



清华大学

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Modeling and Analyzing Latency in the Memcached System

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Memcached System

- Key-Value system
- In-memory caching solution
- Speed up searching applications



facebook



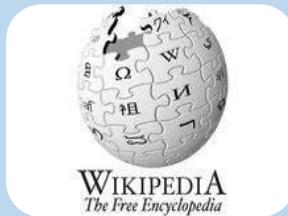
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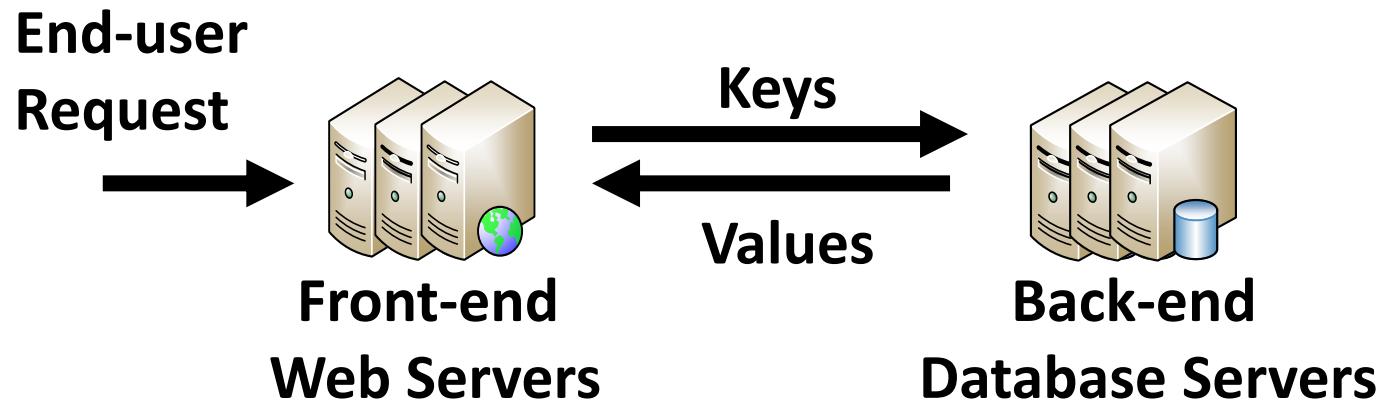


LIVEJOURNAL™

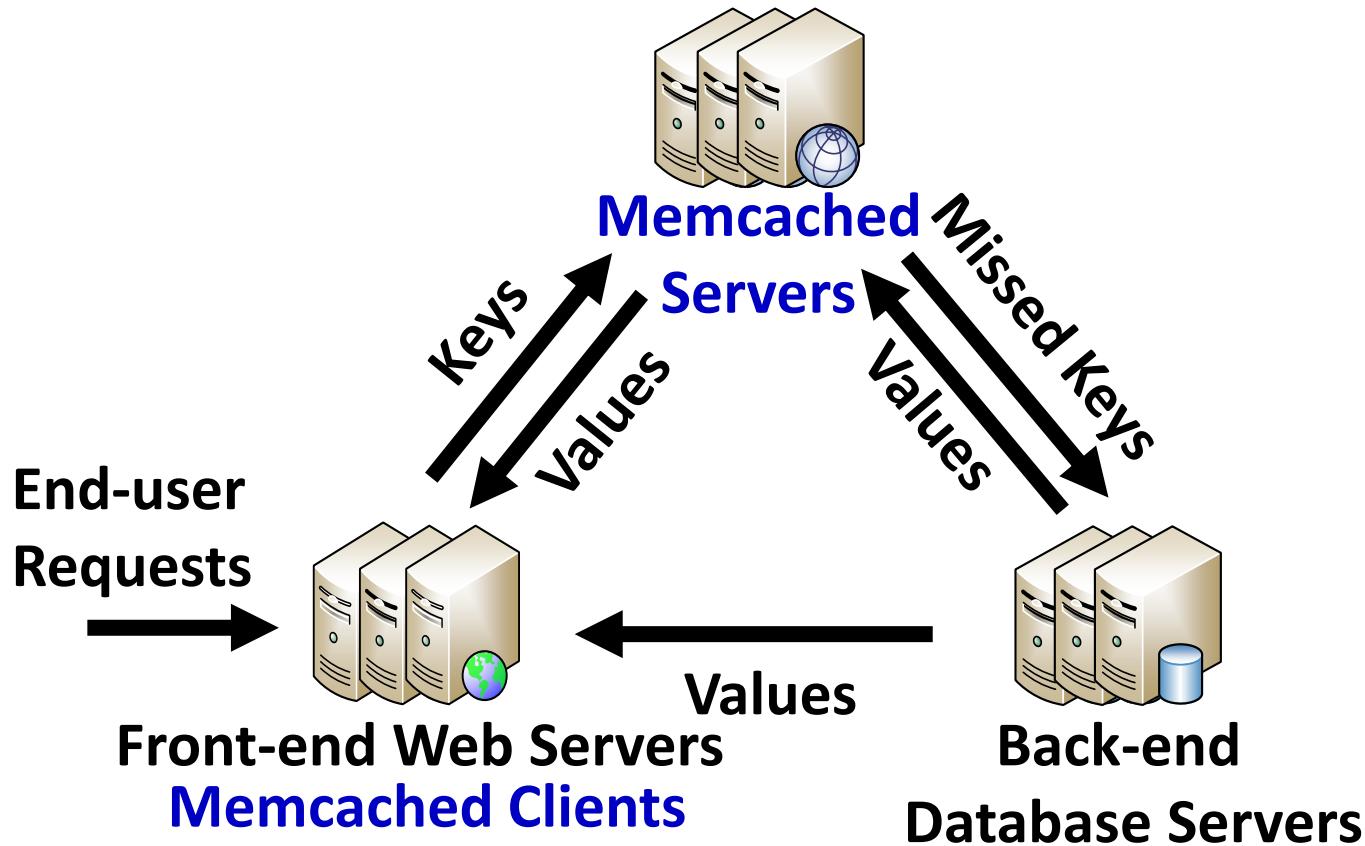


YouTube

Traditional Websites

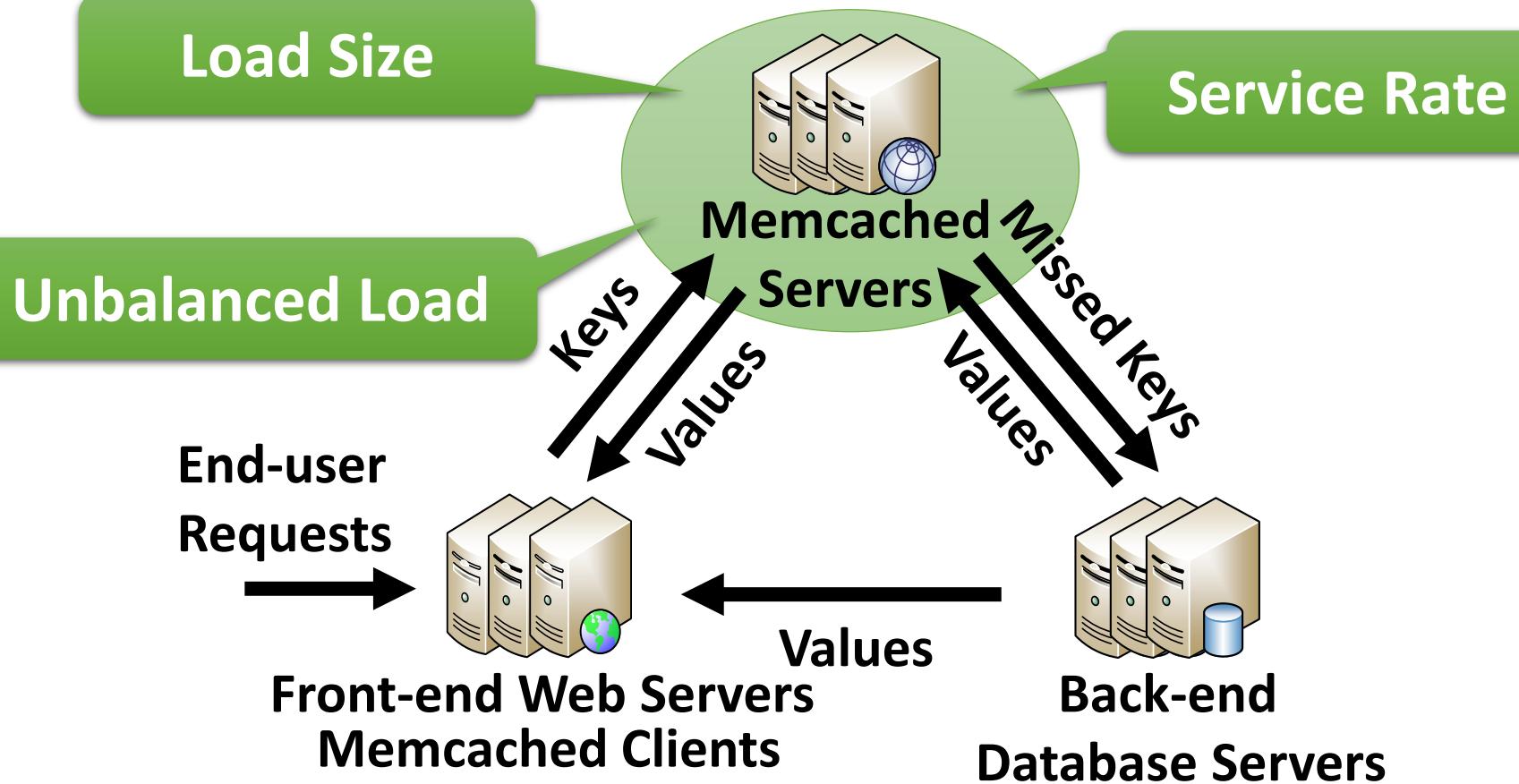


Memcached Architecture

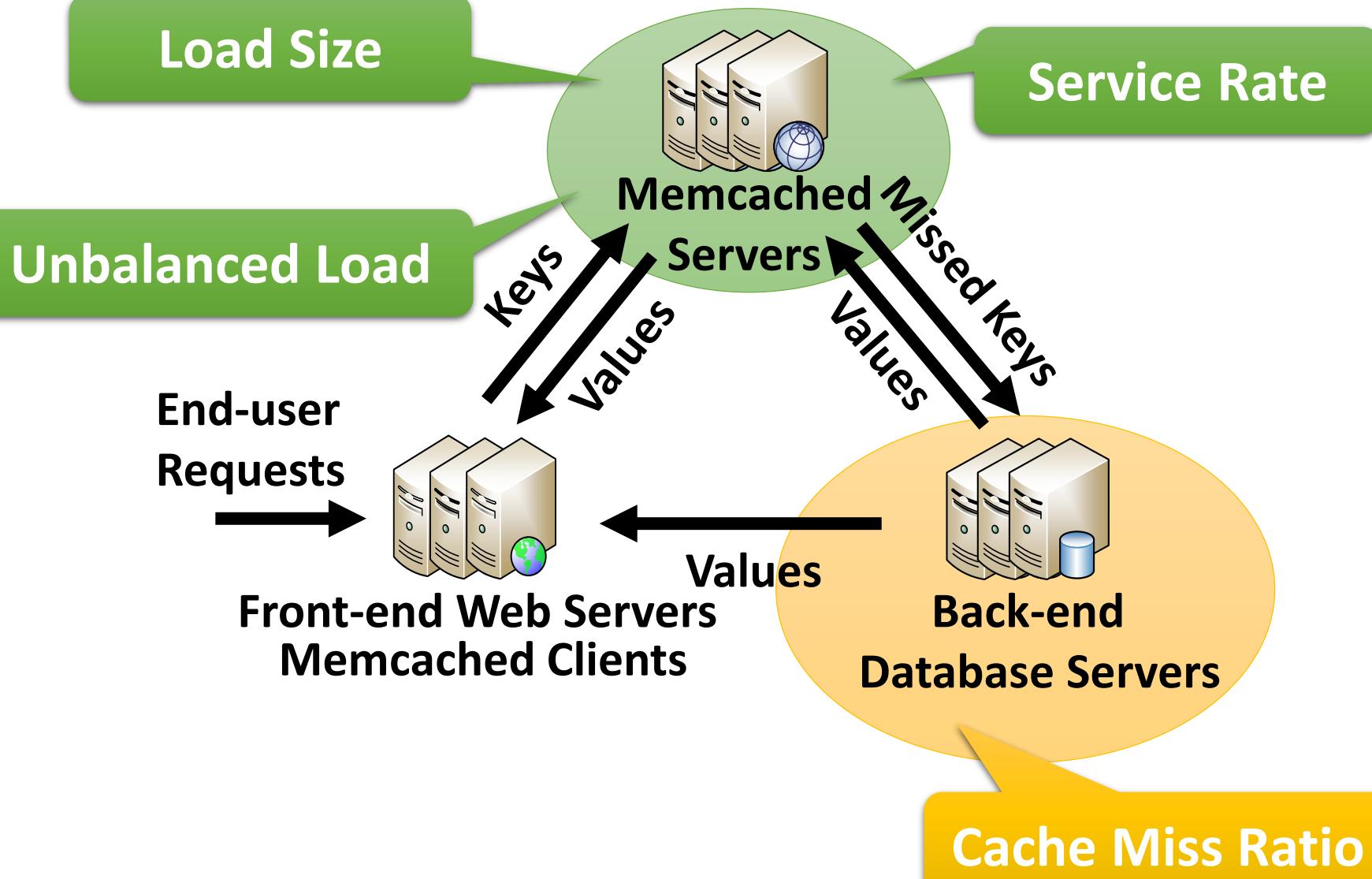


Latency is the most pivotal performance metric.

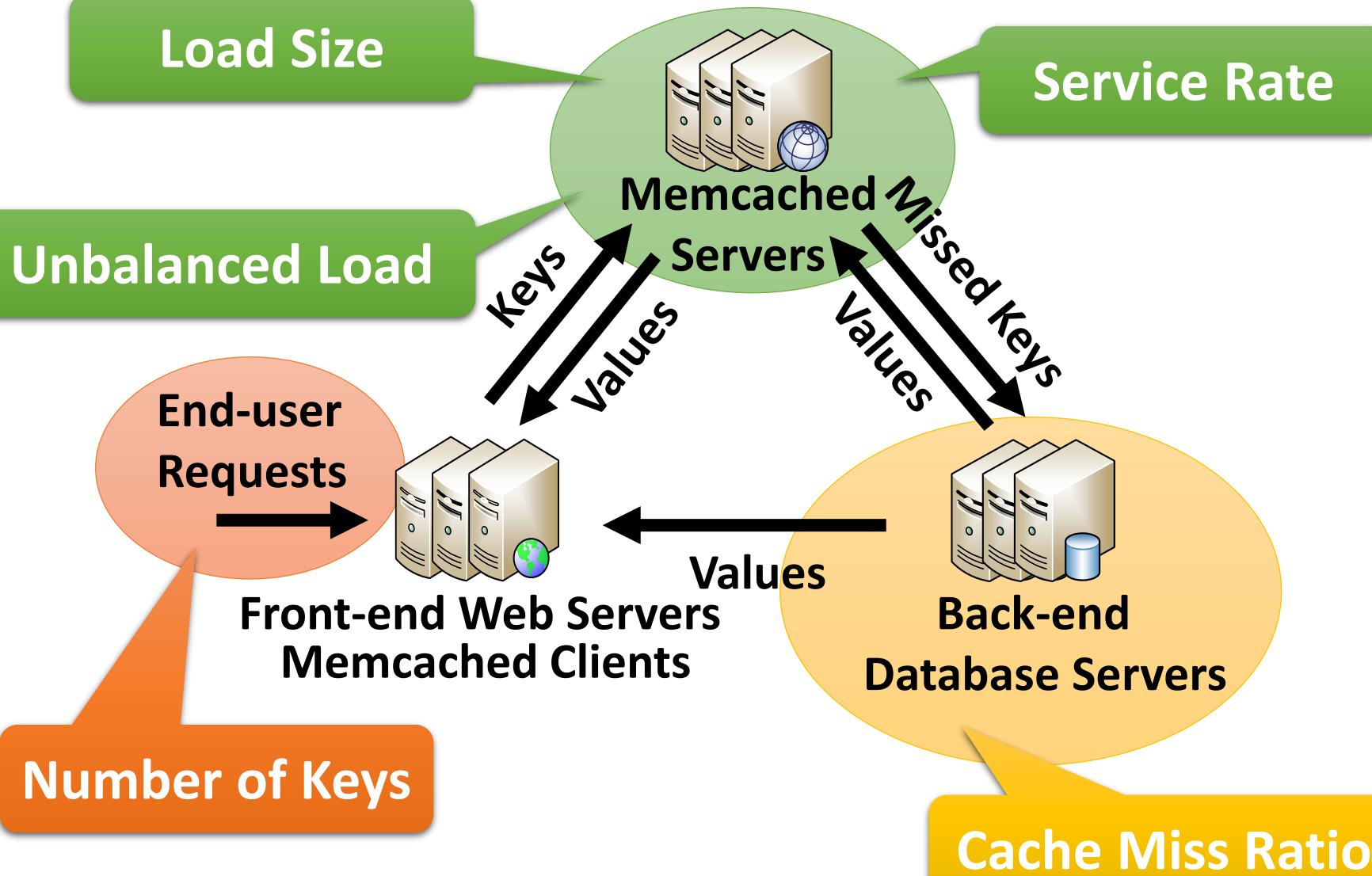
Factors on Latency



Factors on Latency



Factors on Latency



Factors on Latency

Load Size



Service Rate

- How much improvement on latency can be achieved by optimizing each factor?



Values

Front-end Web Servers
Memcached Clients



Back-end
Database Servers

Number of Keys

Cache Miss Ratio

Factors on Latency

Load Size



Service Rate

- How much improvement on latency can be achieved by optimizing each factor?

Theoretical Model & Quantitative Analysis



Front-end Web Servers
Memcached Clients

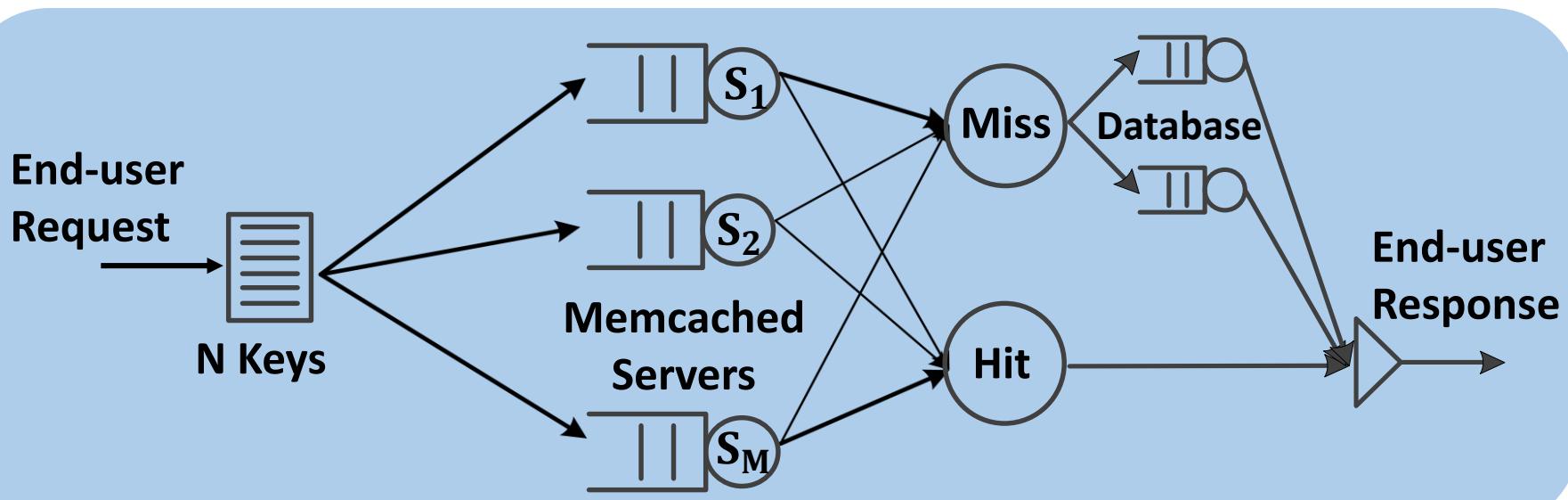
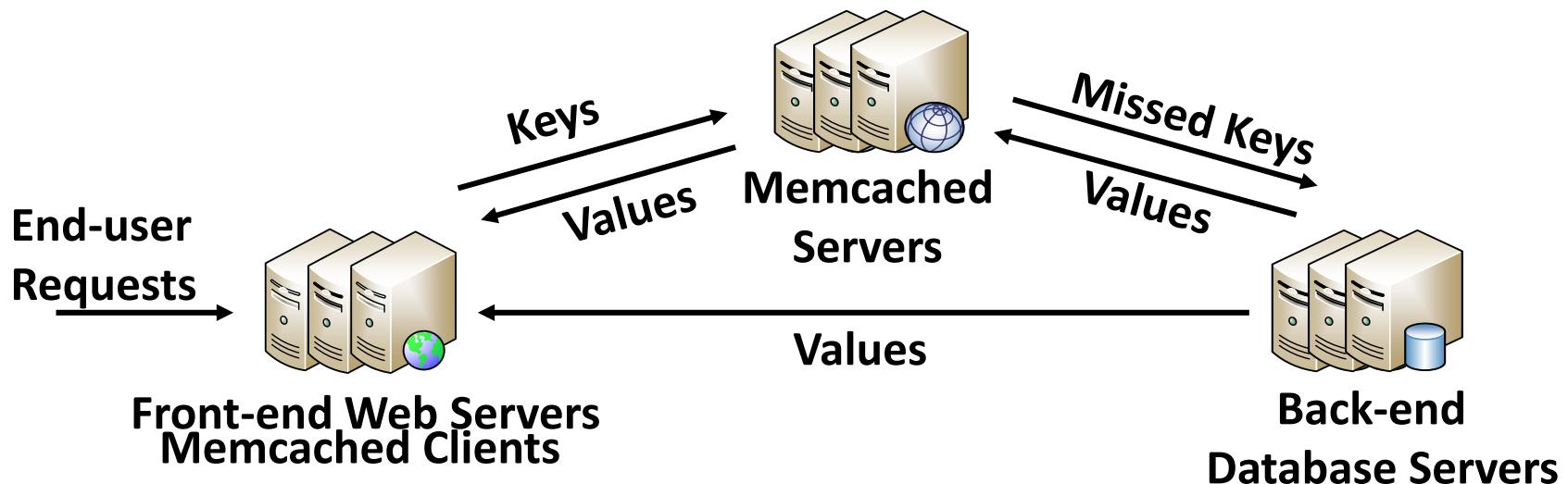
Number of Keys

Values

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Cache Miss Ratio

Model for Memcached



Specific Features

Unbalanced Load

End-user Requests

Front-end Web Servers
Memcached Clients

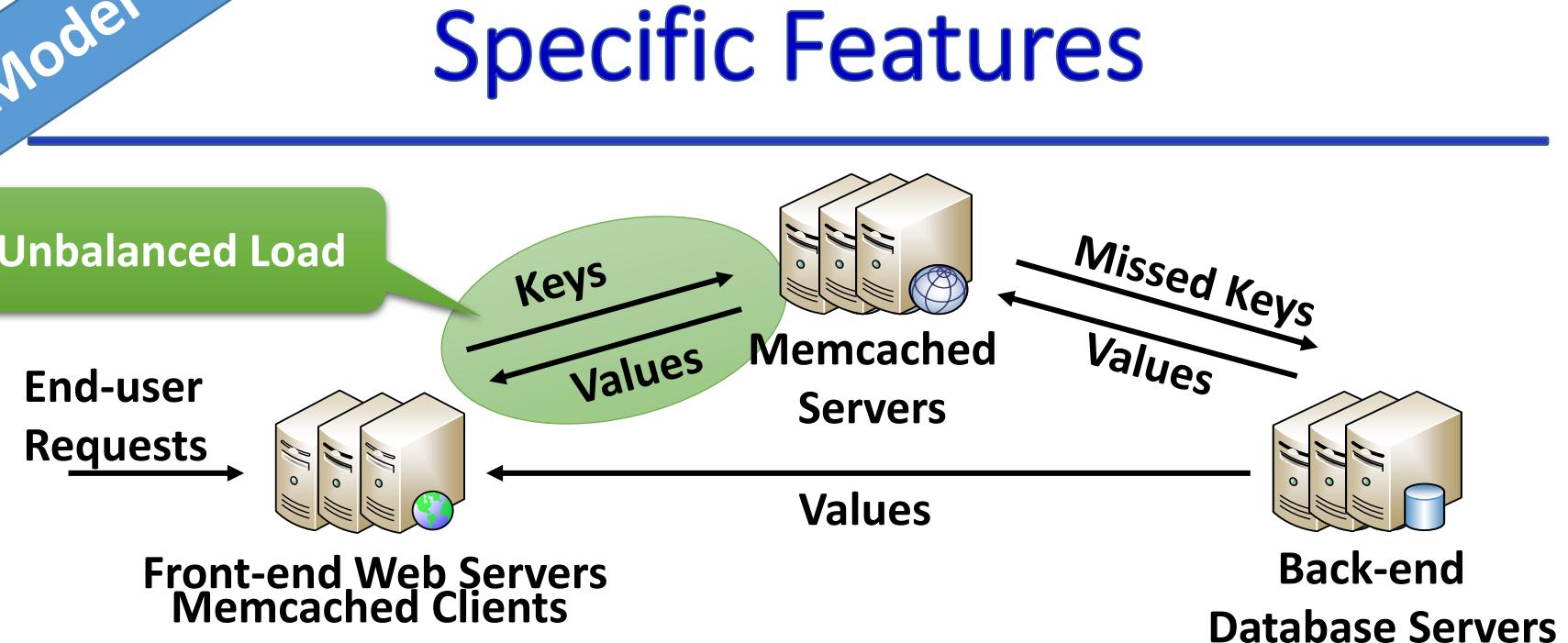


Missed Keys
Values



Memcached
Servers

Values



End-user Request

N Keys

$$P_1 * N$$
$$P_2 * N$$
$$P_M * N$$

Memcached
Servers

S₁

S₂

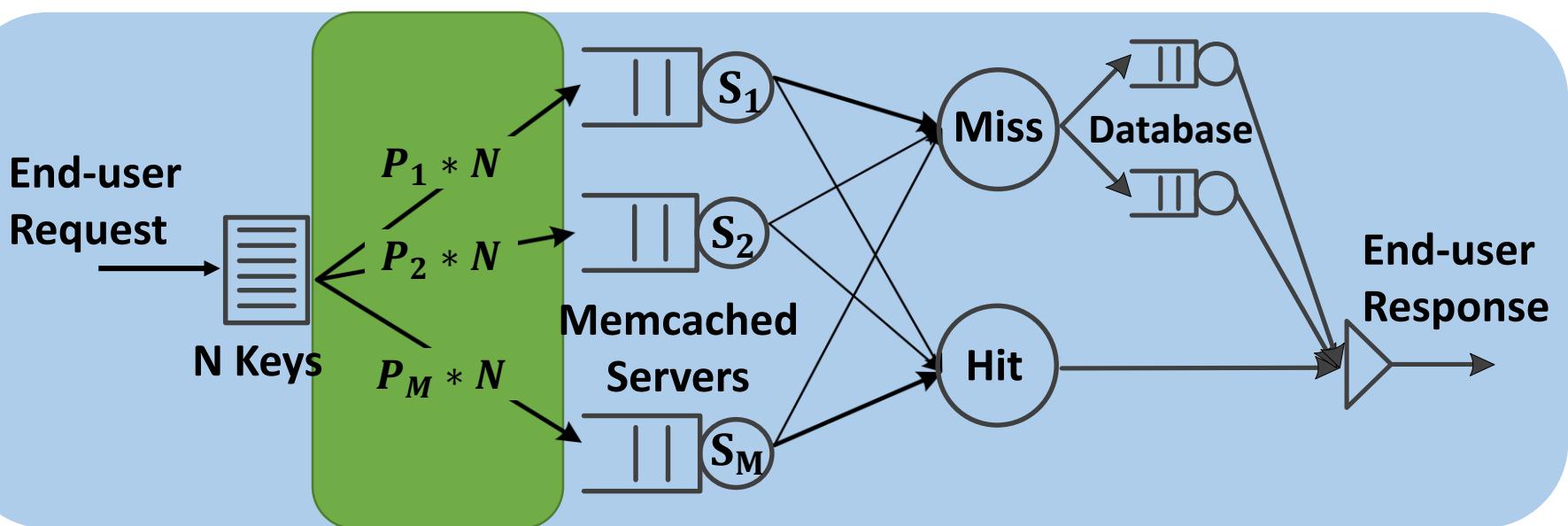
S_M

Miss

Hit

Database

End-user Response



Specific Features

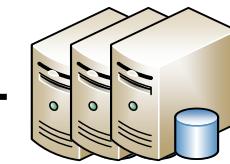
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End-user Requests

Front-end Web Servers
Memcached Clients

Memcached
Servers

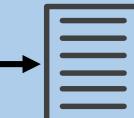
Missed Keys
Values



Back-end
Database Servers

Burst and Concurrent
Key Arrivals

End-user
Request



N Keys

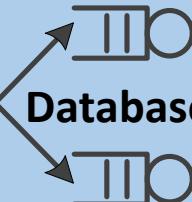
$$P_1 * N \\ P_2 * N \\ \vdots \\ P_M * N$$

Memcached
Servers



Miss

Hit



End-user
Response

Specific Features

Unbalanced Load

End-user Requests

Front-end Web Servers
Memcached Clients

Memcached Servers

Database Processing

Missed Keys

Values

Back-end Database Servers

Burst and Concurrent Key Arrivals

End-user Request

N Keys

$$P_1 * N \\ P_2 * N \\ \vdots \\ P_M * N$$

Memcached Servers



Hit

Miss

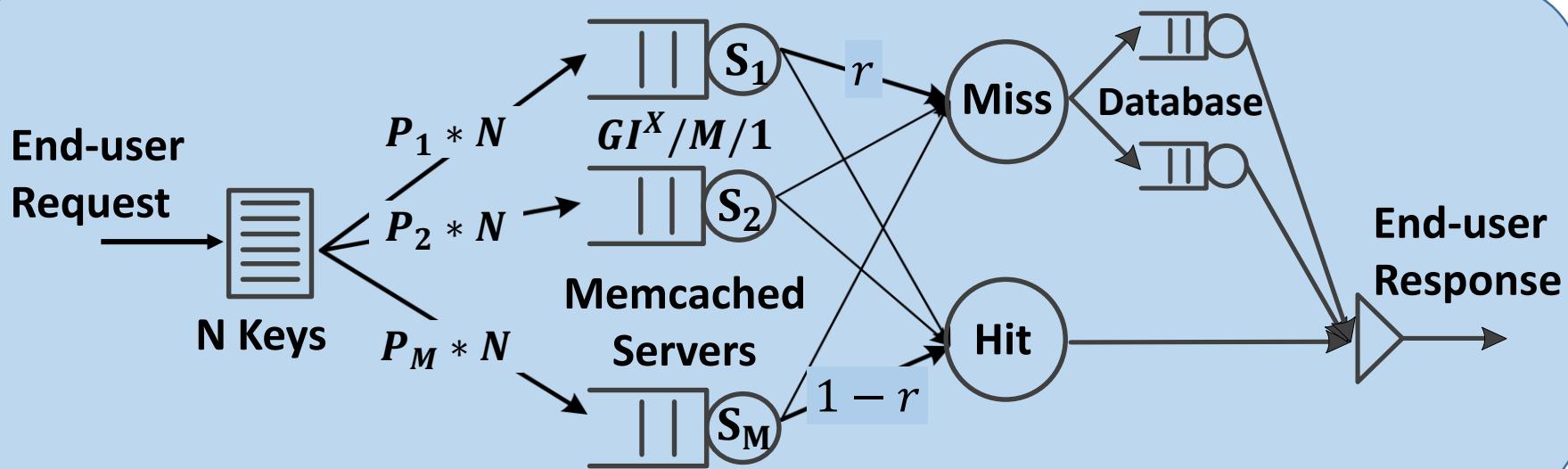
Database

End-user Response

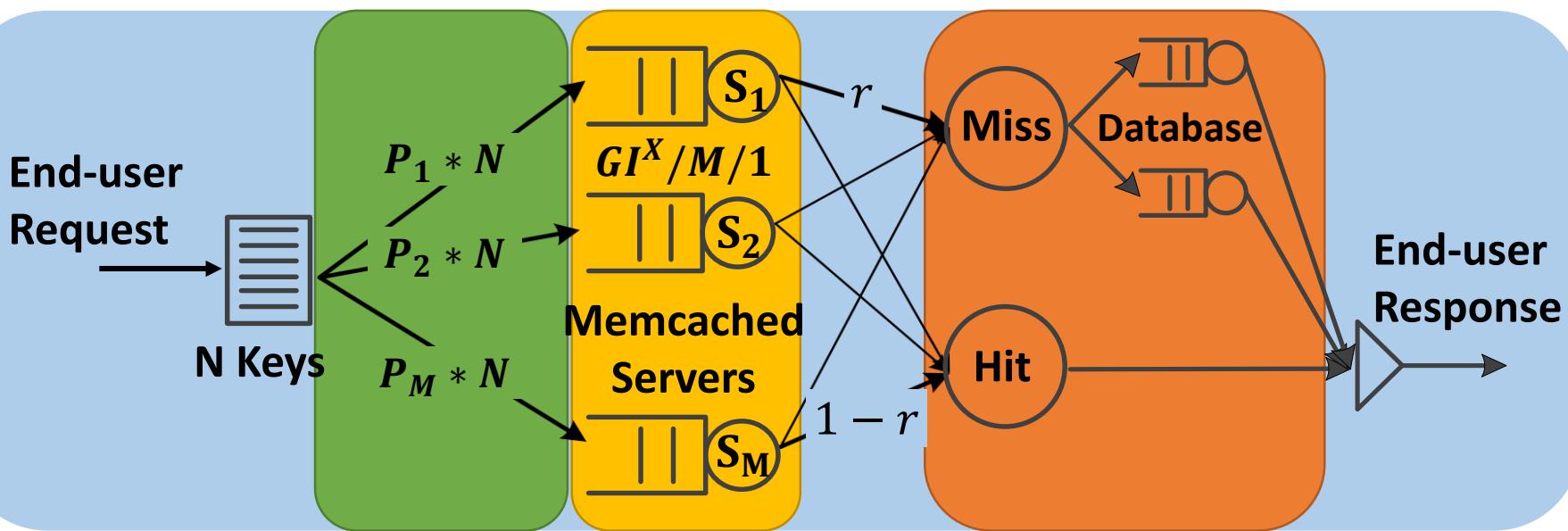
$$1 - r$$

Assumptions

- One end-user request simultaneously generates multiple independent keys.
- Service times are exponential, both at Memcached servers and database. [CoNEXT'13][NSDI'15]

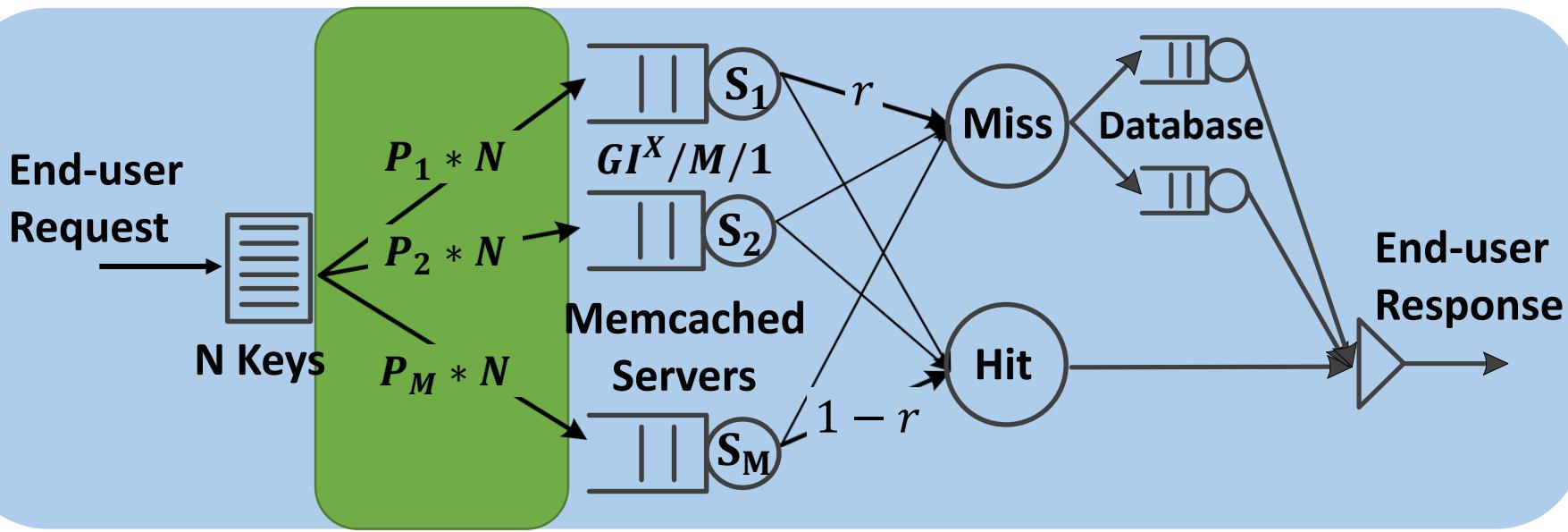


Latency Composition



- Network Latency
- Processing Latency at Memcached Servers
- Processing Latency at Database

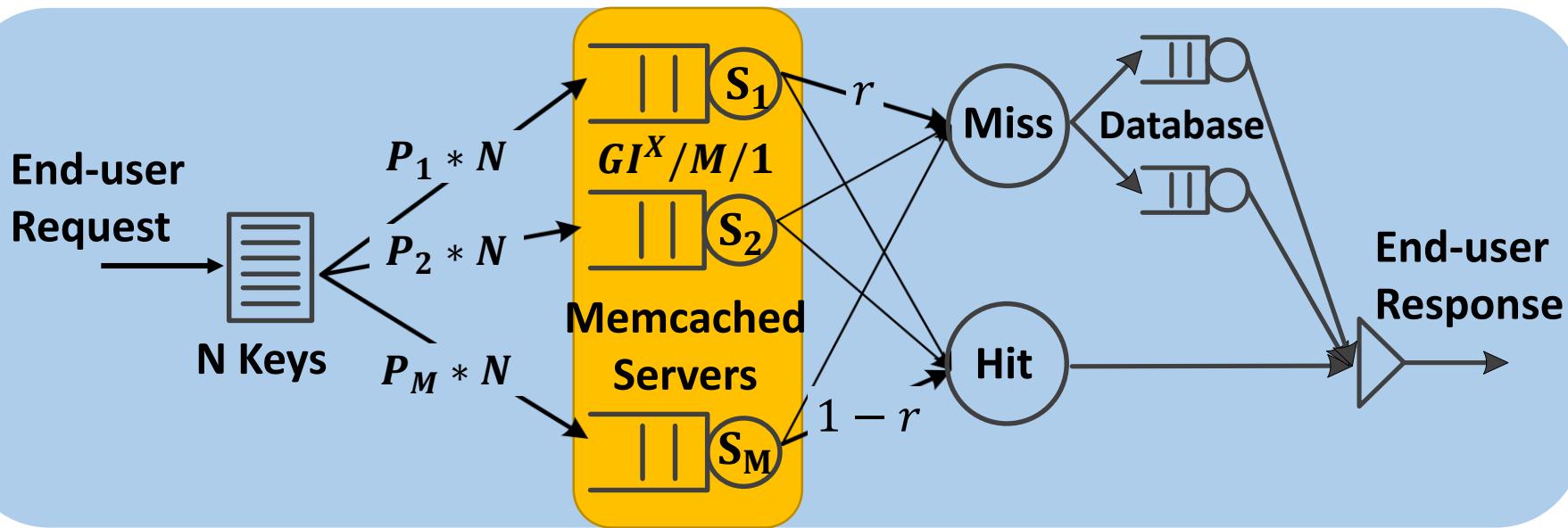
Network Latency



- Network Utilization $\ll 10\% \rightarrow$ No queueing in network transmission
 - Request rate: $10^5 / s$
 - Request size: Key < 200B, Value < 1KB^[SIGMETRICS'12]
 - \rightarrow Workload $\ll 1\text{Gbps}$
 - \rightarrow Network bandwidth 10Gbps

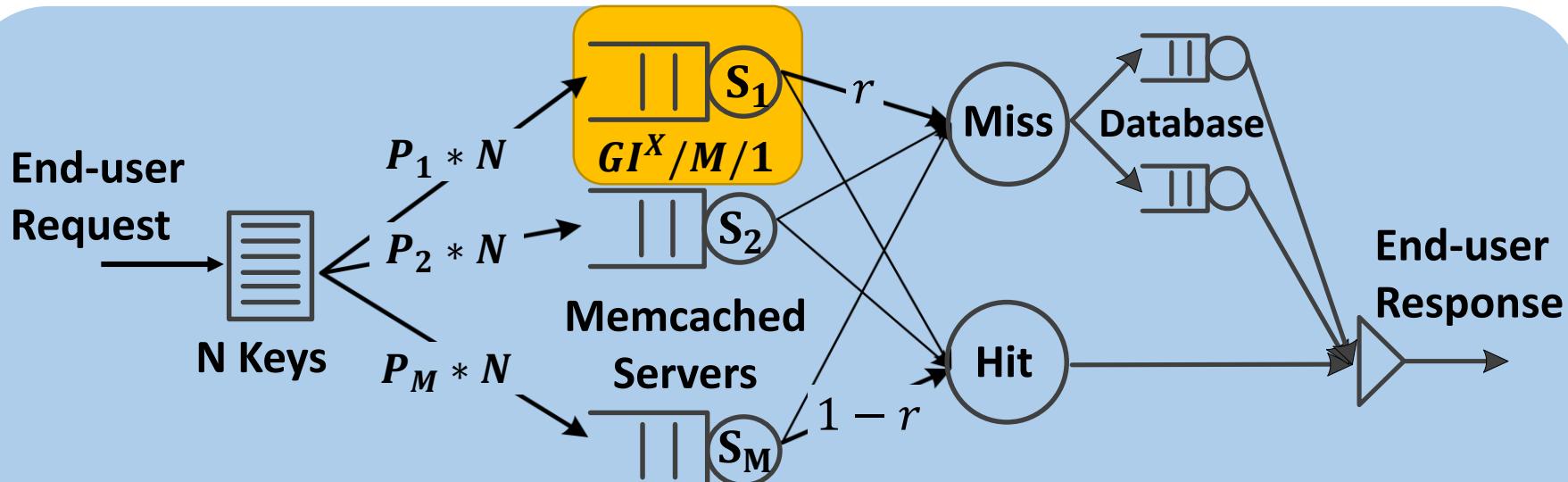
Network latency is almost constant

Latency at Memcached servers



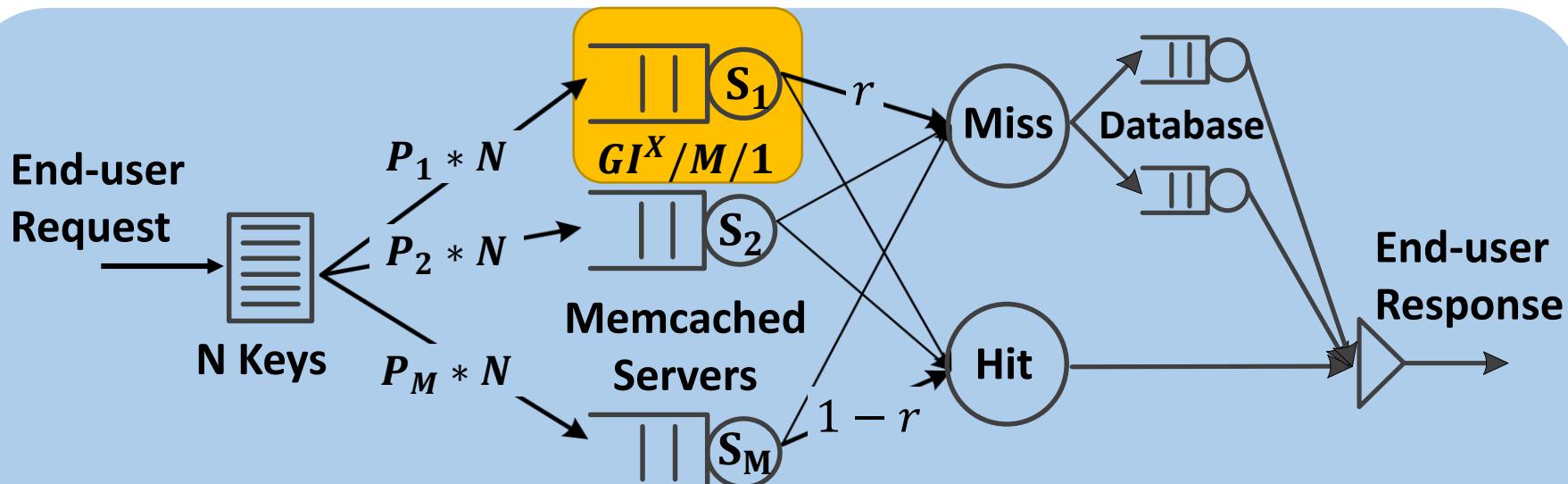
- Latency for 1 Key \rightarrow
- Maximum Latency for N Keys

Latency at Memcached servers



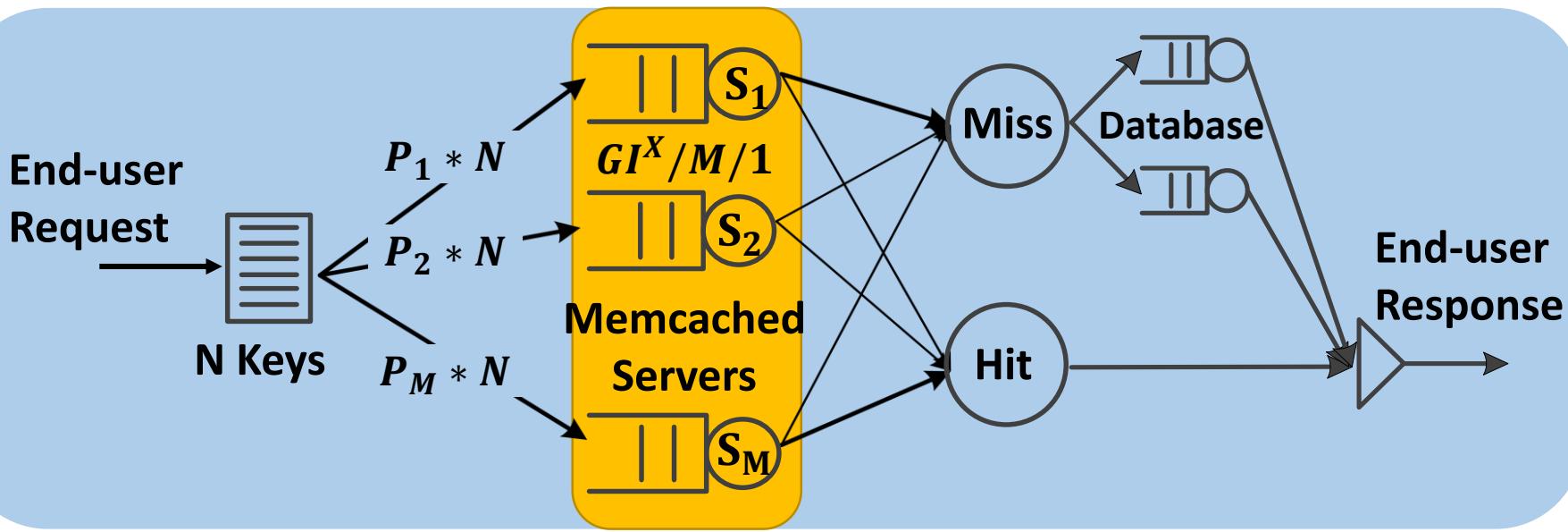
- Latency for 1 Key
 - Queuing at each Memcached server : $GI^X/M/1$
 - GI : Any distribution → Burst arrival of keys
 - X : Batch blocks → Concurrent arrival of keys
 - $GI^X/M/1 \rightarrow GI/M/1$
 - Response time for a batch < Latency for 1 Key \leq Response time for a batch

Latency at Memcached servers



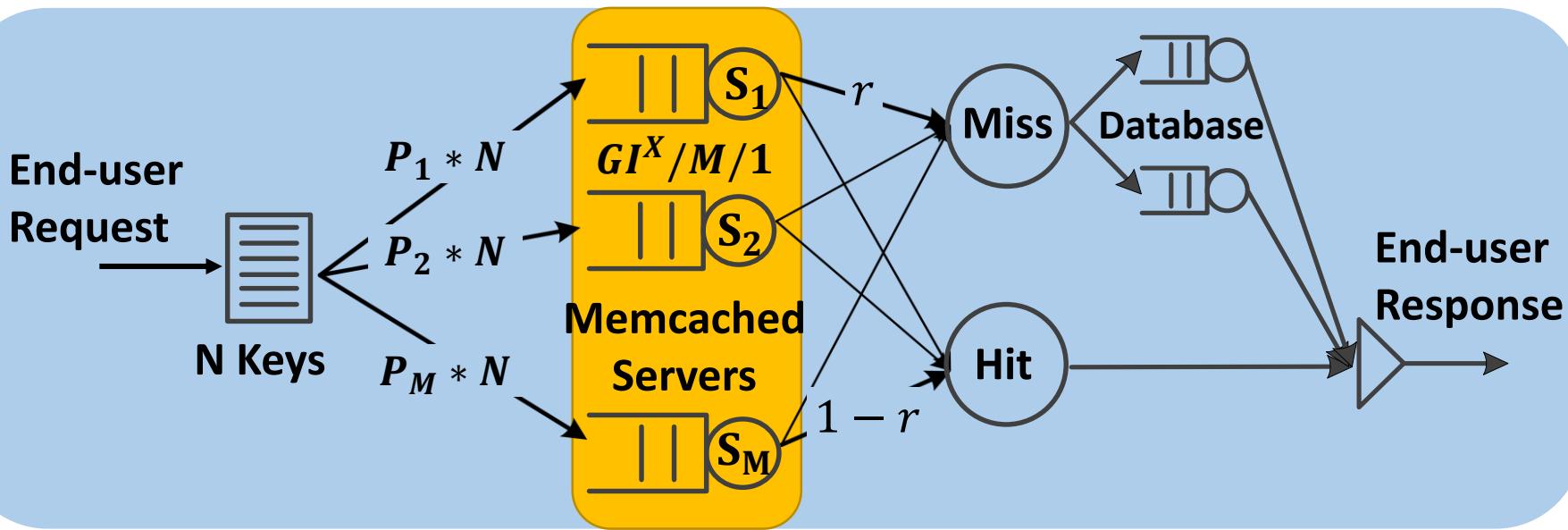
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 - GI : Any distribution → Burst arrival of keys
 - X : Batch blocks → Concurrent arrival of keys
 - $GI^X/M/1 \rightarrow GI/M/1$
 - Waiting time for a batch < Latency for 1 Key \leq Response time for a batch

Latency at Memcached servers



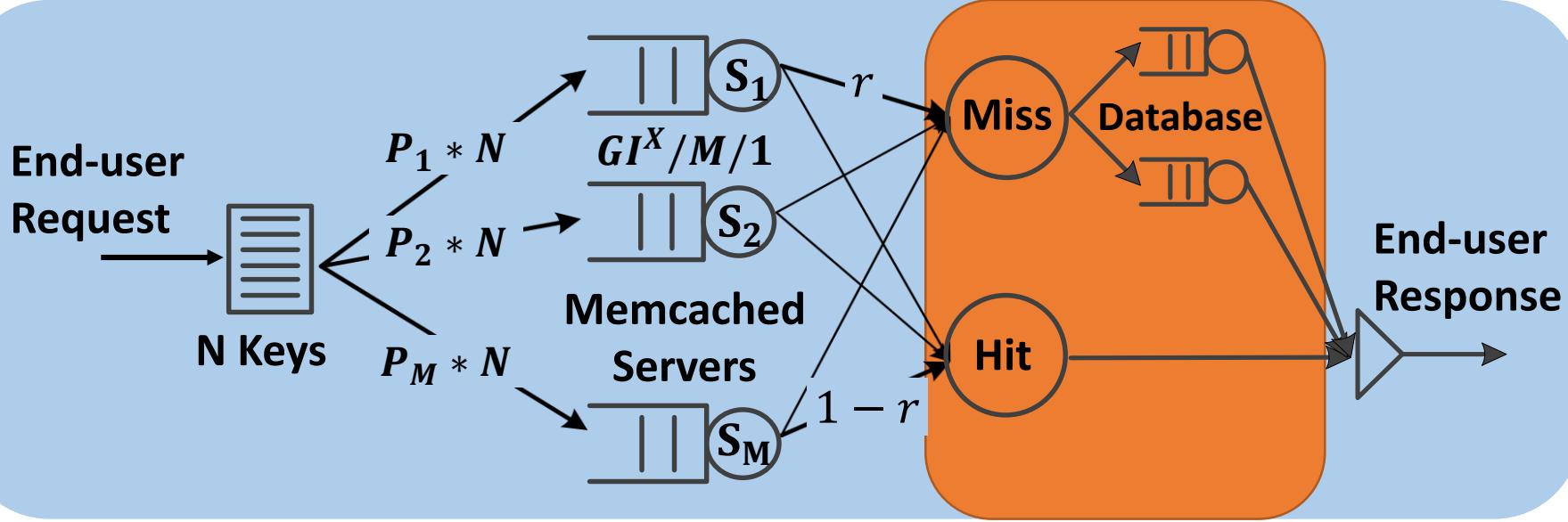
- Maximum Latency for N Keys
 - Balanced load $\rightarrow \approx \left(\frac{N}{N+1}\right)_{th}$ of latency for 1 key
 - Unbalanced load among Memcached servers : $\{P_j\}_{j=1}^M$
 - $\in \left[\left(\frac{N}{N+1}\right)^{\frac{1}{p_{max}}}, \left(\frac{N}{N+1}\right)_{th}\right]$ of latency for a key at the heaviest server

Latency at Memcached servers



- Maximum Latency for N Keys
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 - Unbalanced load among Memcached servers : $\{P_j\}_{j=1}^M$
 - $\in \left[\left(\frac{N}{N+1}\right)^{p_{max}}, \left(\frac{N}{N+1}\right)\right]_{th}$ of latency for 1 key at the heaviest server

Latency at Database



- Depends on the number of keys N and cache miss ratio r
 - $(1 - r)^N$: All-hit cases → Zero latency.
 - $1 - (1 - r)^N$: Miss cases → Large latency.

Symbols

Symbol	Definition
N	Number of keys generated from an end user request.
M	Number of Memcached servers.
$\{p_j\}_{j=1}^M$	Load distribution among Memcached servers.
q	Concurrent probability of keys.
X	Batch size for concurrent keys.
T_X	Inter-arrival gap of batches.
λ	Arrival rate of keys.
μ_S	Service rate of one Memcached server.
r	Cache miss ratio.
μ_D	Service rate at database.

Latency Estimation

Latency $T(N)$ of an end-user request that generates N keys is separately bounded by three parts,

$$\max\{T_N(N), T_S(N), T_D(N)\} \leq T(N) \leq T_N(N) + T_S(N) + T_D(N)$$

- The network latency $T_N(N)$ is almost constant.
- The processing latency at Memcached servers $T_S(N)$ satisfies

$$\max \left\{ \frac{\ln \delta + \frac{1}{p_{max}} \ln(N+1)}{(1-\delta)(1-q)\mu_S}, 0 \right\} \leq E[T_S(N)] \leq \frac{\frac{1}{p_{max}} \ln(N+1)}{(1-\delta)(1-q)\mu_S}$$

- The processing latency at database $T_D(N)$ can be estimated by

$$E[T_D(N)] \approx \frac{1 - (1-r)^N}{\mu_D} * \ln \left(\frac{N * r}{1 - (1-r)^N} + 1 \right)$$

Basic Validation

- Experimental Setup
 - 2 client machines and 4 server machines
 - Intel CoreTM i5-5200U CPU and 8GB memory
 - Same Rank, 1Gbps links
 - Workload
 - Mutilate: 512 connections
 - Measurements of Facebook [Sigmetrics'12]

Latency	Estimation	Experiment	Confidence Interval
Network	$20\mu s$	$20\mu s$	$[18.12\mu s, 21.68\mu s]$
Memcached Servers	$351\mu s \sim 366\mu s$	$368\mu s$	$[362\mu s, 373\mu s]$
Database	$836\mu s$	$867\mu s$	$[855\mu s, 879\mu s]$
Total	$836\mu s \sim 1222\mu s$	$1144\mu s$	$[1128\mu s, 1160\mu s]$

Factors on Latency

Load Size



Service Rate

- How much improvement on latency can be achieved by optimizing each factor?

Quantitative Analysis



Front-end Web Servers
Memcached Clients

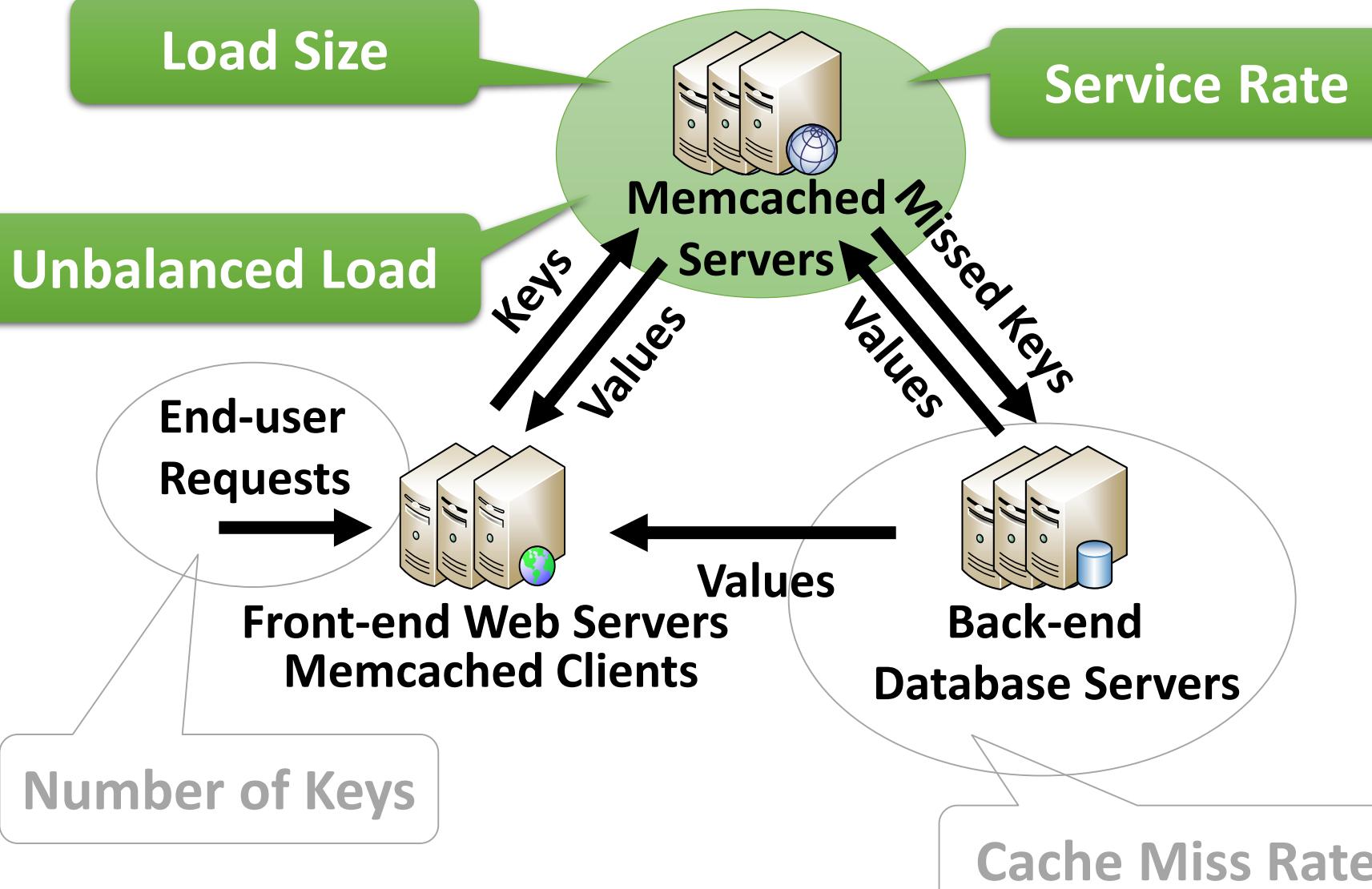
Number of Keys

Values

Back-end
Database Servers

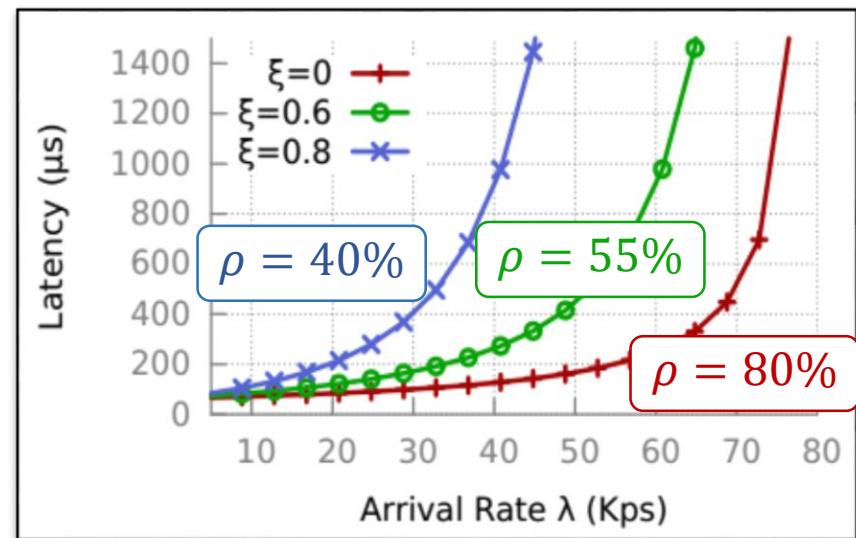
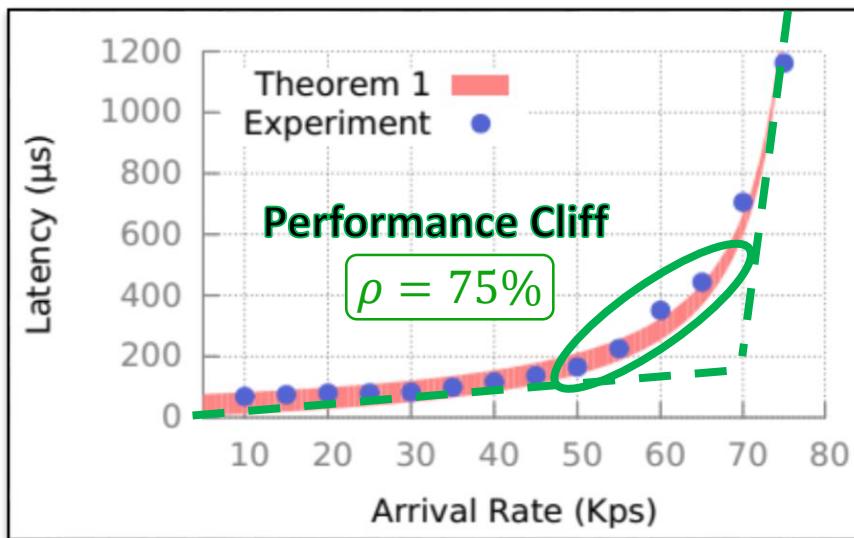
Cache Miss Ratio

Factors on Latency



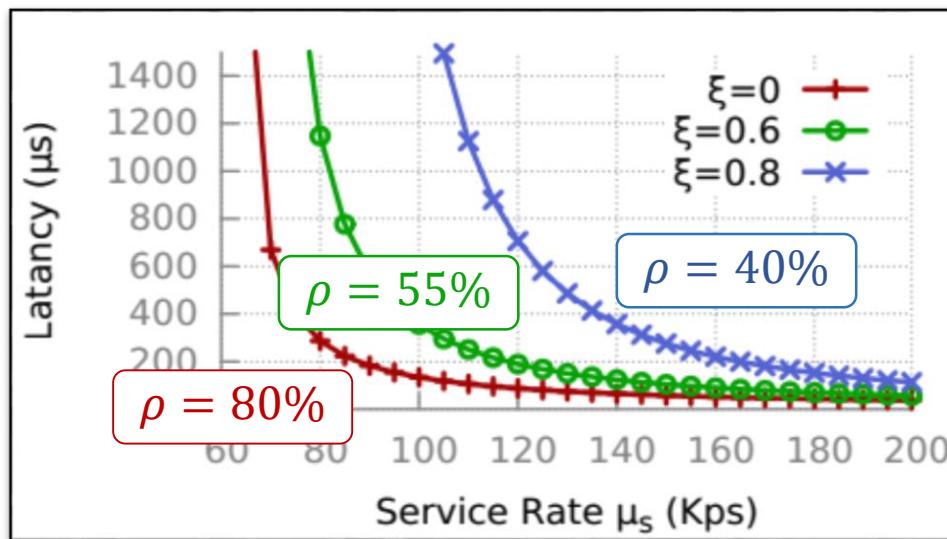
Load Size

- Workload: Generalized Pareto Distribution
 - Arrival rate: λ (default = 62.5Kbps)
 - Burst degree: ξ (default = 0.15)
- Service rate: μ_s (default = 80Kps)

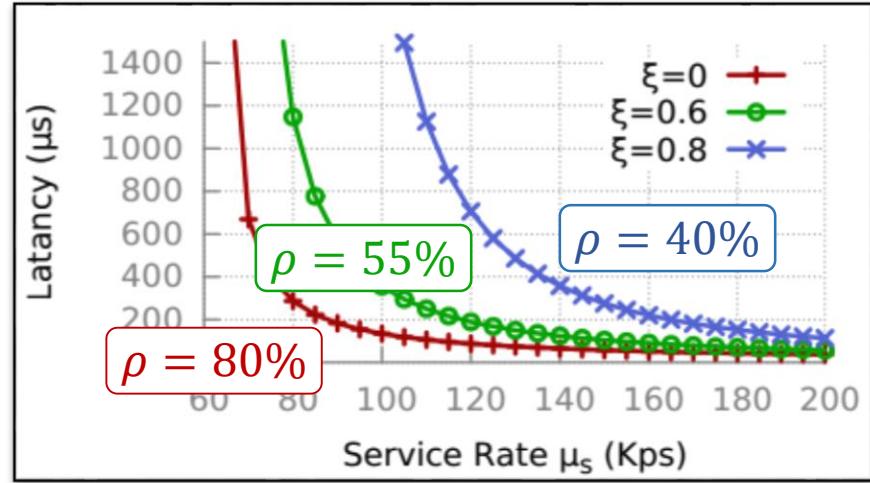
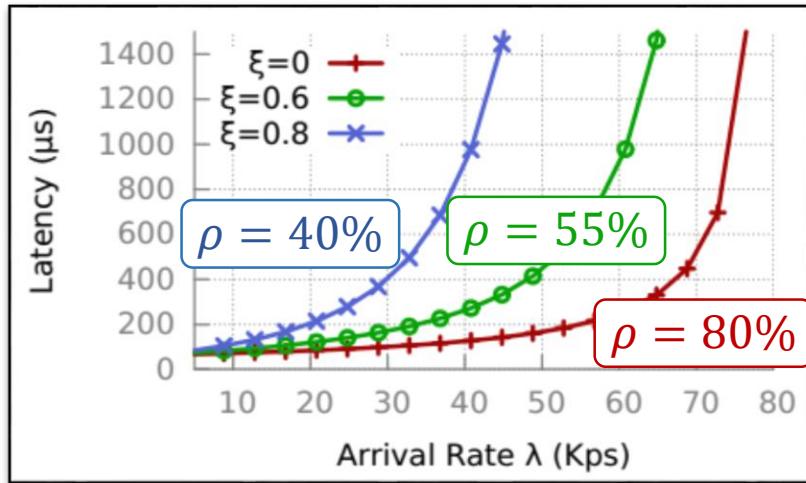


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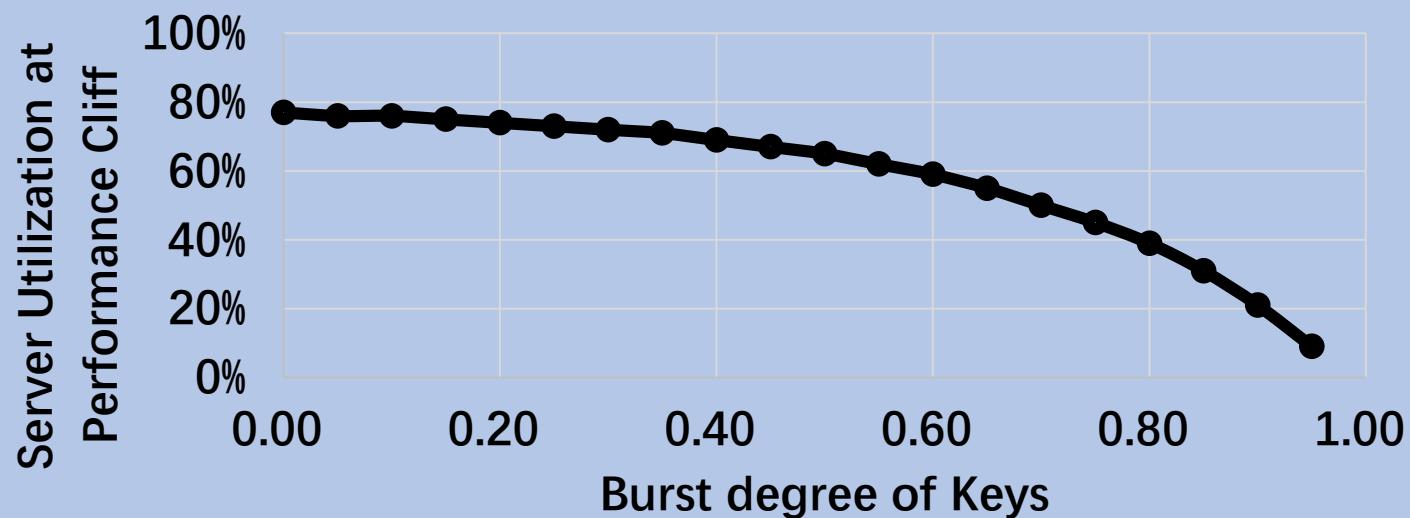


Load Size & Service Rate



- *Latency reaches a **cliff point** when the Memcached **server utilization** gets to a specific value.*
- *This specific utilization is only determined by the **burst degree** of workload.*

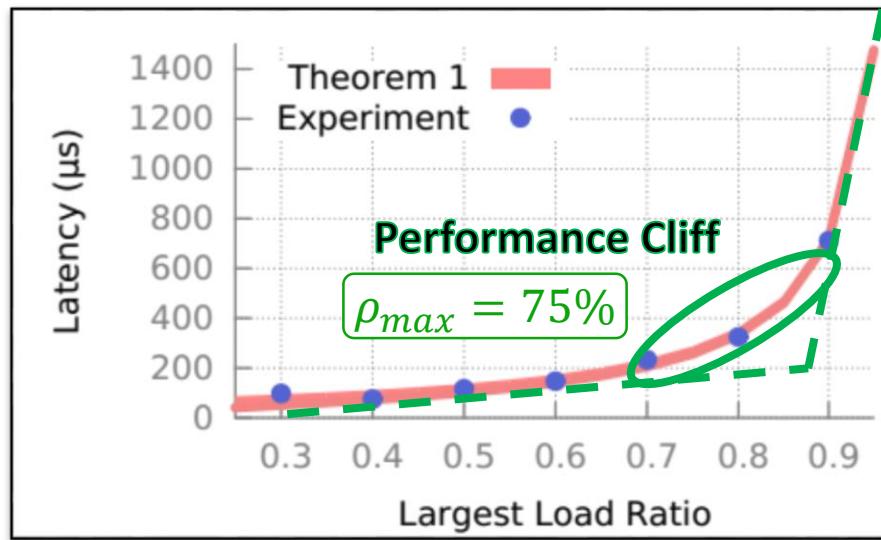
Server Utilization



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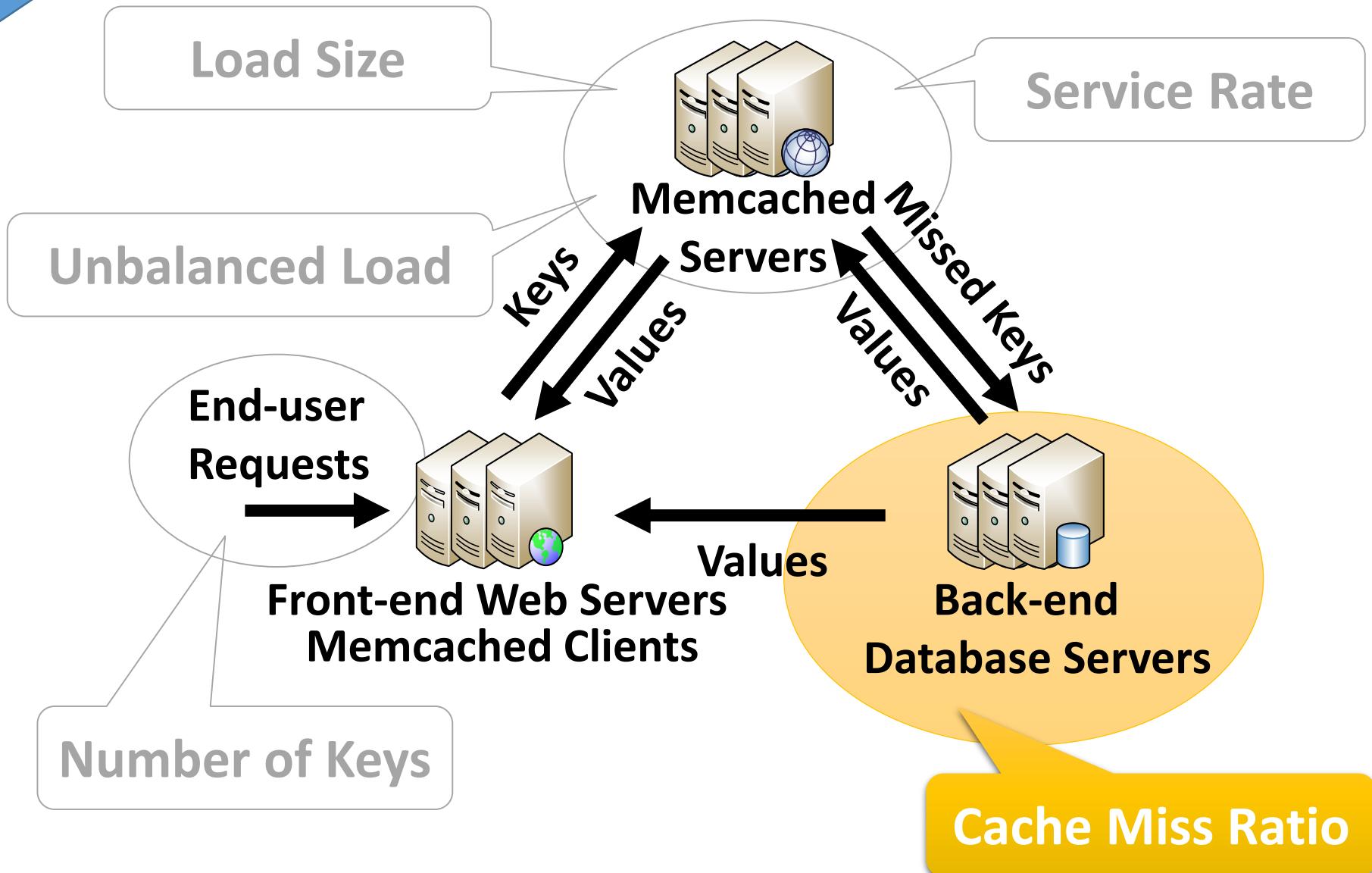
Unbalanced Load

- 4 Memcached servers: Largest load ratio
 - Total load size: $\Lambda = 80Kbps$
 - Burst degree: $\xi = 0.15$
 - Service rate: $\mu_s = 80Kbps$



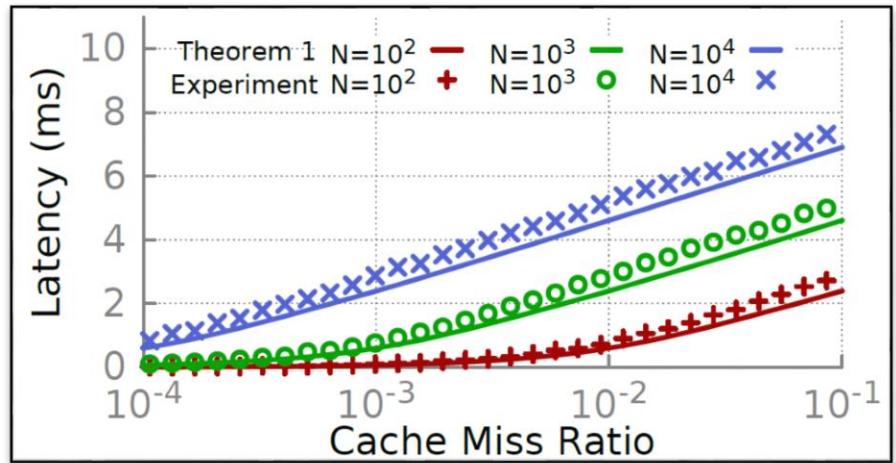
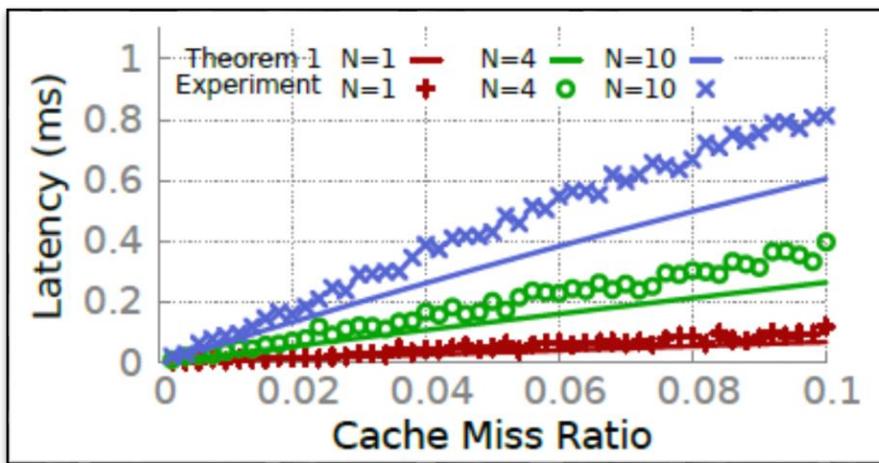
- Latency still reaches a **cliff point** when the **heaviest** server gets to the **specific utilization**.

Factors on Latency



Cache Miss Ratio

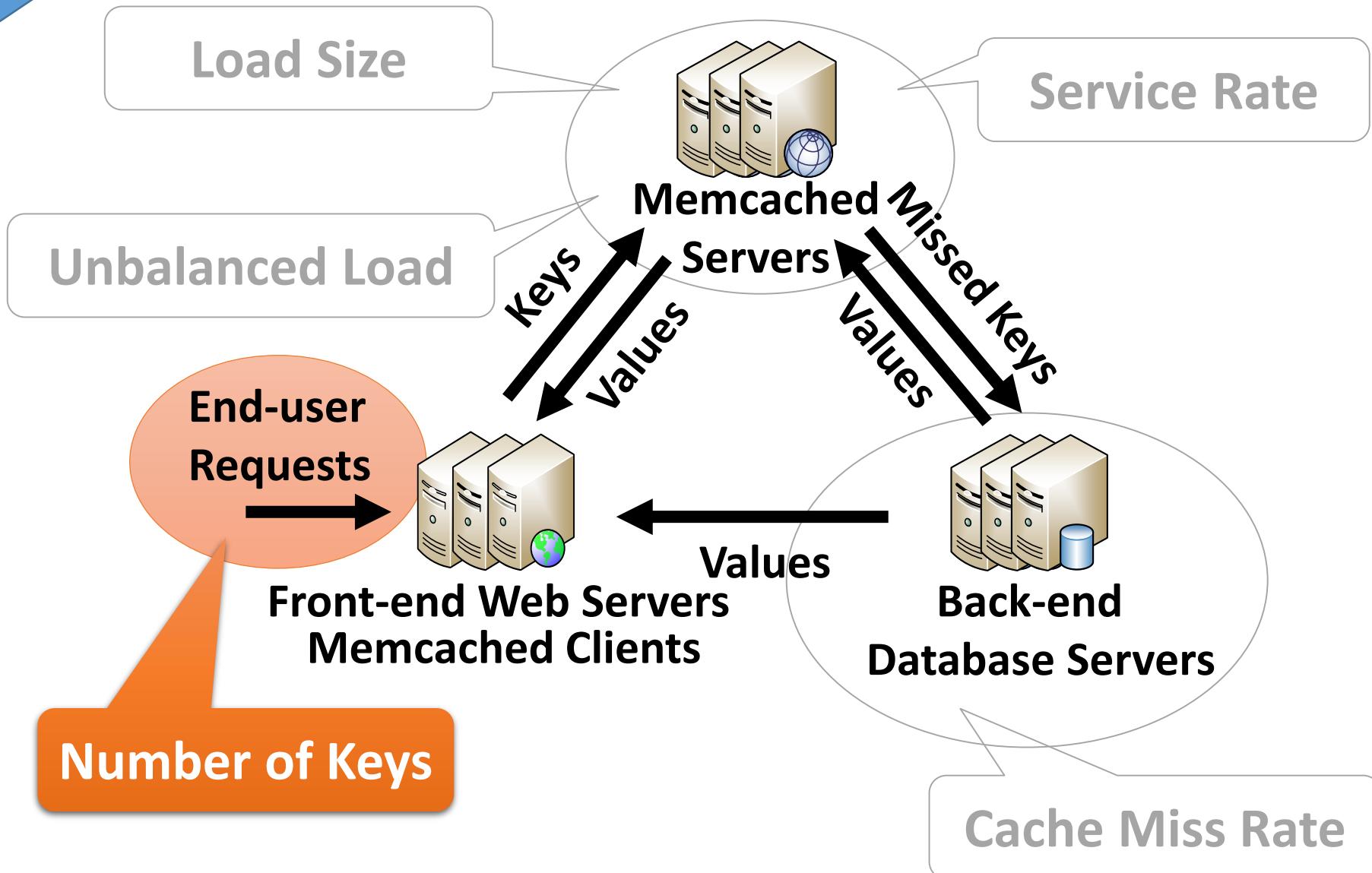
- Average cost at database: $\frac{1}{\mu_D} = 1ms$



- $r = 0.1 \rightarrow 0.01$
- $N = 1 \rightarrow$ latency: 0.8ms \rightarrow 0.08ms, $\downarrow 90\%$
- $N = 10^4 \rightarrow$ latency: 7.2ms \rightarrow 4.3ms, $\downarrow 40\%$

- Latency grows **logarithmically** with the increase of the **cache miss ratio**.

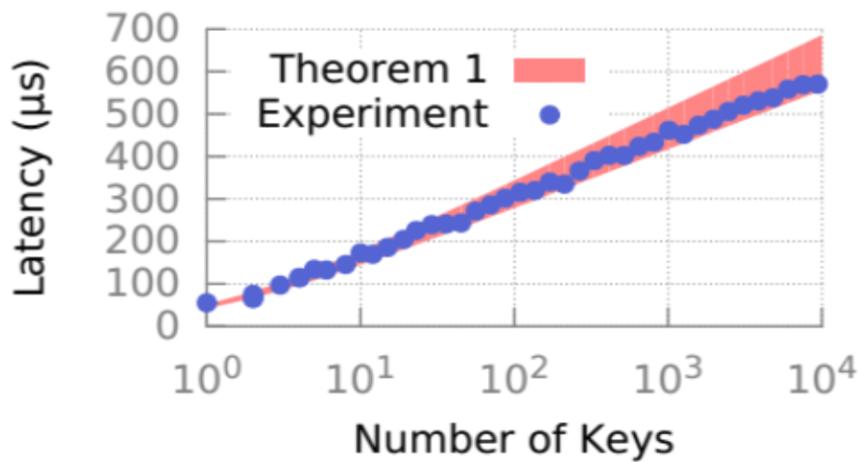
Factors on Latency



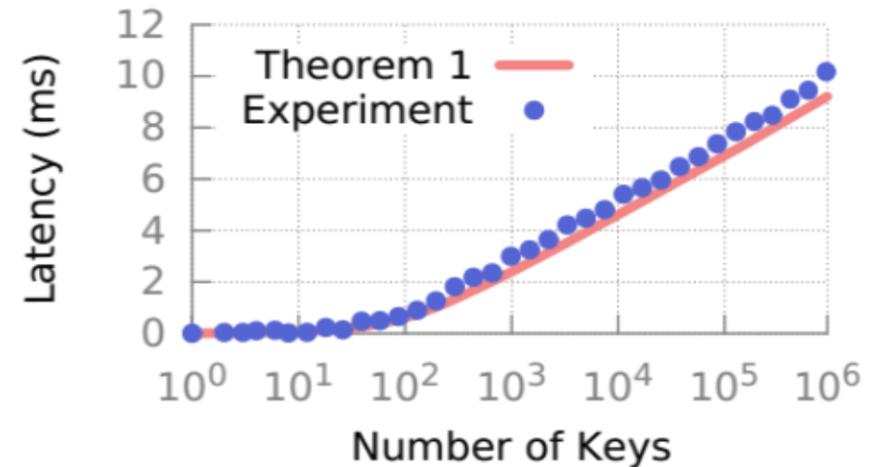
Number of Keys

- Workload burst degree $\xi = 0.15$
- Memcached server utilization $\rho = 0.75$
- Cache miss rate $r = 0.01$

Latency at Memcached Servers



Latency at Database

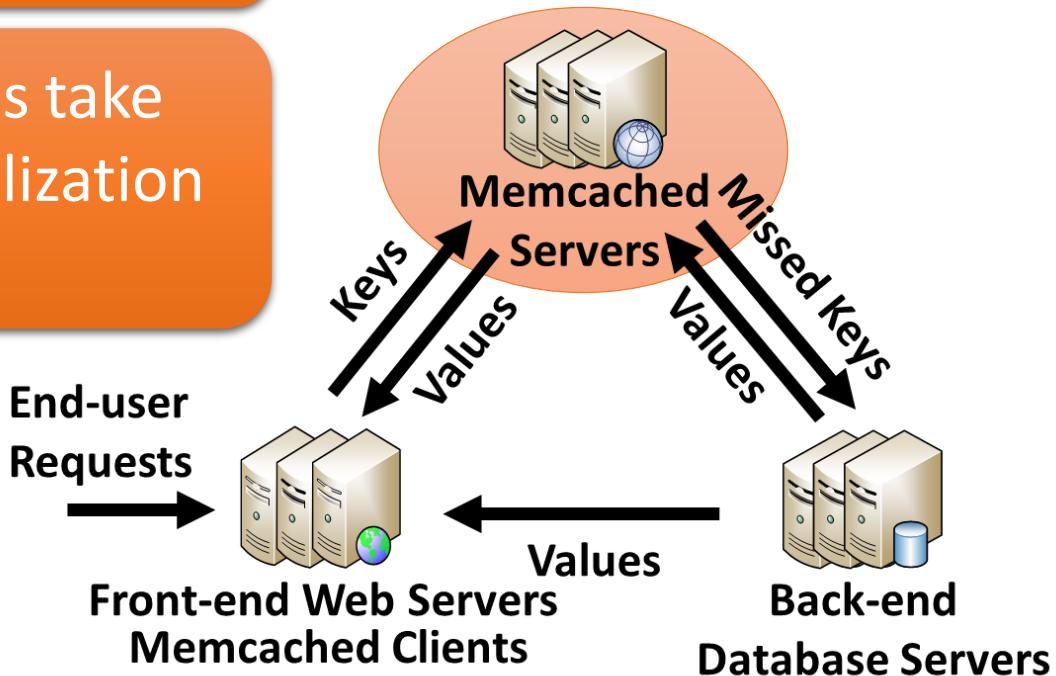


- Latency grows **logarithmically** with the increase of the number of keys.

Recommendations

Server utilization should not exceed a specific value.

Load-balancing mechanisms take effect before the largest utilization extends the specific value.

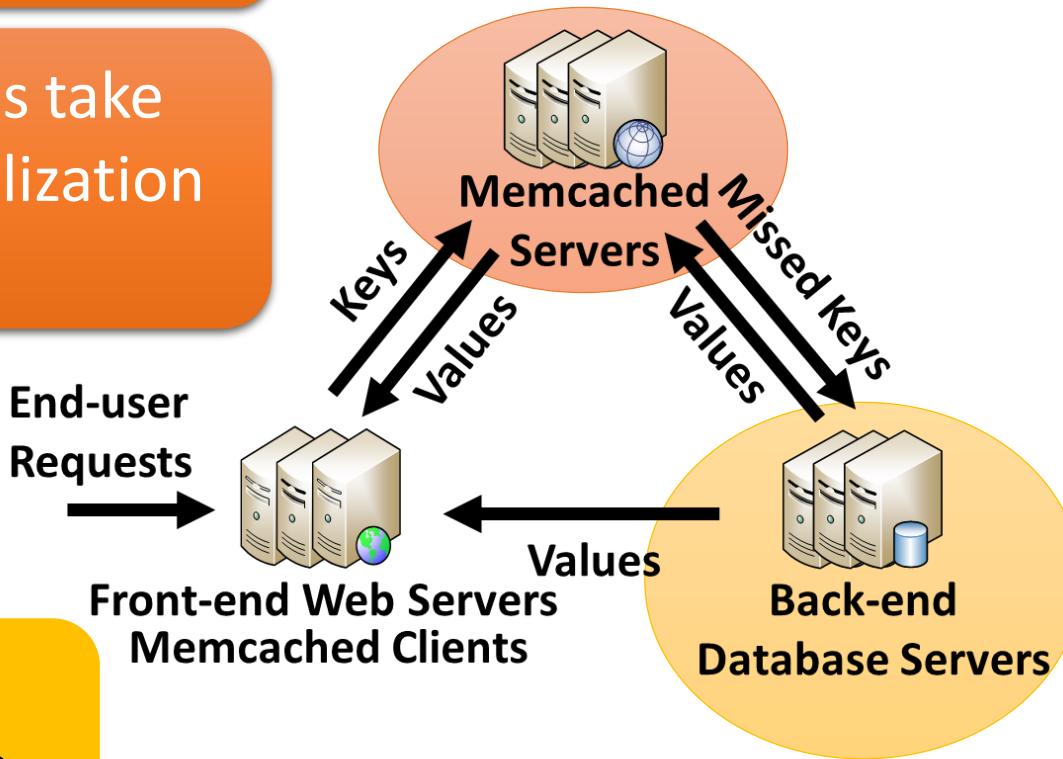


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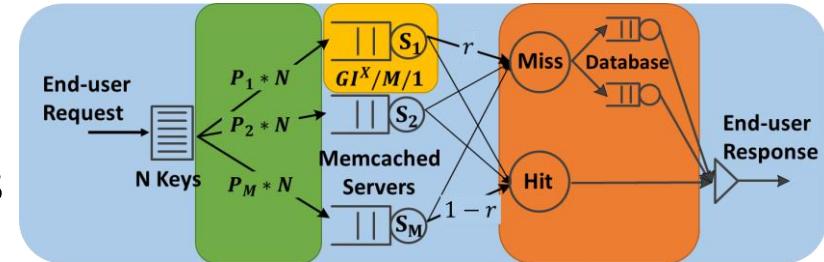
Load-balancing mechanisms take effect before the largest utilization extends the specific value.

Minimizing the number of keys instead of reducing the cache miss ratio.



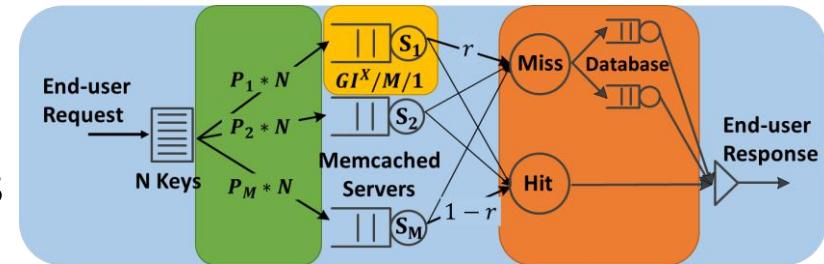
Conclusion

- Specific model for Memcached
 - Unbalanced load distribution
 - Burst and concurrent key arrivals
 - Database processing
- Latency estimation, validation & Analysis
 - Network latency is almost constant
 - Latency at Memcached servers is determined by server utilization
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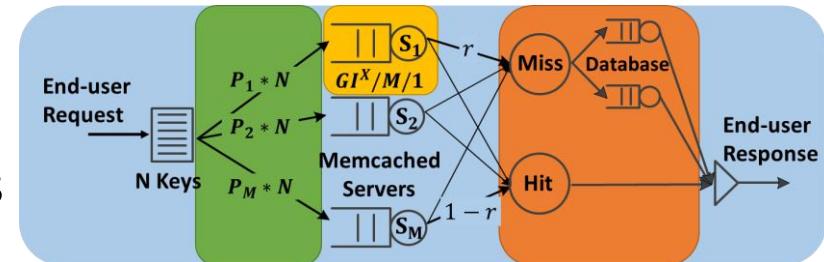
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Thanks!