

第三次作业.

14. (1) 对于任意一个模型 (M, σ)

$$\begin{aligned} \text{由于 } \forall a \in M, \quad (\forall x (x \doteq x))_{M[\sigma]} &= \forall a (x \doteq x)_{M[\sigma][x:=a]} \\ &= \forall a (a = a) \\ &= T \end{aligned}$$

 $\therefore \forall x (x \doteq x)$ 为永真式(2) 对于任意一个模型 (M, σ)

$$\begin{aligned} (\forall x \forall y (x \doteq y \rightarrow y \doteq x))_{M[\sigma]} &= \forall a (\forall y (x \doteq y \rightarrow y \doteq x)_{M[\sigma][x:=a]}) \\ &= \forall a (\forall y (a \doteq y \rightarrow y \doteq a)) \\ &= \forall a \forall b (a \doteq b \rightarrow b \doteq a)_{M[\sigma][y:=b]} \\ &= \forall a \forall b (a = b \rightarrow b = a) \end{aligned}$$

其中 $a, b \in M$ ① 若 $a = b$, 则 $a \doteq b$ 为 T, $b \doteq a$ 为 T, $\forall a \forall b (a \doteq b \rightarrow b \doteq a)$ 为 T② 若 $a \neq b$, 则 $a \doteq b$ 为 F, $b \doteq a$ 为 F, $\forall a \forall b (a \doteq b \rightarrow b \doteq a)$ 为 T \therefore 原式永真(3) 对于任意一个模型 (M, σ)

$$\begin{aligned} &\forall x \forall y \forall z ((x \doteq y) \wedge (y \doteq z) \rightarrow x \doteq z)_{M[\sigma]} \\ &= \forall a \forall b \forall c ((x \doteq y) \wedge (y \doteq z) \rightarrow x \doteq z)_{M[\sigma][x:=a][y:=b][z:=c]} \\ &= \forall a \forall b \forall c (B \rightarrow (B \wedge (a = b, b = c), a = c)) \end{aligned}$$

若 $B \wedge (a = b, b = c)$ 为 T, 为等号传递性知 $a = b = c$ \therefore 上式成立若 $B \wedge (a = b, b = c)$ 为 F, $B \rightarrow (F, a = c)$ 为 T \therefore 上式也成立 \therefore 原式永真

16. 证: $\models (\neg \forall x A) \leftrightarrow (\exists x \neg A)$

\Leftrightarrow 对任何模型 $M \models (\neg \forall x A) \leftrightarrow (\exists x \neg A)$

\Leftrightarrow 对任何模型 $M \models ((\neg \forall x A) \rightarrow (\exists x \neg A)) \wedge ((\exists x \neg A) \rightarrow (\neg \forall x A))$

$\Leftrightarrow B_{\wedge}(\neg \forall x A \rightarrow \exists x \neg A, \exists x \neg A \rightarrow \neg \forall x A) = T$

① $\neg \forall x A \rightarrow \exists x \neg A = B_{\rightarrow}(\neg \forall x A, \exists x \neg A)$

当 $\neg \forall x A$ 为 T 时, $\forall x A$ 为 F

\therefore 存在 $a \in M$, 使得 $A_{M[G[x:=a]]} = F$

即 $\exists x \neg A$ 为 T $\therefore B_{\rightarrow}(\neg \forall x A, \exists x \neg A) = T$

② $\exists x \neg A \rightarrow \neg \forall x A = B_{\rightarrow}(\exists x \neg A, \neg \forall x A)$

当 $\exists x \neg A$ 为 T 时, 存在 $a \in M$, 使得 $\neg A_{M[G[x:=a]]} = T$

$\therefore A_{M[G[x:=a]]} = F$

$\therefore \forall x A$ 为 F, $\neg \forall x A$ 为 T

$\therefore B_{\rightarrow}(\exists x \neg A, \neg \forall x A) = T$

$\therefore B_{\wedge}(B_{\rightarrow}(\neg \forall x A, \exists x \neg A), B_{\rightarrow}(\exists x \neg A, \neg \forall x A)) = T$

$\models (\neg \exists x A) \leftrightarrow (\forall x \neg A)$

$\Leftrightarrow \models (\neg \exists x A \rightarrow \forall x \neg A) \wedge (\forall x \neg A \rightarrow \neg \exists x A)$

\Leftrightarrow 对任何模型 (M, σ) , $M \models (\neg \exists x A \rightarrow \forall x \neg A) \wedge (\forall x \neg A \rightarrow \neg \exists x A)$

$\Leftrightarrow B_{\wedge}(B_{\rightarrow}(\neg \exists x A, \forall x \neg A), B_{\rightarrow}(\forall x \neg A, \neg \exists x A))_{M[G[x:=a]]} = T$ for all $a \in M$.

① $\neg \exists x A$ 为 T, $\exists x A$ 为 F

即任意 a , 使得 $A_{M[G[x:=a]]} = F$

$\therefore \forall x \neg A = T \therefore B_{\rightarrow}(\neg \exists x A, \forall x \neg A) = T$

② $\forall x \neg A$ 为 T, 即任意 $a \in M$, $\neg A_{M[G[x:=a]]} = T$, $A_{M[G[x:=a]]} = F$

$\therefore \exists x A$ 为 F, $\neg \exists x A$ 为 T

$\therefore B_{\rightarrow}(\forall x \neg A, \neg \exists x A) = T$

$\therefore B_{\wedge}(B_{\rightarrow}(\neg \exists x A, \forall x \neg A), B_{\rightarrow}(\forall x \neg A, \neg \exists x A)) = T$ for all $a \in M$.

18. (1) $\models \forall x A \leftrightarrow \forall y A[y/x]$

即对任何模型, 有 $M \models \forall x A \leftrightarrow \forall y A[y/x]$

$$(\forall x A)_{M[G]} = \begin{cases} T, & \text{对所有 } a \in M, A_{M[G][x:=a]} = T \\ F, & \text{否则} \end{cases}$$

$$(\forall y A[y/x])_{M[G]} = \begin{cases} T, & \text{对所有 } a \in M, A_{M[G][y:=a]} = T \\ F, & \text{否则} \end{cases}$$

$\because y$ 是新变元

$$\begin{aligned} \therefore A[y/x]_{M[G][y:=a]} &= A_{M[G][x:=y_{M[G]}, y:=a]} \\ &= A_{M[G][x:=a]} \end{aligned}$$

$$\therefore B(\forall x A, \forall y A[y/x]) = T$$

$$\therefore \models \forall x A \leftrightarrow \forall y A[y/x]$$

(2) $\models \exists x A \leftrightarrow \exists y A[y/x]$

即对任何模型, 有 $M \models \exists x A \leftrightarrow \exists y A[y/x]$

$$(\exists x A)_{M[G]} = \begin{cases} T, & \text{存在 } a \in M, A_{M[G][x:=a]} = T \\ F, & \text{否则} \end{cases}$$

$$(\exists y A[y/x])_{M[G]} = \begin{cases} T, & \text{存在 } a \in M, A_{M[G][y:=a]} = T \\ F, & \text{否则} \end{cases}$$

$\because y$ 是新变元

\therefore 由替换引理,

$$\begin{aligned} A[y/x]_{M[G][y:=a]} &= A_{M[G][x:=y_{M[G]}, y:=a]} \\ &= A_{M[G][x:=a]} \end{aligned}$$

$$\therefore \exists x A \leftrightarrow \exists y A[y/x] \text{ 对任意模型}$$

19. (1) $\forall x A(x) \leftrightarrow \forall t. A(t/x)$ 永真

即对任何模型, $M \models \forall x A \leftrightarrow \forall t A[t/x]$

即 $B \rightarrow (\forall x A, A[t/x])_{M[G]} = T$

即 $B \rightarrow (\forall x A, \forall t A[t/x])_{M[G]} = T$ 且

$B \rightarrow (A[t/x], \forall x A)_{M[G]} = T$

若 $\forall x A = T$, 则所有 $a \in M$, $A_{M[G]}[x:=a] = T$

由引理 3.11, $t_{M[G]} \in M$

$\therefore A_{M[G]}[x:=t_{M[G]}] = T$

由替换引理, $A[\frac{t}{x}]_{M[G]} = A_{M[G]}[x:=t_{M[G]}] = T$

$\therefore B \rightarrow (\forall x A, A[t/x])_{M[G]} = T$

若 $\forall x A[t/x] = T$

即 $A[\frac{t}{x}]_{M[G]} = A_{M[G]}[x:=t_{M[G]}] = T$, 对所有项 t

$\because t_{M[G]} \in M$, t 为任意一个项

$\therefore \forall a \in M, A_{M[G]}[x:=a] = T$

$\therefore B \rightarrow (\forall x A, A[t/x])_{M[G]} = T$

$\therefore \forall x A \leftrightarrow A[t/x]$ 永真.

(2) $\forall t. A[t/x] \rightarrow \exists x A$ 永真

即对任何模型, $M \models \forall t A[t/x] \rightarrow \exists x A$

即 $B \rightarrow (A[t/x], \exists x A)_{M[G]} = T$

当 $A[t/x]_{M[G]} = T$ 时,

即 $A[t/x]_{M[G]} = A_{M[G]}[x:=t_{M[G]}] = T$

由引理 3.11, $t_{M[G]} \in M$, 设 $t_{M[G]} = a$.

所以存在 $a \in M$, $A_{M[G]}[x:=a] = T$

即 $\exists x A = T$

$\therefore B \rightarrow (A[t/x], \exists x A) = T$

$\therefore A[t/x] \rightarrow \exists x A$ 永真.

解: 令 $t_1 = x_1, t_2 = x_2, \dots, t_n = x_n$.

知 $A[\frac{t}{x}]_{M[G]} = A_{M[G]}[x:=t_{M[G]}]$

$= A_{M[G]}[x:=x_1] = T$

\dots

$A[\frac{t_n}{x}]_{M[G]} = A_{M[G]}[x:=t_{M[G]}]$

$= A_{M[G]}[x:=x_n] = T$

20. (1) 若 $A_{M[cx=a]} = T$, $A_{M[cx=b]} = F$ ($a \neq b$)

则 $(\exists x A)_{M[c]} = T$ 因为 $\exists a \in M$, 使得 $A_{M[cx=a]} = T$

$(\forall x A)_{M[c]} = F$ 因为 $\exists b \in M$, 使得 $A_{M[cx=b]} = F$

(2) 若 $A_{M[cx=a]} = F$, $B_{M[cx=a]} = T$

$A_{M[cx=b]} = T$, $B_{M[cx=b]} = F$ ($a \neq b$)

对其他 $c \in M$ 且 $c \neq a, c \neq b$

$A_{M[cx=c]} = T$, $B_{M[cx=c]} = T$

则 $\forall x (A \vee B)_{M[c]} = T$

$(\forall x A)_{M[c]} = F$ 因为 $\exists a \in M$, $A_{M[cx=a]} = F$

$(\forall x B)_{M[c]} = F$ 因为 $\exists b \in M$, $B_{M[cx=b]} = F$

$\therefore \forall x (A \vee B) \rightarrow ((\forall x A) \vee (\forall x B))$ 不是永真

21. (1) 设 A 的一个模型 (M, σ)

$M = N$, $\sigma(x_n) = n$

$(x, y) \in R(x, y)$ 当且仅当 $x < y$

则对所有 $a \in M$, $(a, a) \notin R(x, y)$

$\therefore \forall x \neg R(x, x) = T$

$\forall a, b, c \in M$, 若 $(a, b) \in R(x, y)$, $(b, c) \in R(x, y)$

即 $a < b$ 且 $b < c$, 则 $a < c \Rightarrow (a, c) \in R(x, y)$

$\therefore \forall x \forall y \forall z ((R(x, y) \wedge R(y, z)) \rightarrow R(x, z)) = T$

$\forall a \in M$, $\exists a+1 \in M$

使得 $(a, a+1) \in R(x, y)$

$\therefore \forall x \exists y R(x, y) = T$

$\therefore (M, \sigma)$ 为 A 的一个无穷模型

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(2) 假设 A 有一个无穷模型 (M, σ)

~~若~~ $\sigma(x_n) = a_i \quad (i=0, 1, \dots, k)$

$$M = \{a_i \mid 0 \leq i \leq k\}$$

若 $(a_i, a_j) \in R(x, y)$

记为 $a_i < a_j \quad (i < j)$

$$\therefore \forall x \exists y R(x, y) = T$$

\therefore ~~$\exists a_k$~~ 取 $x_{n[k]} = a_k$

则 $\exists a_j R(a_k, a_j) = T, a_j \in M$

则 $j > k \quad \text{F}$

与 $0 \leq j \leq k$ 矛盾

$\therefore A$ 没有有穷模型

(1) 证: case 1: t 为某个变元 x_i ($1 \leq i$)

$$\because \forall i \leq n, s_1(x_i) = s_2(x_i)$$

$$\therefore t_{M[s_1]} = x_{i[M[s_1]]} = s_1(x_i) = s_2(x_i) = t_{M[s_2]}$$

case 2: t 为某个常元 c

$$\text{则 } t_{M[s_1]} = c_{M[s_1]} = c = c_{M[s_2]} = t_{M[s_2]}$$

case 3: t 为 $f(t_1, t_2, \dots, t_n)$ 其中 t_1, t_2, \dots, t_n 均为项

$$\begin{aligned} \text{则 } t_{M[s_1]} &= f(t_1, t_2, \dots, t_n)_{M[s_1]} = f_M(t_{1M[s_1]}, \dots, t_{nM[s_1]}) \\ &= f_M(t_{1M[s_2]}, \dots, t_{nM[s_2]}) \\ &= f(t_1, t_2, \dots, t_n)_{M[s_2]} \\ &= t_{M[s_2]} \end{aligned}$$

Q.E.D.

(2) case 1: 若 A 呈形 $s \doteq t$ (s 和 t 为项)

$$\text{则 } A_{M[s_1]} = (s \doteq t)_{M[s_1]} = \begin{cases} T, & \text{if } s_{M[s_1]} = t_{M[s_1]} \\ F, & \text{else} \end{cases}$$

$$A_{M[s_2]} = (s \doteq t)_{M[s_2]} = \begin{cases} T, & \text{if } s_{M[s_2]} = t_{M[s_2]} \\ F, & \text{else} \end{cases}$$

由 (1) 知, 对任意项 t , $t_{M[s_1]} = t_{M[s_2]}$

$$\therefore A_{M[s_1]} = A_{M[s_2]}$$

case 2: A 呈形 $R(t_1, t_2, \dots, t_n)$, R 为 n 元谓词, t_1, \dots, t_n 为项.

$$\text{则 } A_{M[s_1]} = R(t_1, t_2, \dots, t_n)_{M[s_1]} = \begin{cases} T, & \text{若 } \langle t_{1M[s_1]}, \dots, t_{nM[s_1]} \rangle \in P_M \\ F, & \text{若 } \langle t_{1M[s_1]}, \dots, t_{nM[s_1]} \rangle \notin P_M \end{cases}$$

$$A_{M[s_2]} = R(t_1, t_2, \dots, t_n)_{M[s_2]} = \begin{cases} T, & \text{若 } \langle t_{1M[s_2]}, \dots, t_{nM[s_2]} \rangle \in P_M \\ F, & \text{若 } \langle t_{1M[s_2]}, \dots, t_{nM[s_2]} \rangle \notin P_M \end{cases}$$

由 (1), 对任意项 t , $t_{M[s_1]} = t_{M[s_2]}$

$$\therefore A_{M[s_1]} = A_{M[s_2]}$$

case 3: A 是形 $\neg B$, 其中 B 为公式

$$\text{则 } A_{MIS_1} = (\neg B)_{MIS_1} = \neg(B_{MIS_1})$$

$$A_{MIS_2} = (\neg B)_{MIS_2} = \neg(B_{MIS_2})$$

$$\text{由 case 1 和 case 2 知 } B_{MIS_1} = B_{MIS_2}$$

$$\therefore A_{MIS_1} = A_{MIS_2}$$

case 4: A 是形 $B * C$, $*$ $\in \{\wedge, \vee, \rightarrow\}$

$$\text{则 } A_{MIS_1} = B * (C_{MIS_1})$$

$$A_{MIS_2} = B * (C_{MIS_2})$$

$$\because B_{MIS_1} = B_{MIS_2}, C_{MIS_1} = C_{MIS_2}$$

$$\therefore A_{MIS_1} = A_{MIS_2}$$

case 5: A 是形 $\forall x. B$ (x 为变元)

$$A_{MIS_1} = (\forall x. B)_{MIS_1} = \begin{cases} T, & \text{若对所有 } a \in M, B_{MIS_1}[x=a] = T \\ F, & \text{否则} \end{cases}$$

$$A_{MIS_2} = (\forall x. B)_{MIS_2} = \begin{cases} T, & \text{若对所有 } a \in M, B_{MIS_2}[x=a] = T \\ F, & \text{否则} \end{cases}$$

$$\because B_{MIS_1} = B_{MIS_2}$$

$$\therefore A_{MIS_1} = A_{MIS_2}$$

case 6: A 是形 $\exists x. B$

$$A_{MIS_1} = (\exists x. B)_{MIS_1} = \begin{cases} T, & \text{若对某个 } a \in M, B_{MIS_1}[x=a] = T \\ F, & \text{否则} \end{cases}$$

$$A_{MIS_2} = (\exists x. B)_{MIS_2} = \begin{cases} T, & \text{若对某个 } a \in M, B_{MIS_2}[x=a] = T \\ F, & \text{否则} \end{cases}$$

$$\because B_{MIS_1} = B_{MIS_2}$$

$$\therefore A_{MIS_1} = A_{MIS_2}$$

Q.E.D.

$$23. \text{ 证: } M \models (A \wedge B) \Leftrightarrow (A \wedge B)_{M[G]}$$

$$\Leftrightarrow B_{\wedge}(A_{M[G]}, B_{M[G]})$$

$$= \begin{cases} T, & A_{M[G]} = T \text{ and } B_{M[G]} = T \\ F, & \text{否则} \end{cases}$$

$$M \models A \text{ and } M \models B = \begin{cases} T, & A_{M[G]} = T \text{ and } B_{M[G]} = T \\ F, & \text{否则} \end{cases}$$

$$\therefore M \models (A \wedge B) \Leftrightarrow (M \models A \text{ and } M \models B)$$

$$24. \text{ 证: } M \models \forall z. A[\frac{z}{x}] \Leftrightarrow (\forall z. A[\frac{z}{x}])_{M[G]} = T$$

$$\Leftrightarrow \text{对任意 } a \in M, (A[\frac{z}{x}])_{M[G][z:=a]} = T$$

$$\Leftrightarrow A_{M[G][x:=z][z:=a]} = A_{M[G][x:=a]} = T$$

(替换引理)

$$\Leftrightarrow M \models \forall x. A$$