## 练习 0: 填写已有实验

本实验依赖实验 2/3/4/5/6/7。请把你做的实验 2/3/4/5/6/7 的代码填入本实验中代码中有 "LAB2"/"LAB3"/"LAB4"/"LAB5"/"LAB6" /"LAB7"的注释相应部分。并确保编译通过。注意:为了能够正确执行 lab8 的测试应用程序,可能需对已完成的实验 2/3/4/5/6/7 的代码进行进一步改进。

## 练习 1: 完成读文件操作的实现(需要编码)

首先了解打开文件的处理流程,然后参考本实验后续的文件读写操作的过程分析,填写在 kern/fs/sfs/sfs\_inode.c 中 的 sfs\_io\_nolock()函数,实现读文件中数据的代码。

- (1) 如果偏移量不与第一个块对齐,从偏移量到第一个块的末尾读取/写入一些内容 读/写大小 = (nblks != 0) ? (SFS\_BLKSIZE blkoff) : (endpos offset)
- (2) 读/写对齐的块
- (3) 如果结束位置不与最后一个块对齐,从开头到最后一个块的(endpos % SFS\_BLKSIZE) 处读取/写入一些内容

```
//LAB8:EXERCISE1 YOUR CODE HINT: call sfs bmap load nolock, sfs rbuf, sfs
  * (1) If offset isn't aligned with the first block, Rd/Wr some content
          NOTICE: useful function: sfs_bmap_load_nolock, sfs_buf_op
                  Rd/Wr size = (nblks != 0) ? (SFS_BLKSIZE - blkoff) : (
   * (2) Rd/Wr aligned blocks
          NOTICE: useful function: sfs bmap load nolock, sfs block op
   * (3) If end position isn't aligned with the last block, Rd/Wr some cor
          NOTICE: useful function: sfs bmap load nolock, sfs buf op
  blkoff=offset%SFS_BLKSIZE;
  if (blkoff != 0) {
     // if(nblks != 0)size=SFS_BLKSIZE-blkoff;
     // else size=endpos-offset;
     size = (nblks != 0) ? (SFS_BLKSIZE - blkoff) : (endpos - offset);
     ret = sfs_bmap_load_nolock(sfs, sin, blkno, &ino);
      if (ret != 0) {
         goto out;
      ret = sfs_buf_op(sfs, buf, size, ino, blkoff);
      if (ret != 0) {
          goto out;
     alen += size;
     buf += size;
     if(nblks == 0)goto out;
     blkno++;
     nblks--;
```

```
if (nblks>0) {
        ret = sfs bmap load nolock(sfs, sin, blkno, &ino);
        if (ret < 0) {
            goto out;
       ret = sfs_block_op(sfs, buf, blkno, nblks);
       if (ret < 0) {
            goto out;
        alen += nblks * SFS BLKSIZE;
        buf += nblks * SFS_BLKSIZE;
        blkno += nblks;
       nblks = 0;
    size = endpos % SFS BLKSIZE;
    if (endpos % SFS_BLKSIZE!=0) {
        ret = sfs_bmap_load_nolock(sfs, sin, blkno, &ino);
        if (ret != 0) {
            goto out;
        }
       ret = sfs_buf_op(sfs, buf, size, ino, 0);
        if (ret != 0) {
            goto out;
        }
       alen += size;
       // }
out:
    *alenp = alen;
    if (offcot , alon & cin \din \cizo) [
```

## 练习 2: 完成基于文件系统的执行程序机制的实现(需要编码)

改写 proc.c 中的 load\_icode 函数和其他相关函数,实现基于文件系统的执行程序机制。执行: make qemu。如果能看看到 sh 用户程序的执行界面,则基本成功了。如果在 sh 用户界面上可以执行"ls","hello"等其他放置在 sfs 文件系统中的其他执行程序,则可以认为本实验基本成功。

proc.c/alloc\_proc

```
//LAB8 YOUR CODE : (update LAB6 steps)
 /*
* below fields(add in LAB6) in proc struct need to be initialized
        struct files_struct * filesp;
                                                      file struct point
*/
   proc->state = PROC_UNINIT;
proc->pid = -1;
proc->runs = 0;
proc->kstack = 0;
proc->need resched = 0;
proc->parent = NULL;
proc->mm = NULL;
//初始化context结构体
memset(&(proc->context), 0, sizeof(struct context));
proc->tf = NULL;
proc->cr3 = boot_cr3;
proc->flags = 0;
proc->wait state = 0;
proc->cptr = NULL;
proc->yptr = NULL;
proc->optr = NULL;
proc->rq=NULL;
list init(&(proc->run link));
list init(&(proc->hash link));
//memset(&(proc->run_link), 0, sizeof(struct context));
proc->time_slice=0;
proc->filesp=NULL;
memset(proc->name, 0, PROC_NAME_LEN+1);
return proc;
```

proc.c/proc\_run

```
proc_run(struct proc_struct *proc) {
    if (proc != current) {
        // LAB4: EXERCISE3 YOUR CODE
        * Some Useful MACROs, Functions and DEFINEs, you can use them in below implementation.
        * MACROs or Functions:
       * local_intr_save():
* local_intr_restore():
                                Disable interrupts
                               Enable Interrupts
       * lcr3():
                                Modify the value of CR3 register
          switch_to():
                                Context switching between two processes
    //检查要切换的进程是否与当前正在运行的进程相同,如果相同则不需要切换
       if (proc->pid == current->pid)
        | return;
| return;
//禁 用 中 断。 你 可 以 使 用/kern/sync/sync.h 中 定 义 好 的 宏 local_intr_save(x) 和local_intr_restore(x) 来实现关、开中断。
       bool intrstate;
        struct proc_struct *currentpointer = current, *procpointer = proc;
       local_intr_save(intrstate);
//切换当前进程为要运行的进程。
       current=proc;
        //切换页表,以便使用新进程的地址空间。/libs/<u>riscy</u>.h 中提供了 lcr3(unsigned int cr3)函数,可实现修改 CR3 寄存器值的功能。
        lcr3(procpointer->cr3);
        flush_tlb();
        switch_to(&(currentpointer->context),&(procpointer->context));
      local_intr_restore(intrstate);
//LABB YOUR CODE : (update LAB4 steps)
     * below fields(add in LAB6) in proc_struct need to be initialized
           before switch_to();you should flush the tlb
            MACROs or Functions:
            flush_tlb():
                              flush the tlb
proc.c/load icode
   if (current->mm != NULL) {
       panic("load_icode: current->mm must be empty.\n");
   int ret = -E NO MEM;
   struct mm_struct *mm;
   //(1) create a new mm for current process
   if ((mm = mm_create()) == NULL) {
       goto bad mm;
   //(2) create a new PDT, and mm->pgdir= kernel virtual addr of PDT
   if (setup pgdir(mm) != 0) {
        goto bad_pgdir_cleanup_mm;
   //(3) copy TEXT/DATA/BSS parts in binary to memory space of process
   struct Page *page;
   struct elfhdr __elf, *elf = &__elf;
   struct proghdr __ph, * ph = &__ph;
   //(3.1) read raw data content in file and resolve elfhdr
   load icode read(fd, (void *)elf, sizeof(struct elfhdr), 0);
   // //(3.2) read raw data content in file and resolve proghdr based on info in elfhdr
   //load icode read(fd, (void *)ph, sizeof(struct proghdr), elf->e phoff );
   //(3.3) This program is valid?
   if (elf->e_magic != ELF_MAGIC) {
       ret = -E INVAL ELF;
        goto bad_elf_cleanup_pgdir;
```

void

```
uint32_t vm_flags, perm;
struct proghdr *ph_end = ph + elf->e_phnum;
// for (; ph < ph_end; ph ++) {
for(int index=0; index<elf->e phnum; index++)
    //(3.4) find every program section headers
   off_t ph_off = elf->e_phoff + sizeof(struct proghdr) * index;
   load_icode_read(fd, (void*)ph, sizeof(struct proghdr), ph_off);
   if (ph->p_type != ELF_PT_LOAD) {
        continue;
   if (ph->p filesz > ph->p memsz) {
       ret = -E_INVAL_ELF;
       goto bad cleanup mmap;
    if (ph->p filesz == 0) {
       // continue;
//(3.5) call mm_map fun to setup the new vma ( ph->p_va, ph->p_memsz)
    vm_flags = 0, perm = PTE_U | PTE_V;
    if (ph->p_flags & ELF_PF_X) vm_flags |= VM_EXEC;
   if (ph->p_flags & ELF_PF_W) vm_flags |= VM_WRITE;
   if (ph->p flags & ELF PF R) vm flags |= VM READ;
    // modify the perm bits here for RISC-V
   if (vm_flags & VM_READ) perm |= PTE_R;
   if (vm_flags & VM_WRITE) perm |= (PTE_W | PTE_R);
   if (vm_flags & VM_EXEC) perm |= PTE_X;
   if ((ret = mm_map(mm, ph->p_va, ph->p_memsz, vm_flags, NULL)) != 0) {
        goto bad_cleanup_mmap;
   size_t from = ph->p_offset;
    size t off, size;
    uintptr t start = ph->p va, end, la = ROUNDDOWN(start, PGSIZE);
    ret = -E NO MEM;
```

```
//(3.6.1) copy TEXT/DATA section of bianry program
  while (start < end) {
      if ((page = pgdir_alloc_page(mm->pgdir, la, perm)) == NULL) {
          goto bad_cleanup_mmap;
      off = start - la, size = PGSIZE - off, la += PGSIZE;
      if (end < la) {
          size -= la - end;
      load_icode_read(fd, page2kva(page) + off, size, from);
      start += size, from += size;
//(3.6.2) build BSS section of binary program
  end = ph->p va + ph->p memsz;
  if (start < la) {
      /* ph->p_memsz == ph->p_filesz */
      if (start == end) {
          continue;
      off = start + PGSIZE - la, size = PGSIZE - off;
      if (end < la) {
          size -= la - end;
      memset(page2kva(page) + off, 0, size);
      start += size:
      assert((end < la && start == end) || (end >= la && start == la));
  while (start < end) {
      if ((page = pgdir alloc page(mm->pgdir, la, perm)) == NULL) {
          goto bad cleanup mmap;
      off = start - la, size = PGSIZE - off, la += PGSIZE;
       if (end < la) {
          size -= la - end;
```

为当前进程创建一个新的内存管理结构(mm); 创建一个新的页目录表(PDT), 并将mm->pgdir 设置为页目录表的内核虚拟地址,将二进制文件中的 TEXT/DATA/BSS 部分复制到进程的内存空间中;读取文件中的原始数据内容并解析 ELF 头部(elfhdr), 根据 ELF 头部信息,读取文件中的原始数据内容并解析程序头部(proghdr),调用 mm\_map 构建与TEXT/DATA 相关的虚拟内存区域(VMA),调用 pgdir\_alloc\_page 为 TEXT/DATA 分配页面,读取文件内容并复制到新分配的页面中,调用 pgdir\_alloc\_page 为 BSS 分配页面,在这些页面中填充零;调用 mm\_map 设置用户栈,并将参数放入用户栈; 设置当前进程的内存管理结构、cr3 寄存器,重置页目录表(使用 1cr3 宏); 在用户栈中设置 uargc 和 uargv; 为用户环境设置陷阱帧(trapframe);如果前面的步骤失败,应清理环境\*/