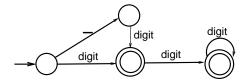
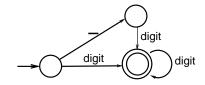
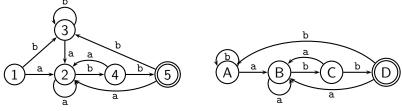
- The DFA constructed by the procedure mentioned earlier does not result in a minimum DFA.
- The DFA that we had seen earlier:



is not minimum. There is another DFA for the same regular expression with lesser number of states.

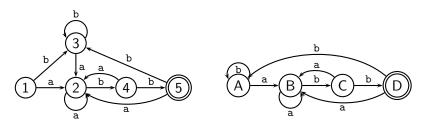


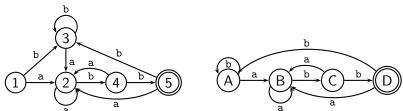
- The DFA constructed for $(b|\epsilon)(a|b)^*abb$.
- There is another DFA for the same regular expression with lesser number of states.



- For a typical language, the number of states of the DFA is in order of hundreds.
- Therefore we should try to minimize the number of states.

- The second DFA has been obtained by merging states 1 and 3 of the first DFA.
- Under what conditions can this merging take place?

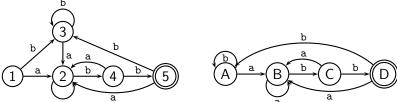




- The string bb takes both states 1 and 3 to a non-final state.
- The string aba takes both states 1 and 3 to a non-final state.
- The string ϵ takes both states 1 and 3 to a non-final state.
- The string bbabb takes both states 1 and 3 to a final state.

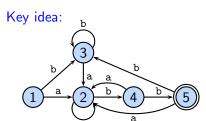
Observation:

Any string that takes state 1 to a final state also takes 3 to a final state. Conversely, any string that takes state 1 to a non-final state also takes 3 to a non-final state.

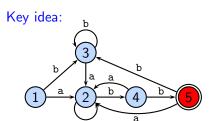


- States 1 and 3 are said to be indistinguishable.
- Minimimization strategy:
 - Find indistinguishable states.
 - Merge them.
- Question: How does one find indistingushable states?

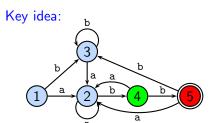
Key idea:



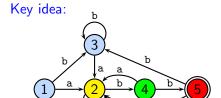
 Initially assume all states to be indistinguishable. Put them in a single set.



• The stri $\hat{\mathbf{n}}$ g ϵ distinguishes between final states and non-final states. Create two partitions.

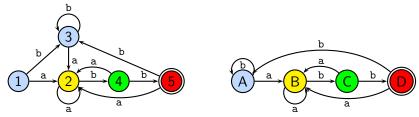


• b takes $\hat{4}$ to a red partition and retains other blue states in blue partition. Put 4 in a separate partition.



• The string b distinguishes 2 from other states in the blue partition.

Key idea:



• No other partition possible. Merge all states in the same partition.

Summary of the Method

- 1. Construct an initial partition $\pi = \{S F, F_1, \dots, F_n, \}$, where $F = F_1 \cup F_2 \cup \dots F_n s$, and each F_i is the set of final states for some token i.
- 2. for each set G in π do partition G into subsets such that two states s and t of G are in the same subset if and only if for all input symbols a, states s and t have transitions onto states in the same set of π ; replace G in π_{new} by the set of all subsets formed
- 3. If $\pi_{new} = \pi$, let $\pi_{final} := \pi$ and continue with step 4. Otherwise repeat step 2 with $\pi := \pi_{new}$.
- 4. Merge states in the same set of the partition.
- 5. Remove any dead states.