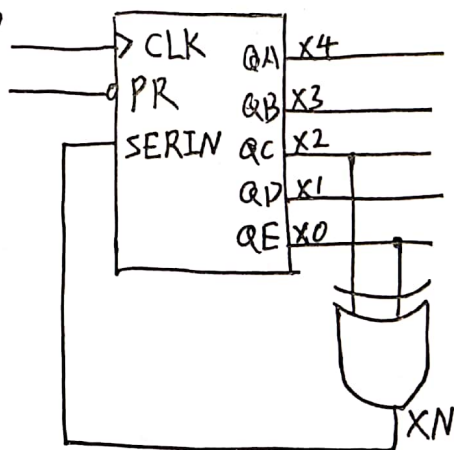


$$\begin{array}{cccccccc}
 15 & & 7 & & 6 & & 5 & & 4 & & 3 & & 2 \\
 1111 & \rightarrow & 0111 & \rightarrow & 0110 & \rightarrow & 0101 & \rightarrow & 0100 & \rightarrow & 0011 & \rightarrow & 0010 \\
 \uparrow & & & & & & & & & & & & \downarrow \\
 14 & 1110 & & & & & & & & & & & 0001 \\
 \uparrow & & & & & & & & & & & & \downarrow \\
 1101 & \leftarrow & 1100 & \leftarrow & 1011 & \leftarrow & 1010 & \leftarrow & 1001 & \leftarrow & 1000 & \leftarrow & 0000 \\
 13 & & 12 & & 11 & & 10 & & 9 & & 8 & & 0
 \end{array}$$

$0000 \rightarrow 0001 \rightarrow 0010 \rightarrow 0011 \rightarrow 0100 \rightarrow 0101$
 $\rightarrow 1000 \rightarrow 1001 \rightarrow 1010 \rightarrow 1011 \rightarrow 1100 \rightarrow 1101$
 $\rightarrow 1110 \rightarrow 1111 \rightarrow 0000$

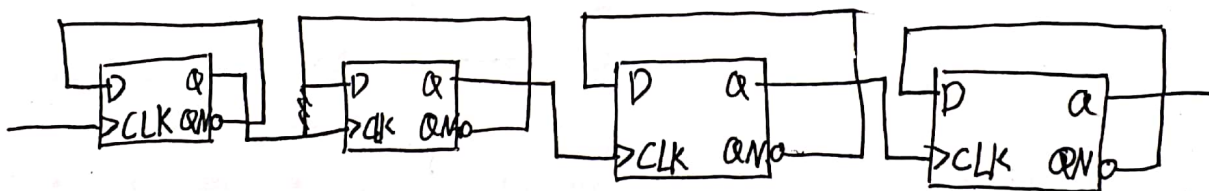


8.16



00001 → 00010 → 00100 → 01001 → 10010
 → 00101 → 01011 → 10110 → 01100 → 11001

8.28



$$t_{\text{最大延迟}} = 4 \times \overset{44}{17} = \overset{44}{68} \text{ ns} = 176 \text{ ns}$$

对于 74AHCT74, $t_{\text{最大延迟}} = 4 \times 10 = 40 \text{ ns}$

对于 74LS74, $t_{\text{最大延迟}} = 4 \times 40 = 160 \text{ ns}$

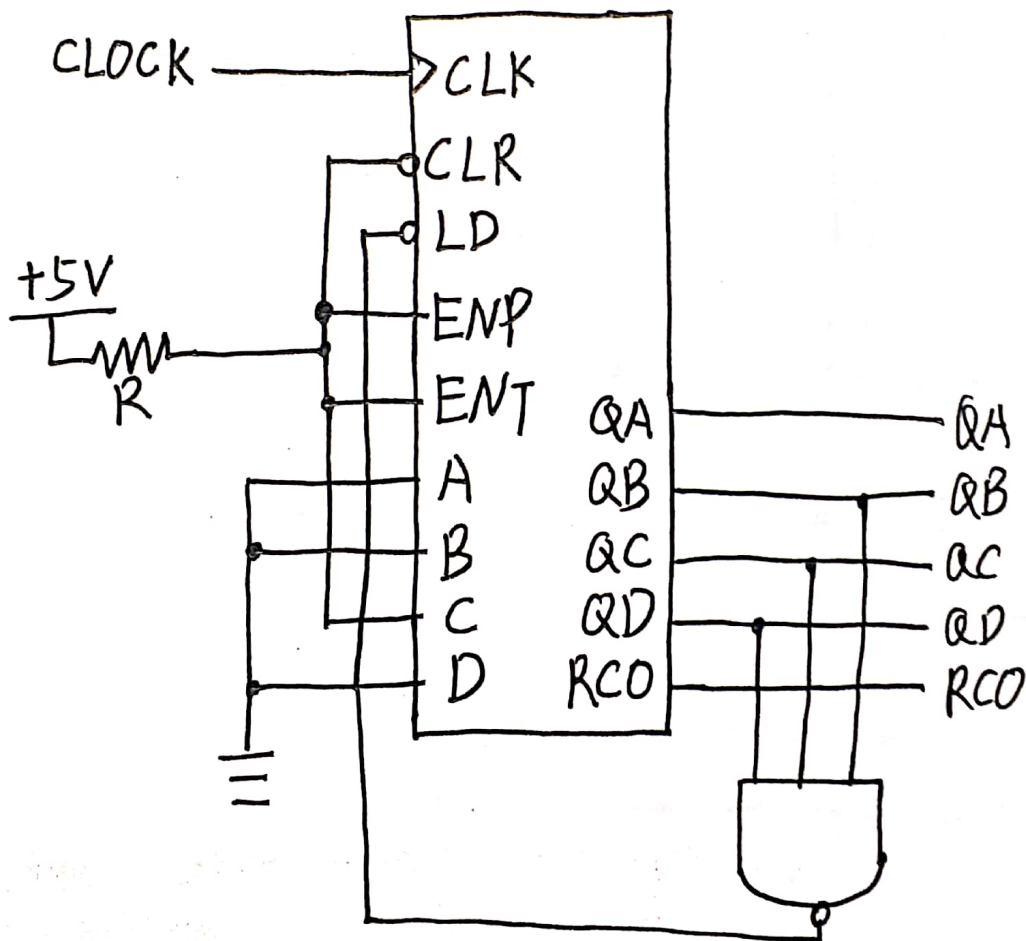
电路图均如上所示

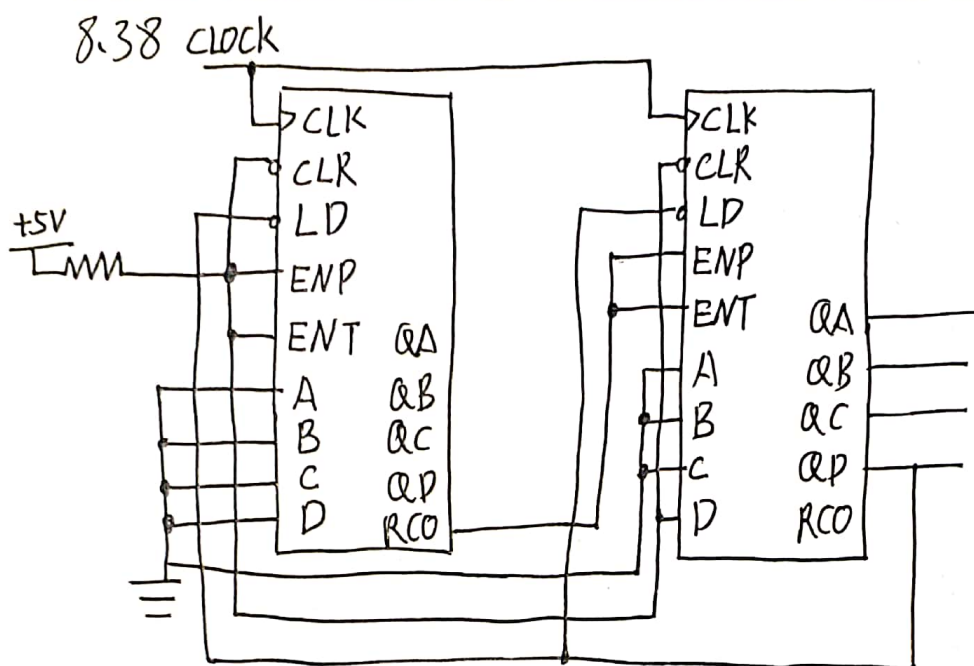
$$8.31. f_{\text{max}} = \frac{1}{T_{\text{min}}} = \frac{1}{3t_{\text{AND}} + 4(t_{\text{setup}} + t_{\text{TQ}})}$$



8.35 $14 = \cancel{1100_2} 1110_2$

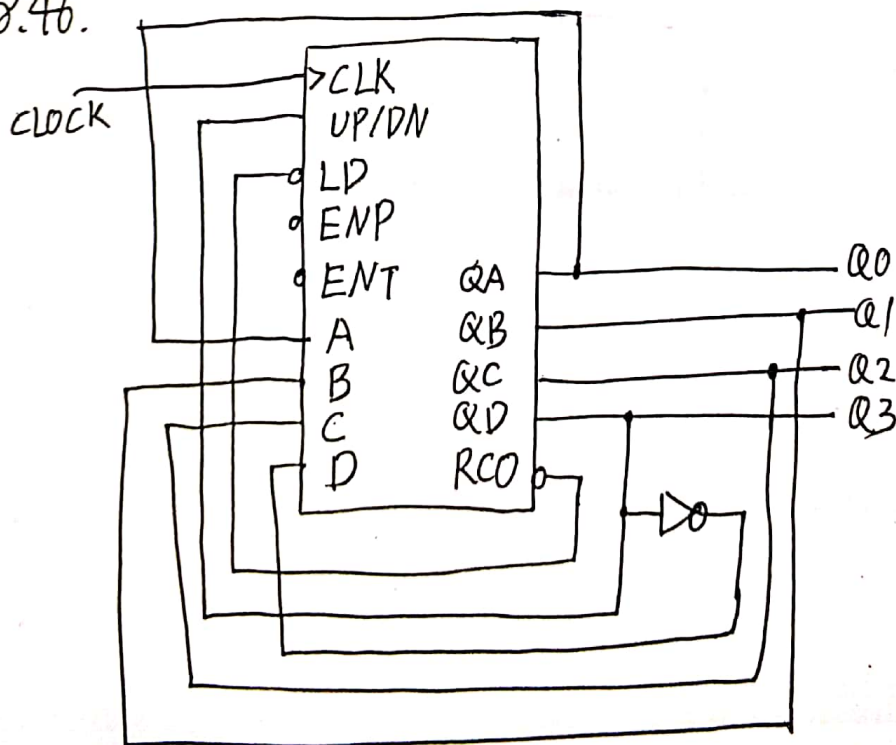
$$LD_L = \overline{Q_D \cdot Q_C \cdot Q_B}$$



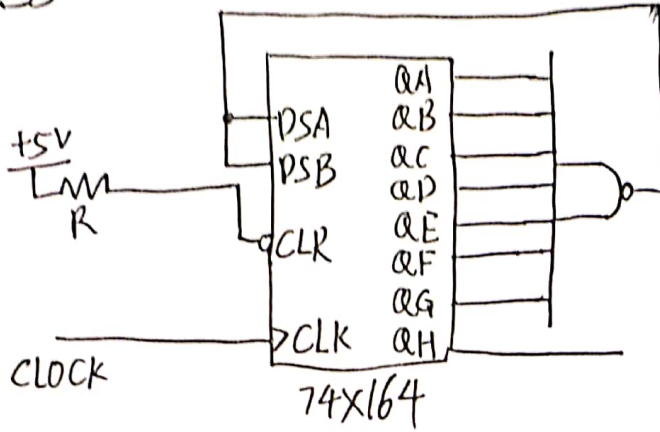


状态从 1000 0000 递增至 1111 1111, 接下来是 0000 0000.
此时 LD 信号有效, 下一状态为 1000 0000.

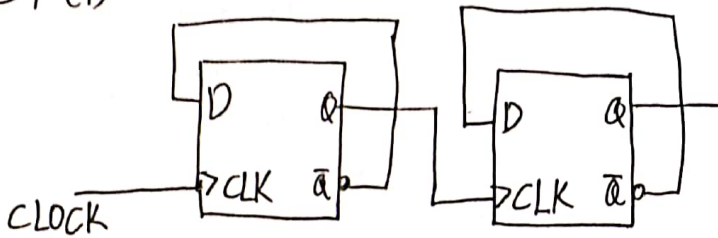
8.46.



8.55



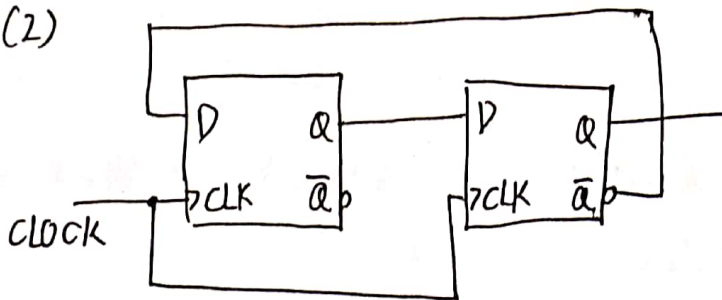
8.57 (1)



状态:

00 → 01 → 10 → 01
↑

(2)



状态:

00 → 01 → 11 → 10
↑



8.63 $x_4 = x_3 \oplus x_0$

	x_0	x_1	x_2	x_3	x_4
s_1	1	1	1	1	0
s_2	1	1	1	0	1
s_3	1	1	0	1	0
s_4	1	0	1	0	1
s_5	0	1	0	1	1
s_6	1	0	1	1	0
s_7	0	1	1	0	0
s_8	1	1	0	0	1
s_9	1	0	0	1	0
s_{10}	0	0	1	0	0
s_{11}	0	1	0	0	0
s_{12}	1	0	0	0	1
s_{13}	0	0	0	1	1
s_{14}	0	0	1	1	1
s_{15}	0	1	1	1	1

8.64 如图 8.5.2 的 n 位 LFSR 计数器, 反馈方程为 $x_n = x_0 \oplus x_1$

当 $x_0 x_1 \dots x_{n-1}$ 为 $1000 \dots 0$ 时, 下一状态为 $x_0 x_1 \dots x_{n-1} = 000 \dots 001$

加入 $n-1$ 输入或非门与一个异或门后:

当 $x_0 x_1 \dots x_{n-1} = 1000 \dots 0$ 时,

下一状态为 $x_0 x_1 \dots x_{n-1} = 00 \dots 0$

下一状态为 $x_0 x_1 \dots x_{n-1} = 00 \dots 01$

对于其他状态, $x_1 \dots x_{n-1}$ 中必存在 1

此时有 $x_0 \oplus x_1 = (x_0 \oplus x_1) \oplus 0 = (x_0 \oplus x_1) \oplus (x_1 + x_2 + \dots + x_{n-1})$

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所以只加入了一种状态(全零), 共 2^n 种状态

