
   memcpy(ord, nord, sizeof(int) * n);
   p *= 2;
 for(int i=0: i<n: i++) rev[sfx[i]] = i:</pre>
 int h = 0;
 for(int i=0; i<n; i++){</pre>
   if(rev[i]){
     int prv = sfx[rev[i] - 1];
     while(str[prv + h] == str[i + h]) h++;
     lcp[rev[i]] = h;
   }
   h = \max(h-1, 0);
 }
3.3 Manacher's Algorithm
const int MAXN = 1000005:
int aux[2 * MAXN - 1];
void solve(int n, int *str, int *ret){
// *ret : number of nonobvious palindromic character pair
 for(int i=0; i<n; i++){
   aux[2*i] = str[i];
   if(i != n-1) aux[2*i+1] = -1:
 int p = 0, c = 0;
 for(int i=0; i<2*n-1; i++){
   int cur = 0;
   if(i \le p) cur = min(ret[2 * c - i], p - i);
   while(i - cur - 1 >= 0 && i + cur + 1 < 2*n-1 && aux[i-cur-1] == aux[i+cur+1]){
   }
   ret[i] = cur;
   if(i + ret[i] > p){
     p = i + ret[i];
     c = i;
   }
 }
3.4 Suffix Array (Linear time)
 Should be revised.
class SuffixArray {
public:
 int A[7 * N / 10], B[7 * N / 10], cnt[N + 2], SAV[N];
 int mem[5 * N]; int* mem_pt = mem;
 void clear(int n){
   int *ptr = mem;
```

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```
while(ptr != mem pt){
     *ptr = 0;
     ptr++;
   mem_pt = mem;
   for(int i=0: i<n+10 && i < 7 * N / 10: i++) A[i] = B[i] = 0:
   for(int i=0; i< n+2; i++) cnt[i] = 0;
   for(int i=0; i<n; i++) SAV[i] = 0;
 inline int* mloc(size_t sz) {
   int* ret = mem_pt; mem_pt = mem_pt + sz;
   return ret:
 void rsort(int* a, int* b, int* dat, int n, int k) {
   for (int i = 0; i <= k; i++) cnt[i] = 0;
   for (int i = 0; i < n; i++) SAV[i] = dat[a[i]], cnt[SAV[i]]++;
   for (int i = 1; i <= k; i++) cnt[i] += cnt[i - 1];
   for (int i = n - 1; i \ge 0; i--) b[--cnt[SAV[i]]] = a[i];
 }
#define I(x) ((x)\%3==1)?((x)/3):((x)/3+num1)
#define I2(x) (x<num1)?(3*x+1):(3*(x-num1)+2)
 static int cmp(int x, int y, int str[], int A[], int num1) {
   if (x \% 3 == 1) {
     if (y \% 3 == 2) return A[I(x)] < A[I(y)];
     else return str[x] < str[y] \mid \mid str[x] == str[y] && A[I(x + 1)] < A[I(y + 1)];
   }
     return str[x] < str[y] \mid \mid str[x] == str[y] && cmp(x + 1, y + 1, str, A, num1);
 }
 void make(int* str, int* sa, int n, int k) {
   if (n == 0) return:
   int num1 = (n + 2) / 3, num2 = n / 3;
   int num = num1 + num2:
   str[n] = str[n + 1] = str[n + 2] = 0:
   int *nsa = mloc(num), *nstr = mloc(num + 3);
   for (int i = 0, j = 0; i < n; i++) if (i % 3) A[j++] = i;
   if (n \% 3 == 1) A[num - 1] = n;
   rsort(A, B, str + 2, num, k); rsort(B, A, str + 1, num, k); rsort(A, B, str, num, k);
   int cnt = 1:
   nstr[I(B[0])] = 1:
   for (int i = 1; i < num; i++) {
     int c = B[i], p = B[i - 1];
     if (str[p] != str[c] || str[p + 1] != str[c + 1] || str[p + 2] != str[c + 2]) cnt++;
     nstr[I(c)] = cnt;
   if (cnt == num) for (int i = 0; i < num; i++) nsa[nstr[i] - 1] = i;
   else make(nstr, nsa, num, cnt);
   for (int i = 0, j = 0; i < num; i++) if (nsa[i] < num1) A[j++] = 3 * nsa[i];
   rsort(A. B. str. num1, k):
   for (int i = 0; i < num; i++) A[nsa[i]] = i, nsa[i] = I2(nsa[i]);</pre>
   A[num] = -1:
```

```
merge(B, B + num1, nsa + (n \% 3 == 1), nsa + num, sa, [&](int x, int y) {
     return cmp(x, y, str, A, num1);
   });
   return;
 }
}sa:
3.5 eertree
int nxt[MAXN][26];
int par[MAXN], len[MAXN], slink[MAXN], ptr[MAXN], diff[MAXN], series[MAXN], piv;
void clear(int n = MAXN){
 memset(par, 0, sizeof(int) * n);
 memset(len, 0, sizeof(int) * n);
 memset(slink, 0, sizeof(int) * n);
 memset(nxt, 0, sizeof(int) * 26 * n);
 piv = 0;
void init(int n, char *a){
 par[0] = 0:
 par[1] = 1;
 a[0] = -1;
 len[0] = -1:
 piv = 1;
  int cur = 1:
  for(int i=1; i<=n; i++){
    while(a[i] != a[i - len[cur] - 1]) cur = slink[cur];
    if(!nxt[cur][a[i]]){
     nxt[cur][a[i]] = ++piv;
      par[piv] = cur:
      len[piv] = len[cur] + 2;
      int lnk = slink[cur];
      while (a[i] != a[i - len[lnk] - 1])
       lnk = slink[lnk];
      if(nxt[lnk][a[i]]) lnk = nxt[lnk][a[i]]:
      if(len[piv] == 1 || lnk == 0) lnk = 1;
      slink[piv] = lnk;
      diff[piv] = len[piv] - len[lnk];
      if(diff[piv] == diff[lnk]) series[piv] = series[lnk];
      else series[piv] = piv;
    cur = nxt[cur][a[i]];
   ptr[i] = cur;
 }
int query(int s. int e){
 int pos = ptr[e];
  while(len[pos] >= e - s + 1){
        if(len[pos] % diff[pos] == (e - s + 1) % diff[pos] &&
           len[series[pos]] <= e - s + 1) return true;</pre>
        pos = series[pos]:
        pos = slink[pos];
 }
 return false;
vector<pi> minimum_partition(int n){ // (odd min, even min)
```

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int solve(string a, string b){

```
vector < pi > dp(n + 1):
 vector<pi> series_ans(n + 10);
 dp[0] = pi(1e9 + 1, 0);
 for(int i=1; i<=n; i++){</pre>
   dp[i] = pi(1e9 + 1, 1e9);
   for(int j=ptr[i]; len[j] > 0;){
     int slv = slink[series[j]];
     series_ans[j] = dp[i - (len[slv] + diff[j])];
     if(diff[j] == diff[slink[j]]){
       series_ans[j].first = min(series_ans[j].first, series_ans[slink[j]].first);
       series_ans[j].second = min(series_ans[j].second, series_ans[slink[j]].second);
      auto val = series_ans[i];
     dp[i].first = min(dp[i].first, val.second + 1):
     dp[i].second = min(dp[i].second, val.first + 1);
     j = slv;
 }
 return dp;
3.6 Circular LCS
string s1. s2:
int dp[4005][2005];
int nxt[4005][2005]:
int n, m;
void reroot(int px){
 int py = 1;
 while(py <= m && nxt[px][py] != 2) py++;</pre>
 nxt[px][py] = 1;
 while(px < 2 * n && py < m){
   if(nxt[px+1][py] == 3){
     px++;
     nxt[px][py] = 1;
   else if(nxt[px+1][py+1] == 2){
     px++;
     py++;
     nxt[px][py] = 1;
    else py++;
 while(px < 2 * n && nxt[px+1][py] == 3){
   px++;
   nxt[px][py] = 1;
int track(int x, int y, int e){ // use this routine to find LCS as string
 int ret = 0:
 while(y != 0 \&\& x != e){
   if(nxt[x][v] == 1) v--:
   else if (nxt[x][y] == 2) ret += (s1[x] == s2[y]), x--, y--;
   else if(nxt[x][v] == 3) x--:
 return ret;
```

```
n = a.size(), m = b.size();
 s1 = "#" + a + a:
 s1 = '#' + b;
 for(int i=0: i<=2*n: i++){
   for(int j=0; j<=m; j++){</pre>
     if(i == 0){
       nxt[i][j] = 3;
       continue;
     if(i == 0){
       nxt[i][j] = 1;
       continue:
     dp[i][j] = -1;
     if(dp[i][j] < dp[i][j-1]){
       dp[i][j] = dp[i][j-1];
       nxt[i][j] = 1;
     if(dp[i][j] < dp[i-1][j-1] + (s1[i] == s2[j])){
       dp[i][j] = dp[i-1][j-1] + (s1[i] == s2[j]);
       nxt[i][j] = 2;
     if(dp[i][j] < dp[i-1][j]){
       dp[i][j] = dp[i-1][j];
       nxt[i][j] = 3;
   }
 int ret = dp[n][m];
 for(int i=1; i<n; i++){
   reroot(i), ret = max(ret, track(n+i, m, i));
 }
 return ret:
  Geometry
4.1 Smallest Enclosing Circle / Sphere
namespace cover_2d{
 double eps = 1e-9;
 using Point = complex<double>;
 struct Circle{ Point p; double r; };
 double dist(Point p, Point q){ return abs(p-q); }
 double area2(Point p, Point q){ return (conj(p)*q).imag(); }
 bool in(const Circle& c. Point p) { return dist(c.p. p) < c.r + eps; }
 Circle INVAL = Circle{Point(0, 0), -1};
 Circle mCC(Point a, Point b, Point c){
   b -= a; c -= a;
   double d = 2*(conj(b)*c).imag(); if(abs(d)<eps) return INVAL;</pre>
   Point ans = (c*norm(b) - b*norm(c)) * Point(0, -1) / d:
   return Circle{a + ans, abs(ans)};
 Circle solve(vector<Point> p) {
   mt19937 gen(0x94949); shuffle(p.begin(), p.end(), gen);
   Circle c = INVAL:
```

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```
for(int i=0: i<p.size(): ++i) if(c.r<0 |\cdot| in(c. p[i])){
      c = Circle{p[i], 0};
      for(int j=0; j<=i; ++j) if(!in(c, p[j])){</pre>
        Circle ans{(p[i]+p[j])*0.5, dist(p[i], p[j])*0.5};
        if(c.r == 0) { c = ans; continue; }
        Circle 1. r: 1 = r = INVAL:
        Point pq = p[j]-p[i];
        for(int k=0; k<=j; ++k) if(!in(ans, p[k])) {</pre>
          double a2 = area2(pq, p[k]-p[i]);
          Circle c = mCC(p[i], p[j], p[k]);
          if(c.r<0) continue;</pre>
          else if(a2 > 0 && (1.r<0||area2(pq, c.p-p[i]) > area2(pq, l.p-p[i]))) 1 = c;
          else if(a2 < 0 && (r.r<0||area2(pq, c.p-p[i]) < area2(pq, r.p-p[i]))) r = c;
        if(1.r<0\&\&r.r<0) c = ans;
        else if(1.r<0) c = r;
        else if(r.r<0) c = 1:
        else c = 1.r<=r.r?1:r;
     }
   }
    return c;
};
namespace cover 3d{
  double enclosing_sphere(vector<double> x, vector<double> y, vector<double> z){
    int n = x.size():
    auto hyp = [](double x, double y, double z){
     return x * x + y * y + z * z;
   };
    double px = 0, py = 0, pz = 0;
    for(int i=0; i<n; i++){</pre>
      px += x[i];
     py += y[i];
      pz += z[i]:
    px *= 1.0 / n;
    pv *= 1.0 / n:
    pz *= 1.0 / n;
    double rat = 0.1, maxv;
    for(int i=0: i<10000: i++){
      maxv = -1;
      int maxp = -1;
      for(int j=0; j<n; j++){
        double tmp = hyp(x[j] - px, y[j] - py, z[j] - pz);
        if(maxv < tmp){</pre>
          maxv = tmp;
          maxp = j;
        }
      px += (x[maxp] - px) * rat;
      py += (y[maxp] - py) * rat;
      pz += (z[maxp] - pz) * rat;
      rat *= 0.998:
    return sqrt(maxv);
```

```
}
};
4.2 3D Convex Hull
struct vec3{
 11 x, y, z;
 vec3(): x(0), y(0), z(0) {}
 vec3(11 a, 11 b, 11 c): x(a), y(b), z(c) {}
 vec3 operator*(const vec3& v) const{ return vec3(y*v.z-z*v.y, z*v.x-x*v.z, x*v.y-y*v.x); }
 vec3 operator-(const vec3& v) const{ return vec3(x-v.x, y-v.y, z-v.z); }
 vec3 operator-() const{ return vec3(-x, -y, -z); }
 11 dot(const vec3 &v) const{ return x*v.x+v*v.v+z*v.z: }
};
struct twoset {
 int a, b;
 void insert(int x) { (a == -1 ? a : b) = x : }
 bool contains(int x) { return a == x || b == x; }
 void erase(int x) { (a == x ? a : b) = -1; }
 int size() { return (a != -1) + (b != -1); }
} E[MAXN] [MAXN]; // i < j</pre>
struct face{
 vec3 norm:
 ll disc:
 int I[3];
};
face make_face(int i, int j, int k, int ii, vector<vec3> &A){ // p^T * norm < disc</pre>
 E[i][j].insert(k); E[i][k].insert(j); E[j][k].insert(i);
 face f; f.I[0]=i, f.I[1]=j, f.I[2]=k;
 f.norm = (A[i]-A[i])*(A[k]-A[i]);
 f.disc = f.norm.dot(A[i]):
 if(f.norm.dot(A[ii])>f.disc){
   f.norm = -f.norm;
   f.disc = -f.disc:
 }
 return f;
vector<face> get_hull(vector<vec3> &A){
 int N = A.size();
 vector<face> faces; memset(E, -1, sizeof(E));
 faces.push_back(make_face(0,1,2,3,A));
 faces.push_back(make_face(0,1,3,2,A));
 faces.push_back(make_face(0,2,3,1,A));
 faces.push back(make face(1,2,3,0,A)):
 for(int i=4; i<N; ++i){</pre>
   for(int j=0; j<faces.size(); ++j){</pre>
     face f = faces[i];
      if(f.norm.dot(A[i])>f.disc){
       E[f.I[0]][f.I[1]].erase(f.I[2]):
        E[f.I[0]][f.I[2]].erase(f.I[1]);
       E[f.I[1]][f.I[2]].erase(f.I[0]);
       faces[j--] = faces.back();
        faces.pop_back();
```

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```
}
    int nf = faces.size();
    for(int j=0; j<nf; ++i){</pre>
      face f=faces[j];
      for(int a=0; a<3; ++a) for(int b=a+1; b<3; ++b){
        int c=3-a-b:
        if(E[f.I[a]][f.I[b]].size()==2) continue;
        faces.push_back(make_face(f.I[a], f.I[b], i, f.I[c], A));
   }
 }
  return faces:
     Dynamic Convex Hull Trick
using line_t = double;
const line_t is_query = -1e18;
struct Line {
 line t m. b:
  mutable function<const Line*()> succ;
  bool operator<(const Line& rhs) const {
   if (rhs.b != is_query) return m < rhs.m;</pre>
    const Line* s = succ();
    if (!s) return 0:
   line t x = rhs.m:
   return b - s->b < (s->m - m) * x;
 }
};
struct HullDynamic : public multiset<Line> { // will maintain upper hull for maximum
  bool bad(iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
      return v->m == z->m && v->b <= z->b:
   }
    auto x = prev(y);
    if (z == end()) return y->m == x->m && y->b <= x->b;
    return (x-b - y-b)*(z-m - y-m) >= (y-b - z-b)*(y-m - x-m);
  void insert_line(line_t m, line_t b) {
    auto v = insert({ m, b }):
    y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
    if (bad(y)) { erase(y); return; }
    while (next(y) != end() && bad(next(y))) erase(next(y));
    while (y != begin() && bad(prev(y))) erase(prev(y));
  line t querv(line t x) {
    auto 1 = *lower_bound((Line) { x, is_query });
    return 1.m * x + 1.b:
 }
}H;
4.4 Half-plane Intersection
const double eps = 1e-8;
typedef pair < long double, long double > pi;
```

```
bool z(long double x) { return fabs(x) < eps: }
struct line{
 long double a, b, c;
 bool operator<(const line &1)const{</pre>
   bool flag1 = pi(a, b) > pi(0, 0);
   bool flag2 = pi(1.a, 1.b) > pi(0, 0);
   if(flag1 != flag2) return flag1 > flag2;
   long double t = ccw(pi(0, 0), pi(a, b), pi(l.a, l.b));
   return z(t) ? c * hypot(1.a, 1.b) < 1.c * hypot(a, b) : t > 0;
 pi slope(){ return pi(a, b); }
pi cross(line a, line b){
 long double det = a.a * b.b - b.a * a.b;
 return pi((a.c * b.b - a.b * b.c) / det, (a.a * b.c - a.c * b.a) / det);
bool bad(line a, line b, line c){
 if(ccw(pi(0, 0), a.slope(), b.slope()) <= 0) return false;</pre>
 pi crs = cross(a, b):
 return crs.first * c.a + crs.second * c.b >= c.c;
bool solve(vector<line> v, vector<pi> &solution){ // ax + by <= c;</pre>
 sort(v.begin(), v.end());
 deque<line> dq;
 for(auto &i : v){
   if(!dq.empty() && z(ccw(pi(0, 0), dq.back().slope(), i.slope()))) continue;
   while(dq.size() >= 2 && bad(dq[dq.size()-2], dq.back(), i)) dq.pop_back();
   while(dq.size() >= 2 \&\& bad(i, dq[0], dq[1])) dq.pop_front();
   dq.push_back(i);
 while(dq.size() > 2 && bad(dq[dq.size()-2], dq.back(), dq[0])) dq.pop_back();
 while(dq.size() > 2 && bad(dq.back(), dq[0], dq[1])) dq.pop_front();
 vector<pi> tmp;
 for(int i=0; i<dq.size(); i++){</pre>
   line cur = dq[i], nxt = dq[(i+1)%dq.size()];
   if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps) return false;</pre>
   tmp.push_back(cross(cur, nxt));
 solution = tmp;
 return true;
4.5 Point-in-polygon test / Point-to-polygon tangent
// C : counter_clockwise(C[0] == C[N]), N >= 3
// return highest point in C <- P(clockwise) or -1 if strictly in P
// polygon is strongly convex, C[i] != P
int convex_tangent(vector<pi> &C, pi P, int up = 1){
 auto sign = [&](lint c){ return c > 0 ? up : c == 0 ? 0 : -up; };
 auto local = [&](pi P, pi a, pi b, pi c) {
   return sign(ccw(P, a, b)) \le 0 \&\& sign(ccw(P, b, c)) >= 0;
 };
 int N = C.size()-1, s = 0, e = N, m:
 if( local(P, C[1], C[0], C[N-1]) ) return 0;
 while(s+1 < e){
   m = (s+e) / 2:
   if( local(P, C[m-1], C[m], C[m+1]) ) return m;
   if ( sign(ccw(P, C[s], C[s+1])) < 0) { // up}
```

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```
if (sign(ccw(P, C[m], C[m+1])) > 0) e = m:
      else if( sign(ccw(P, C[m], C[s])) > 0) s = m;
      else e = m:
   }
    else{ // down
      if (sign(ccw(P, C[m], C[m+1])) < 0) s = m;
      else if( sign(ccw(P, C[m], C[s])) < 0) s = m;
      else e = m:
   }
  if( s && local(P, C[s-1], C[s], C[s+1]) ) return s;
 if( e != N && local(P. C[e-1], C[e], C[e+1]) ) return e:
  return -1;
4.6 kd-tree
typedef pair<int, int> pi;
struct node{
 pi pnt;
  int spl, sx, ex, sy, ey;
}tree[270000];
pi a[100005];
int n, ok[270000];
lint sqr(int x){ return 111 * x * x; }
bool cmp1(pi a, pi b){ return a < b; }</pre>
bool cmp2(pi a, pi b){ return pi(a.second, a.first) < pi(b.second, b.first); }</pre>
// init(0, n-1, 1) : Initialize kd-tree
// set dap = INF, and call solve(1, P). dap = (closest point from P)
void init(int s, int e, int p){ // Initialize kd-tree
  int minx = 1e9, maxx = -1e9, miny = 1e9, maxy = -1e9;
  int m = (s+e)/2;
  for(int i=s; i<=e; i++){</pre>
   minx = min(minx, a[i].first):
   miny = min(miny, a[i].second);
   maxx = max(maxx, a[i].first);
   maxv = max(maxv. a[i].second):
  tree[p].spl = (maxx - minx < maxy - miny);</pre>
  sort(a+s, a+e+1, [&](const pi &a, const pi &b){
   return tree[p].spl ? cmp2(a, b) : cmp1(a, b);
 }):
  ok[p] = 1;
  tree[p] = {a[m], tree[p].spl, minx, maxx, miny, maxy};
 if (s \le m-1) init(s, m-1, 2*p):
  if(m+1 \le e) init(m+1, e, 2*p+1);
lint dap = 3e18;
void solve(int p, pi x){ // find closest point from point x (L^2)
  if(x != tree[p].pnt) dap = min(dap, sqr(x.first - tree[p].pnt.first) + sqr(x.second -
  tree[p].pnt.second));
  if(tree[p].spl){
   if(!cmp2(tree[p].pnt, x)){
```

```
if(ok[2*p]) solve(2*p, x):
     if(ok[2*p+1] && sqr(tree[2*p+1].sy - x.second) < dap) solve(2*p+1, x);
   }
   else{
     if(ok[2*p+1]) solve(2*p+1, x);
     if(ok[2*p] && sqr(tree[2*p].ey - x.second) < dap) solve(2*p, x);
 }
 else{
   if(!cmp1(tree[p].pnt, x)){
     if(ok[2*p]) solve(2*p, x);
     if(ok[2*p+1] && sqr(tree[2*p+1].sx - x.first) < dap) solve(2*p+1, x);
   }
   else{
     if(ok[2*p+1]) solve(2*p+1, x);
      if(ok[2*p] && sqr(tree[2*p].ex - x.first) < dap) solve(2*p, x);
 }
5 Math
5.1 FFT / NTT
typedef complex<double> base;
void fft(vector<base> &a. bool inv){
 int n = a.size(), j = 0;
 vector<base> roots(n/2);
 for(int i=1: i<n: i++){
   int bit = (n >> 1);
   while(j >= bit){
     j -= bit;
     bit >>= 1;
   j += bit;
   if(i < j) swap(a[i], a[j]);</pre>
 double ang = 2 * acos(-1) / n * (inv ? -1 : 1);
 for(int i=0: i<n/2: i++){
   roots[i] = base(cos(ang * i), sin(ang * i));
 }
 /* In NTT, let prr = primitive root. Then,
 int ang = ipow(prr, (mod - 1) / n);
 if(inv) ang = ipow(ang, mod - 2);
 for(int i=0; i<n/2; i++){
   roots[i] = (i ? (111 * roots[i-1] * ang % mod) : 1);
 Others are same. If there is /= n, do *= ipow(n, mod - 2).
 In XOR convolution, roots[*] = 1.
 for(int i=2; i<=n; i<<=1){
   int step = n / i;
   for(int i=0: i<n: i+=i){
     for(int k=0; k<i/2; k++){
       base u = a[j+k], v = a[j+k+i/2] * roots[step * k];
       a[j+k] = u+v;
       a[j+k+i/2] = u-v;
```

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```
}
 }
 if(inv) for(int i=0; i<n; i++) a[i] /= n;
vector<lint> multiplv(vector<lint> &v. vector<lint> &w){
 vector<base> fv(v.begin(), v.end()), fw(w.begin(), w.end());
 int n = 2; while(n < v.size() + w.size()) n <<= 1;</pre>
 fv.resize(n); fw.resize(n);
 fft(fv, 0); fft(fw, 0);
 for(int i=0; i<n; i++) fv[i] *= fw[i];</pre>
 fft(fv. 1):
 vector<lint> ret(n);
 for(int i=0: i<n: i++) ret[i] = (lint)round(fv[i].real()):
 return ret;
vector<lint> multiply(vector<lint> &v, vector<lint> &w, lint mod){
 int n = 2; while(n < v.size() + w.size()) n <<= 1;</pre>
 vector<base> v1(n), v2(n), r1(n), r2(n);
 for(int i=0; i<v.size(); i++){</pre>
   v1[i] = base(v[i] >> 15, v[i] & 32767);
 for(int i=0; i<w.size(); i++){</pre>
   v2[i] = base(w[i] >> 15, w[i] & 32767);
 fft(v1, 0);
 fft(v2, 0);
 for(int i=0: i<n: i++){
   int j = (i ? (n - i) : i);
   base ans1 = (v1[i] + conj(v1[j])) * base(0.5, 0);
   base ans2 = (v1[i] - conj(v1[i])) * base(0, -0.5);
   base ans3 = (v2[i] + conj(v2[j])) * base(0.5, 0);
   base ans4 = (v2[i] - conj(v2[j])) * base(0, -0.5);
   r1[i] = (ans1 * ans3) + (ans1 * ans4) * base(0, 1);
   r2[i] = (ans2 * ans3) + (ans2 * ans4) * base(0, 1);
 fft(r1, 1);
 fft(r2, 1):
 vector<lint> ret(n);
  for(int i=0; i<n; i++){</pre>
   lint av = (lint)round(r1[i].real());
   lint bv = (lint)round(r1[i].imag()) + (lint)round(r2[i].real());
   lint cv = (lint)round(r2[i].imag());
   av %= mod, bv %= mod, cv %= mod;
   ret[i] = (av << 30) + (bv << 15) + cv;
   ret[i] %= mod;
   ret[i] += mod;
   ret[i] %= mod;
 return ret;
5.2 Hell-Joseon style FFT
#include <smmintrin.h>
#include <immintrin.h>
#pragma GCC target("avx2")
#pragma GCC target("fma")
```

```
m256d mult( m256d a, m256d b){
  _{\rm m256d} c = _{\rm mm256\_movedup\_pd(a)};
 _{m256d} d = _{mm256\_shuffle\_pd(a, a, 15)};
 _{m256d cb} = _{mm256_{mul_pd(c, b)}}
 _{\rm m256d\ db} = _{\rm mm256\_mul\_pd(d, b)};
 m256d = mm256  shuffle pd(db, db, 5):
 _{\rm m256d} r = _{\rm mm256\_addsub\_pd(cb, e)};
 return r;
void fft(int n, __m128d a[], bool invert){
 for(int i=1, j=0; i<n; ++i){
   int bit = n >> 1:
   for(; j>=bit; bit>>=1) j -= bit;
    i += bit:
   if(i<j) swap(a[i], a[j]);</pre>
 for(int len=2; len<=n; len<<=1){</pre>
    double ang = 2*3.14159265358979/len*(invert?-1:1);
    _{\text{m256d wlen; wlen[0]}} = \cos(\text{ang}), \text{ wlen[1]} = \sin(\text{ang});
   for(int i=0: i<n: i += len){</pre>
      _{\rm m}256d w; w[0] = 1; w[1] = 0;
      for(int j=0; j<len/2; ++j){</pre>
        w = _{mm256_permute2f128_pd(w, w, 0)};
        wlen = _{mm256_{insertf128_{pd}(wlen, a[i+j+len/2], 1)};
        w = mult(w, wlen):
        _{m128d} vw = _{mm256}extractf128_{pd}(w, 1);
        _{m128d} u = a[i+i];
        a[i+j] = _mm_add_pd(u, vw);
        a[i+j+len/2] = _mm_sub_pd(u, vw);
   }
 }
 if(invert){
    _{m128d inv; inv[0] = inv[1] = 1.0/n;
    for(int i=0; i<n; ++i) a[i] = _mm_mul_pd(a[i], inv);</pre>
vector<int64 t> multiplv(vector<int64 t>& v. vector<int64 t>& w){
 int n = 2; while(n < v.size()+w.size()) n<<=1;</pre>
 _{m128d*} fv = new _{m128d[n]};
 for(int i=0: i<n: ++i) fv[i][0] = fv[i][1] = 0:
 for(int i=0; i<v.size(); ++i) fv[i][0] = v[i];</pre>
 for(int i=0; i<w.size(); ++i) fv[i][1] = w[i];</pre>
 fft(n, fv, 0); // (a+bi) is stored in FFT
 for(int i=0; i < n; i += 2){
   m256d a:
   a = _mm256_insertf128_pd(a, fv[i], 0);
   a = _mm256_insertf128_pd(a, fv[i+1], 1);
   a = mult(a, a):
   fv[i] = _mm256_extractf128_pd(a, 0);
   fv[i+1] = _mm256_extractf128_pd(a, 1);
 }
 fft(n, fv, 1);
 vector<int64 t> ret(n):
 for(int i=0; i<n; ++i) ret[i] = (int64_t)round(fv[i][1]/2);</pre>
 delete[] fv:
```

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```
return ret:
5.3 NTT Polynomial Division
vector<lint> get_inv(int n, const vector<lint> &p){
 vector<lint> q = \{ipow(p[0], mod - 2)\};
 for(int i=2; i<=n; i<<=1){
   vector<lint> res:
   vector<lint> fq(q.begin(), q.end()); fq.resize(2*i);
   vector<lint> fp(p.begin(), p.begin() + i); fp.resize(2*i);
   fft(fq, 0); fft(fp, 0);
   for(int j=0; j<2*i; j++){
     fp[j] *= fq[j] * fq[j] % mod;
     fp[j] %= mod;
   fft(fp, 1);
   res.resize(i):
   for(int j=0; j<i; j++){
     res[j] = mod - fp[j];
     if(j < i/2) res[j] += 2 * q[j];
     res[i] %= mod;
   q = res;
 return q;
vector<lint> poly_divide(const vector<lint> &a, const vector<lint> &b){
 assert(b.back() != 0); // please trim leading zero
 int n = a.size(), m = b.size();
 int k = 2; while (k < n-m+1) k <<= 1;
 vector<lint> rb(k), ra(k);
 for(int i=0; i<m && i<k; ++i) rb[i] = b[m-i-1];</pre>
 for(int i=0; i<n && i<k; ++i) ra[i] = a[n-i-1];</pre>
 vector<lint> rbi = get_inv(k, rb);
 vector<lint> res = multiply(rbi, ra);
 res.resize(n - m + 1):
 reverse(res.begin(), res.end());
 return res:
5.4 Black Box Linear Algebra + Kitamasa
vector<int> berlekamp_massey(vector<int> x){
 vector<int> ls, cur;
 int lf. ld:
 for(int i=0; i<x.size(); i++){</pre>
   lint t = 0:
   for(int j=0; j<cur.size(); j++){</pre>
     t = (t + 111 * x[i-j-1] * cur[j]) \% mod;
   if((t - x[i]) \% mod == 0) continue;
   if(cur.empty()){
     cur.resize(i+1):
     lf = i;
     ld = (t - x[i]) \% mod:
     continue;
   lint k = -(x[i] - t) * ipow(ld, mod - 2) % mod;
```

```
vector<int> c(i-lf-1):
    c.push_back(k);
    for(auto &j : ls) c.push_back(-j * k % mod);
    if(c.size() < cur.size()) c.resize(cur.size());</pre>
    for(int j=0; j<cur.size(); j++){</pre>
      c[j] = (c[j] + cur[j]) \% mod;
    if(i-lf+(int)ls.size()>=(int)cur.size()){
      tie(ls, lf, ld) = make_tuple(cur, i, (t - x[i]) % mod);
    cur = c:
  for(auto &i : cur) i = (i % mod + mod) % mod;
  return cur:
int get_nth(vector<int> rec, vector<int> dp, lint n){
  int m = rec.size():
  vector<int> s(m), t(m);
  s[0] = 1:
  if(m != 1) t[1] = 1;
  else t[0] = rec[0];
  auto mul = [&rec](vector<int> v, vector<int> w){
   int m = v.size();
    vector<int> t(2 * m);
    for(int j=0; j<m; j++){</pre>
      for(int k=0; k<m; k++){</pre>
        t[j+k] += 111 * v[j] * w[k] % mod;
        if(t[i+k] >= mod) t[i+k] -= mod:
    for(int j=2*m-1; j>=m; j--){
      for(int k=1; k<=m; k++){</pre>
        t[j-k] += 111 * t[j] * rec[k-1] % mod;
        if(t[j-k] >= mod) t[j-k] -= mod;
    }
    t.resize(m);
    return t:
  };
  while(n){
   if(n \& 1) s = mul(s, t):
   t = mul(t, t);
   n >>= 1:
 }
  lint ret = 0:
  for(int i=0; i<m; i++) ret += 111 * s[i] * dp[i] % mod;</pre>
 return ret % mod;
int guess_nth_term(vector<int> x, lint n){
 if(n < x.size()) return x[n];</pre>
 vector<int> v = berlekamp_massey(x);
 if(v.empty()) return 0;
 return get_nth(v, x, n);
struct elem{int x, y, v_i}; // A_i(x, y) \leftarrow v_i 0-based. no duplicate please..
vector<int> get_min_poly(int n, vector<elem> M){
```

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```
// smallest polv P such that A^i = sum {i < i} {A^i \times P i}
  vector<int> rnd1, rnd2;
 mt19937 rng(0x14004);
  auto randint = [&rng](int lb, int ub){
   return uniform_int_distribution<int>(lb, ub)(rng);
  for(int i=0; i<n; i++){</pre>
   rnd1.push_back(randint(1, mod - 1));
   rnd2.push_back(randint(1, mod - 1));
  vector<int> gobs;
  for(int i=0: i<2*n+2: i++){
   int tmp = 0;
   for(int j=0; j<n; j++){</pre>
     tmp += 111 * rnd2[j] * rnd1[j] % mod;
      if(tmp >= mod) tmp -= mod;
   gobs.push_back(tmp);
   vector<int> nxt(n):
   for(auto &i : M){
     nxt[i.x] += 111 * i.v * rnd1[i.v] % mod;
     if(nxt[i.x] >= mod) nxt[i.x] -= mod;
   }
   rnd1 = nxt;
  auto sol = berlekamp_massey(gobs);
 reverse(sol.begin(), sol.end());
  return sol:
lint det(int n. vector<elem> M){
 vector<int> rnd;
 mt19937 rng(0x14004);
 auto randint = [&rng](int lb, int ub){
   return uniform_int_distribution<int>(lb, ub)(rng);
 };
  for(int i=0: i<n: i++) rnd.push back(randint(1, mod - 1)):
  for(auto &i : M){
   i.v = 111 * i.v * rnd[i.v] % mod:
  auto sol = get_min_poly(n, M)[0];
 if (n \% 2 == 0) sol = mod - sol:
 for(auto &i : rnd) sol = 111 * sol * ipow(i, mod - 2) % mod;
 return sol:
5.5 Gaussian Elimination
int n. inv:
vector<int> basis[505];
lint gyesu = 1;
void insert(vector<int> v){
 for(int i=0: i<n: i++){</pre>
    if(basis[i].size()) inv ^= 1; // inversion num increases
   if(v[i] && basis[i].empty()){
     basis[i] = v;
      return;
   }
```

```
if(v[i]){
      lint minv = ipow(basis[i][i], mod - 2) * v[i] % mod;
      for(auto &j : basis[i]) j = (j * minv) % mod;
      gvesu *= minv:
     gyesu %= mod;
     for(int j=0; j<basis[i].size(); j++){</pre>
        v[j] += mod - basis[i][j];
        while(v[j] >= mod) v[j] -= mod;
   }
 }
 puts("0");
 exit(0);
// Sample: Calculates Determinant in Z_p Field
int main(){
 scanf("%d",&n);
 for(int i=0: i<n: i++){
   vector<int> v(n);
   for(int j=0; j<n; j++) scanf("%d",&v[j]);
   if(i % 2 == 1) inv ^= 1:
   insert(v);
 }
 if(inv) gyesu = mod - gyesu;
 gyesu = ipow(gyesu, mod - 2);
 for(int i=0; i<n; i++) gyesu = gyesu * basis[i][i] % mod;</pre>
 cout << gvesu % mod << endl:
5.6 Simplex Algorithm
using T = long double;
const int N = 410, M = 30010;
const T eps = 1e-7;
int n, m;
int Left[M]. Down[N]:
// time complexity: exponential. fast $O(MN^2)$ in experiment. dependent on the modeling.
// Ax <= b, max c^T x. 최댓값: v, 답 추적: sol[i]. 1 based
Ta[M][N], b[M], c[N], v, sol[N]:
bool eq(T a, T b) { return fabs(a - b) < eps; }</pre>
bool ls(T a, T b) { return a < b && !eq(a, b); }
void init(int p, int q) {
 n = p; m = q; v = 0;
 for(int i = 1; i <= m; i++){
   for(int j = 1; j <= n; j++) a[i][j]=0;
 for(int i = 1; i \le m; i++) b[i]=0;
 for(int i = 1; i <= n; i++) c[i]=sol[i]=0;
void pivot(int x,int y) {
 swap(Left[x], Down[y]);
 T k = a[x][y]; a[x][y] = 1;
 vector<int> nz;
 for(int i = 1; i <= n; i++){
   a[x][i] /= k;
    if(!eq(a[x][i], 0)) nz.push_back(i);
```

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```
b[x] /= k:
                                                                                                int aux[10000000]; // >= k*n
  for(int i = 1; i <= m; i++){
                                                                                                  if(k == 1) {
   if(i == x || eq(a[i][y], 0)) continue;
                                                                                                    res[0] = 0;
   k = a[i][y]; a[i][y] = 0;
                                                                                                    return 1;
   b[i] -= k*b[x]:
   for(int j : nz) a[i][j] -= k*a[x][j];
                                                                                                  for(int i = 0; i < k * n; i++)
                                                                                                    aux[i] = 0;
  if(eq(c[y], 0)) return;
                                                                                                  int sz = 0:
 k = c[v]; c[v] = 0;
 v += k*b[x]:
                                                                                                    if(t > n)  {
 for(int i : nz) c[i] -= k*a[x][i];
                                                                                                      if(n \% p == 0)
                                                                                                        for(int i = 1; i <= p; i++)
// 0: found solution, 1: no feasible solution, 2: unbounded
                                                                                                          res[sz++] = aux[i]:
int solve() {
                                                                                                    }
  for(int i = 1; i <= n; i++) Down[i] = i;</pre>
                                                                                                    else {
  for(int i = 1; i <= m; i++) Left[i] = n+i;</pre>
                                                                                                      aux[t] = aux[t - p];
  while(1) { // Eliminating negative b[i]
                                                                                                      db(t + 1, p);
   int x = 0, y = 0;
                                                                                                      for(int i = aux[t - p] + 1; i < k; i++) {
   for(int i = 1; i \le m; i++) if (ls(b[i], 0) && (x == 0 || b[i] < b[x])) x = i;
                                                                                                        aux[t] = i;
   if(x == 0) break;
                                                                                                        db(t + 1, t);
   for(int i = 1; i \le n; i +++) if (ls(a[x][i], 0) && (y == 0 || a[x][i] \le a[x][y])) y = i;
   if(y == 0) return 1;
                                                                                                    }
                                                                                                  };
   pivot(x, y);
                                                                                                  db(1, 1);
  while(1) {
                                                                                                  return sz;
   int x = 0, y = 0;
   for(int i = 1: i <= n: i++)
                                                                                                5.9 Range Prime Counting
     if (ls(0, c[i]) && (!v || c[i] > c[v])) v = i;
                                                                                                // Primes up to 10^12 can be counted in ~1 second.
   if(y == 0) break;
   for(int i = 1; i <= m; i++)
                                                                                                bool prime[MAXN];
     if (ls(0, a[i][y]) && (!x || b[i]/a[i][y] < b[x]/a[x][y])) x = i;
                                                                                                int prec[MAXN];
   if(x == 0) return 2;
                                                                                                vector<int> P;
   pivot(x, y);
                                                                                                void init() {
 for(int i = 1: i <= m: i++) if(Left[i] <= n) sol[Left[i]] = b[i]:</pre>
                                                                                                  prime[2] = true;
 return 0;
                                                                                                  for (int i = 3: i*i < MAXN: i += 2){
5.7 Pentagonal Number Theorem for Partition Number Counting
                                                                                                    if (prime[i]){
vector<pair<int, int>> gp;
                                                                                                    }
lint P[MAXN+1] = \{\};
                                                                                                  }
gp.emplace_back(0, 0);
                                                                                                  for(int i=1; i<MAXN; i++){</pre>
for(int i = 1; gp.back().second <= MAXN; i++) {</pre>
                                                                                                    if (prime[i]) P.push_back(i);
  gp.emplace_back(i \% 2 ? 1 : -1, i * (3*i - 1) / 2);
                                                                                                    prec[i] = prec[i-1] + prime[i];
 gp.emplace_back(i % 2 ? 1 : -1, i * (3*i + 1) / 2);
                                                                                                 }
}
                                                                                                }
P[1] = 1;
for(int n = 2: n \le MAXN: n++) {
                                                                                                lint rec(lint N, int K) {
 for(auto it : gp) if(n >= it.second) P[n] += P[n - it.second] * it.first + MOD;
                                                                                                  if (N <= 1 | | K < 0) return 0;
 P[n] \% = MOD;
                                                                                                  if (N \leq P[K]) return N-1;
}
5.8 De Bruin Sequence
                                                                                                  const int LIM = 250:
// Create cyclic string of length k^n that contains every length n string as substring.
                                                                                                  static int memo[LIM*LIM][LIM];
alphabet = [0, k - 1]
                                                                                                  bool ok = N < LIM*LIM;</pre>
int res[10000000]; // >= k^n
```

```
int de_bruijn(int k, int n) { // Returns size (k^n)
 function<void(int, int)> db = [&](int t, int p) {
const int MAXN = 1000005; // MAXN is the maximum value of sqrt(N) + 2
 for (int i = 3; i < MAXN; i += 2) prime[i] = true;</pre>
     for (int j = i*i; j < MAXN; j += i+i) prime[j] = false;</pre>
 if (N < MAXN && 111 * P[K]*P[K] > N) return N-1 - prec[N] + prec[P[K]];
 if (ok && memo[N][K]) return memo[N][K];
```

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```
lint ret = N/P[K] - rec(N/P[K], K-1) + rec(N, K-1):
 if (ok) memo[N][K] = ret;
 return ret;
lint count_primes(lint N) { // less than or equal to
 if (N < MAXN) return prec[N];</pre>
 int K = prec[(int)sqrt(N) + 1];
 return N-1 - rec(N, K) + prec[P[K]];
5.10 Discrete Kth root
* Solve x for x^P = A mod Q
* (P, Q-1) = 1 \rightarrow P^-1 \mod (Q-1) exists
 * x has solution iff A^{(Q-1)}/P = 1 mod Q
 * PP | (Q-1) \rightarrow P < sqrt(Q), solve lgQ rounds of discrete log
* else -> find a s.t. s | (Pa - 1) -> ans = A^a */
using LL = long long;
LL mul(LL x, LL y, LL mod) { return (__int128) x * y % mod; }
LL add(LL x, LL y, LL mod) { return (x + y) % mod; }
LL pw(LL x, LL y, LL mod){
 LL ret = 1, piv = x;
 while(y){
   if(y & 1) ret = mul(ret, piv, mod);
   piv = mul(piv, piv, mod);
   y >>= 1;
 return ret % mod;
void gcd(LL a, LL b, LL &x, LL &y, LL &g){
 if (b == 0) {
   x = 1, y = 0, g = a;
   return;
 LL tx, ty;
  gcd(b, a%b, tx, ty, g);
 x = ty; y = tx - ty * (a / b);
LL P, A, Q, g; // x^P = A \mod Q
const int X = 1e5:
LL base, ae[X], aXe[X], iaXe[X];
unordered_map<LL, LL> ht;
#define FOR(i, c) for (int i = 0; i < (c); ++i)
#define REP(i, 1, r) for (int i = (1); i <= (r); ++i)
void build(LL a) { // ord(a) = P < sqrt(Q)</pre>
 base = a:
 ht.clear();
 ae[0] = 1; ae[1] = a; aXe[0] = 1; aXe[1] = pw(a, X, Q);
 iaXe[0] = 1; iaXe[1] = pw(aXe[1], Q-2, Q);
 REP(i, 2, X-1) {
   ae[i] = mul(ae[i-1], ae[1], Q):
   aXe[i] = mul(aXe[i-1], aXe[1], Q);
   iaXe[i] = mul(iaXe[i-1], iaXe[1], Q);
 FOR(i, X) ht[ae[i]] = i;
```

```
LL dis_log(LL x) {
 FOR(i, X) {
   LL iaXi = iaXe[i];
   LL rst = mul(x, iaXi, Q);
   if (ht.count(rst)) return i*X + ht[rst]:
 }
LL main2() {
   cin >> P >> A >> Q;
 LL t = 0, s = Q-1;
 while (s % P == 0) {
   ++t;
   s /= P:
 }
 if (A == 0) return 0;
  if (t == 0) {
   // a^{P^-1 mod phi(Q)}
   LL x, y, _;
   gcd(P, Q-1, x, y, _);
   if (x < 0) {
     x = (x \% (Q-1) + Q-1) \% (Q-1);
   LL ans = pw(A, x, Q);
   if (pw(ans, P, Q) != A) while(1);
   return ans;
 // A is not P-residue
 if (pw(A, (Q-1) / P, Q) != 1) return -1;
 for (g = 2; g < Q; ++g) {
   if (pw(g, (Q-1) / P, Q) != 1)
     break:
 }
 LL alpha = 0;
   LL y, _;
    gcd(P, s, alpha, y, _);
   if (alpha < 0) alpha = (alpha \% (Q-1) + Q-1) \% (Q-1);
 if (t == 1) {
   LL ans = pw(A, alpha, Q);
   return ans;
 LL a = pw(g, (Q-1) / P, Q);
 build(a):
 LL b = pw(A, add(mul(P%(Q-1), alpha, Q-1), Q-2, Q-1), Q);
 LL c = pw(g, s, Q);
 LL h = 1;
 LL e = (Q-1) / s / P: // r^{t-1}
 REP(i, 1, t-1) {
   e /= P;
   LL d = pw(b, e, Q);
   LL j = 0;
   if (d != 1) {
     j = -dis_log(d);
      if (j < 0) j = (j % (Q-1) + Q-1) % (Q-1);
```

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```
}
   b = mul(b, pw(c, mul(P%(Q-1), j, Q-1), Q), Q);
   h = mul(h, pw(c, j, Q), Q);
   c = pw(c, P, Q);
 }
 return mul(pw(A, alpha, Q), h, Q);
5.11 Miller-Rabin Test + Pollard Rho Factorization
namespace miller_rabin{
   lint mul(lint x, lint y, lint mod){ return (__int128) x * y % mod; }
 lint ipow(lint x, lint y, lint p){
   lint ret = 1, piv = x \% p;
   while(y){
     if(y&1) ret = mul(ret, piv, p);
     piv = mul(piv, piv, p);
     y >>= 1;
   }
   return ret;
 bool miller_rabin(lint x, lint a){
   if(x \% a == 0) return 0;
   lint d = x - 1:
   while(1){
     lint tmp = ipow(a, d, x);
     if(d&1) return (tmp != 1 && tmp != x-1);
     else if(tmp == x-1) return 0;
     d >>= 1:
   }
 }
 bool isprime(lint x){
   for(auto &i : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
     if(x == i) return 1:
     if (x > 40 \&\& miller_rabin(x, i)) return 0;
   if(x \le 40) return 0:
   return 1;
 }
namespace pollard_rho{
 lint f(lint x, lint n, lint c){
   return (c + miller_rabin::mul(x, x, n)) % n;
 void rec(lint n, vector<lint> &v){
   if(n == 1) return;
   if(n \% 2 == 0){
     v.push_back(2);
     rec(n/2, v);
     return;
   if(miller_rabin::isprime(n)){
     v.push_back(n);
     return:
   lint a, b, c;
   while(1){
```

```
a = rand() \% (n-2) + 2:
     b = a;
     c = rand() \% 20 + 1;
       a = f(a, n, c);
       b = f(f(b, n, c), n, c);
     \frac{1}{2} while \left( \gcd(abs(a-b), n) == 1 \right);
     if(a != b) break;
   lint x = gcd(abs(a-b), n);
   rec(x, v);
   rec(n/x, v);
 vector<lint> factorize(lint n){
   vector<lint> ret;
   rec(n, ret);
   sort(ret.begin(), ret.end());
   return ret;
 }
};
5.12 Highly Composite Numbers, Large Prime
                           divisors 2 3 5 71113171923293137
 < 10^k
                number
                     6
 2
                    60
                                 12 2 1 1
 3
                   840
                                 32 3 1 1 1
                  7560
                                64 3 3 1 1
 5
                 83160
                                128 3 3 1 1 1
                720720
                                240 4 2 1 1 1 1
 7
                               448 6 3 1 1 1 1
                8648640
 8
               73513440
                               768 5 3 1 1 1 1 1
 9
              735134400
                               1344 6 3 2 1 1 1 1
 10
            6983776800
                               2304 5 3 2 1 1 1 1 1
 11
            97772875200
                               4032 6 3 2 2 1 1 1 1
 12
                               6720 6 4 2 1 1 1 1 1 1
           963761198400
 13
          9316358251200
                              10752 6 3 2 1 1 1 1 1 1 1
                              17280 5 4 2 2 1 1 1 1 1 1
 14
         97821761637600
 15
        866421317361600
                              26880 6 4 2 1 1 1 1 1 1 1 1
 16
       8086598962041600
                              41472 8 3 2 2 1 1 1 1 1 1 1
      74801040398884800
                              64512 6 3 2 2 1 1 1 1 1 1 1 1
                             103680 8 4 2 2 1 1 1 1 1 1 1 1
 18 897612484786617600
          prime # of prime
                                      < 10^k
 1
              7
                          4
                                     10
                                                  999999967
 2
             97
                         25
                                     11
                                                 9999999977
                                     12
            997
                         168
                                                99999999989
 4
            9973
                        1229
                                     13
                                           999999999971
 5
           99991
                        9592
                                     14
                                           9999999999973
 6
          999983
                       78498
                                     15
                                             99999999999989
 7
         9999991
                      664579
                                            99999999999937
        9999989
                     5761455
                                   17
                                           999999999999997
 9
       99999937
                    50847534
                                          999999999999999
NTT Prime:
 998244353 = 119 \times 2^{23} + 1. Primitive root: 3.
```

 $985661441 = 235 \times 2^{22} + 1$ . Primitive root: 3.

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 $1012924417 = 483 \times 2^{21} + 1$ . Primitive root: 5.

### Miscellaneous

#### 6.1 Mathematics

- Tutte Matrix. For a simple undirected graph G, Let M be a matrix with entries  $A_{i,j} = 0$  if  $(i,j) \notin E$  and  $A_{i,j} = -A_{j,i} = X$  if  $(i,j) \in E$ . X could be any random value. If the determinants are non-zero, then a perfect matching exists, while other direction might not hold for very small probability.
- Cayley's Formula. Given a degree sequence  $d_1, d_2 \cdots, d_n$  for each labeled vertices, there exists  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees. Summing this for every possible degree sequence gives  $n^{n-2}$ .
- Kirchhoff's Theorem. For a multigraph G with no loops, define Laplacian matrix as L = D A. D is a diagonal matrix with  $D_{i,i} = deg(i)$ , and A is an adjacency matrix. If you remove any row and column of L, the determinant gives a number of spanning trees.
- Green's Theorem. Let C is positive, smooth, simple curve. D is region bounded by C.  $\oint_C (Ldx + Mdy) = \iint_D (\frac{\partial M}{\partial x} \frac{\partial L}{\partial y})$

To calculate area,  $\frac{\partial M}{\partial x} - \frac{\partial L}{\partial y} = 1$ , common selection is  $M = \frac{1}{2}x$ ,  $L = -\frac{1}{2}y$ .

Line integral of circle parametrized by  $(x,y) = (x_C + r_C \cos \theta, \ y_C + r_C \sin \theta)$ , when  $\theta = t\theta_i + (1-t)\theta_f$ , is given as follows:  $\frac{1}{2}(r_C(x_C(\sin \theta_f - \sin \theta_i) - y_C(\cos \theta_f - \cos \theta_i)) + (\theta_f - \theta_i)r_C^2)$ .

Line integral of line parametrized by  $(x, y) = t(x_1, y_1) + (1 - t)(x_2, y_2)$  is given as follows::  $\frac{1}{2}(x_1y_2 - x_2y_1)$ .

• Burnside's lemma / Pólya enumeration theorem. let G and H be groups of permutations of finite sets X and Y. Let  $c_m(g)$  denote the number of cycles of length m in  $g \in G$  when permuting X. The number of colorings of X into |Y| = n colors with exactly  $r_i$  occurrences of the i-th color is the coefficient of  $w_1^{r_1} \dots w_n^{r_n}$  in the following polynomial:

$$P(w_1, \dots, w_n) = \frac{1}{|H|} \sum_{h \in H} \frac{1}{|G|} \sum_{g \in G} \prod_{m > 1} (\sum_{h^m(b) = b} (w_b^m))^{c_m(g)}$$

When  $H = \{I\}$  (No color permutation):  $P(w_1, \dots, w_n) = \frac{1}{|G|} \sum_{g \in G} \prod_{m > 1} (w_1^m + \dots + w_n^m)^{c_m(g)}$ 

Without the occurrence restriction:

 $P(1,...,1) = \frac{1}{|G|} \sum_{g \in G} n^{c(g)}$ 

where c(g) could also be interpreted as the number of elements in X that are fixed up to g.

Pick's Theorem. A = i + b/2 - 1, where: P is a simple polygon whose vertices are grid points, A is area of P, i is # of grid points in the interior of P, and b is # of grid points on the boundary of P. If h is # of holes of P (h + 1 simple closed curves in total), A = i + b/2 + h - 1.

```
// number of (x, y) : (0 <= x < n && 0 < y < k/d x + b/d)
ll count_solve(ll n, ll k, ll b, ll d) {
   if (k == 0) {
     return (b / d) * n;
   }
   if (k >= d || b >= d) {
     return ((k / d) * (n - 1) + 2 * (b / d)) * n / 2 + count_solve(n, k % d, b % d, d);
   }
   return count_solve((k * n + b) / d, d, (k * n + b) % d, k);
}
```

```
• Xudyh Sieve. F(n) = \sum_{d|n} f(d)

S(n) = \sum_{i \leq n} f(i) = \sum_{i \leq n} F(i) - \sum_{d=2}^n S\left(\left\lfloor \frac{n}{d} \right\rfloor\right)

Preprocess S(1) to S(M) (Set M = n^{\frac{2}{3}} for complexity)

S(n) = \sum f(i) = \sum_{i \leq n} \left[F(i) - \sum_{j|i,j \neq i} f(j)\right] = \sum F(i) - \sum_{i/j=d=2}^n \sum_{dj \leq n} f(j)

S(n) = \sum i f(i) = \sum_{i \leq n} i \left[F(i) - \sum_{j|i,j \neq i} f(j)\right] = \sum i F(i) - \sum_{i/j=d=2}^n \sum_{dj \leq n} dj f(j)

\sum_{d|n} \varphi(d) = n \sum_{d|n} \mu(d) = \text{if } (n > 1) \text{ then } 0 \text{ else } 1 \sum_{d|n} (\mu(\frac{n}{d}) \sum_{e|d} f(e)) = f(n)
```

### 6.2 Popular Optimization Technique

• CHT. DnC optimization. Mo's algorithm trick (on tree). IOI 2016 Aliens trick. IOI 2009 Regions trick.

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- Knuth's  $O(n^2)$  Optimal BST: minimize  $D_{i,j} = Min_{i \leq k < j}(D_{i,k} + D_{k+1,j}) + C_{i,j}$ . Quadrangle Inequality :  $C_{a,c} + C_{b,d} \leq C_{a,d} + C_{b,c}$ ,  $C_{b,c} \leq C_{a,d}$ . Now monotonicity holds.
- Sqrt batch processing Save queries in buffer, and update in every sqrt steps (cf : IOI 2011 Elephant. hyea calls it "ainta technique")
- Dynamic insertion in static set (Make  $O(\log n)$  copy. Merge like binomial heap.)
- Offline insertion / deletion in insert-only set (Pair insertion-deletion operation, and regard it as range query)
- Accoder Median Pyramid: Reduce the input to binary, and solve the easier problem.
- LP Duality, max  $c^T x$  sit to  $Ax \leq b$ . Dual problem is min  $b^T x$  sit to  $A^T x \geq c$ . By strong duality, min max value coincides.

## 6.3 Bit Twiddling Hack

```
int __builtin_clz(int x);// number of leading zero
int __builtin_ctz(int x);// number of trailing zero
int __builtin_clzll(long long x);// number of leading zero
int __builtin_ctzll(long long x);// number of trailing zero
int __builtin_popcount(int x);// number of 1-bits in x
int __builtin_popcountll(long long x);// number of 1-bits in x
lsb(n): (n & -n): // last bit (smallest)
floor(log2(n)): 31 - __builtin_clz(n | 1);
floor(log2(n)): 63 - __builtin_clzll(n | 1);
// compute next perm. ex) 00111, 01011, 01101, 01110, 10011, 10101...
long long next_perm(long long v){
 long long t = v \mid (v-1);
 return (t + 1) \mid (((^t \& -^t) - 1) >> (_builtin_ctz(v) + 1));
6.4 Fast Integer IO
static char buf[1 << 19]: // size : any number geg than 1024
static int idx = 0;
static int bytes = 0;
static inline int _read() {
 if (!bytes || idx == bytes) {
   bytes = (int)fread(buf, sizeof(buf[0]), sizeof(buf), stdin);
   idx = 0;
 return buf[idx++];
static inline int readInt() {
```

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```
int x = 0, s = 1:
  int c = _read();
  while (c \le 32) c = read();
  if (c == '-') s = -1, c = _read();
  while (c > 32) x = 10 * x + (c - '0'), c = _read();
 if (s < 0) x = -x:
 return x;
6.5 OSRank in g++
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
typedef
tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update> ordered_set;
ordered_set X;
X.insert(1): X.insert(2): X.insert(4): X.insert(8): X.insert(16):
cout<<*X.find_by_order(1)<<end1; // 2</pre>
cout<<*X.find by order(2)<<endl: // 4
cout<<*X.find_by_order(4)<<endl; // 16</pre>
cout<<(end(X)==X.find_by_order(6))<<endl; // true</pre>
cout<<X.order_of_key(-5)<<endl; // 0</pre>
cout<<%.order of kev(1)<<endl: // 0
cout<<X.order_of_key(3)<<endl; // 2</pre>
cout<<X.order_of_key(4)<<endl; // 2</pre>
cout<<X.order_of_key(400)<<endl; // 5</pre>
6.6 Nasty Stack Hacks
// 64bit ver.
int main2(){ return 0: }
int main(){
  size_t sz = 1 << 29; // 512MB
  void* newstack = malloc(sz):
  void* sp_dest = newstack + sz - sizeof(void*);
  asm __volatile__("movq %0, %%rax\n\t"
  "movq %%rsp , (%%rax)\n\t"
  "movq %0, %%rsp\n\t": : "r"(sp_dest): );
 main2():
  asm __volatile__("pop %rsp\n\t");
 return 0;
6.7 C++ / Environment Overview
// vimrc : set nu sc ci si ai sw=4 ts=4 bs=2 mouse=a svntax on
// compile : g++ -o PROB PROB.cpp -std=c++11 -Wall -02
// options : -fsanitize=address -Wfatal-errors
struct StupidGCCCantEvenCompileThisSimpleCode{
  pair<int, int> array[1000000];
}; // https://gcc.gnu.org/bugzilla/show_bug.cgi?id=68203
// how to use rand (in 2018)
mt19937 rng(0x14004);
```

```
int randint(int lb, int ub){    return uniform int distribution<int>(lb, ub)(rng): }
// comparator overload
auto cmp = [](seg a, seg b){ return a.func() < b.func(); };</pre>
set<seg, decltype(cmp)> s(cmp);
map<seg, int, decltype(cmp)> mp(cmp);
priority_queue<seg, vector<seg>, decltype(cmp)> pq(cmp); // max heap
// hash func overload
struct point{
   int x, y;
   bool operator == (const point &p)const{ return x == p.x && y == p.y; }
};
struct hasher {
    size_t operator()(const point &p)const{ return p.x * 2 + p.y * 3; }
unordered_map<point, int, hasher> hsh;
6.8 Credits
 • cki86201, zigui, PavelKunyavskiy
 • https://gist.github.com/msg555/4963794
 • https://github.com/niklasb/contest-algos/blob/master/convex_hull/dynamic.cpp
 • https://github.com/jaehyunp/stanfordacm
 • https://github.com/stjepang/snippets/blob/master/count_primes.cpp
 • https://github.com/tzupengwang/PECaveros/blob/master/codebook/graph/
   BorrowedGeneralWeightedMatching.cpp
 • https://github.com/tzupengwang/PECaveros/blob/master/codebook/math/DiscreteKthsqrt.cpp
 • http://www-math.mit.edu/~etingof/groups.pdf
```