Team Note of Deobureo Minkyu Party

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ALL BELOW HERE ARE USELESS IF YOU READ THE STATEMENT WRONG

1 Quote

```
럭스를 럭스답게 든든한 연습헬팟 더불어민규당 hyea: Prove by solving. koosaga: Locality to the rescue! alex9801: Isshoman Beenzino
```

2 Flows, Matching

2.1 Hopcroft-Karp Bipartite Matching

```
const int MAXN = 50005, MAXM = 50005;
vector<int> gph[MAXN];
int dis[MAXN], 1[MAXN], r[MAXM], vis[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int 1, int r){ gph[1].push_back(r); }
bool bfs(int n){
  queue<int> que;
 bool ok = 0;
  memset(dis, 0, sizeof(dis));
  for(int i=0: i<n: i++){
   if(l[i] == -1 && !dis[i]){
      que.push(i);
      dis[i] = 1;
  while(!que.empty()){
    int x = que.front();
    que.pop();
    for(auto &i : gph[x]){
     if(r[i] == -1) ok = 1;
      else if(!dis[r[i]]){
        dis[r[i]] = dis[x] + 1;
        que.push(r[i]);
   }
 }
  return ok;
bool dfs(int x){
  for(auto &i : gph[x]){
    if(r[i] == -1 \mid | (!vis[r[i]] && dis[r[i]] == dis[x] + 1 && dfs(r[i]))){
      vis[r[i]] = 1; l[x] = i; r[i] = x;
      return 1:
   }
 }
```

```
return 0;
int match(int n){
 memset(1, -1, sizeof(1));
 memset(r, -1, sizeof(r));
 int ret = 0;
 while(bfs(n)){
   memset(vis, 0, sizeof(vis));
   for(int i=0; i<n; i++) if(l[i] == -1 && dfs(i)) ret++;
 return ret;
bool chk[MAXN + MAXM]:
void rdfs(int x, int n){
 if(chk[x]) return;
 chk[x] = 1:
 for(auto &i : gph[x]){
   chk[i + n] = 1;
   rdfs(r[i], n);
 }
vector<int> getcover(int n, int m){ // solve min. vertex cover
 match(n);
 memset(chk, 0, sizeof(chk));
 for(int i=0; i<n; i++) if(l[i] == -1) rdfs(i, n);
 vector<int> v:
 for(int i=0: i<n: i++) if(!chk[i]) v.push back(i):</pre>
 for(int i=n; i<n+m; i++) if(chk[i]) v.push_back(i);</pre>
 return v:
2.2 Dinic's Algorithm
const int MAXN = 505:
struct edg{int pos, cap, rev;};
vector<edg> gph[MAXN];
void clear(){for(int i=0; i<MAXN; i++) gph[i].clear();}</pre>
void add_edge(int s, int e, int x){
 gph[s].push_back({e, x, (int)gph[e].size()});
 gph[e].push_back({s, 0, (int)gph[s].size()-1});
int dis[MAXN], pnt[MAXN];
bool bfs(int src, int sink){
 memset(dis, 0, sizeof(dis));
 memset(pnt, 0, sizeof(pnt));
 queue<int> que;
 que.push(src):
 dis[src] = 1;
 while(!que.empty()){
   int x = que.front();
   que.pop();
```

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```
for(auto &e : gph[x]){
      if(e.cap > 0 && !dis[e.pos]){
        dis[e.pos] = dis[x] + 1;
        que.push(e.pos);
     }
   }
  return dis[sink] > 0:
int dfs(int x, int sink, int f){
  if(x == sink) return f;
  for(; pnt[x] < gph[x].size(); pnt[x]++){</pre>
    edg e = gph[x][pnt[x]];
    if(e.cap > 0 \&\& dis[e.pos] == dis[x] + 1){
      int w = dfs(e.pos, sink, min(f, e.cap));
      if(w){
        gph[x][pnt[x]].cap -= w;
        gph[e.pos][e.rev].cap += w;
       return w:
   }
  }
  return 0;
}
lint match(int src, int sink){
 lint ret = 0:
  while(bfs(src. sink)){
   int r;
    while((r = dfs(src, sink, 2e9))) ret += r:
  return ret;
}
     Min Cost Max Flow
const int MAXN = 100:
struct mincostflow{
  struct edg{ int pos, cap, rev, cost; };
  vector<edg> gph[MAXN];
  void clear(){
    for(int i=0; i<MAXN; i++) gph[i].clear();</pre>
  void add_edge(int s, int e, int x, int c){
    gph[s].push_back({e, x, (int)gph[e].size(), c});
    gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
  int dist[MAXN], pa[MAXN], pe[MAXN];
  bool inque[MAXN];
  bool spfa(int src, int sink){
   memset(dist, 0x3f, sizeof(dist));
    memset(inque, 0, sizeof(inque));
```

```
queue<int> que;
   dist[src] = 0;
   inque[src] = 1:
   que.push(src);
   bool ok = 0:
   while(!que.empty()){
     int x = que.front();
     que.pop();
     if(x == sink) ok = 1;
     inque[x] = 0;
     for(int i=0; i<gph[x].size(); i++){</pre>
        edg e = gph[x][i];
        if(e.cap > 0 \&\& dist[e.pos] > dist[x] + e.cost){
         dist[e.pos] = dist[x] + e.cost;
         pa[e.pos] = x;
         pe[e.pos] = i;
          if(!inque[e.pos]){
           inque[e.pos] = 1;
           que.push(e.pos);
   }
   return ok;
 int match(int src, int sink){
   int ret = 0:
   while(spfa(src, sink)){
     int cap = 1e9:
     for(int pos = sink; pos != src; pos = pa[pos]){
        cap = min(cap, gph[pa[pos]][pe[pos]].cap);
     ret += dist[sink] * cap;
     for(int pos = sink; pos != src; pos = pa[pos]){
       int rev = gph[pa[pos]][pe[pos]].rev;
        gph[pa[pos]][pe[pos]].cap -= cap;
       gph[pos][rev].cap += cap;
   }
   return ret:
 }
}mcmf;
2.4 Hell-Joseon style MCMF
const int MAXN = 100;
struct mincostflow{
 struct edg{ int pos, cap, rev, cost; };
 vector<edg> gph[MAXN];
 void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
 void add_edge(int s, int e, int x, int c){
```

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```
gph[s].push_back({e, x, (int)gph[e].size(), c});
 gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
int phi[MAXN], inque[MAXN], dist[MAXN];
void prep(int src. int sink){
 memset(phi, 0x3f, sizeof(phi));
 memset(dist, 0x3f, sizeof(dist));
 queue<int> que;
 que.push(src);
 inque[src] = 1;
  while(!que.empty()){
   int x = que.front();
   que.pop();
   inque[x] = 0;
   for(auto &i : gph[x]){
     if(i.cap > 0 && phi[i.pos] > phi[x] + i.cost){
        phi[i.pos] = phi[x] + i.cost;
        if(!inque[i.pos]){
         inque[i.pos] = 1;
          que.push(i.pos);
     }
   }
 }
  for(int i=0; i<MAXN; i++){</pre>
   for(auto &j : gph[i]){
     if(j.cap > 0) j.cost += phi[i] - phi[j.pos];
   }
 }
 priority_queue<pi, vector<pi>, greater<pi> > pq;
 pq.push(pi(0, src));
  dist[src] = 0:
  while(!pq.empty()){
   auto 1 = pq.top();
   pq.pop();
   if(dist[1.second] != 1.first) continue;
   for(auto &i : gph[l.second]){
     if(i.cap > 0 && dist[i.pos] > l.first + i.cost){
        dist[i.pos] = l.first + i.cost;
        pq.push(pi(dist[i.pos], i.pos));
   }
 }
bool vis[MAXN]:
int ptr[MAXN];
int dfs(int pos, int sink, int flow){
 vis[pos] = 1;
 if(pos == sink) return flow;
 for(; ptr[pos] < gph[pos].size(); ptr[pos]++){</pre>
```

```
auto &i = gph[pos][ptr[pos]];
     if(!vis[i.pos] && dist[i.pos] == i.cost + dist[pos] && i.cap > 0){
       int ret = dfs(i.pos, sink, min(i.cap, flow));
       if(ret != 0){
         i.cap -= ret:
         gph[i.pos][i.rev].cap += ret;
         return ret;
     }
   }
   return 0;
 int match(int src, int sink, int sz){
   prep(src, sink);
   for(int i=0; i<sz; i++) dist[i] += phi[sink] - phi[src];</pre>
   int ret = 0:
   while(true){
     memset(ptr, 0, sizeof(ptr));
     memset(vis, 0, sizeof(vis));
     int tmp = 0;
     while((tmp = dfs(src, sink, 1e9))){
       ret += dist[sink] * tmp;
       memset(vis, 0, sizeof(vis));
     tmp = 1e9;
     for(int i=0: i<sz: i++){</pre>
       if(!vis[i]) continue:
       for(auto &j : gph[i]){
         if(j.cap > 0 && !vis[j.pos]){
            tmp = min(tmp, (dist[i] + j.cost) - dist[j.pos]);
       }
     if(tmp > 1e9 - 200) break:
     for(int i=0: i<sz: i++){</pre>
       if(!vis[i]) dist[i] += tmp;
     }
   return ret;
}mcmf;
2.5 Circulation Problem
struct circ{
 maxflow mf;
 lint lsum:
 void clear(){
   lsum = 0:
   mf.clear();
```

```
void add_edge(int s, int e, int l, int r){
   lsum += 1:
    mf.add\_edge(s + 2, e + 2, r - 1);
    mf.add_edge(0, e + 2, 1);
   mf.add edge(s + 2, 1, 1):
  bool solve(int s, int e){
    mf.add_edge(e+2, s+2, 1e9); // to reduce as maxflow with lower bounds, in
    circulation problem skip this line
    return lsum == mf.match(0, 1):
    // to get maximum LR flow, run maxflow from s+2 to e+2 again
}circ:
     Min Cost Circulation
// Cycle canceling (Dual of successive shortest path)
// Time complexity is ridiculously high (F * maxC * nm^2). But runs reasonably in
practice (V = 70 \text{ in } 1s)
struct cost_circ{
  struct edg{ int pos, cap, rev, cost; };
  vector<edg> gph[MAXN];
  void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
  void add_edge(int s, int e, int x, int c){
    gph[s].push_back({e, x, (int)gph[e].size(), c});
    gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
  int dist[MAXN], par[MAXN], pae[MAXN];
  int negative_cycle(int n){
   bool mark[MAXN] = {};
    memset(dist, 0, sizeof(dist));
    int upd = -1;
    for(int i=0: i<=n: i++){
     for(int j=0; j<n; j++){</pre>
       int idx = 0:
        for(auto &k : gph[j]){
          if(k.cap > 0 && dist[k.pos] > dist[j] + k.cost){
            dist[k.pos] = dist[i] + k.cost:
            par[k.pos] = j;
            pae[k.pos] = idx;
```

 $if(i == n){$

upd = j;

return upd;

}

idx++;

}

while(!mark[upd]){

mark[upd] = 1;

upd = par[upd];

```
}
   }
   return -1;
 int match(int n){
   int rt = -1;
   int ans = 0:
   while(~(rt = negative_cycle(n))){
     bool mark[MAXN] = {}:
     vector<pi> cyc;
     while(!mark[rt]){
       cyc.push_back(pi(par[rt], pae[rt]));
       mark[rt] = 1;
       rt = par[rt];
     reverse(cyc.begin(), cyc.end());
     int capv = 1e9;
     for(auto &i : cyc){
       auto e = &gph[i.first][i.second];
        capv = min(capv, e->cap);
     for(auto &i : cvc){
        auto e = &gph[i.first][i.second];
       e->cap -= capv;
       gph[e->pos][e->rev].cap += capv;
       ans += e->cost * capv:
   }
   return ans;
 }
}circ:
2.7 Gomory-Hu Tree
struct edg{ int s, e, x; };
vector<edg> edgs;
maxflow mf;
void clear(){edgs.clear():}
void add_edge(int s, int e, int x){edgs.push_back({s, e, x});}
bool vis[MAXN]:
void dfs(int x){
 if(vis[x]) return;
 vis[x] = 1:
 for(auto &i : mf.gph[x]) if(i.cap > 0) dfs(i.pos);
vector<pi> solve(int n){ // i - j cut : i - j minimum edge cost. 0 based.
 vector<pi> ret(n); // if i > 0, stores pair(parent,cost)
 for(int i=1; i<n; i++){
   for(auto &j : edgs){
     mf.add_edge(j.s, j.e, j.x);
```

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```
mf.add_edge(j.e, j.s, j.x);
    ret[i].first = mf.match(i, ret[i].second);
    memset(vis, 0, sizeof(vis));
    dfs(i):
    for(int j=i+1; j<n; j++){</pre>
      if(ret[i].second == ret[i].second && vis[i]){
        ret[i].second = i:
    }
    mf.clear();
  return ret;
}
```

Blossom Algorithm for General Matching

```
const int MAXN = 2020 + 1;
struct GM { // 1-based Vertex index
 int vis[MAXN], par[MAXN], orig[MAXN], match[MAXN], aux[MAXN], t, N;
 vector<int> conn[MAXN];
  queue<int> Q;
  void addEdge(int u. int v) {
   conn[u].push_back(v); conn[v].push_back(u);
  void init(int n) {
   N = n; t = 0;
   for(int i=0; i<=n; ++i) {</pre>
     conn[i].clear();
     match[i] = aux[i] = par[i] = 0;
  void augment(int u, int v) {
   int pv = v. nv:
   do {
     pv = par[v]: nv = match[pv]:
     match[v] = pv; match[pv] = v;
     v = nv;
   } while(u != pv);
  int lca(int v. int w) {
   ++t.:
    while(true) {
     if(v) {
        if(aux[v] == t) return v; aux[v] = t;
       v = orig[par[match[v]]];
      swap(v, w);
 }
  void blossom(int v, int w, int a) {
```

```
while(orig[v] != a) {
     par[v] = w; w = match[v];
     if(vis[w] == 1) Q.push(w), vis[w] = 0;
     orig[v] = orig[w] = a;
     v = par[w]:
   }
 }
 bool bfs(int u) {
   fill(vis+1, vis+1+N, -1); iota(orig + 1, orig + N + 1, 1);
   Q = queue<int> (); Q.push(u); vis[u] = 0;
    while(!Q.empty()) {
     int v = Q.front(); Q.pop();
     for(int x: conn[v]) {
       if(vis[x] == -1) {
          par[x] = v; vis[x] = 1;
         if(!match[x]) return augment(u, x), true;
          Q.push(match[x]); vis[match[x]] = 0;
        else if(vis[x] == 0 && orig[v] != orig[x]) {
         int a = lca(orig[v], orig[x]);
         blossom(x, v, a): blossom(v, x, a):
   }
   return false;
  int Match() {
   int ans = 0;
   //find random matching (not necessary, constant improvement)
   vector<int> V(N-1); iota(V.begin(), V.end(), 1);
    shuffle(V.begin(), V.end(), mt19937(0x94949));
   for(auto x: V) if(!match[x]){
     for(auto y: conn[x]) if(!match[y]) {
        match[x] = y, match[y] = x;
        ++ans; break;
   }
   for(int i=1; i<=N; ++i) if(!match[i] && bfs(i)) ++ans;</pre>
    return ans;
 }
};
     Blossom Algorithm for Weighted General Matching
// https://github.com/tzupengwang/PECaveros/blob/master/
// codebook/graph/BorrowedGeneralWeightedMatching.cpp
// N^3 (but fast in practice)
struct WeightGraph {
 static const int INF = INT_MAX;
 static const int N = 514:
```

```
struct edge{
```

```
int u,v,w; edge(){}
edge(int ui,int vi,int wi)
  :u(ui).v(vi).w(wi){}
int n.n x:
edge g[N*2][N*2];
int lab[N*2];
int match[N*2],slack[N*2],st[N*2],pa[N*2];
int flo_from[N*2][N+1],S[N*2],vis[N*2];
vector<int> flo[N*2]:
queue<int> q;
int e_delta(const edge &e){
return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
void update_slack(int u,int x){
if(!slack[x]||e_delta(g[u][x])<e_delta(g[slack[x]][x]))slack[x]=u;</pre>
void set_slack(int x){
slack[x]=0:
for(int u=1;u<=n;++u)</pre>
  if(g[u][x].w>0&&st[u]!=x&&S[st[u]]==0)
  update_slack(u,x);
void q_push(int x){
if(x<=n)q.push(x);</pre>
else for(size_t i=0;i<flo[x].size();i++)</pre>
  q_push(flo[x][i]);
void set st(int x.int b){
st[x]=b:
if(x>n)for(size_t i=0;i<flo[x].size();++i)</pre>
  set st(flo[x][i].b):
int get pr(int b.int xr){
int pr=find(flo[b].begin(),flo[b].end(),xr)-flo[b].begin();
if(pr%2==1){
 reverse(flo[b].begin()+1,flo[b].end());
  return (int)flo[b].size()-pr;
}else return pr;
}
void set_match(int u,int v){
match[u]=g[u][v].v;
if(u<=n) return:
edge e=g[u][v];
int xr=flo_from[u][e.u],pr=get_pr(u,xr);
for(int i=0;i<pr;++i)set_match(flo[u][i],flo[u][i^1]);</pre>
set match(xr.v):
rotate(flo[u].begin(),flo[u].begin()+pr,flo[u].end());
void augment(int u,int v){
```

```
for(;;){
  int xnv=st[match[u]];
  set match(u.v):
  if(!xnv)return;
  set match(xnv.st[pa[xnv]]):
  u=st[pa[xnv]],v=xnv;
}
int get_lca(int u,int v){
static int t=0:
for(++t;u||v;swap(u,v)){
  if(u==0)continue;
  if(vis[u]==t)return u;
  vis[u]=t;
  u=st[match[u]];
  if(u)u=st[pa[u]]:
}
return 0;
}
void add_blossom(int u,int lca,int v){
int b=n+1:
while(b \le n_x \& t[b])++b;
if(b>n_x)++n_x;
lab[b]=0,S[b]=0;
match[b]=match[lca];
flo[b].clear():
flo[b].push back(lca):
for(int x=u,y;x!=lca;x=st[pa[y]])
  flo[b].push_back(x),flo[b].push_back(y=st[match[x]]),q_push(y);
reverse(flo[b].begin()+1,flo[b].end());
for(int x=v,v;x!=lca;x=st[pa[v]])
  flo[b].push_back(x),flo[b].push_back(y=st[match[x]]),q_push(y);
set_st(b,b);
for(int x=1; x \le n_x; ++x)g[b][x].w=g[x][b].w=0;
for(int x=1;x<=n;++x)flo_from[b][x]=0;</pre>
for(size_t i=0;i<flo[b].size();++i){</pre>
  int xs=flo[b][i];
  for(int x=1;x<=n_x;++x)</pre>
  if(g[b][x].w==0||e_delta(g[xs][x]) < e_delta(g[b][x]))
    g[b][x]=g[xs][x],g[x][b]=g[x][xs];
  for(int x=1;x\leq n;++x)
  if(flo_from[xs][x])flo_from[b][x]=xs;
set_slack(b);
void expand_blossom(int b){
for(size_t i=0;i<flo[b].size();++i)</pre>
  set_st(flo[b][i],flo[b][i]);
int xr=flo_from[b][g[b][pa[b]].u],pr=get_pr(b,xr);
for(int i=0;i<pr;i+=2){</pre>
```

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```
int xs=flo[b][i],xns=flo[b][i+1];
  pa[xs]=g[xns][xs].u;
  S[xs]=1.S[xns]=0:
  slack[xs]=0,set_slack(xns);
  q_push(xns);
S[xr]=1,pa[xr]=pa[b];
for(size_t i=pr+1;i<flo[b].size();++i){</pre>
  int xs=flo[b][i];
 S[xs]=-1, set_slack(xs);
}
st[b]=0;
}
bool on_found_edge(const edge &e){
int u=st[e.u],v=st[e.v];
if(S[v]==-1){
 pa[v]=e.u,S[v]=1;
  int nu=st[match[v]];
  slack[v]=slack[nu]=0;
  S[nu]=0,q_push(nu);
}else if(S[v]==0){
  int lca=get_lca(u,v);
  if(!lca)return augment(u,v),augment(v,u),true;
  else add_blossom(u,lca,v);
return false:
bool matching(){
memset(S+1,-1,sizeof(int)*n x):
memset(slack+1,0,sizeof(int)*n_x);
q=queue<int>();
for(int x=1;x<=n_x;++x)</pre>
  if (st[x] == x\&\&! match[x])pa[x] = 0, S[x] = 0, q_push(x);
if(q.empty())return false;
for(;;){
  while(q.size()){
  int u=q.front();q.pop();
  if(S[st[u]]==1)continue;
  for(int v=1; v<=n;++v)</pre>
    if(g[u][v].w>0&&st[u]!=st[v]){
    if(e_delta(g[u][v])==0){
      if(on_found_edge(g[u][v]))return true;
    }else update_slack(u,st[v]);
    }
  }
  int d=INF;
  for(int b=n+1;b \le n_x;++b)
  if(st[b]==b&&S[b]==1)d=min(d,lab[b]/2);
  for(int x=1;x<=n_x;++x)</pre>
  if(st[x] == x&&slack[x]){
```

```
if(S[x]==-1)d=min(d,e_delta(g[slack[x]][x]));
    else if(S[x]==0)d=min(d,e_delta(g[slack[x]][x])/2);
  }
  for(int u=1;u<=n;++u){
  if(S[st[u]]==0){
    if(lab[u] <= d) return 0;</pre>
    lab[u]-=d;
  }else if(S[st[u]]==1)lab[u]+=d:
  }
  for(int b=n+1;b<=n_x;++b)
  if(st[b]==b){
    if(S[st[b]]==0)lab[b]+=d*2;
    else if(S[st[b]]==1)lab[b]-=d*2;
  }
  q=queue<int>();
  for(int x=1:x<=n x:++x)
  if(st[x]==x&&slack[x]&&st[slack[x]]!=x&&e_delta(g[slack[x]][x])==0)
    if(on_found_edge(g[slack[x]][x]))return true;
  for(int b=n+1:b \le n x:++b)
  if (st[b]==b\&\&S[b]==1\&\&lab[b]==0) expand_blossom(b);
}
return false;
pair<long long,int> solve(){
memset(match+1,0,sizeof(int)*n);
n x=n:
int n matches=0:
long long tot_weight=0;
for(int u=0:u<=n:++u)st[u]=u.flo[u].clear():</pre>
int w max=0:
for(int u=1;u<=n;++u)</pre>
 for(int v=1:v<=n:++v){
  flo_from[u][v]=(u==v?u:0);
  w max=max(w max.g[u][v].w):
  }
for(int u=1;u<=n;++u)lab[u]=w_max;</pre>
while(matching())++n_matches;
for(int u=1;u<=n;++u)</pre>
  if (match[u] &&match[u] < u)</pre>
  tot weight+=g[u][match[u]].w:
return make_pair(tot_weight,n_matches);
void add_edge( int ui , int vi , int wi ){
g[ui][vi].w = g[vi][ui].w = wi;
void init( int _n ){
n = _n;
for(int u=1:u<=n:++u)
 for(int v=1;v<=n;++v)</pre>
  g[u][v]=edge(u,v,0);
```

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```
} graph;
3 Graph
3.1 2-SAT
  Should be revised.
strongly_connected scc;
int n: // = number of clauses
void init(int _n){ scc.clear(); n = _n; }
int NOT(int x) { return x \ge n ? (x - n) : (x + n); }
void add_edge(int x, int y){ // input ~x to denote NOT
 if((x >> 31) & 1) x = (~x) + n;
 if((y >> 31) & 1) y = (~y) + n;
  scc.add_edge(x, y), scc.add_edge(NOT(y), NOT(x));
bool satisfy(vector<bool> &res){
  res.resize(n);
  scc.get_scc(2*n);
  for(int i=0; i<n; i++){</pre>
    if(scc.comp[i] == scc.comp[NOT(i)]) return 0;
    if(scc.comp[i] < scc.comp[NOT(i)]) res[i] = 0;</pre>
    else res[i] = 1;
 }
  return 1;
}
3.2 BCC
void color(int x, int p){
  if(p){
    bcc[p].push_back(x);
    cmp[x].push_back(p);
  for(auto &i : gph[x]){
    if(cmp[i].size()) continue;
    if(low[i] >= dfn[x]){
     bcc[++c].push_back(x);
      cmp[x].push_back(c);
      color(i, c);
    else color(i, p);
      Splay Tree + Link-Cut Tree
// Checklist 1. Is it link cut, or splay?
// Checklist 2. In link cut, is son always root?
```

```
void rotate(node *x){
 if(!x->p) return;
 push(x->p); // if there's lazy stuff
 push(x);
 node *p = x->p;
 bool is_left = (p->l == x);
 node *b = (is_left ? x->r : x->1);
 x->p = p->p;
 if(x-p \&\& x-p-1 == p) x-p-1 = x;
 if(x-p \&\& x-p-r == p) x-p-r = x;
 if(is_left){
   if(b) b->p = p;
   p->1 = b;
   p->p = x;
   x->r = p;
 }
 else{
   if(b) b->p = p;
   p->r = b;
   p->p = x;
   x->1 = p;
 pull(p); // if there's something to pull up
 pull(x);
 if(!x->p) root = x; // IF YOU ARE SPLAY TREE
 if(p->pp){ // IF YOU ARE LINK CUT TREE
   x->pp = p->pp;
   p->pp = NULL;
 }
void splay(node *x){
 while(x->p){
   node *p = x->p;
   node *g = p \rightarrow p;
   if(g){
     if((p->1 == x) ^ (g->1 == p)) rotate(x);
     else rotate(p);
   rotate(x);
 }
void access(node *x){
 splay(x);
 push(x);
 if(x->r){
   x->r->pp = x;
   x->r->p = NULL;
   x->r = NULL;
 pull(x);
```

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```
while(x->pp){
   node *nxt = x->pp;
    splay(nxt);
    push(nxt);
    if(nxt->r){}
     nxt->r->pp = nxt;
      nxt->r->p = NULL;
     nxt->r = NULL:
    }
    nxt->r = x;
    x->p = nxt;
    x->pp = NULL;
    pull(nxt);
    splay(x);
}
node *root(node *x){
  access(x):
  while(x->1){
   push(x);
   x = x -> 1:
  access(x);
  return x;
node *par(node *x){
  access(x):
 if(!x->1) return NULL;
 push(x):
  x = x->1;
  while (x->r) {
   push(x);
   x = x->r;
  access(x);
  return x;
}
node *lca(node *s, node *t){
  access(s):
  access(t):
  splay(s);
  if(s->pp == NULL) return s;
  return s->pp;
void link(node *par, node *son){
  access(par);
  access(son);
  son->rev ^= 1: // remove if needed
  push(son);
  son->1 = par;
```

```
par->p = son;
 pull(son);
void cut(node *p){
 access(p):
 push(p);
 if(p->1){
   p->1->p = NULL;
   p->1 = NULL;
 pull(p);
3.4 Offline Dynamic MST
 Should be revised.
int n, m, q;
int st[MAXN], ed[MAXN], cost[MAXN], chk[MAXN];
pi qr[MAXN];
bool cmp(int &a, int &b){ return pi(cost[a], a) < pi(cost[b], b);}
void contract(int s, int e, vector<int> v, vector<int> &must_mst, vector<int>
&maybe_mst){
 sort(v.begin(), v.end(), cmp);
 vector<pi> snapshot;
 for(int i=s; i<=e; i++) disj.uni(st[qr[i].first], ed[qr[i].first], snapshot);</pre>
 for(auto &i : v) if(disj.uni(st[i], ed[i], snapshot)) must_mst.push_back(i);
 disj.revert(snapshot);
 for(auto &i : must_mst) disj.uni(st[i], ed[i], snapshot);
 for(auto &i : v) if(disj.uni(st[i], ed[i], snapshot)) maybe_mst.push_back(i);
 disj.revert(snapshot);
void solve(int s. int e. vector<int> v. lint cv){
 if(s == e){}
    cost[qr[s].first] = qr[s].second;
   if(st[qr[s].first] == ed[qr[s].first]){
     printf("%lld\n", cv);
     return:
   7
   int minv = qr[s].second;
   for(auto &i : v) minv = min(minv, cost[i]);
   printf("%lld\n",minv + cv);
   return;
 int m = (s+e)/2;
 vector<int> lv = v, rv = v;
 vector<int> must_mst, maybe_mst;
 for(int i=m+1; i<=e; i++){</pre>
```

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```
chk[qr[i].first]--;
    if(chk[qr[i].first] == 0) lv.push_back(qr[i].first);
  }
  vector<pi> snapshot;
  contract(s. m. lv. must mst. maybe mst):
  lint lcv = cv;
  for(auto &i : must_mst) lcv += cost[i], disj.uni(st[i], ed[i], snapshot);
  solve(s, m, maybe_mst, lcv);
  disj.revert(snapshot);
  must_mst.clear(); maybe_mst.clear();
  for(int i=m+1; i<=e; i++) chk[qr[i].first]++;</pre>
  for(int i=s; i<=m; i++){</pre>
    chk[qr[i].first]--;
    if(chk[qr[i].first] == 0) rv.push_back(qr[i].first);
  lint rcv = cv:
  contract(m+1, e, rv, must_mst, maybe_mst);
  for(auto &i : must_mst) rcv += cost[i], disj.uni(st[i], ed[i], snapshot);
  solve(m+1, e, maybe_mst, rcv);
  disj.revert(snapshot);
  for(int i=s: i<=m: i++) chk[qr[i].first]++:</pre>
}
int main(){
  scanf("%d %d",&n,&m);
  vector<int> ve:
  for(int i=0: i<m: i++){</pre>
    scanf("%d %d %d", &st[i], &ed[i], &cost[i]);
 }
  scanf("%d",&q);
  for(int i=0; i<q; i++){</pre>
    scanf("%d %d",&qr[i].first,&qr[i].second);
    gr[i].first--;
    chk[ar[i].first]++:
  disj.init(n);
  for(int i=0; i<m; i++) if(!chk[i]) ve.push_back(i);</pre>
  solve(0, q-1, ve, 0);
      Dominator Tree
namespace dtree{ // by cki86201
  vector<int> E[MAXN], RE[MAXN], rdom[MAXN];
  int S[MAXN], RS[MAXN], cs;
  int par[MAXN], val[MAXN], sdom[MAXN], rp[MAXN], dom[MAXN];
  void clear(int n) {
    cs = 0:
    for(int i=0;i<=n;i++) {</pre>
```

```
par[i] = val[i] = sdom[i] = rp[i] = dom[i] = S[i] = RS[i] = 0;
     E[i].clear(); RE[i].clear(); rdom[i].clear();
   }
 }
 void add edge(int x, int v) { E[x].push back(v): }
 void Union(int x, int y) { par[x] = y; }
  int Find(int x, int c = 0) {
   if(par[x] == x) return c ? -1 : x:
   int p = Find(par[x], 1);
   if(p == -1) return c ? par[x] : val[x];
   if(sdom[val[x]] > sdom[val[par[x]]]) val[x] = val[par[x]];
   par[x] = p;
   return c ? p : val[x];
 }
 void dfs(int x) {
   RS[S[x] = ++cs] = x:
   par[cs] = sdom[cs] = val[cs] = cs;
   for(int e : E[x]) {
     if(S[e] == 0) dfs(e), rp[S[e]] = S[x];
     RE[S[e]].push_back(S[x]);
   }
 }
 int solve(int s, int *up) { // Calculate idoms
   dfs(s):
   for(int i=cs;i;i--) {
     for(int e : RE[i]) sdom[i] = min(sdom[i], sdom[Find(e)]);
     if(i > 1) rdom[sdom[i]].push back(i):
     for(int e : rdom[i]) {
       int p = Find(e):
       if(sdom[p] == i) dom[e] = i;
       else dom[e] = p;
     if(i > 1) Union(i, rp[i]);
   for(int i=2:i<=cs:i++) if(sdom[i] != dom[i]) dom[i] = dom[dom[i]]:</pre>
   for(int i=2;i<=cs;i++) up[RS[i]] = RS[dom[i]];</pre>
   return cs:
 }
}
3.6 Global Min-Cut.
namespace stoer_wagner{
 int minimum_cut_phase(int n, int &s, int &t, vector<vector<int>> &adj, vector<int>
 vis){
   vector<int> dist(n);
   int mincut = 1e9:
   while(true){
     int pos = -1, cur = -1e9;
     for(int i=0: i<n: i++){
       if(!vis[i] && dist[i] > cur){
```

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```
cur = dist[i];
          pos = i;
      if(pos == -1) break:
      s = t;
      t = pos;
      mincut = cur:
      vis[pos] = 1;
     for(int i=0; i<n; i++){</pre>
        if(!vis[i]) dist[i] += adj[pos][i];
     }
    return mincut; // optimal s-t cut here is, {t} and V \ {t}
  int solve(int n. vector<vector<int>> adi){
    if(n <= 1) return 0;
    vector<int> vis(n);
    int ans = 1e9:
    for(int i=0; i<n-1; i++){</pre>
     int s. t:
      ans = min(ans, minimum_cut_phase(n, s, t, adj, vis));
      vis[t] = 1;
      for(int j=0; j<n; j++){</pre>
       if(!vis[i]){
          adj[s][j] += adj[t][j];
          adj[j][s] += adj[j][t];
       }
     }
      adj[s][s] = 0;
    return ans;
}:
      Edmond's Directed MST
// starts from node 0. assumes there exists at least one dmst.
// edge is reversed : if there is edge s -> e, INSERT IN gph[e]
struct edge { int to, cost, id; };
using elist = vector<edge>;
void dmst(vector<elist> &g, vector<int>& res) {
  const int n = g.size();
  vector<edge*> to(n);
  vector<int> u(n, 0);
  for (int i = 1; i < n; ++i) {
   int mn = g[i][0].cost:
   for (int j = 0; j < g[i].size(); ++j)
     mn = min(mn, g[i][j].cost);
    for (int j = 0; j < g[i].size(); ++j) {
```

if (g[i][j].cost == mn)

```
to[i] = &g[i][i];
    g[i][j].cost -= mn;
 }
for (int i = 1: i < n: ++i) {
 if (u[i]) continue;
 int x = i;
 vector<int> order(1, x);
  u[x] = 1;
  while (to[x]->to > 0 && u[to[x]->to] == 0) {
   x = to[x]->to;
   u[x] = 1:
    order.push_back(x);
  int y = to[x] \rightarrow to;
  vector<int> cycle(find(order.begin(), order.end(), y), order.end());
  if (cycle.size() == 0) continue;
  vector<int> in_cycle(n);
  for (int j = 0; j < cycle.size(); ++j) in_cycle[cycle[j]] = 1;</pre>
 vector<int> nw_id(n);
 for (int j = 1; j < n; ++j) nw_id[j] = nw_id[j - 1] + !in_cycle[j];
  int nn = n - cycle.size();
  vector<elist> gn(nn + 1);
 for (int j = 1; j < n; ++j) {
   if (in_cycle[j]) {
     for (int k = 0; k < g[j].size(); ++k)
        if (!in cvcle[g[i][k].to])
          gn[nn].push_back({nw_id[g[j][k].to], g[j][k].cost, g[j][k].id});
      for (int k = 0; k < g[j].size(); ++k)</pre>
        if (in_cycle[g[j][k].to])
          gn[nw_id[j]].push_back({nn, g[j][k].cost, g[j][k].id});
          gn[nw_id[j]].push_back({nw_id[g[j][k].to], g[j][k].cost, g[j][k].id});
   }
 }
  dmst(gn, res);
  set<int> used_e(res.begin(), res.end());
  for (int j = 0; j < cycle.size(); ++j) {
   bool found = false:
    for (int k = 0; k < g[cycle[j]].size(); ++k)
      found |= used_e.count(g[cycle[j]][k].id);
    if (found) {
      for (int k = 0; k < cycle.size(); ++k) {
        if (k != i)
          res.push_back(to[cycle[k]]->id);
      }
      return;
 }
```

```
for (int i = 1; i < n; ++i) res.push_back(to[i]->id);
}
```

4 Strings

4.1 Aho-Corasick Algorithm

```
const int MAXN = 100005, MAXC = 26:
int trie[MAXN][MAXC], fail[MAXN], term[MAXN], piv;
void init(vector<string> &v){
 memset(trie, 0, sizeof(trie));
 memset(fail, 0, sizeof(fail));
 memset(term, 0, sizeof(term));
 piv = 0;
 for(auto &i : v){
   int p = 0;
   for(auto &j : i){
     if(!trie[p][j]) trie[p][j] = ++piv;
     p = trie[p][j];
   }
   term[p] = 1;
  queue<int> que;
  for(int i=0: i<MAXC: i++){</pre>
   if(trie[0][i]) que.push(trie[0][i]);
  while(!que.empty()){
   int x = que.front();
   que.pop();
   for(int i=0; i<MAXC; i++){</pre>
     if(trie[x][i]){
       int p = fail[x]:
       while(p && !trie[p][i]) p = fail[p];
       p = trie[p][i];
       fail[trie[x][i]] = p;
       if(term[p]) term[trie[x][i]] = 1;
       que.push(trie[x][i]);
   }
 }
bool query(string &s){
 int p = 0;
 for(auto &i : s){
   while(p && !trie[p][i]) p = fail[p];
   p = trie[p][i];
   if(term[p]) return 1;
 }
 return 0;
```

```
4.2 Suffix Array
 Should be revised.
const int MAXN = 500005;
int ord[MAXN], nord[MAXN], cnt[MAXN], aux[MAXN];
void solve(int n, char *str, int *sfx, int *rev, int *lcp){
 int p = 1;
 memset(ord, 0, sizeof(ord));
 for(int i=0; i<n; i++){
   sfx[i] = i;
   ord[i] = str[i];
 int pnt = 1;
 while(1){
   memset(cnt, 0, sizeof(cnt));
   for(int i=0; i<n; i++) cnt[ord[min(i+p, n)]]++;</pre>
   for(int i=1; i<=n || i<=255; i++) cnt[i] += cnt[i-1];
   for(int i=n-1; i>=0; i--)
     aux[--cnt[ord[min(i+p, n)]]] = i:
   memset(cnt, 0, sizeof(cnt));
   for(int i=0; i<n; i++) cnt[ord[i]]++;</pre>
   for(int i=1; i<=n || i<=255; i++) cnt[i] += cnt[i-1];
   for(int i=n-1; i>=0; i--)
     sfx[--cnt[ord[aux[i]]]] = aux[i];
   if(pnt == n) break;
   pnt = 1;
   nord[sfx[0]] = 1;
   for(int i=1; i<n; i++){</pre>
     if(ord[sfx[i-1]] != ord[sfx[i]] || ord[sfx[i-1] + p] != ord[sfx[i] + p]){
       pnt++;
     nord[sfx[i]] = pnt:
   memcpy(ord, nord, sizeof(int) * n);
 for(int i=0; i<n; i++) rev[sfx[i]] = i;</pre>
 int h = 0:
 for(int i=0; i<n; i++){
   if(rev[i]){
     int prv = sfx[rev[i] - 1];
     while(str[prv + h] == str[i + h]) h++;
     lcp[rev[i]] = h:
   }
   h = \max(h-1, 0);
```

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4.3 Manacher's Algorithm

```
const int MAXN = 1000005;
int aux[2 * MAXN - 1]:
void solve(int n. int *str. int *ret){
 // *ret : number of nonobvious palindromic character pair
 for(int i=0: i<n: i++){
   aux[2*i] = str[i];
   if(i != n-1) aux[2*i+1] = -1;
 int p = 0, c = 0;
 for(int i=0: i<2*n-1: i++){
   int cur = 0:
   if(i \le p) cur = min(ret[2 * c - i], p - i);
   while(i - cur - 1 >= 0 && i + cur + 1 < 2*n-1 && aux[i-cur-1] == aux[i+cur+1]){
     cur++;
   }
   ret[i] = cur:
   if(i + ret[i] > p){
     p = i + ret[i];
     c = i:
   }
 }
4.4
     eertree
namespace eertree{
 int nxt[MAXN][26];
 int par[MAXN], len[MAXN], slink[MAXN], ptr[MAXN], diff[MAXN], series[MAXN], piv;
  void clear(int n = MAXN){
   memset(par, 0, sizeof(int) * n);
   memset(len, 0, sizeof(int) * n);
   memset(slink, 0, sizeof(int) * n);
   memset(nxt, 0, sizeof(int) * 26 * n);
   piv = 0;
  void init(int n, char *a){
   par[0] = 0:
   par[1] = 1;
   a[0] = -1:
   len[0] = -1:
   piv = 1;
   int cur = 1:
   for(int i=1; i<=n; i++){
     while(a[i] != a[i - len[cur] - 1]) cur = slink[cur];
     if(!nxt[cur][a[i]]){
       nxt[cur][a[i]] = ++piv;
       par[piv] = cur;
       len[piv] = len[cur] + 2;
        int lnk = slink[cur];
```

```
while(a[i] != a[i - len[lnk] - 1]){
          lnk = slink[lnk]:
        if(nxt[lnk][a[i]]) lnk = nxt[lnk][a[i]];
        if(len[piv] == 1 || lnk == 0) lnk = 1:
        slink[piv] = lnk;
        diff[piv] = len[piv] - len[lnk];
        if(diff[piv] == diff[lnk]) series[piv] = series[lnk];
        else series[piv] = piv;
      cur = nxt[cur][a[i]];
      ptr[i] = cur;
 }
  int query(int s, int e){
    int pos = ptr[e]:
    while(len[pos] \geq e - s + 1){
            if(len[pos] % diff[pos] == (e - s + 1) % diff[pos] &&
               len[series[pos]] <= e - s + 1) return true;</pre>
            pos = series[pos];
            pos = slink[pos]:
   }
    return false;
  vector<pi> minimum_partition(int n){ // (odd min, even min)
   vector < pi > dp(n + 1):
   vector<pi> series ans(n + 10):
   dp[0] = pi(1e9 + 1, 0);
   for(int i=1: i<=n: i++){
      dp[i] = pi(1e9 + 1, 1e9);
     for(int j=ptr[i]; len[j] > 0;){
        int slv = slink[series[j]];
        series_ans[j] = dp[i - (len[slv] + diff[j])];
        if(diff[i] == diff[slink[i]]){
          series_ans[j].first = min(series_ans[j].first,
          series_ans[slink[j]].first);
          series_ans[j].second = min(series_ans[j].second,
          series_ans[slink[j]].second);
        auto val = series ans[i]:
        dp[i].first = min(dp[i].first, val.second + 1);
        dp[i].second = min(dp[i].second, val.first + 1);
        i = slv:
   }
   return dp;
 }
};
```

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4.5 Circular LCS

```
string s1, s2;
int dp[4005][2005];
int nxt[4005][2005];
int n, m;
void reroot(int px){
  int py = 1;
  while(py <= m && nxt[px][py] != 2) py++;</pre>
  nxt[px][py] = 1;
  while(px < 2 * n \&\& py < m){
    if(nxt[px+1][py] == 3){
      px++;
      nxt[px][py] = 1;
    else if(nxt[px+1][py+1] == 2){
      px++;
      py++;
      nxt[px][py] = 1;
    else py++;
  while(px < 2 * n \&\& nxt[px+1][py] == 3){
    nxt[px][py] = 1;
}
int track(int x, int y, int e){ // use this routine to find LCS as string
  int ret = 0;
  while(y != 0 \&\& x != e){
    if(nxt[x][v] == 1) v--:
    else if(nxt[x][y] == 2) ret += (s1[x] == s2[y]), x--, y--;
    else if(nxt[x][v] == 3) x--:
 }
  return ret;
int solve(string a, string b){
  n = a.size(), m = b.size();
  s1 = "#" + a + a:
  s1 = '#' + b:
  for(int i=0; i<=2*n; i++){
    for(int j=0; j<=m; j++){</pre>
      if(i == 0){
        nxt[i][j] = 3;
        continue;
      }
      if(i == 0){
```

```
nxt[i][j] = 1;
      continue;
    dp[i][j] = -1;
    if(dp[i][j] < dp[i][j-1]){
      dp[i][j] = dp[i][j-1];
      nxt[i][j] = 1;
    if(dp[i][j] < dp[i-1][j-1] + (s1[i] == s2[j])){
      dp[i][j] = dp[i-1][j-1] + (s1[i] == s2[j]);
      nxt[i][j] = 2;
    if(dp[i][j] < dp[i-1][j]){
      dp[i][j] = dp[i-1][j];
      nxt[i][j] = 3;
 }
int ret = dp[n][m];
for(int i=1; i<n; i++){
 reroot(i), ret = max(ret, track(n+i, m, i));
return ret;
```

6 Geometry

5.1 Green's Theorem

Let C is positive, smooth, simple curve. D is region bounded by C.

$$\oint_C (Ldx + Mdy) = \iint_D (\frac{\partial M}{\partial x} - \frac{\partial L}{\partial y})$$

To calculate area, $\frac{\partial M}{\partial x} - \frac{\partial L}{\partial y} = 1$, common selection is $M = \frac{1}{2}x$, $L = -\frac{1}{2}y$.

Line integral of circle parametrized by $(x,y)=(x_C+r_C\cos\theta,\ y_C+r_C\sin\theta)$ is given as follows.: $\frac{1}{2}(r_C(x_C(\sin\theta_f-\sin\theta_i)-y_C(\cos\theta_f-\cos\theta_i))+(\theta_f-\theta_i)r_C^2)$.

Line integral of line parametrized by $(x,y)=t(x_1,y_1)+(1-t)(x_2,y_2)$ is given as follows:: $\frac{1}{2}(x_1y_2-x_2y_1)$.

inline double arc_area(double x, double y, double r, double s, double e){ //s and e
are line integral theta value.
 return (r*(x*(sin(e)-sin(s))-y*(cos(e)-cos(s)))+(e-s)*r*r)*0.5;

inline double polygon_area(double x1, double y1, double x2, double y2){ //Shoelace formula

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```
return (x1*y2-x2*y1)*0.5;
```

Smallest Enclosing Circle / Sphere

```
namespace cover_2d{
  double eps = 1e-9;
  using Point = complex<double>;
  struct Circle{ Point p; double r; };
  double dist(Point p, Point q){ return abs(p-q); }
  double area2(Point p, Point q){ return (conj(p)*q).imag();}
  bool in(const Circle& c. Point p) { return dist(c.p. p) < c.r + eps: }
  Circle INVAL = Circle{Point(0, 0), -1};
  Circle mCC(Point a, Point b, Point c){
   b -= a: c -= a:
   double d = 2*(conj(b)*c).imag(); if(abs(d)<eps) return INVAL;</pre>
   Point ans = (c*norm(b) - b*norm(c)) * Point(0, -1) / d:
    return Circle{a + ans, abs(ans)};
  Circle solve(vector<Point> p) {
    mt19937 gen(0x94949); shuffle(p.begin(), p.end(), gen);
    Circle c = INVAL;
    for(int i=0; i<p.size(); ++i) if(c.r<0 ||!in(c, p[i])){</pre>
     c = Circle{p[i], 0};
     for(int j=0; j<=i; ++j) if(!in(c, p[j])){
       Circle ans{(p[i]+p[j])*0.5, dist(p[i], p[j])*0.5};
        if(c.r == 0) {c = ans; continue;}
        Circle 1. r: 1 = r = INVAL:
        Point pq = p[i]-p[i];
        for(int k=0; k<=j; ++k) if(!in(ans, p[k])) {</pre>
          double a2 = area2(pq, p[k]-p[i]);
          Circle c = mCC(p[i], p[j], p[k]);
          if(c.r<0) continue:
          else if(a2 > 0 && (1.r<0||area2(pg, c.p-p[i]) > area2(pg, 1.p-p[i]))) 1 =
          else if(a2 < 0 && (r.r<0||area2(pq, c.p-p[i]) < area2(pq, r.p-p[i]))) r =
        if(1.r<0\&\&r.r<0) c = ans:
        else if(1.r<0) c = r;
        else if(r.r<0) c = 1:
        else c = 1.r<=r.r?1:r:
      }
    }
    return c;
}:
namespace cover_3d{
  double enclosing_sphere(vector<double> x, vector<double> y, vector<double> z){
    int n = x.size();
```

```
auto hyp = [](double x, double y, double z){
     return x * x + y * y + z * z;
   };
   double px = 0, py = 0, pz = 0;
   for(int i=0: i<n: i++){
     px += x[i];
     pv += v[i];
     pz += z[i];
    px *= 1.0 / n:
   py *= 1.0 / n;
   pz *= 1.0 / n;
   double rat = 0.1, maxv:
   for(int i=0; i<10000; i++){
     maxv = -1;
     int maxp = -1:
     for(int j=0; j<n; j++){
       double tmp = hyp(x[j] - px, y[j] - py, z[j] - pz);
        if(maxv < tmp){</pre>
         maxv = tmp;
         maxp = j;
      px += (x[maxp] - px) * rat;
     py += (y[maxp] - py) * rat;
     pz += (z[maxp] - pz) * rat;
     rat *= 0.998:
   return sqrt(maxv);
};
5.3 3D Convex Hull
// code credit : https://gist.github.com/msg555/4963794
struct vec3{
 11 x, y, z;
 vec3(): x(0), y(0), z(0) {}
 vec3(11 a, 11 b, 11 c): x(a), y(b), z(c) {}
 vec3 operator*(const vec3& v) const{ return vec3(y*v.z-z*v.y, z*v.x-x*v.z,
 x*v.v-v*v.x): }
 vec3 operator-(const vec3& v) const{ return vec3(x-v.x, y-v.y, z-v.z); }
 vec3 operator-() const{ return vec3(-x, -y, -z); }
 11 dot(const vec3 &v) const{ return x*v.x+y*v.y+z*v.z; }
};
struct twoset {
 int a, b;
 void insert(int x) { (a == -1 ? a : b) = x; }
 bool contains(int x) { return a == x || b == x: }
 void erase(int x) { (a == x ? a : b) = -1; }
```

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```
int size() { return (a != -1) + (b != -1); }
} E[MAXN][MAXN]; // i < j</pre>
struct face{
  vec3 norm:
 ll disc;
  int I[3];
}:
face make_face(int i, int j, int k, int ii, vector<vec3> &A){ // p^T * norm < disc</pre>
  E[i][j].insert(k); E[i][k].insert(j); E[j][k].insert(i);
  face f; f.I[0]=i, f.I[1]=j, f.I[2]=k;
  f.norm = (A[j]-A[i])*(A[k]-A[i]);
  f.disc = f.norm.dot(A[i]);
  if(f.norm.dot(A[ii])>f.disc){
   f.norm = -f.norm:
    f.disc = -f.disc;
  return f;
}
vector<face> get_hull(vector<vec3> &A){
  int N = A.size():
  vector<face> faces; memset(E, -1, sizeof(E));
  faces.push_back(make_face(0,1,2,3,A));
  faces.push_back(make_face(0,1,3,2,A));
  faces.push back(make face(0.2.3.1.A)):
  faces.push_back(make_face(1,2,3,0,A));
  for(int i=4: i<N: ++i){</pre>
   for(int j=0; j<faces.size(); ++j){</pre>
      face f = faces[i];
      if(f.norm.dot(A[i])>f.disc){
        E[f.I[0]][f.I[1]].erase(f.I[2]);
        E[f.I[0]][f.I[2]].erase(f.I[1]);
        E[f.I[1]][f.I[2]].erase(f.I[0]);
        faces[j--] = faces.back();
        faces.pop_back();
    }
    int nf = faces.size():
    for(int j=0; j<nf; ++j){</pre>
      face f=faces[i]:
      for(int a=0: a<3: ++a) for(int b=a+1: b<3: ++b){
        int c=3-a-b:
        if(E[f.I[a]][f.I[b]].size()==2) continue;
        faces.push_back(make_face(f.I[a], f.I[b], i, f.I[c], A));
     }
    }
  return faces:
```

5.4 Dynamic Convex Hull Trick

```
// code credit : https://github.com/niklasb/contest-algos/
// blob/master/convex_hull/dynamic.cpp
using line t = double:
const line_t is_query = -1e18;
struct Line {
 line_t m, b;
  mutable function<const Line*()> succ:
  bool operator<(const Line& rhs) const {
   if (rhs.b != is_query) return m < rhs.m;</pre>
   const Line* s = succ():
   if (!s) return 0:
   line t x = rhs.m:
   return b - s->b < (s->m - m) * x;
 }
};
struct HullDynamic : public multiset<Line> { // will maintain upper hull for maximum
 bool bad(iterator v) {
   auto z = next(y);
   if (v == begin()) {
     if (z == end()) return 0:
     return y->m == z->m && y->b <= z->b;
   auto x = prev(v):
   if (z == end()) return y->m == x->m && y->b <= x->b;
   return (x-b - y-b)*(z-m - y-m) >= (y-b - z-b)*(y-m - x-m);
 }
  void insert_line(line_t m, line_t b) {
   auto v = insert({ m, b }):
   v->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
   if (bad(v)) { erase(v): return: }
   while (next(y) != end() && bad(next(y))) erase(next(y));
    while (y != begin() && bad(prev(y))) erase(prev(y));
 line_t query(line_t x) {
   auto 1 = *lower_bound((Line) { x, is_query });
   return 1.m * x + 1.b:
 }
}H:
     Half-plane Intersection
const double eps = 1e-8;
```

```
typedef pair < long double, long double > pi;
namespace hpi{
bool z(long double x){ return fabs(x) < eps; }</pre>
```

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```
struct line{
    long double a, b, c;
    bool operator<(const line &1)const{</pre>
                                                                                          };
     bool flag1 = pi(a, b) > pi(0, 0);
     bool flag2 = pi(1.a, 1.b) > pi(0, 0):
      if(flag1 != flag2) return flag1 > flag2;
                                                                                          while(s+1 < e){
      long double t = ccw(pi(0, 0), pi(a, b), pi(l.a, l.b));
                                                                                            m = (s+e) / 2;
      return z(t) ? c * hypot(l.a, l.b) < l.c * hypot(a, b) : <math>t > 0;
    pi slope(){ return pi(a, b);}
  };
  pi cross(line a, line b){
                                                                                              else e = m:
    long double det = a.a * b.b - b.a * a.b;
    return pi((a.c * b.b - a.b * b.c) / det, (a.a * b.c - a.c * b.a) / det);
                                                                                             else{ // down
  bool bad(line a, line b, line c){
    if(ccw(pi(0, 0), a.slope(), b.slope()) <= 0) return false;</pre>
   pi crs = cross(a, b);
                                                                                            }
                                                                                          }
   return crs.first * c.a + crs.second * c.b >= c.c:
  bool solve(vector<line> v, vector<pi> &solution){ // ax + by <= c;</pre>
    sort(v.begin(), v.end());
                                                                                          return -1;
    deque<line> dq;
    for(auto &i : v){
      if(!dq.empty() && z(ccw(pi(0, 0), dq.back().slope(), i.slope()))) continue;
                                                                                         5.7 kd-tree
      while(dq.size() >= 2 && bad(dq[dq.size()-2], dq.back(), i)) dq.pop_back();
      while(dq.size() >= 2 && bad(i, dq[0], dq[1])) dq.pop_front();
      dq.push_back(i);
                                                                                         struct node{
                                                                                          pi pnt;
    while(dq.size() > 2 && bad(dq[dq.size()-2], dq.back(), dq[0])) dq.pop_back();
    while(dq.size() > 2 && bad(dq.back(), dq[0], dq[1])) dq.pop_front();
                                                                                         }tree[270000];
    vector<pi> tmp:
    for(int i=0; i<dq.size(); i++){</pre>
                                                                                        pi a[100005];
                                                                                         int n. ok[270000]:
     line cur = dq[i], nxt = dq[(i+1)%dq.size()];
      if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps) return false;</pre>
      tmp.push_back(cross(cur, nxt));
    }
    solution = tmp;
    return true;
};
      Point-in-polygon test / Point-to-polygon tangent
// by zigui
                                                                                          int m = (s+e)/2;
// C : counter_clockwise(C[0] == C[N]), N >= 3
// return highest point in C <- P(clockwise) or -1 if strictly in P
// polygon is strongly convex, C[i] != P
int convex_tangent(vector<pi> &C, pi P, int up = 1){
  auto sign = [&](lint c){ return c > 0 ? up : c == 0 ? 0 : -up; };
```

```
auto local = [&](pi P, pi a, pi b, pi c) {
   return sign(ccw(P, a, b)) \le 0 && sign(ccw(P, b, c)) >= 0;
 int N = C.size()-1, s = 0, e = N, m;
 if( local(P, C[1], C[0], C[N-1]) ) return 0:
   if( local(P, C[m-1], C[m], C[m+1]) ) return m;
   if ( sign(ccw(P, C[s], C[s+1])) < 0) { // up}
     if (sign(ccw(P, C[m], C[m+1])) > 0) e = m;
     else if( sign(ccw(P, C[m], C[s])) > 0) s = m;
      if (sign(ccw(P, C[m], C[m+1])) < 0) s = m;
     else if( sign(ccw(P, C[m], C[s])) < 0) s = m;
 if( s && local(P, C[s-1], C[s], C[s+1]) ) return s;
 if( e != N && local(P, C[e-1], C[e], C[e+1]) ) return e;
typedef pair<int, int> pi;
 int spl, sx, ex, sy, ey;
lint sqr(int x){ return 111 * x * x: }
bool cmp1(pi a, pi b){ return a < b; }</pre>
bool cmp2(pi a, pi b){ return pi(a.second, a.first) < pi(b.second, b.first); }
// init(0, n-1, 1) : Initialize kd-tree
// set dap = INF, and call solve(1, P). dap = (closest point from P)
void init(int s, int e, int p){ // Initialize kd-tree
 int minx = 1e9, maxx = -1e9, miny = 1e9, maxy = -1e9;
 for(int i=s; i<=e; i++){</pre>
   minx = min(minx, a[i].first);
   miny = min(miny, a[i].second);
   maxx = max(maxx, a[i].first);
   maxv = max(maxv, a[i].second):
```

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```
tree[p].spl = (maxx - minx < maxy - miny);</pre>
  sort(a+s, a+e+1, [&](const pi &a, const pi &b){
   return tree[p].spl ? cmp2(a, b) : cmp1(a, b);
  });
  ok[p] = 1:
  tree[p] = {a[m], tree[p].spl, minx, maxx, miny, maxy};
  if(s <= m-1) init(s, m-1, 2*p);
  if(m+1 \le e) init(m+1, e, 2*p+1);
}
lint dap = 3e18;
void solve(int p, pi x){ // find closest point from point x (L^2)
  if(x != tree[p].pnt) dap = min(dap, sqr(x.first - tree[p].pnt.first) +
  sqr(x.second - tree[p].pnt.second));
  if(tree[p].spl){
    if(!cmp2(tree[p].pnt, x)){
      if(ok[2*p]) solve(2*p, x);
      if (ok[2*p+1] \&\& sqr(tree[2*p+1].sy - x.second) < dap) solve(2*p+1, x);
    }
    else{
      if(ok[2*p+1]) solve(2*p+1, x);
      if(ok[2*p] && sqr(tree[2*p].ey - x.second) < dap) solve(2*p, x);</pre>
    }
  }
  else{
    if(!cmp1(tree[p].pnt, x)){
      if(ok[2*p]) solve(2*p, x);
      if (ok[2*p+1] \&\& sqr(tree[2*p+1].sx - x.first) < dap) solve(2*p+1, x);
    }
    else{
      if(ok[2*p+1]) solve(2*p+1, x);
      if(ok[2*p] && sqr(tree[2*p].ex - x.first) < dap) solve(2*p, x);</pre>
 }
}
    Math
6.1 FFT / NTT
namespace fft{
  typedef complex<double> base;
  void fft(vector<base> &a, bool inv){
    int n = a.size(), j = 0;
    vector<base> roots(n/2):
    for(int i=1; i<n; i++){</pre>
     int bit = (n \gg 1);
```

while(j >= bit){
 j -= bit;

```
bit >>= 1;
    i += bit:
    if(i < j) swap(a[i], a[j]);</pre>
  double ang = 2 * acos(-1) / n * (inv ? -1 : 1);
  for(int i=0; i<n/2; i++){
    roots[i] = base(cos(ang * i), sin(ang * i));
  /* In NTT, let prr = primitive root. Then,
  int ang = ipow(prr, (mod - 1) / n);
  if(inv) ang = ipow(ang, mod - 2);
  for(int i=0: i<n/2: i++){
    roots[i] = (i ? (111 * roots[i-1] * ang % mod) : 1);
  Others are same. If there is /= n, do *= ipow(n, mod - 2).
  In XOR convolution, roots[*] = 1.
  for(int i=2: i<=n: i<<=1){
    int step = n / i;
    for(int j=0; j<n; j+=i){</pre>
      for(int k=0; k<i/2; k++){
        base u = a[j+k], v = a[j+k+i/2] * roots[step * k];
        a[j+k] = u+v;
        a[j+k+i/2] = u-v;
    }
  if(inv) for(int i=0; i<n; i++) a[i] /= n;
vector<lint> multiply(vector<lint> &v, vector<lint> &w){
  vector<base> fv(v.begin(), v.end()), fw(w.begin(), w.end());
  int n = 2; while(n < v.size() + w.size()) n <<= 1;</pre>
  fv.resize(n): fw.resize(n):
  fft(fv, 0); fft(fw, 0);
  for(int i=0; i<n; i++) fv[i] *= fw[i];</pre>
  fft(fv, 1);
  vector<lint> ret(n);
  for(int i=0: i<n: i++) ret[i] = (lint)round(fv[i].real()):
  return ret;
7
vector<lint> multiply(vector<lint> &v, vector<lint> &w, lint mod){
  int n = 2; while(n < v.size() + w.size()) n <<= 1;</pre>
  vector<base> v1(n), v2(n), r1(n), r2(n);
  for(int i=0; i<v.size(); i++){</pre>
    v1[i] = base(v[i] >> 15, v[i] & 32767);
  for(int i=0; i<w.size(); i++){</pre>
    v2[i] = base(w[i] >> 15, w[i] & 32767);
```

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```
}
    fft(v1, 0);
    fft(v2, 0):
    for(int i=0; i<n; i++){</pre>
      int i = (i ? (n - i) : i):
      base ans1 = (v1[i] + conj(v1[j])) * base(0.5, 0);
      base ans2 = (v1[i] - conj(v1[j])) * base(0, -0.5);
      base ans3 = (v2[i] + conj(v2[j])) * base(0.5, 0);
      base ans4 = (v2[i] - conj(v2[j])) * base(0, -0.5);
      r1[i] = (ans1 * ans3) + (ans1 * ans4) * base(0, 1);
      r2[i] = (ans2 * ans3) + (ans2 * ans4) * base(0, 1);
    fft(r1, 1):
    fft(r2, 1);
    vector<lint> ret(n);
    for(int i=0: i<n: i++){</pre>
      lint av = (lint)round(r1[i].real());
      lint bv = (lint)round(r1[i].imag()) + (lint)round(r2[i].real());
      lint cv = (lint)round(r2[i].imag());
      av %= mod, bv %= mod, cv %= mod;
      ret[i] = (av << 30) + (bv << 15) + cv:
      ret[i] %= mod;
      ret[i] += mod:
      ret[i] %= mod;
    return ret:
 }
}
     Hell-Joseon style FFT
#include <smmintrin.h>
#include <immintrin.h>
#pragma GCC target("avx2")
#pragma GCC target("fma")
m256d mult( m256d a, m256d b){
  _{\rm m256d} c = _{\rm mm256\_movedup\_pd(a)};
  _{m256d} d = _{mm256\_shuffle\_pd(a, a, 15)};
  m256d cb = mm256 mul pd(c, b):
  _{m256d} db = _{mm256_{mul_pd(d, b)}};
  _{\rm m256d} e = _{\rm mm256\_shuffle\_pd(db, db, 5)};
  _{\rm m256d} r = _{\rm mm256\_addsub\_pd(cb, e)};
  return r:
void fft(int n, __m128d a[], bool invert){
  for(int i=1, j=0; i<n; ++i){</pre>
    int bit = n >> 1:
   for(; j>=bit; bit>>=1) j -= bit;
    j += bit;
    if(i<j) swap(a[i], a[j]);</pre>
```

```
for(int len=2; len<=n; len<<=1){</pre>
    double ang = 2*3.14159265358979/len*(invert?-1:1);
    m256d wlen; wlen[0] = cos(ang), wlen[1] = sin(ang);
   for(int i=0: i<n: i += len){</pre>
     m256d w: w[0] = 1: w[1] = 0:
     for(int j=0; j<len/2; ++j){</pre>
        w = _mm256_permute2f128_pd(w, w, 0);
        wlen = _{mm256_{insertf128_{pd}(wlen, a[i+j+len/2], 1)};
        w = mult(w, wlen);
        _{m128d} vw = _{mm256} extractf128 pd(w, 1);
        _{m128d} u = a[i+j];
        a[i+j] = _mm_add_pd(u, vw);
        a[i+j+len/2] = _mm_sub_pd(u, vw);
   }
 7
 if(invert){
    _{m128d inv; inv[0] = inv[1] = 1.0/n;
   for(int i=0; i<n; ++i) a[i] = _mm_mul_pd(a[i], inv);</pre>
 }
vector<int64_t> multiply(vector<int64_t>& v, vector<int64_t>& w){
 int n = 2; while(n < v.size()+w.size()) n<<=1;</pre>
  _{m128d*} fv = new _{m128d[n]};
 for(int i=0; i<n; ++i) fv[i][0] = fv[i][1] = 0;
 for(int i=0; i<v.size(); ++i) fv[i][0] = v[i];</pre>
 for(int i=0: i<w.size(): ++i) fv[i][1] = w[i]:</pre>
 fft(n, fv, 0); //(a+bi) is stored in FFT
 for(int i=0: i<n: i += 2){
    __m256d a:
   a = _mm256_insertf128_pd(a, fv[i], 0);
   a = _mm256_insertf128_pd(a, fv[i+1], 1);
   a = mult(a, a);
   fv[i] = mm256 extractf128 pd(a, 0):
   fv[i+1] = _mm256_extractf128_pd(a, 1);
 fft(n, fv, 1):
  vector<int64_t> ret(n);
 for(int i=0; i<n; ++i) ret[i] = (int64_t)round(fv[i][1]/2);
 delete[] fv:
 return ret;
6.3 NTT Polynomial Division
vector<lint> get_inv(int n, const vector<lint> &p){
 vector<lint> q = \{ipow(p[0], mod - 2)\}:
 for(int i=2; i<=n; i<<=1){
   vector<lint> res:
    vector<lint> fq(q.begin(), q.end()); fq.resize(2*i);
    vector<lint> fp(p.begin(), p.begin() + i); fp.resize(2*i);
```

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```
fft(fq, 0); fft(fp, 0);
    for(int j=0; j<2*i; j++){
     fp[j] *= fq[j] * fq[j] % mod;
     fp[j] %= mod;
    fft(fp, 1);
    res.resize(i);
    for(int j=0; j<i; j++){</pre>
     res[i] = mod - fp[i];
     if(j < i/2) res[j] += 2 * q[j];
     res[j] %= mod;
    q = res;
  return q;
vector<lint> poly_divide(const vector<lint> &a, const vector<lint> &b){
  assert(b.back() != 0); // please trim leading zero
  int n = a.size(), m = b.size();
  int k = 2; while (k < n-m+1) k <<= 1;
  vector<lint> rb(k), ra(k);
  for(int i=0; i<m && i<k; ++i) rb[i] = b[m-i-1];</pre>
  for(int i=0; i<n && i<k; ++i) ra[i] = a[n-i-1];</pre>
  vector<lint> rbi = get_inv(k, rb);
  vector<lint> res = multiply(rbi, ra);
  res.resize(n - m + 1):
  reverse(res.begin(), res.end()):
  return res;
}
     Simplex Algorithm
/* Ax <= b. max c^T x
* Usage : Simplex(VVD A, VD b, VD c).solve(VD ans)
* not feasible : -INF; unbounded : INF
* accuracy ~ (size of ans) * EPS
* EPS recommended 1e-9 on double, 1e-12 on long double
* expected n ~ 100, 10ms. worst case is exponential */
using real_t = double;
using VD = vector<real_t>;
using VVD = vector<VD>;
const real_t EPS = 1e-9;
struct Simplex{
  int m, n;
  vector<int> B. N:
  VVD D:
  Simplex(const VVD& A, const VD& b, const VD &c)
   : m(b.size()), n(c.size()), N(n+1), B(m), D(m+2, VD(n+2)){
      for(int i=0; i<m; ++i) for(int j=0; j<n; ++j) D[i][j] = A[i][j];
```

```
for(int i=0; i<m; ++i) B[i] = n+i, D[i][n] = -1, D[i][n+1] = b[i];
      for(int j=0; j<n; ++j) N[j] = j, D[m][j] = -c[j];</pre>
     N[n] = -1: D[m+1][n] = 1:
 void Pivot(int r. int s) {
   real_t inv = 1/D[r][s];
   for(int i=0; i<m+2; ++i){
     for(int j=0; j<n+2; ++j){
        if(i != r && j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
     }
   }
   for(int i=0; i<m+2; ++i) if(i != r) D[i][s] *= -inv;
   for(int j=0; j<n+2; ++j) if(j != s) D[r][j] *= inv;
   D[r][s] = inv; swap(B[r], N[s]);
 bool Phase(bool p) {
    int x = m + p;
   while(true) {
      int s = -1:
     for(int j=0; j<=n; ++j){
        if(!p && N[j] == -1) continue;
        if(s == -1 || D[x][j] < D[x][s]) s = j;
      if(D[x][s] > -EPS) return true;
      int r = -1;
      for(int i=0: i<m: ++i){
        if(D[i][s] <= EPS) continue:</pre>
        if(r == -1 || D[i][n+1] / D[i][s] < D[r][n+1] / D[r][s]) r = i;
     if(r == -1) return false;
     Pivot(r, s);
   }
 real t solve(VD &x) {
    int r = 0:
   for(int i=1; i<m; ++i) if(D[i][n+1] < D[r][n+1]) r=i;
    if(D[r][n+1] < -EPS) {
     Pivot(r, n);
      if(!Phase(1) || D[m+1][n+1] < -EPS) return -1/0.0;
     for(int i=0; i<m; ++i) if(B[i] == -1) {
        int s = min_element(D[i].begin(), D[i].end() - 1) - D[i].begin();
        Pivot(i, s);
   }
   if(!Phase(0)) return 1/0.0:
   x = VD(n):
   for(int i=0; i<m; ++i) if(B[i] < n) x[B[i]] = D[i][n+1];
   return D[m][n+1]:
 }
};
```

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6.5 Range Prime Counting

```
// credit : https://github.com/stjepang/snippets/blob/master/count_primes.cpp
// Primes up to 10^12 can be counted in ~1 second.
const int MAXN = 1000005; // MAXN is the maximum value of sqrt(N) + 2
bool prime[MAXN]:
int prec[MAXN];
vector<int> P;
void init() {
 prime[2] = true;
 for (int i = 3; i < MAXN; i += 2) prime[i] = true;
 for (int i = 3: i*i < MAXN: i += 2){
    if (prime[i]){
     for (int i = i*i: i < MAXN: i += i+i) prime[i] = false:
   }
 }
  for(int i=1: i<MAXN: i++){</pre>
    if (prime[i]) P.push_back(i);
   prec[i] = prec[i-1] + prime[i];
}
lint rec(lint N, int K) {
 if (N <= 1 || K < 0) return 0;
 if (N <= P[K]) return N-1:
  if (N < MAXN && 111 * P[K]*P[K] > N) return N-1 - prec[N] + prec[P[K]];
  const int LIM = 250:
  static int memo[LIM*LIM][LIM];
  bool ok = N < LIM*LIM;</pre>
  if (ok && memo[N][K]) return memo[N][K]:
 lint ret = N/P[K] - rec(N/P[K], K-1) + rec(N, K-1);
  if (ok) memo[N][K] = ret;
 return ret;
}
lint count primes(lint N) { //less than or equal to
  if (N < MAXN) return prec[N];</pre>
 int K = prec[(int)sqrt(N) + 1];
  return N-1 - rec(N, K) + prec[P[K]];
}
     Discrete Kth root
/* code source : https://github.com/tzupengwang/PECaveros/blob/
master/codebook/math/DiscreteKthsqrt.cpp
* Solve x for x^P = A mod Q
* (P, Q-1) = 1 -> P^-1 \mod (Q-1) exists
* x has solution iff A^{(Q-1)} / P = 1 \mod Q
```

```
* PP | (Q-1) -> P < sqrt(Q), solve lgQ rounds of discrete log
* else \rightarrow find a s.t. s | (Pa - 1) \rightarrow ans = A^a */
```

```
using LL = long long;
LL mul(LL x, LL y, LL mod) { return (__int128) x * y % mod; }
LL add(LL x, LL v, LL mod) { return (x + v) % mod: }
LL pw(LL x, LL v, LL mod){
 LL ret = 1, piv = x:
 while(v){
   if(y & 1) ret = mul(ret, piv, mod);
   piv = mul(piv, piv, mod);
   y >>= 1;
 return ret % mod;
void gcd(LL a, LL b, LL &x, LL &y, LL &g){
 if (b == 0) {
   x = 1, y = 0, g = a;
   return:
 LL tx, ty;
 gcd(b, a%b, tx, ty, g);
 x = ty; y = tx - ty * (a / b);
LL P, A, Q, g; // x^P = A \mod Q
const int X = 1e5;
LL base, ae[X], aXe[X], iaXe[X];
unordered_map<LL, LL> ht;
#define FOR(i, c) for (int i = 0; i < (c); ++i)
#define REP(i, 1, r) for (int i = (1): i \le (r): ++i)
void build(LL a) { // ord(a) = P < sqrt(Q)</pre>
 base = a:
 ht.clear():
 ae[0] = 1; ae[1] = a; aXe[0] = 1; aXe[1] = pw(a, X, Q);
 iaXe[0] = 1; iaXe[1] = pw(aXe[1], Q-2, Q);
 REP(i, 2, X-1) {
   ae[i] = mul(ae[i-1], ae[1], Q):
   aXe[i] = mul(aXe[i-1], aXe[1], 0):
   iaXe[i] = mul(iaXe[i-1], iaXe[1], Q);
 FOR(i, X) ht[ae[i]] = i;
LL dis_log(LL x) {
 FOR(i, X) {
   LL iaXi = iaXe[i]:
   LL rst = mul(x, iaXi, Q);
   if (ht.count(rst)) return i*X + ht[rst]:
 }
LL main2() {
 cin >> P >> A >> Q;
 LL t = 0, s = 0-1:
```

```
while (s % P == 0) {
  ++t:
  s /= P:
if (A == 0) return 0:
if (t == 0) {
 // a^{P^-1 mod phi(Q)}
 LL x, y, _;
  gcd(P, Q-1, x, y, _);
  if (x < 0) {
    x = (x \% (Q-1) + Q-1) \% (Q-1);
  LL ans = pw(A, x, Q);
  if (pw(ans, P, Q) != A) while(1);
  return ans;
}
// A is not P-residue
if (pw(A, (Q-1) / P, Q) != 1) return -1;
for (g = 2; g < Q; ++g) {
  if (pw(g, (Q-1) / P, Q) != 1)
    break:
}
LL alpha = 0;
 LL y, _;
  gcd(P, s, alpha, y, _);
  if (alpha < 0) alpha = (alpha \% (Q-1) + Q-1) \% (Q-1);
if (t == 1) {
 LL ans = pw(A, alpha, Q);
  return ans;
}
LL a = pw(g, (Q-1) / P, Q);
build(a):
LL b = pw(A, add(mul(P%(Q-1), alpha, Q-1), Q-2, Q-1), Q);
LL c = pw(g, s, Q);
I.I. h = 1:
LL e = (Q-1) / s / P; // r^{t-1}
REP(i, 1, t-1) {
 e /= P:
 LL d = pw(b, e, Q);
 LL j = 0;
  if (d != 1) {
   j = -dis_log(d);
    if (j < 0) j = (j % (Q-1) + Q-1) % (Q-1);
  b = mul(b, pw(c, mul(P%(Q-1), j, Q-1), Q), Q);
  h = mul(h, pw(c, j, Q), Q);
  c = pw(c, P, Q);
}
```

```
return mul(pw(A, alpha, Q), h, Q);
6.7 Miller-Rabin Test + Pollard Rho Factorization
namespace miller_rabin{
   lint mul(lint x, lint y, lint mod){ return (__int128) x * y % mod; }
 lint ipow(lint x, lint y, lint p){
   lint ret = 1, piv = x \% p;
   while(v){
     if(y&1) ret = mul(ret, piv, p);
     piv = mul(piv, piv, p);
     y >>= 1;
   }
   return ret;
 bool miller rabin(lint x. lint a){
   if(x % a == 0) return 0;
   lint d = x - 1;
   while(1){
     lint tmp = ipow(a, d, x);
     if(d&1) return (tmp != 1 && tmp != x-1);
     else if(tmp == x-1) return 0;
     d >>= 1;
   }
 }
 bool isprime(lint x){
   for(auto &i : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
     if(x == i) return 1;
     if (x > 40 \&\& miller_rabin(x, i)) return 0;
   if(x <= 40) return 0;
   return 1:
 }
namespace pollard_rho{
 lint f(lint x, lint n, lint c){
   return (c + miller_rabin::mul(x, x, n)) % n;
 void rec(lint n, vector<lint> &v){
   if(n == 1) return:
   if(n \% 2 == 0){
   v.push_back(2);
   rec(n/2, v);
   return:
 if(miller_rabin::isprime(n)){
   v.push_back(n);
   return;
```

```
lint a, b, c;
  while(1){
    a = rand() \% (n-2) + 2;
    c = rand() \% 20 + 1;
     a = f(a, n, c);
     b = f(f(b, n, c), n, c);
    }while(gcd(abs(a-b), n) == 1);
    if(a != b) break;
  lint x = gcd(abs(a-b), n);
  rec(x, v);
  rec(n/x, v);
vector<lint> factorize(lint n){
  vector<lint> ret;
  rec(n, ret);
  sort(ret.begin(), ret.end());
  return ret:
}
};
```

6.8 Highly Composite Numbers, Large Prime

< 10	^k nu	mber	di	visors	2	3	5	7:	11:	13:	17:	192	232	29:	31	37
1		6		4	1	1										
2		60		12	2	1	1									
3		840		32	3	1	1	1								
4		7560		64	3	3	1	1								
5	8	3160		128	3	3	1	1	1							
6	72	0720		240	4	2	1	1	1	1						
7	864	8640		448	6	3	1	1	1	1						
8	7351	3440		768	5	3	1	1	1	1	1					
9	73513	4400		1344	6	3	2	1	1	1	1					
10	698377	6800		2304	5	3	2	1	1	1	1	1				
11	9777287	5200		4032	6	3	2	2	1	1	1	1				
12	96376119	8400		6720	6	4	2	1	1	1	1	1	1			
13	931635825	1200		10752	6	3	2	1	1	1	1	1	1	1		
14	9782176163	7600		17280	5	4	2	2	1	1	1	1	1	1		
15	86642131736	1600		26880	6	4	2	1	1	1	1	1	1	1	1	
16	808659896204	1600		41472	8	3	2	2	1	1	1	1	1	1	1	
17	7480104039888	4800		64512	6	3	2	2	1	1	1	1	1	1	1	1
18	89761248478661	7600		103680	8	4	2	2	1	1	1	1	1	1	1	1
< 10	^k prime	# of	prime		<	10)^ <u>]</u>	K						pı	riı	me
1	7		4	10	10						9999999967					
2	97		25		1	1					99	999	999	999	99,	77

```
168
                                                               99999999989
                9973
                                1229
                                                  13
                                                              999999999971
               99991
                                9592
                                                  14
                                                             9999999999973
             999983
                              78498
                                                  15
                                                            99999999999989
 7
                              664579
                                                           99999999999937
 8
                            5761455
                                                  17
           99999989
                                                         999999999999997
          99999937
                           50847534
                                                        9999999999999989
NTT Prime:
 998244353 = 119 \times 2^{23} + 1. Primitive root: 3.
 985661441 = 235 \times 2^{22} + 1. Primitive root: 3. 1012924417 = 483 \times 2^{21} + 1. Primitive root: 5.
```

7 Miscellaneous

7.1 Popular Optimization Technique

- Add or remove the constraints / Simplify the problem / Solve similar problem
- "Colliding ants problem": Find a better way to describe the situation
- CHT. DnC optimization. Mo's algorithm trick (on tree). IOI 2016 Aliens trick. IOI 2009 Regions trick.
- Knuth's $O(n^2)$ Optimal BST : minimize $D_{i,j} = Min_{i \leq k < j}(D_{i,k} + D_{k+1,j}) + C_{i,j}$. Quadrangle Inequality : $C_{a,c} + C_{b,d} \leq C_{a,d} + C_{b,c}$, $C_{b,c} \leq C_{a,d}$. Now monotonicity holds.
- Sqrt batch processing Save queries in buffer, and update in every sqrt steps (cf : IOI 2011 Elephant. hyea calls it "ainta technique")
- Dynamic insertion in static set (Make O(lqn) copy. Merge like binomial heap.)
- Offline insertion / deletion in insert-only set (Pair insertion-deletion operation, and regard it as range query)
- Atcoder Median Pyramid: Reduce the input to binary, and solve the easier problem.

7.2 Bit Twiddling Hack

```
int __builtin_clz(int x);//number of leading zero
int __builtin_ctz(int x);//number of trailing zero
int __builtin_clzll(long long x);//number of leading zero
int __builtin_ctzll(long long x);//number of trailing zero
int __builtin_popcount(int x);// number of 1-bits in x
int __builtin_popcountll(long long x);//number of 1-bits in x

lsb(n): (n & -n); // last bit (smallest)
floor(log2(n)): 31 - __builtin_clz(n | 1);
floor(log2(n)): 63 - __builtin_clzll(n | 1);
//compute next perm. ex) 00111, 01011, 01101, 01110, 10011, 10101...
long long next perm(long long y){
```

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```
long long t = v \mid (v-1);
  return (t + 1) \mid (((^t \& -^t) - 1) >> (_builtin_ctz(v) + 1));
}
     Fast Integer IO
static char buf[1 << 19]; // size : any number geq than 1024
static int idx = 0;
static int bytes = 0:
static inline int _read() {
  if (!bvtes || idx == bvtes) {
    bytes = (int)fread(buf, sizeof(buf[0]), sizeof(buf), stdin);
  return buf[idx++];
static inline int _readInt() {
  int x = 0, s = 1;
  int c = _read();
  while (c \le 32) c = read():
  if (c == '-') s = -1, c = _read();
  while (c > 32) x = 10 * x + (c - '0'), c = read():
  if (s < 0) x = -x;
  return x;
7.4 OSRank in g++
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update>
ordered_set;
ordered set X:
X.insert(1); X.insert(2); X.insert(4); X.insert(8); X.insert(16);
cout<<*X.find_by_order(1)<<endl; // 2</pre>
cout<<*X.find_by_order(2)<<endl; // 4</pre>
cout<<*X.find_by_order(4)<<endl; // 16</pre>
cout<<(end(X)==X.find_by_order(6))<<end1; // true</pre>
cout<<X.order of kev(-5)<<endl: // 0
cout<<X.order_of_key(1)<<endl; // 0</pre>
cout<<X.order_of_key(3)<<endl; // 2</pre>
cout<<X.order_of_key(4)<<endl; // 2</pre>
cout<<X.order_of_key(400)<<endl; // 5</pre>
```

7.5 Nasty Stack Hacks

```
int main2(){ return 0;}
int main(){
 size_t sz = 1<<29; //512MB
 void* newstack = malloc(sz);
 void* sp_dest = newstack + sz - sizeof(void*);
  asm volatile ("movg %0, %%rax\n\t"
  "movg %%rsp , (%%rax)\n\t"
  "movq %0, %%rsp\n\t": : "r"(sp_dest): );
 asm __volatile__("pop %rsp\n\t");
 return 0:
7.6 C++ / Environment Overview
// vimrc : set nu sc ci si ai sw=4 ts=4 bs=2 mouse=a syntax on
// compile : g++ -o PROB PROB.cpp -std=c++11 -Wall -02
// options : -fsanitize=address -Wfatal-errors
#include <bits/stdc++.h> // magic header
using namespace std; // magic namespace
struct StupidGCCCantEvenCompileThisSimpleCode{
 pair<int, int> array[1000000];
}; // https://gcc.gnu.org/bugzilla/show_bug.cgi?id=68203
// how to use rand (in 2017)
mt19937 rng(0x14004);
int randint(int lb, int ub){ return uniform_int_distribution<int>(lb, ub)(rng); }
// comparator overload
auto cmp = [](seg a, seg b){return a.func() < b.func(); };</pre>
set<seg, decltype(cmp)> s(cmp);
map<seg, int, decltype(cmp)> mp(cmp);
priority_queue<seg, vector<seg>, decltype(cmp)> pq(cmp); // max heap
// hash func overload
struct point{
bool operator == (const point &p)const{ return x == p.x && y == p.y: }
struct hasher {
size_t operator()(const point &p)const{ return p.x * 2 + p.y * 3; }
unordered_map<point, int, hasher> hsh;
```