

## 第五次作业.

1. 证: (1) 对于任意的集合  $\Delta$ , 若  $\Delta$  为  $\Sigma \cap \bar{\Sigma}$  的子集

则  $\Delta \subseteq \Sigma$  且  $\Delta \subseteq \bar{\Sigma}$ ,

$\therefore \text{Con}(\Sigma), \text{Con}\bar{\Sigma}$ .

$\therefore \Delta \vdash$  不可证.

$\therefore$  由定义知  $\text{Con}(\Sigma \cap \bar{\Sigma})$

(2) 令  $\Sigma = \{A, B\}$

$\bar{\Sigma} = \{\neg A, C\}$

则  $\Sigma \cup \bar{\Sigma} = \{A, \neg A, B, C\}$

对于公式  $A$ , 存在  $\Sigma \cup \bar{\Sigma}$  的有穷子集  $\Delta = \{A, \neg A\}$

使得  $\Delta \vdash A$  和  $\Delta \vdash \neg A$  可证

由命题 6.2 知  $\text{Incon}(\Sigma \cup \bar{\Sigma})$

3. 证: 对 $\mathcal{L}$ 的任何协调公式集 $\Sigma$ ,

设 $\Sigma$ 的全体公式为 $\varphi_0, \varphi_1, \dots, \varphi_n, \dots (n \in \mathbb{N})$ ,

$$\text{令 } \begin{cases} \Sigma_0 = \Sigma \\ \Sigma_{n+1} = \begin{cases} \Sigma_n, & \text{若 } \text{Incon}(\Sigma_n \cup \{\varphi_n\}) \\ \Sigma_n \cup \{\varphi_n\}, & \text{若 } \text{Con}(\Sigma_n \cup \{\varphi_n\}) \end{cases} \end{cases}$$

$$\text{令 } \Sigma = \bigcup \{\Sigma_n \mid n \in \mathbb{N}\}$$

现证 $\Sigma$ 是极大协调的.

反设  $\text{Incon}(\Sigma)$ , 从而存在 $\Sigma$ 的有穷子集 $\Delta$ 使 $\Delta \vdash \perp$ 可证.

$\because \Delta$ 有穷, 不妨设  $\Delta = \{A_1, \dots, A_k\}$

$\therefore A_i (i=1, 2, \dots, k) \in \Sigma = \bigcup \{\Sigma_n \mid n \in \mathbb{N}\}$

故对每个 $i \leq k$ , 有 $n_i$ 使  $A_i \in \Sigma_{n_i}$

故有 $l$ 对每个 $i \leq k$ ,  $A_i \in \Sigma_l$

$\therefore \Delta \subseteq \Sigma_l$ , 这与  $\text{Con}(\Sigma_l)$  矛盾.



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证明极大性.

只需证若  $\text{Con}(\bar{\Psi}_n \cup \{\varphi_n\})$ , 则  $\varphi_n \in \bar{\Psi}$

设  $\text{Con}(\bar{\Psi}_n \cup \{\varphi_n\})$ , 从而  $\text{Con}(\bar{\Psi}_n \cup \{\varphi_n\})$

$\therefore \varphi_n \in \bar{\Psi}_{n+1} \quad \therefore \varphi_n \in \bar{\Psi}.$

4. (2) 证:

$$\frac{s \doteq t, p(s) \vdash p(t)}{\vdash (s \doteq t) \rightarrow (p(s) \rightarrow p(t))} \rightarrow R$$

定义  $P, m \in P \Leftrightarrow m \doteq s$

则有  $s \doteq t, s \doteq s \vdash t \doteq s$

$$\begin{array}{c} \therefore \quad s \doteq t, s \doteq s \vdash t \doteq s \quad \vdash s \doteq s \\ \hline s \doteq t \vdash t \doteq s \quad \text{Cut} \\ \vdash s \doteq t \rightarrow t \doteq s \rightarrow R \end{array}$$

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$$(3) \quad \frac{t \doteq s, t \doteq u \vdash s \doteq u \quad s \doteq t \vdash t \doteq s}{\text{Cot}}$$

$$\frac{s \doteq t, t \doteq u \vdash s \doteq u}{\rightarrow R}$$

$$\frac{s \doteq t \vdash (t \doteq u \rightarrow s \doteq u)}{\vdash s \doteq t \rightarrow (t \doteq u \rightarrow s \doteq u)} \rightarrow R$$



7. (1)  $M = \{a, b, c\}$

$\delta: \delta(P) = \{a\}$

易知  $P(a) \rightarrow \forall y P(y)$  为假  $\therefore$  有反例模型  $\langle M, \delta \rangle$ .

(2)

$$\begin{array}{l}
 \frac{P(a) \vdash P(a), \exists x P(x), Q(a)}{P(a) \vdash \exists x P(x), Q(a)} \exists R \quad \frac{P(a), Q(y), \forall y Q(y) \vdash Q(a)}{P(a), \forall y Q(y) \vdash Q(a)} \forall L \\
 \frac{P(a) \vdash \exists x P(x), Q(a) \quad P(a), \forall y Q(y) \vdash Q(a)}{\exists x P(x) \rightarrow \forall y Q(y), P(a) \vdash Q(a)} \rightarrow L \\
 \frac{\exists x P(x) \rightarrow \forall y Q(y), P(a) \vdash Q(a)}{\exists x P(x) \rightarrow \forall y Q(y) \vdash P(a) \rightarrow Q(a)} \rightarrow R \\
 \frac{\exists x P(x) \rightarrow \forall y Q(y) \vdash P(a) \rightarrow Q(a)}{\exists x P(x) \rightarrow \forall y Q(y) \vdash \forall z (P(z) \rightarrow Q(z))} \forall R \\
 \frac{\exists x P(x) \rightarrow \forall y Q(y) \vdash \forall z (P(z) \rightarrow Q(z))}{\vdash (\exists x P(x) \rightarrow \forall y Q(y)) \rightarrow \forall z (P(z) \rightarrow Q(z))} \rightarrow R
 \end{array}$$

(3)  $M: \{a, b\}$

$\delta(P) = \{a\}$

$\delta(Q) = \{a\}$

则  $\forall z (P(z) \rightarrow Q(z)) = T$

$\exists x P(x) \rightarrow \forall y Q(y) = F$

$\therefore \forall z (P(z) \rightarrow Q(z)) \rightarrow (\exists x P(x) \rightarrow \forall y Q(y)) = F$

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(4)

$$\begin{array}{l}
 \text{Axiom} \\
 \frac{R(z,z), \exists y \forall x (\dots) \vdash R(z,z) \neg R}{\exists y \forall x (\dots) \vdash \neg R(z,z), R(z,z)} \\
 \text{Axiom} \\
 \frac{\exists y \forall x (\dots), \neg R(z,z) \vdash \neg R(z,z)}{R(z,z) \rightarrow \neg R(z,z), \exists y \forall x (R(x,y) \leftrightarrow \neg R(x,x)) \vdash \neg R(z,z)} \rightarrow I \\
 \text{Axiom} \\
 \frac{\exists y \forall x (\dots), R(z,z) \vdash R(z,z)}{R(z,z) \rightarrow R(z,z), \exists y \forall x (\dots), R(z,z) \vdash} \rightarrow I \\
 \text{Axiom} \\
 \frac{R(z,z), \exists y \forall x (\dots) \vdash R(z,z) \neg R}{\exists y \forall x (\dots), R(z,z), \neg R(z,z) \vdash} \neg I \\
 \frac{R(z,z) \rightarrow \neg R(z,z), \neg R(z,z) \rightarrow R(z,z), \exists y \forall x (R(x,y) \leftrightarrow \neg R(x,x)) \vdash}{R(z,z) \rightarrow \neg R(z,z) \wedge \neg R(z,z) \rightarrow R(z,z), \exists y \forall x (R(x,y) \leftrightarrow \neg R(x,x)) \vdash} \wedge I \\
 \frac{R(z,z) \leftrightarrow \neg R(z,z), \exists y \forall x (R(x,y) \leftrightarrow \neg R(x,x)) \vdash}{\forall x (R(x,z) \leftrightarrow \neg R(x,x)), \exists y \forall x (R(x,y) \leftrightarrow \neg R(x,x)) \vdash} \forall I \\
 \frac{\forall x (R(x,z) \leftrightarrow \neg R(x,x)), \exists y \forall x (R(x,y) \leftrightarrow \neg R(x,x)) \vdash}{\exists y \forall x (R(x,y) \leftrightarrow \neg R(x,x)) \vdash} \exists I \\
 \frac{\exists y \forall x (R(x,y) \leftrightarrow \neg R(x,x)) \vdash}{\vdash \neg \exists y \forall x (R(x,y) \leftrightarrow \neg R(x,x))} \neg R
 \end{array}$$